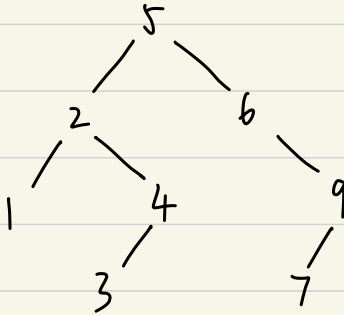


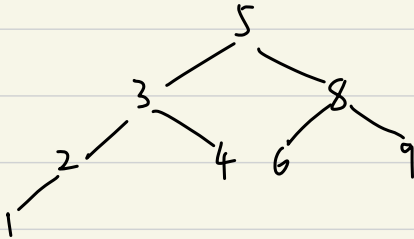
## Assignment 2

### Problem 1

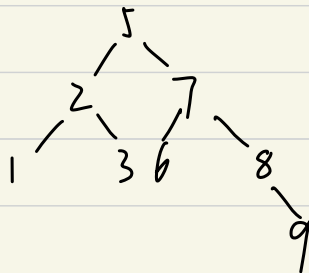
1.  $S = \{2, 6, 4, 9, 7, 3, 1\}$



2.  $S = \{8, 3, 9, 2, 6, 1, 4\}$



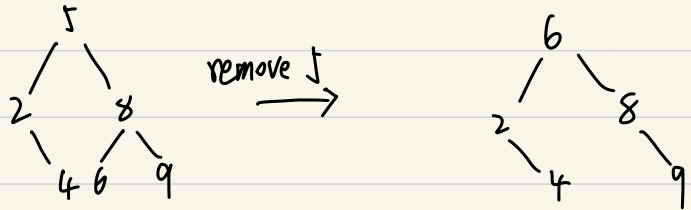
3.  $S = \{7, 2, 1, 8, 3, 6, 9\}$



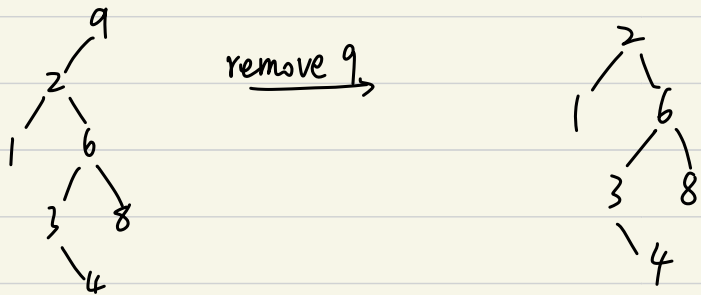
## Problem 2

A is NOT a BST.

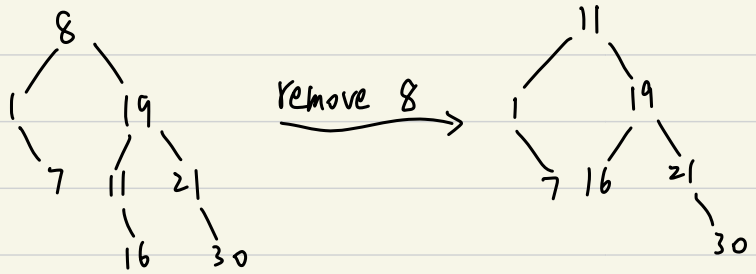
B.



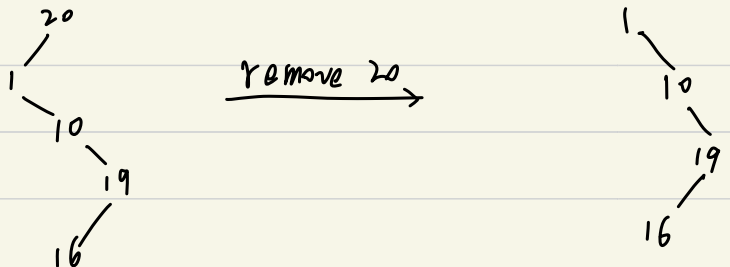
C.



D.



E.



### Problem 3

2 is true.

### Problem 4

1. We can record tree structure first. Then we traverse all the elements in the tree and choose a sorting algorithm applying on them. In the end, put the sorted element back to the tree with in-order traversal.

Time complexity:  $O(n \log n)$

best approach:  $O(n)$

pseudo code:

```
def inorder(node):  
    inorder(node.leftchild)  
    array.append(node)  
    inorder(node.rightchild)
```

```
inorder(root)
```

```
if isincreasing(array) == True: return 0
```

```
sorted_array = merge sort (array) // use merge sort algorithm
```

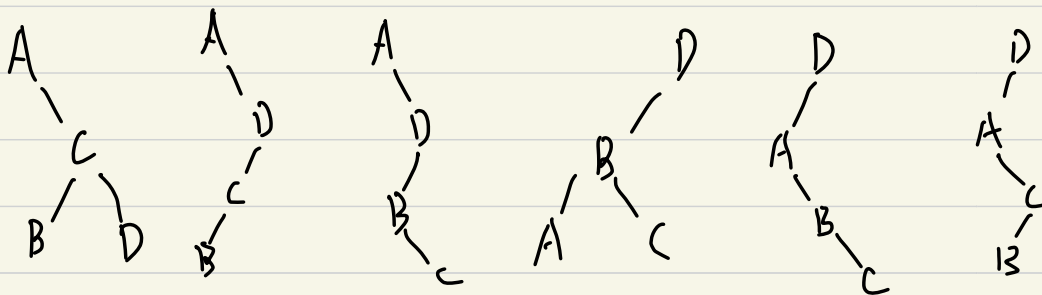
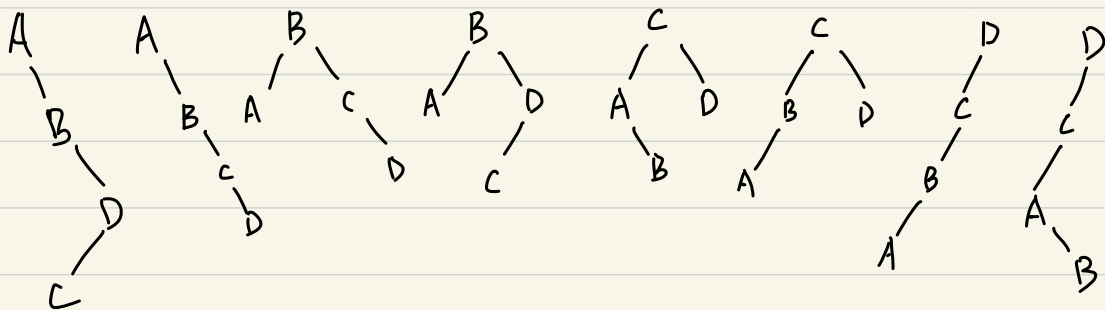
```

def construct_tree(node, i)
    construct_tree(node.leftchild)
    node = array[i]
    i = i + 1
    construct_tree(node.rightchild)
construct_tree(root, 0)

```

Problem 5

Suppose 4 keys are :  $A < B < C < D$

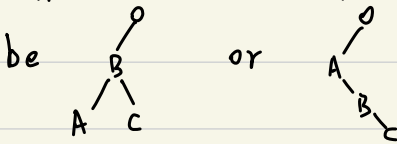


14 BSTs.

## Problem 6

1. True. In-order traversal prints out left child node before root node and then right child node for each layer. For a BST, left child node  $<$  root node  $<$  right child node. The sequence is always in increasing order.

2. False. It's NOT sufficient to build a BST when only be given one traversal result. We can't tell the left child from the right child. It could be



for example.

## Problem 7

1. Downheap for each node

{9, 2, 8, 5, 6, 1, 3}

0 1 2 3 4 5 6

pseudocode:

```
def minheap(heap, i):
```

```
    if heap[left] < heap[i] and left < length(heap):
```

```
        s = left    else: s = i
```

```
    if heap[right] < heap[s] and right < length(heap):
```

```
        s = right
```

```
    if s != i:
```

```
        swap(heap[s], heap[i])
```

```
        minheap(heap, s)
```

```
heap = [9, 2, 8, 5, 6, 1, 3]
```

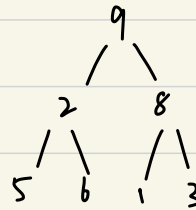
```
n = length(heap) // 2 - 1
```

```
for i in range(n, -1, -1):
```

```
    minheap(heap, i)
```

Time Complexity:  $O(n \log n)$

Best approach:  $O(n)$



test:

9 2 (8) 5 6 (1) 3

9 (2) 1 (5) (6) 8 3

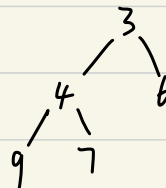
(9) (2) (1) 5 6 8 3

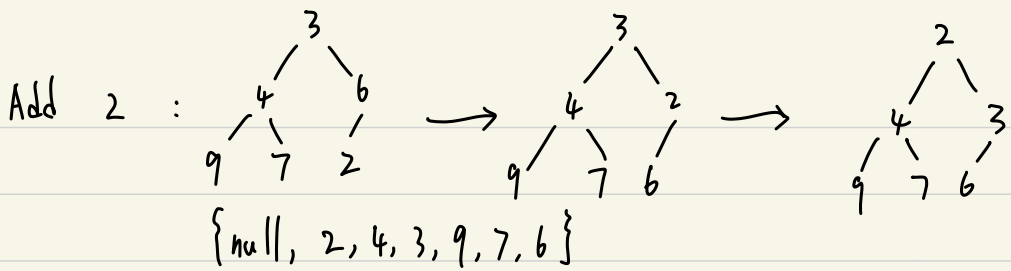
1 2 (9) 5 6 (8) 3

1 2 3 5 6 8 9

## Problem 8

{null, 3, 4, 6, 9, 7}





## Problem 9

1.  $\{9, 12, 14, 3, 4, 21, 18\}$

2.  $\{9, 14, 4, 18, 12, 3, 21\}$

3.  $\{12, 14, 3, 9, 4, 18, 21\}$

4.  $\{12, 3, 14, 18, 4, 9, 21\}$

5.  $\{12, 9, 18, 3, 14, 21, 4\}$

mod 9

	0	1	2	3	4	5	6	7	8	
1.	9	18		12	3	14	4	21		✓
2.	9	18		12	4	14	3	21		✗
3.	9	18		12	3	14	4	21		✓
4.	18	9		12	3	14	4	21		✗
5.	9	18		12	3	14	21	4		✗
Given	9	18		12	3	14	4	21		

Answer : 1, 3

Problem 10

$\{12, 44, 13, 88, 23, 94, 11, 39, 20, 16\}$

29, 93, 31, 181, 51 | 93 27 83 45 37

$$h(k) = 2k + 5 \mod 11$$

0	1	2	3	4	5	6	7	8	9	10
	20			16	11	94	23		13	
					↑	↑	↑			
					88	39	12			
					↑					
					44					

Problem 11

$\{12, 44, 13, 88, 23, 94, 11, 39, 20\}$

$$h(k, i) = (h_0(k) + i + i^2) \mod 11$$

where  $h_0(k) = k \mod 11$

0	1	2	3	4	5	6	7	8	9	10
44	12	13	23	20		88	39	94	11	