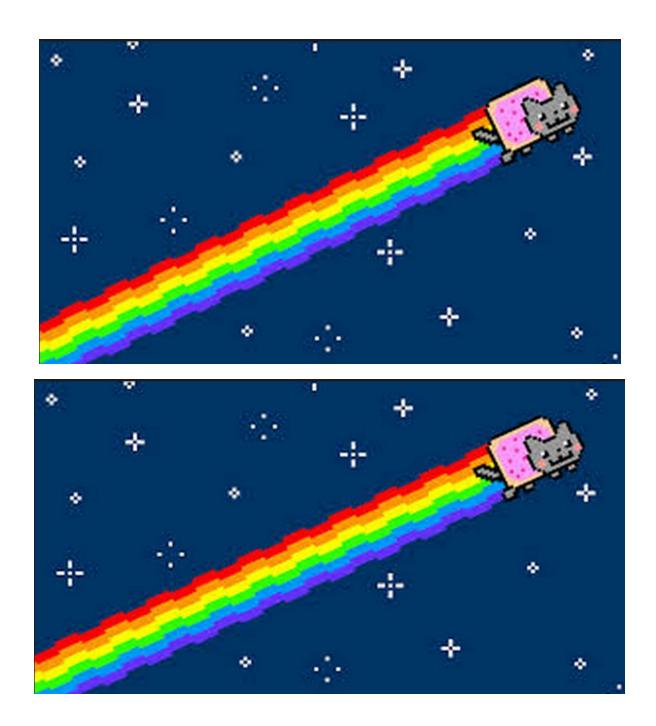
Xi Chen 4626350 DB HW4



В1

- (a) false
- (b) true
- (c) false

B2

(a) R1(X), W2(X), W2(Y), W3(Y), W1(Y), C1, C2, C3

T1 R(X) W(Y) C1

T2 W(X) W(Y) C2

T3 W(Y) C3

Steps

T3 X lock on y

T3 commits

T2 being blocked

T1 S lock on X

T1 X lock on Y

T1 commits

T2 X lock on X, X lock on Y

No Deadlock

(b) R1(X), W2(Y), W2(X), W3(Y), W1(Y), C1, C2, C3

T1 R(X) W(Y) C1

T2 W(Y) W(X) C2

T3 W(Y) C3

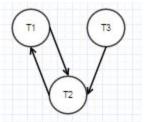
Step:

T1 S lock on X

T2 X lock on Y

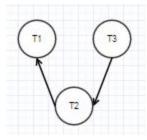
T3 block and wait Y

deadlock

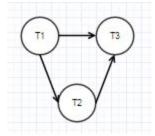


В3

(a) T3T2T1

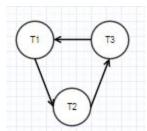


(b) T1T2T3

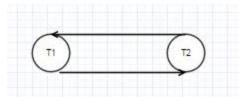


(c)

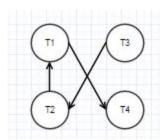
no



(d) no



(e) T4T1T2T3



B4

T1:W(A),W(B),C T2:R(A).R(B),C

that is a conflict serializable schedule. While it is inconsistent with 2PL

C1

a false

b true

c true

d false

C2

(a)

LSN LOG

00 Update T1 writes P2

10 Update T1 writes P1

20 Update T2 writes P5

30 Update T3 writes P3

40 T3 commit

50 Update T2 writes P5

60 Update T2 writes P3

70 T2 abort

(b)

Avort at LSN 70, and previous LSN is 60. SO first undone LSN 6, and then read prevLSN 50, then undone LSN50. Then read previous LSN value 20, undone LSN20, then there is no more prove LSN on LSN20. So Undo is done. T2 ends.

(c)		
LSN	LOG	prvLSN
00	Update T1 writes P2	
10	Update T1 writes P1	00
20	Update T2 writes P5	
30	Update T3 writes P3	
40	T3 commit	30
50	Update T2 writes P5	20
60	Update T2 writes P3	50
70	T2 abort	60
80	CLR undo T2 LSN 60	50
90	CLR undo T2 LSN 50	20
100	CLR undo T2 LSN 20, T2 end	

C3

(a)

transaction table

transaction ID Status LastLSN T1 IN progress 40 T2 Aborted 110

Dirty Page Table

pageIDRecLSN

P1 40 P2 30 P3 60 P5 50

(b)

as same as part a

(c)

LSN LOG

120 UNDO T2 LSN 110
130 UNDO T2 LSN 90
140 UNDO T2 LSN 80

150 UNDO T2 LSN 50 T2 ends

160 UNDO T1 LSN 40

170 UNDO T1 LSN 30 T1 ends