

Introduction to Stacks

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June 8, 2016

We shall introduce stacks from the perspective of moduli spaces following a mixture of [2], [3], and [4]. (For more background, see [1].) The claim is that stacks are a more natural class of objects to consider in moduli problems. Here are two examples.

1. Over \mathbb{C} , two elliptic curves E and E' are isomorphic iff $j(E) = j(E')$. But putting $\mathbb{A}_{\mathbb{C}}^1 := \text{Spec } \mathbb{C}[j]$ as the moduli space is “too coarse” in that the moduli space has forgotten about many non-trivial automorphisms, e.g. the $[-1]$ automorphism possessed by every elliptic curve, or the extra automorphisms of $j = 0$ and $j = 1728$.
2. If X is a variety and G a group acting on X , the quotient X/G may not exist in the category of varieties or even schemes. But it does as a stack. Quotients in general are more natural for stacks.

1 Construction

Definition 1.1. Let S be a scheme and Sch/S be the category of S -schemes. A **category over S** is a category \mathcal{F} equipped with a **projection functor** $p: \mathcal{F} \rightarrow \text{Sch}/S$; we refer to \mathcal{F} as an **S -category**. We view Sch/S as our **base category**.

Definition 1.2. A category \mathcal{F} over S is **fibered in groupoids** if:

1. (arrow lifting) every arrow $\phi: U \rightarrow V$ in Sch/S **lifts** to an arrow $f: x \rightarrow y$ in \mathcal{F} , i.e. $p(f) = \phi$, for every choice of $y \in p^{-1}(V)$;
2. (diagram lifting) for all diagrams $x \xrightarrow{f} y \xleftarrow{g} z$ in \mathcal{F} with image $U \xrightarrow{\phi} W \xleftarrow{\psi} V$ under p , for every $\chi: U \rightarrow V$ completing the image diagram there exists a unique **lift** $h: x \rightarrow y$ of χ completing the original diagram.

Remark 1.3. Note that condition 1 does not seem to specify any conditions on x . However, condition 2 implies the $f: x \rightarrow y$ in condition 1 is unique up to unique isomorphism, since if there were another such arrow $f': x' \rightarrow y$ then we can form the diagram $x \xrightarrow{f} y \xleftarrow{f'} x'$. The image diagram $U \xrightarrow{p(f)} V \xleftarrow{p(f')} U$ is completed by the identity id_U , so by condition 2 there exists a unique lift $x \rightarrow x'$. By swapping x and x' we get its inverse.

Remark 1.4. Condition 2 implies that an arrow f in \mathcal{F} is an isomorphism iff $p(f)$ is an isomorphism. The non-trivial direction: let $\phi := p(f): U \rightarrow V$ be an isomorphism and $\psi: V \rightarrow U$ its left inverse, so that the diagram $V \xrightarrow{\text{id}} V \xleftarrow{\phi} U$ is completed by ψ , which lifts to a left inverse $g: y \rightarrow x$ of f in \mathcal{F} . Repeating the argument with the right inverse shows that f is an isomorphism.

Definition 1.5. Let U be a scheme over S and \mathcal{F} be an S -category. Then $\mathcal{F}(U)$, the **fiber of \mathcal{F} over U** , is the category with objects $p^{-1}(U)$ and arrows $f: x \rightarrow y$ with $x, y \in p^{-1}(U)$ and $p(f) = \text{id}_U$. From the preceding remark, every arrow in $\mathcal{F}(U)$ is an isomorphism since id_U is, i.e. $\mathcal{F}(U)$ is a **groupoid**.

Definition 1.6. Recall that a **Grothendieck topology** on a category \mathcal{C} with fiber products is an assignment to each object U of \mathcal{C} a collection of sets of arrows $\{U_i \rightarrow U\}$, called **coverings** of U , such that

1. (isomorphisms) if $V \rightarrow U$ is an isomorphism, then $\{V \rightarrow U\}$ is a covering,
2. (stability under base change) if $\{U_i \rightarrow U\}_i$ is a covering and $V \rightarrow U$ is any arrow, then $\{U_i \times_U V \rightarrow V\}_i$ is a covering, and
3. (local character) if $\{U_i \rightarrow U\}_i$ is a covering where for each i we have a covering $\{V_{ij} \rightarrow U_i\}_j$, then the composite $\{V_{ij} \rightarrow U_i \rightarrow U\}$ is a covering.

A category with a Grothendieck topology is a **site**.

Example 1.7. Let X be a topological space and X_{cl} be the category of open subsets of X with inclusion maps. The **site of the topological space X** defines a covering of an open subset $U \subset X$ to be an open covering of U , i.e. a set $\{U_i \rightarrow U\}_i$ with U_i open and $U \subset \bigcup_i U_i$. Note that in this case, $U_1 \times_U U_2 = U_1 \cap U_2$.

Example 1.8. We say a set of maps $\{X_i \rightarrow X\}$ is **jointly surjective** if the set-theoretic union of their images is X . The **(global) étale topology** on Sch/S defines a covering to be a jointly surjective set of étale morphisms. (Being **étale** is the algebraic analogue of being a local isomorphism, e.g. a morphism of smooth varieties is étale at a point iff the differential there is an isomorphism.)

Definition 1.9. Assume that the base category Sch/S has a Grothendieck topology: the étale topology, unless otherwise specified. Write $U_{ij} := U_i \times_U U_j$. A **stack** over a scheme S is a category \mathcal{F} fibered in groupoids over S such that the assignment

$$\text{Sch}/S \rightarrow \text{Set}, \quad U \mapsto \mathcal{F}(U) = p^{-1}(U)$$

is a **sheaf of groupoids** (in the topology of Sch/S). In other words:

1. for all $U \in \text{Sch}/S$ and all $x, y \in \mathcal{F}(U)$, the functor

$$\begin{aligned} \text{Isom}_U(x, y): \text{Sch}/U &\rightarrow \text{Set} \\ V &\mapsto \{\alpha: x|_V \rightarrow y|_V \text{ is an isomorphism in } \mathcal{F}(V)\} \end{aligned}$$

is a sheaf (in the topology of Sch/S), i.e. it satisfies

- **(sheaf axiom)** for all $x, y \in \mathcal{F}(U)$, all open covers $\{U_i \rightarrow U\}_i$ of U , and all isomorphisms $\alpha_i: x|_{U_i} \rightarrow y|_{U_i}$ such that $\alpha_i|_{U_{ij}} = \alpha_j|_{U_{ij}}$, there exists a unique isomorphism $\alpha: x \rightarrow y$ such that $\alpha|_{U_i} = \alpha_i$;
- 2. the **descent datum** is **effective**, i.e. for all open covers $\{U_i \rightarrow U\}_i$, all $x_i \in \mathcal{F}(U_i)$, and all $\alpha_{ij}: x_i|_{U_{ij}} \rightarrow x_j|_{U_{ij}}$ such that $\alpha_{ik} = \alpha_{jk} \circ \alpha_{ij}$ over U_{ijk} , there exists an $x \in \mathcal{F}(U)$ and $\alpha_i: x|_{U_i} \rightarrow x_i$ in $\mathcal{F}(U_i)$ such that $\alpha_{ij} = \alpha_j|_{U_{ij}} \circ (\alpha_i|_{U_{ij}})^{-1}$.

If \mathcal{F} satisfies only condition 1, it is a **prestack**. (There are quite a few undefined terms here; we shall define them later.)

Example 1.10 (Vector bundles). Let \mathbf{Bund}_r/S be the category of line bundles of rank r on an S -scheme U . The morphisms are pullbacks. The **forgetful functor** mapping a line bundle $L \rightarrow U$ to its underlying S -scheme U puts an S -category structure on \mathbf{Bund}_r/S . We check that \mathbf{Bund}_r/S is fibered in groupoids:

1. (arrow lifting) every $\phi: U \rightarrow V$ induces a pullback ϕ^* , giving the map $\phi^*L \rightarrow L$ for any line bundle L over V ;
2. (diagram lifting) given line bundles L' and L'' over U' and U'' respectively that map into a line bundle $L \rightarrow U$, a commuting morphism $\chi: U' \rightarrow U''$ induces a pullback χ^* which is unique up to unique isomorphism, completing the diagram of line bundles.

Now we verify that the category \mathbf{Bund}_r/S is a stack over S :

1. condition 1 says that local (wrt an open cover) isomorphisms between two bundles agreeing on overlaps can be glued together to get a (global) bundle isomorphism;
2. condition 2 says that local (wrt an open cover) line bundles can be glued together by specifying transition maps α_{ij} that satisfy the cocycle condition.

Example 1.11 (Schemes and functors). A scheme is a stack via its functor of points. Let X be an S -scheme. The morphism

$$p: \mathbf{Sch}/X \rightarrow \mathbf{Sch}/S, \quad (U \rightarrow X) \mapsto (U \rightarrow X \rightarrow S)$$

gives \mathbf{Sch}/X the structure of an S -category. The structure of the category $\mathcal{F}(U)$ for $U \in \mathbf{Sch}/S$ is particularly simple:

1. the objects are $\mathcal{F}(U) = \mathbf{Hom}_S(U, X)$ since an object in $\mathcal{F}(U)$ is an X -scheme Z such that $Z \rightarrow X \rightarrow S = U \rightarrow S$, i.e. $Z = U$;
2. the only morphisms in $\mathcal{F}(U)$ are the identity morphisms, since an X -morphism $U \rightarrow U$ projects under p to an S -morphism $U \rightarrow U$, which is the identity iff the original X -morphism is also the identity, by 1.4.

Hence we can view $\mathcal{F}(U) = \text{Hom}_S(U, X)$ as a set. Now we check the stack conditions:

1. $\text{Isom}_U(x, y)(V)$ has either one element or none, depending on whether the objects $x, y: U \rightarrow X$ satisfy $x|_V = y|_V$, so locally it is either the empty sheaf or the constant sheaf and trivially satisfies the sheaf condition;
2. fact: the **functor of points** $\text{Hom}_S(-, X)$ is a sheaf in the étale topology on S .

Henceforth we use X to represent both the S -scheme and the stack $p: \text{Sch}/X \rightarrow \text{Sch}/S$.

2 Representability

Definition 2.1. A **morphism of stacks** is a functor $F: \mathcal{F} \rightarrow \mathcal{G}$ such that $p_{\mathcal{F}} = p_{\mathcal{G}} \circ F$. The morphisms from a stack \mathcal{F} to a stack \mathcal{G} form a category: the arrows in $\text{Hom}_S(\mathcal{F}, \mathcal{G})$ are natural transformations of functors. Hence we say the category of stacks is a **2-category**.

Definition 2.2. A stack \mathcal{F} is **representable** if there exists an S -scheme X and a stack isomorphism $\mathcal{F} \rightarrow X$. (Hence when we say isomorphism here, we mean in the sense of an equivalence of categories; the composition of the two functors need not be the identity functor.)

Proposition 2.3. Let $F: \mathcal{F} \rightarrow \mathcal{H}$ and $G: \mathcal{G} \rightarrow \mathcal{H}$ be morphisms of stacks over S . The fiber product $\mathcal{F} \times_{\mathcal{H}} \mathcal{G}$ (defined categorically, as a limit) consists of:

1. objects (x, y, α) where $x \in \mathcal{F}$, $y \in \mathcal{G}$, and $\alpha: F(x) \rightarrow G(y)$ is a morphism in a fiber of \mathcal{H} , i.e. $p_{\mathcal{H}}(F(x)) = p_{\mathcal{H}}(G(y))$ is the same element of Sch/S as $p_{\mathcal{H}}(\alpha) = \text{id}$;
2. morphisms $(x, y, \alpha) \rightarrow (x', y', \alpha')$ given by pairs $(\phi: x \rightarrow x', \psi: y \rightarrow y')$ in fibers of \mathcal{F} and \mathcal{G} such that

$$G(\psi) \circ \alpha = \alpha' \circ F(\phi): F(x) \rightarrow G(y').$$

The fiber of $\mathcal{F} \times_{\mathcal{H}} \mathcal{G}$ over U is precisely pairs $(x, y) \in \mathcal{F}(U) \times \mathcal{G}(U)$ such that $F(x) \cong G(y)$ in $\mathcal{H}(U)$.

Definition 2.4. Let $F: \mathcal{F} \rightarrow \mathcal{G}$ be a morphism of stacks. Then F is **representable** if for all S -schemes X and all morphisms of stacks $X \rightarrow \mathcal{G}$, the fiber product $\mathcal{F} \times_{\mathcal{G}} X$ is a scheme.

Example 2.5. A morphism of schemes is representable, because given schemes X, Y, Z , the fiber product $X \times_Z Y$ is still a scheme. More generally, a morphism $\mathcal{F} \rightarrow X$ with X a scheme is representable iff \mathcal{F} is a representable stack.

3 Deligne–Mumford Stacks

Definition 3.1. Let \mathcal{P} be a property of morphisms of schemes which is stable under base change. A representable morphism of stacks $F: \mathcal{F} \rightarrow \mathcal{G}$ has **property** \mathcal{P} if for each scheme X and morphism of stacks $X \rightarrow \mathcal{G}$, the morphism of schemes $\mathcal{F} \times_{\mathcal{G}} X \rightarrow X$ has property \mathcal{P} .

Definition 3.2. Let \mathcal{F} be a stack and $\Delta: \mathcal{F} \rightarrow \mathcal{F} \times_S \mathcal{F}$ be the **diagonal** given by the two identity morphisms $\text{id}_{\mathcal{F}}$. Then \mathcal{F} is **Deligne–Mumford (DM)** (resp. **Artin** or **algebraic**) if:

1. the diagonal Δ is representable, quasi-compact, and separated;
2. there exists an S -scheme U , called an **atlas**, and an étale (resp. smooth) and surjective morphism $U \rightarrow \mathcal{F}$. (Think orbifold cover.)

By 3.5, the atlas morphism $U \rightarrow \mathcal{F}$ is automatically representable, so it makes sense to talk about it being étale or smooth.

Proposition 3.3 ([3, Prop. 7.15]). *The diagonal Δ of a Deligne–Mumford stack is unramified.*

Corollary 3.4 ([3, pg. 666]). *If \mathcal{F} is a Deligne–Mumford stack and $U \in \text{Sch}/S$ is quasi-compact, then any $x \in \mathcal{F}(U)$ has only finitely many automorphisms.*

The idea is that a DM stack is the algebraic analogue of an orbifold. The rest of this section is for interpreting and exploring the conditions of the definition of being a DM stack.

Proposition 3.5. *The diagonal Δ is representable iff every morphism from a scheme to \mathcal{F} is representable.*

Proof. Suppose Δ is representable. Let $F: X \rightarrow \mathcal{F}$ be a morphism; we must show that for every $G: Y \rightarrow \mathcal{F}$, the stack $X \times_{\mathcal{F}} Y$ is a scheme. But

$$\begin{array}{ccc} X \times_{\mathcal{F}} Y & \longrightarrow & X \times_S Y \\ \downarrow & & (F, G) \downarrow \\ \mathcal{F} & \xrightarrow{\Delta} & \mathcal{F} \times_S \mathcal{F} \end{array}$$

is a cartesian diagram, i.e. $X \times_{\mathcal{F}} Y = (X \times_S Y) \times_{\mathcal{F} \times_S \mathcal{F}} \mathcal{F}$, by checking the universal property. (This diagram is sometimes called the **magic diagram** and holds in any category.) Since Δ is representable, $X \times_{\mathcal{F}} Y$ is a scheme, as desired.

Conversely, suppose every morphism from a scheme to \mathcal{F} is representable. Let $H: X \rightarrow \mathcal{F} \times_S \mathcal{F}$ be a morphism; we must show that $\mathcal{F} \times_{\mathcal{F} \times_S \mathcal{F}} X$ is a scheme. By the universal property of fiber products, $H = (F, G) \circ \Delta_X$, so

$$\begin{array}{ccccc} \mathcal{F} \times_{\mathcal{F} \times_S \mathcal{F}} X & \longrightarrow & X \times_{\mathcal{F}} X & \longrightarrow & \mathcal{F} \\ \downarrow & & \downarrow & & \Delta_{\mathcal{F}} \downarrow \\ X & \xrightarrow{\Delta_X} & X \times_S X & \xrightarrow{(F, G)} & \mathcal{F} \times_S \mathcal{F} \end{array}$$

consists of two cartesian squares using the magic diagram. Clearly $X \times_{\mathcal{F}} X$ is a scheme, since the morphism $X \rightarrow \mathcal{F}$ is representable by hypothesis. But Δ_X in the left square is just a morphism of schemes, which is representable, so since $X \times_{\mathcal{F}} X$ is a scheme, so is $\mathcal{F} \times_{\mathcal{F} \times_S \mathcal{F}} X$, as desired. \square

Example 3.6. Let X be a scheme. Then X is a DM stack, since clearly every morphism from a scheme to X is representable, and we can let X be its own atlas, with the identity map $X \rightarrow X$ both étale and surjective.

Example 3.7. The stack Bund_r/S is not a DM stack for $r \geq 1$, because in general the group of automorphisms of a vector bundle is not finite, violating 3.4.

References

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