

EECS 340: Assignment 4

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Problem 1

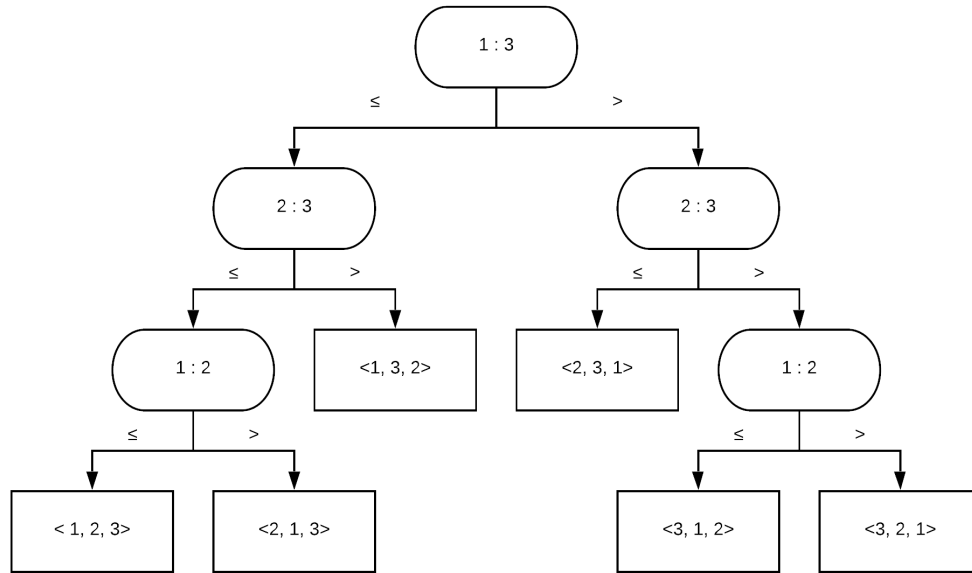


Figure 1: Decision Tree of QUICKSORT

Problem 2

Procedure PREPROCESS(A)

Algorithm 1 PreProcess(A)

```
1: procedure PREPROCESS(A, p, r)
2:   Let  $C[0, 1, \dots, k]$  be an array of 0.
3:   for  $i \leftarrow 0$  to  $n$  do
4:      $C[A[i]] = C[A[i]] + 1$ 
5:   for  $j \leftarrow 1$  to  $k$  do
6:      $C[j] = C[j] + C[j - 1]$ 
7:   return  $C$ 
```

Procedure QUERY(A, a, b)

Algorithm 2 Query(A, a, b)

```
1: procedure QUERY( $A, a, b$ )
2:   if  $a == 0$  then
3:     return  $A[b]$ 
4:   else
5:     return  $A[b] - A[a - 1]$ 
```

Problem 3

Algorithm 3 Sparse-Transpose(R, C, V, m, n, k)

```
1: procedure SPARSE-TRANSPOSE( $A, a, b$ )
2:   Let  $CH[0, 1, \dots, n]$  be an array of 0.
3:   Let  $VH[0, 1, \dots, n]$  be an array of 0.
4:   for  $i \leftarrow 0$  to  $n$  do
5:      $start \leftarrow R[i]$ 
6:      $end \leftarrow R[i + 1] - 1$ 
7:     Let  $CPR[ ]$  be an empty array.
8:     Let  $VPR[ ]$  be an empty array.
9:     for  $j \leftarrow start$  to  $end$  do
10:       $CPR[i].append(C[j])$ 
11:       $VPR[i].append(V[j])$ 
12:       $CH[i].append(CPR)$ 
13:       $VH[i].append(VPR)$ 
14:   Let  $R[ ]$  be an empty array.
15:   Let  $C[ ]$  be an empty array.
16:   Let  $V[ ]$  be an empty array.
17:    $V.append(0)$ 
18:   Let  $len\_max$  to be the maximum length among all elements in  $CH$ .
19:   for  $i \leftarrow 0$  to  $len\_max$  do
20:     POP-SMALLEST( $R, C, V, CH, VH, len\_max$ )
21:   for  $i \leftarrow 1$  to  $n$  do
22:      $R[i] = R[i] + R[i - 1]$ 
23:   return  $R, C, V$ 
```

Algorithm 4 Pop-Smallest($R, C, V, CH, VH, \text{len_max}$)

```
1: procedure POP-SMALLEST( $R, C, V, CH, VH, \text{len\_max}$ )
2:    $\text{min\_value} \leftarrow CH[0][0]$ 
3:    $\text{min\_i} \leftarrow 0$ 
4:    $\text{min\_j} \leftarrow 0$ 
5:   for  $i \leftarrow 0$  to  $\text{len\_max}$  do
6:      $R\_counter \leftarrow 0$ 
7:     for  $j \leftarrow 0$  to  $n$  do
8:       if  $CH[j][i] < \text{min\_value}$  or  $(CH[j][i] == \text{min\_value}$  and  $j < \text{min\_j})$  then
9:          $\text{min\_value} \leftarrow CH[j][i]$ 
10:         $\text{min\_i} \leftarrow i$ 
11:         $\text{min\_j} \leftarrow j$ 
12:        continue
13:       if  $CH[j][i] == \text{min\_value}$  then
14:          $R\_counter \leftarrow R\_counter + 1$ 
15:         if  $j < \text{min\_j}$  then
16:            $\text{min\_value} \leftarrow CH[j][i]$ 
17:            $\text{min\_i} \leftarrow i$ 
18:            $\text{min\_j} \leftarrow j$ 
19:        $CH[\text{min\_j}].\text{remove}(\text{min\_i})$ 
20:        $C.\text{append}(\text{min\_j})$ 
21:        $V.\text{append}(VH[\text{min\_j}][\text{min\_i}])$ 
22:        $VH[\text{min\_j}].\text{remove}(\text{min\_i})$ 
23:        $R.\text{append}(R\_counter)$ 
```
