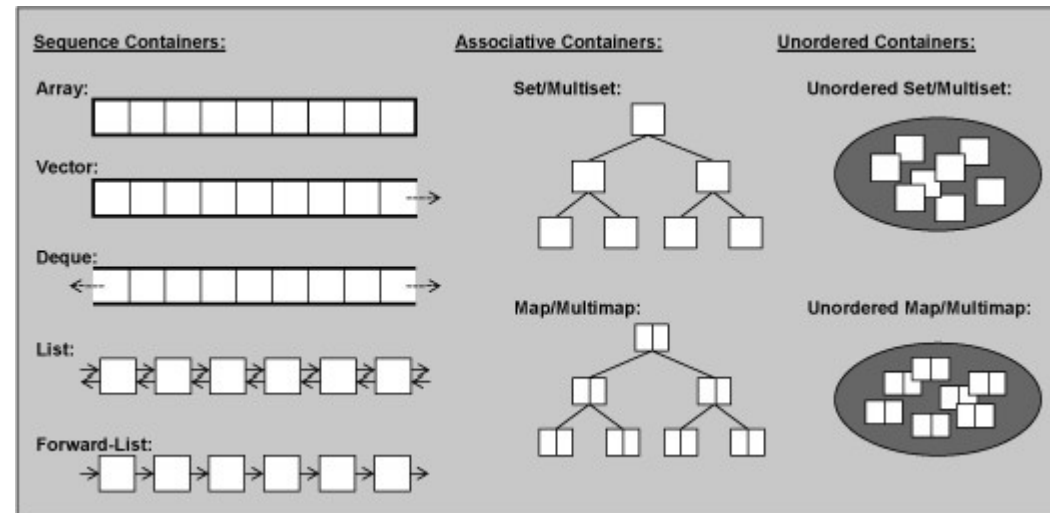


# CPSC 131 – Data Structures

## Singly and Doubly Linked List Abstract Data Types



*Professor T. L. Bettens*  
*Fall 2020*

# Concepts & Interface

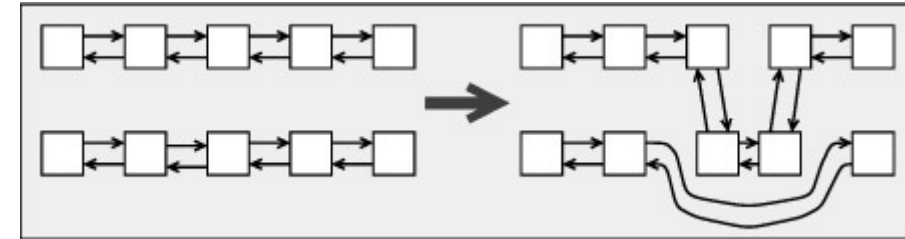
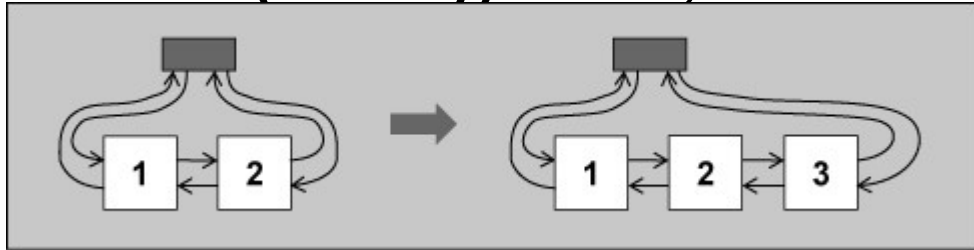
- Jusuttis,  
[The C++ Standard Library](#)
  - [6.2. Containers](#)
  - [7.1. Common Container Abilities and Operations](#)
  - [7.6. Forward \(Singly Linked\) Lists](#)
    - [7.6.4. Examples of Using Forward Lists](#)
  - [7.5. \(Doubly Linked\) Lists](#)
    - [7.5.4. Examples of Using Lists](#)
- [Cplusplus.com](#)
  - [Containers library](#)
  - [std::forward\\_list](#)
  - [std::list](#)
- zyBook
  - tbd



# List Abstract Data Type

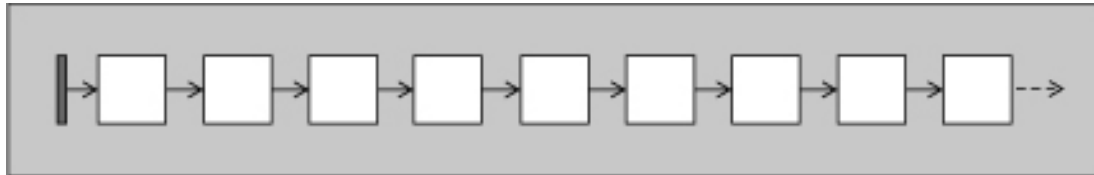
## Common Implementation choices

### Sentinel (Dummy) Nodes, or not



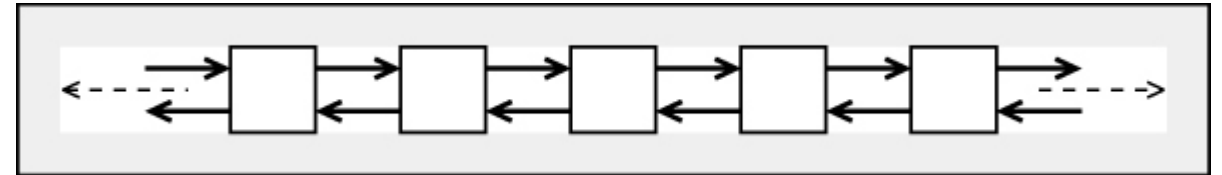
### Singly Linked List

- Can move only forward
- Can insert and delete the “next” element, never the “current” element

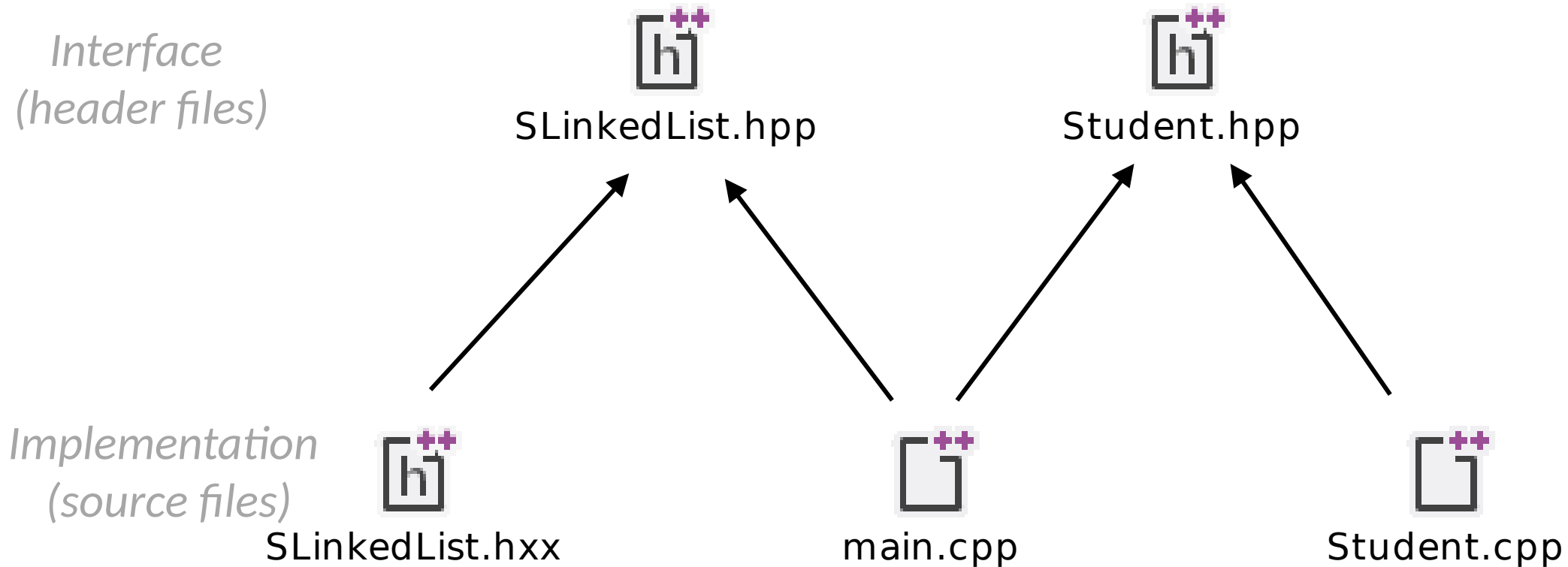


### Doubly Linked List

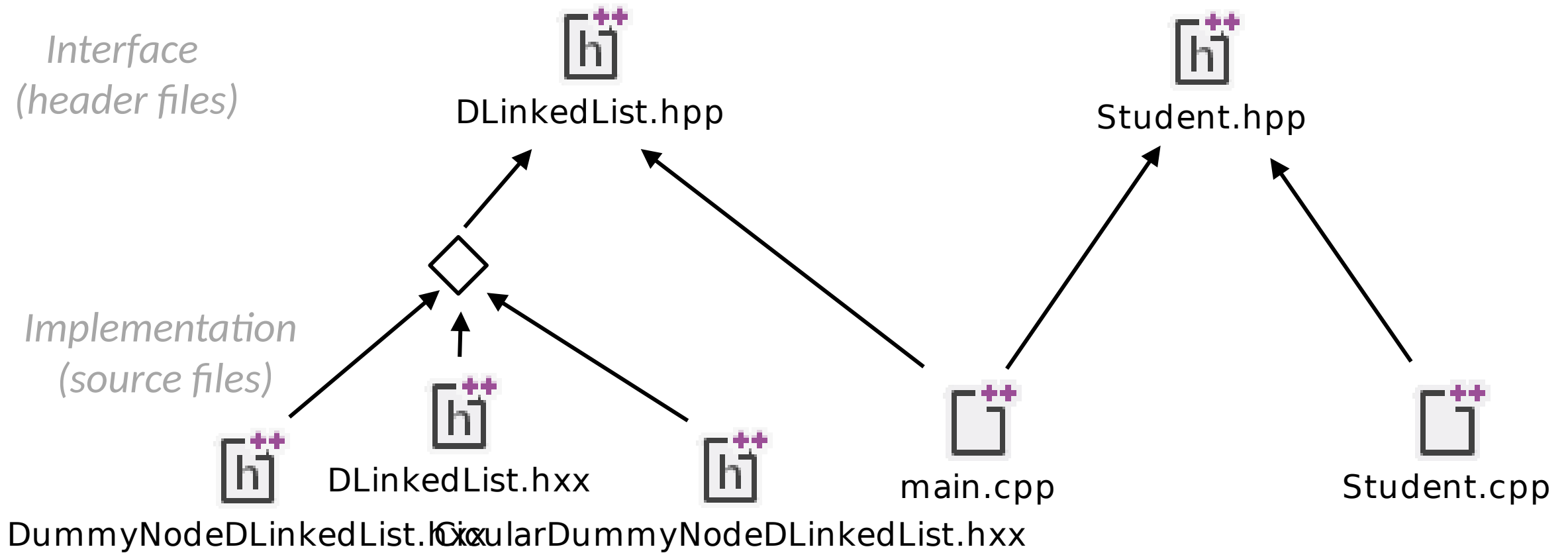
- Can move in both directions
- Can delete “current” element
- Can insert before or after “current”



# Singly Linked List Implementation Example



# Doubly Linked List Implementation Example



# Analysis of the Vector Abstract Data Type

## Complexity Analysis (1)

Function	Analysis – <code>std::forward_list&lt;T&gt;</code> (Singly Linked List)	Analysis – <code>std::list&lt;T&gt;</code> (Doubly Linked List)
<code>at()</code>	Not available	Not available
<code>size()</code>	$O(n)$ <code>std::forward_list&lt;T&gt;</code> calculates size on demand	same
<code>empty()</code>	$O(1)$	same
<code>clear()</code>	$O(n)$ All elements are destroyed and size set to zero	same

# Analysis of the Vector Abstract Data Type

## Complexity Analysis (2)

Function	Analysis – <code>std::forward_list&lt;T&gt;</code> (Singly Linked List)	Analysis – <code>std::list&lt;T&gt;</code> (Doubly Linked List)
<code>push_back()</code>	Not available <code>std::forward_list&lt;T&gt;</code> has no tail pointer	$O(1)$ has tail pointer
<code>erase()</code>	$O(1)$ assumes you have erasure point Can erase only <u>after</u> erasure point	$O(1)$ assumes you have erasure point Can erase only <u>before</u> erasure point
<code>splice</code>	Not available <code>std::forward_list&lt;T&gt;</code> has no tail pointer	$O(1)$

# Analysis of the Vector Abstract Data Type

## Complexity Analysis (3)

Function	Analysis – <code>std::forward_list&lt;T&gt;</code> (Singly Linked List)	Analysis – <code>std::list&lt;T&gt;</code> (Doubly Linked List)
<code>insert()</code>	$O(1)$ assumes you have insertion point Can insert only <u>after</u> insertion point	$O(1)$ assumes you have insertion point Can insert only <u>before</u> insertion point
default construction	$O(1)$ creates an empty container	same
Equality $C_1 == C_2$	$O(n)$	same



# Analysis of the Vector Abstract Data Type

## Complexity Analysis (4)

Function	Analysis – <code>std::forward_list&lt;T&gt;</code> (Singly Linked List)	Analysis – <code>std::list&lt;T&gt;</code> (Doubly Linked List)
<code>push_front</code>	$O(1)$	same
<code>resize</code>	$O(n)$	same
<code>find</code>	$O(n)$ linear search from <code>begin()</code> to <code>end()</code> (i.e. <i>head to tail</i> )	same



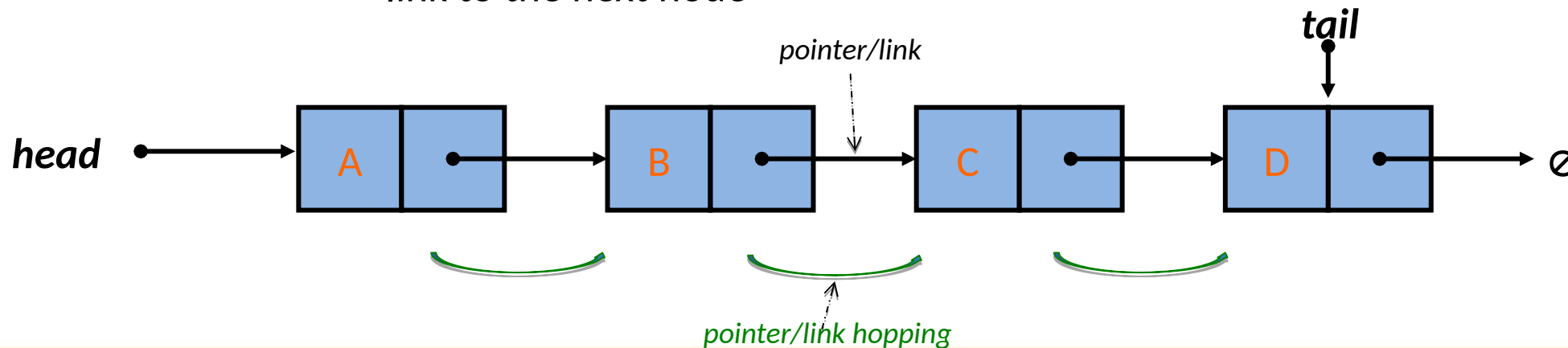
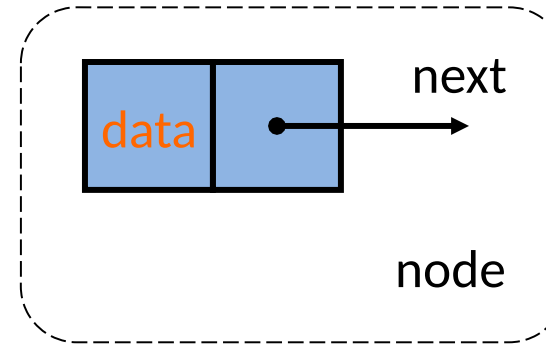
# Analysis of the Vector Abstract Data Type

## Complexity Analysis (5)

Function	Analysis – <code>std::forward_list&lt;T&gt;</code> (Singly Linked List)	Analysis – <code>std::list&lt;T&gt;</code> (Doubly Linked List)
Visit every element e.g. <code>print()</code>	$O(n)$ Visiting every node from <code>begin()</code> to <code>end()</code>	same
Visit in reverse e.g. <code>print_reverse()</code>	not possible	$O(n)$ Visiting every node from <code>rbegin()</code> to <code>rend()</code> Direction doesn't matter

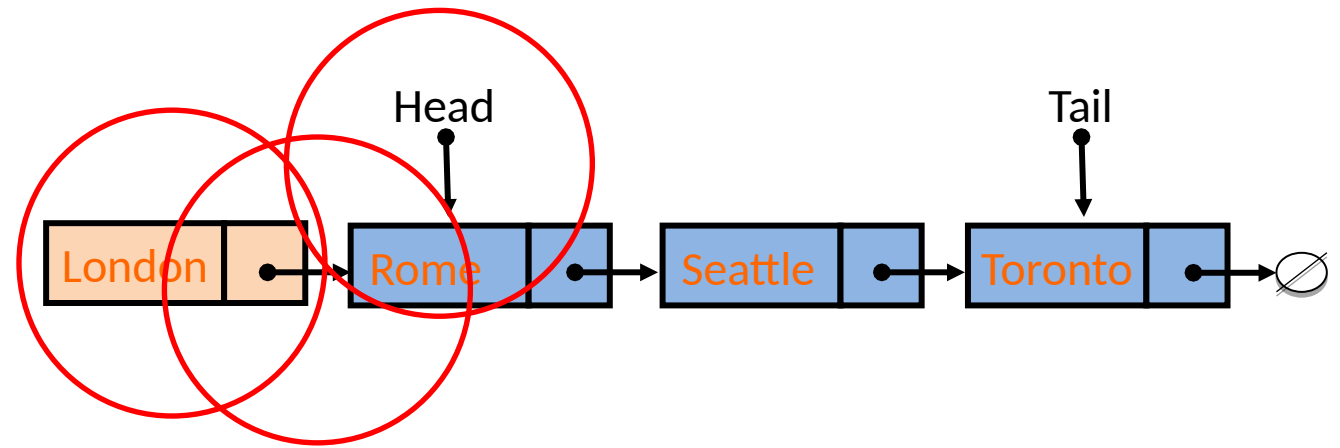
# Singly Linked Lists

- ❖ A singly linked list is a concrete **data structure** consisting of a sequence of nodes
- ❖ Each node stores
  - Data element
  - link to the next node



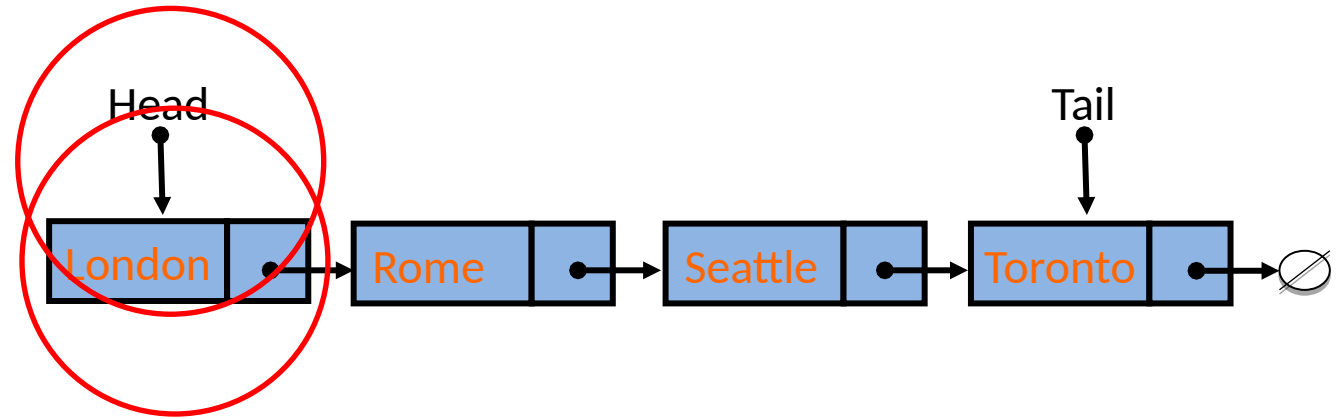
# Inserting at the Head

1. Allocate and populate a new node
2. Have new node point to old head
3. Update head to point to new node



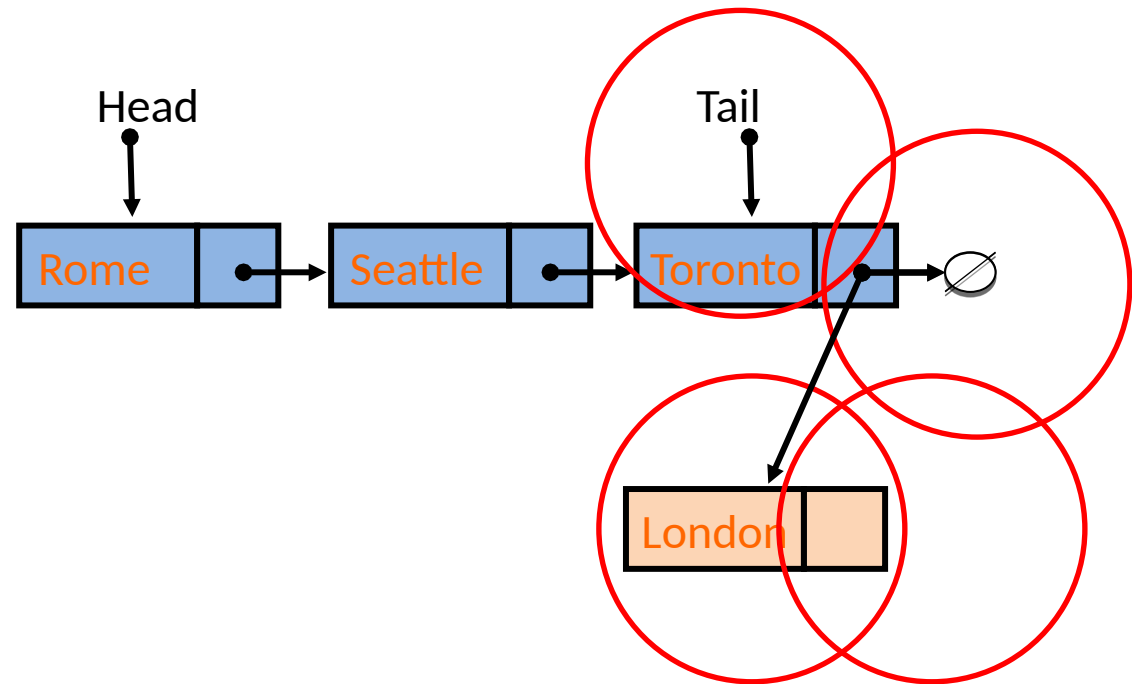
# Deleting at the Head

1. Update head to point to the next node in the list
2. Delete the former first node



# Inserting at the Tail


1. Allocate and populate a new node
2. Have new node point to whatever the tail node pointed to, namely nullptr
3. Point the old Tail node to the new node
4. Update tail to point to new node



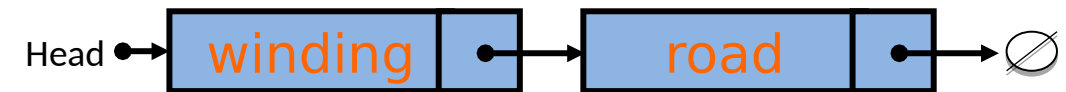
# Draw data structure for this code

prepend == push\_front  
removeFront == pop\_front

```
SLinkedList<string> ds;  
cout << ds.size();
```

Head → 

```
ds.prepend("road");  
ds.prepend("winding");  
cout << ds.front();
```




```
ds.prepend("and");  
ds.removeFront();  
ds.prepend("long");  
cout << ds.front();
```



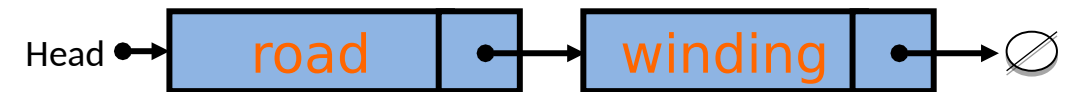
# Draw data structure for this code

append == push\_back  
removeFront == pop\_front

```
SLinkedList<string> ds;  
cout << ds.size();
```

Head → 

```
ds.append("road");  
ds.append("winding");  
cout << ds.front();
```



```
ds.append("and");  
ds.removeFront();  
ds.append("long");  
cout << ds.front();
```





# Nodes

To create a linked list using dynamic storage, we need a class which has two data members:

- one to hold information
- one to point to another object of the *same* class

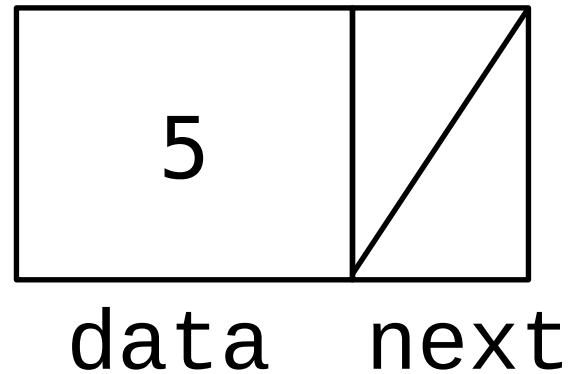
# Singly Linked List Node

```
template <typename Data_t>
struct SLinkedList<Data_t>::Node // singly linked list node
{
    Node() = delete;                // Intentionally no default constructor
    Node( const Data_t & element ) : data( element ) {}

    Data_t data;                    // linked list element value
    Node * next = nullptr;          // next item in the list
};
```

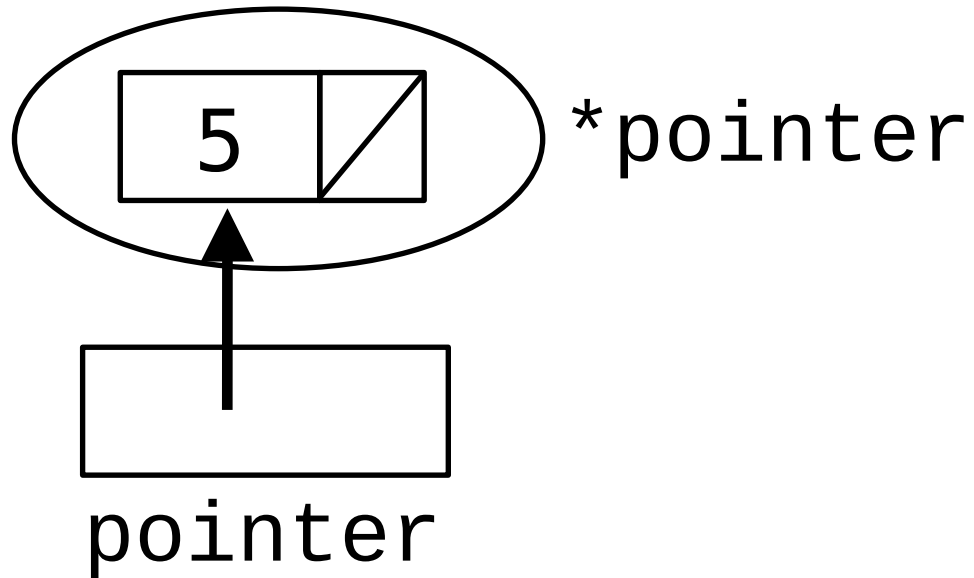
# Picture of a Node

```
SLinkedList<int>::Node node(5);  
node.next = nullptr;
```



# Creating a node as a dynamic variable

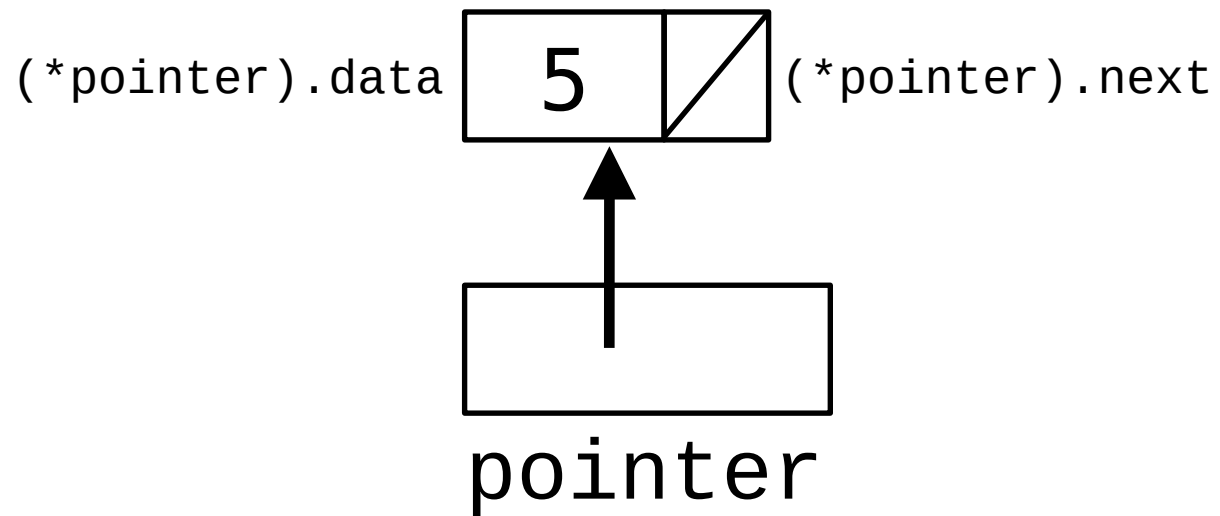
```
SLinkedList<int>::Node * pointer;  
pointer = new SLinkedList<int>::Node(5)
```



# Accessing the fields of the node

`(*pointer).data`

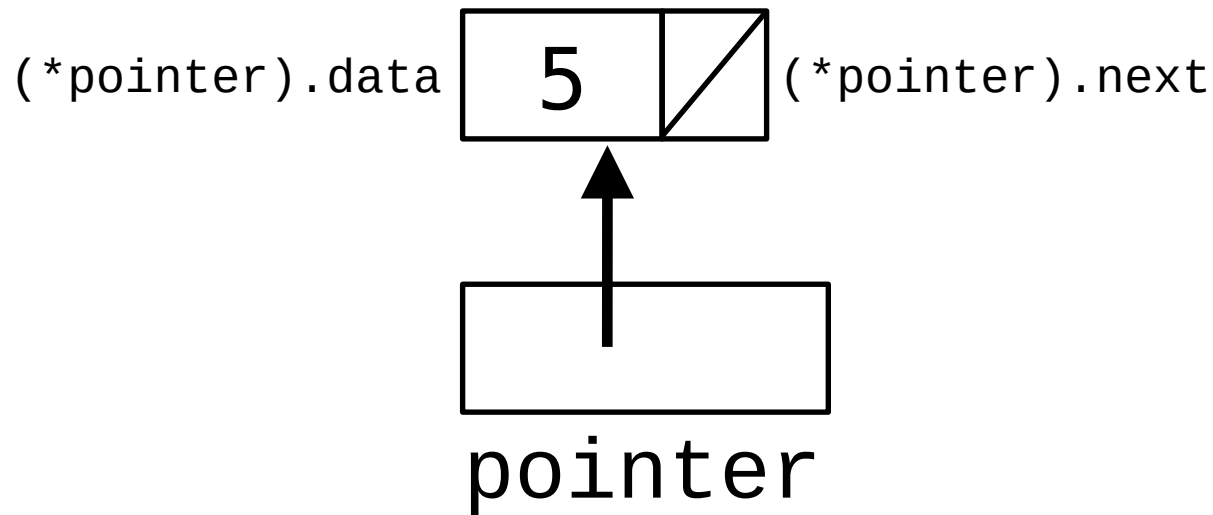
`(*pointer).next`



# Accessing the fields of the node

`(*pointer).data` can also be written as `pointer->data`

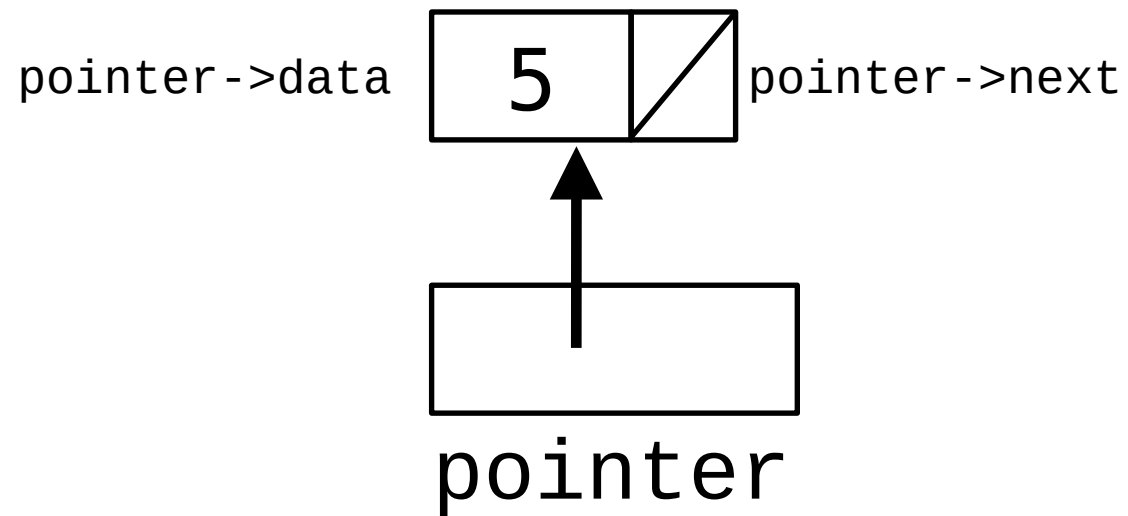
`(*pointer).next` can also be written as `pointer->next`



# Accessing the fields of the node

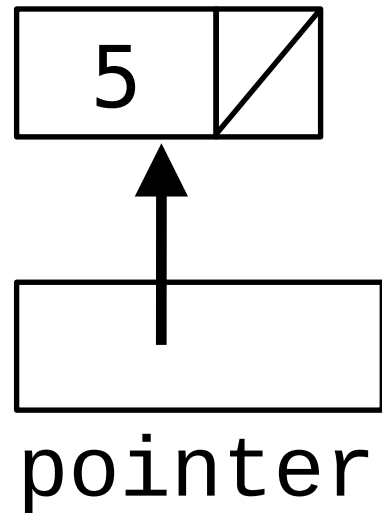
```
pointer->data == 5
```

```
pointer->next == nullptr
```



# Review: delete

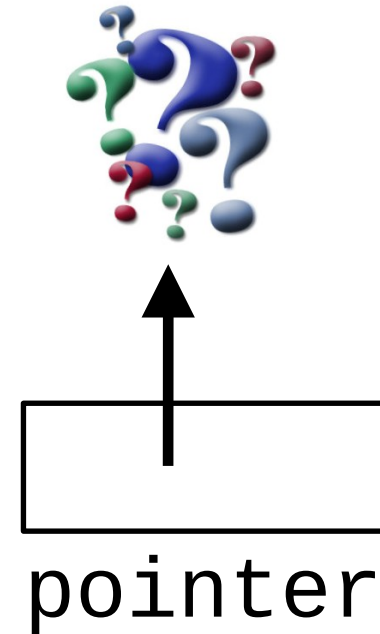
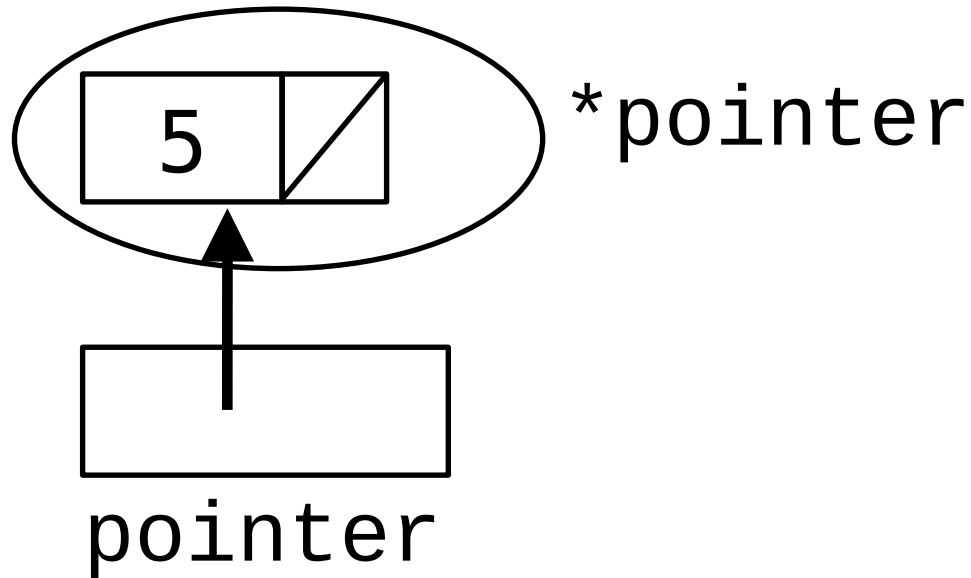
What does  
`delete pointer;`  
do?





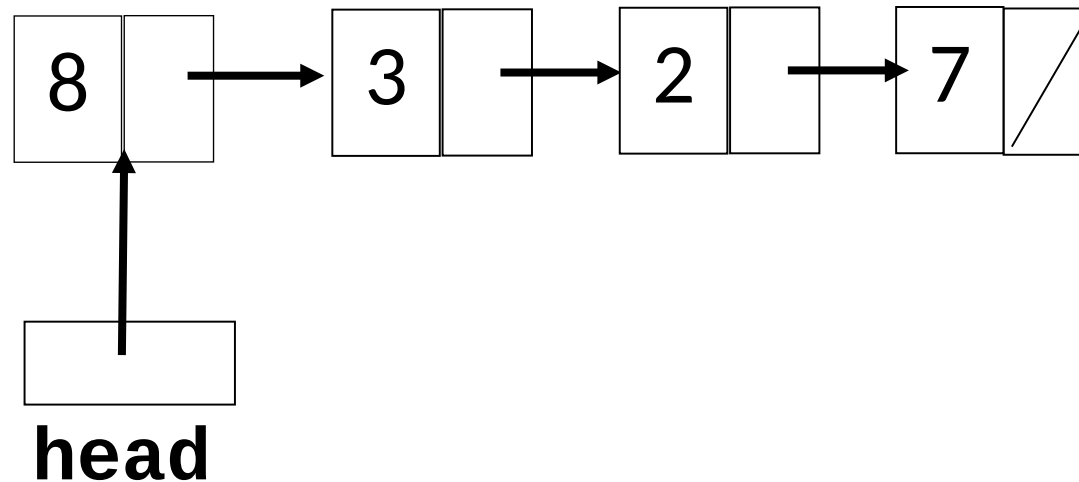
# Answer

It deletes the node, but  
leaves the pointer dangling.



# What about a linked list?

Since the `next` field can point to another node, we can link nodes together like this:



# Creating a linked list

Start out with two NULL pointers to NodeType.

Code for this??

```
SLinkedList<int>::Node *pointer = nullptr;
```

```
SLinkedList<int>::Node *head = nullptr;
```



pointer

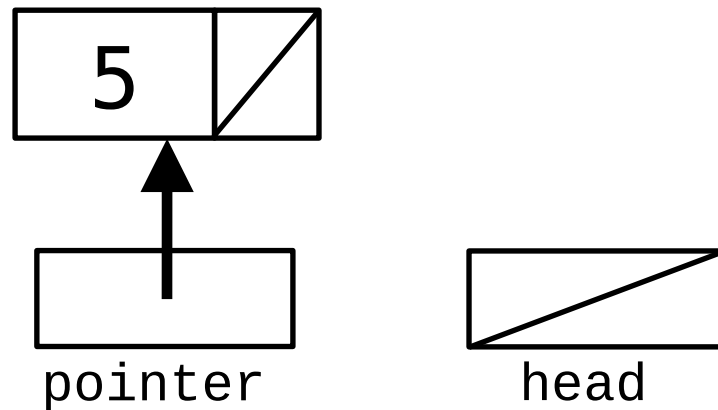


head

# Creating a linked list

Now create a new SinglyLinkedListNode using pointer.  
Code for this??

```
pointer = new SLinkedList<int>::Node(5);
```

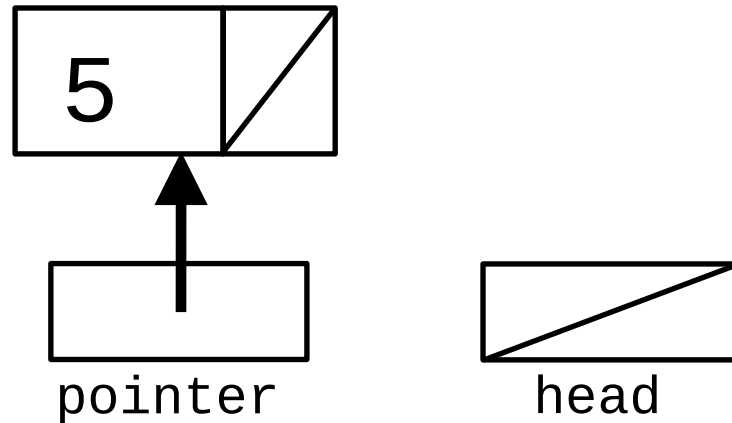


# Creating a linked list

Now what happens if we do

```
pointer->next = head;
```

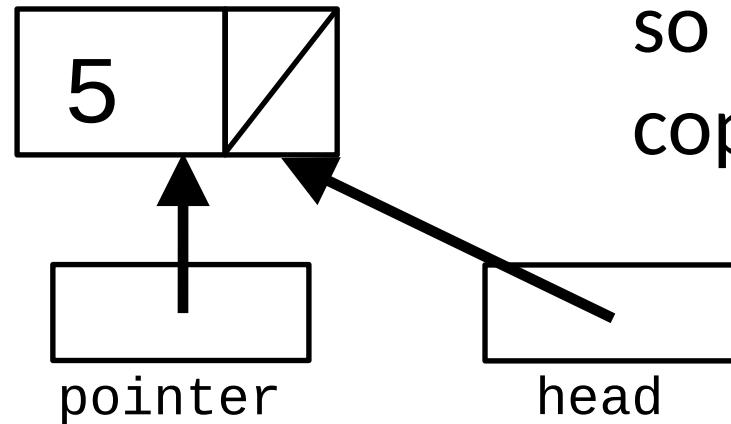
head contains nullptr, which gets copied to pointer->next



# Creating a linked list

Now we want head to point to the new node,  
i.e. head should contain the address of the new node.

What already has that address? **pointer**



so head = pointer will  
copy the address to head .

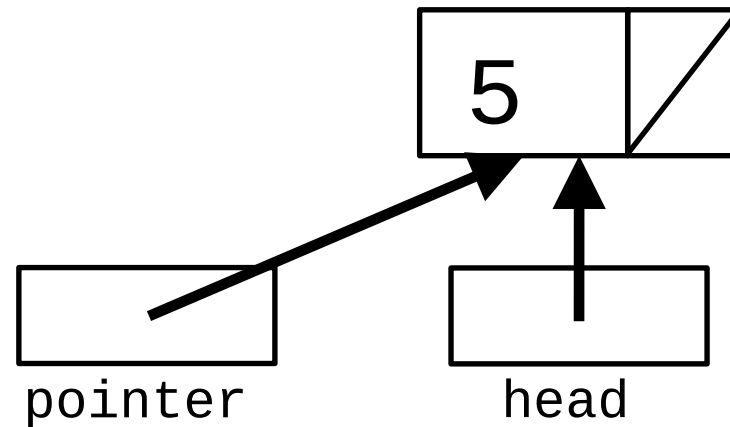
# Creating a linked list

Putting the code together:

```
pointer          = new SLinkedList<int>::Node(5);  
pointer->next    = head;  
head             = pointer;
```

# Going from code to picture

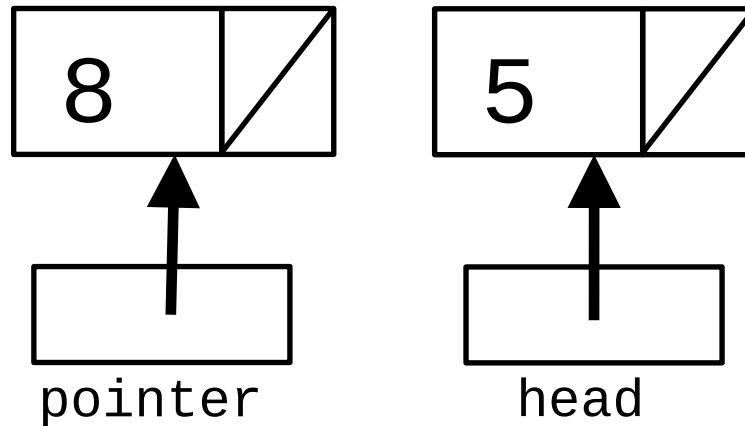
Now try seeing what happens if we repeat the same code with data now set to 8, starting with this picture.





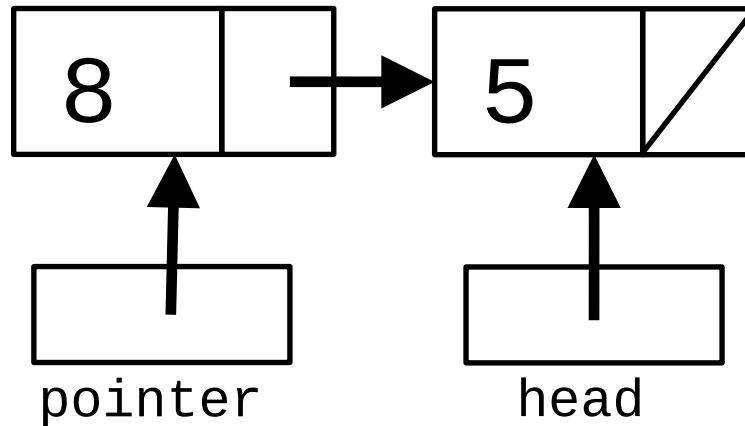
# Going from code to picture

```
pointer      = new SLinkedList<int>::Node(8);  
pointer->next = head;  
head        = pointer;
```



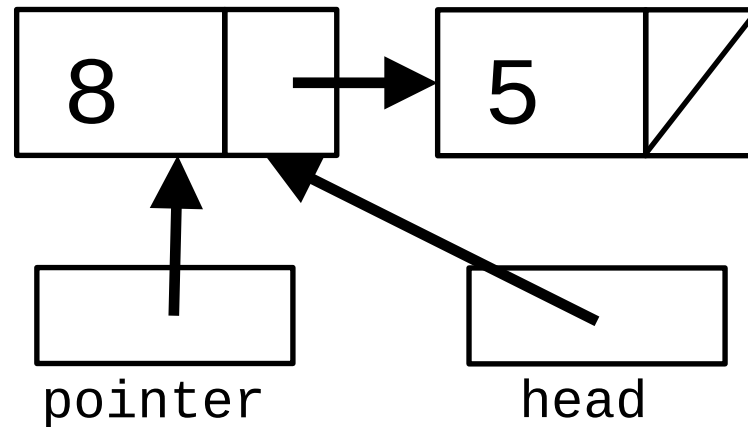
# Going from code to picture

```
pointer      = new SLinkedList<int>::Node(8);  
pointer->next = head;  
head        = pointer;
```



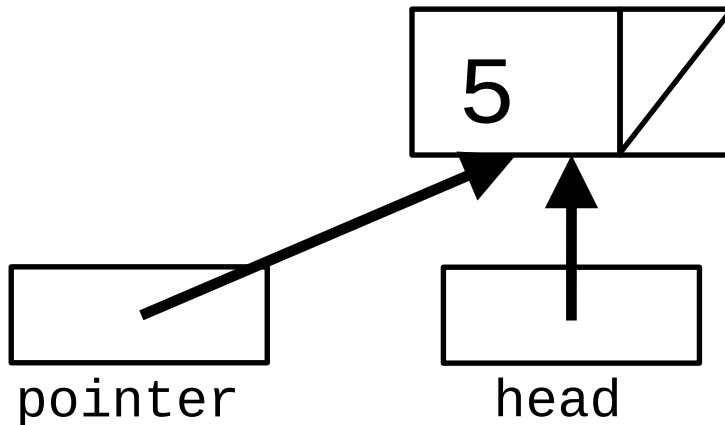
# Going from code to picture

```
pointer      = new SLinkedList<int>::Node(8);  
pointer->next = head;  
head        = pointer;
```



# So the following code will go from

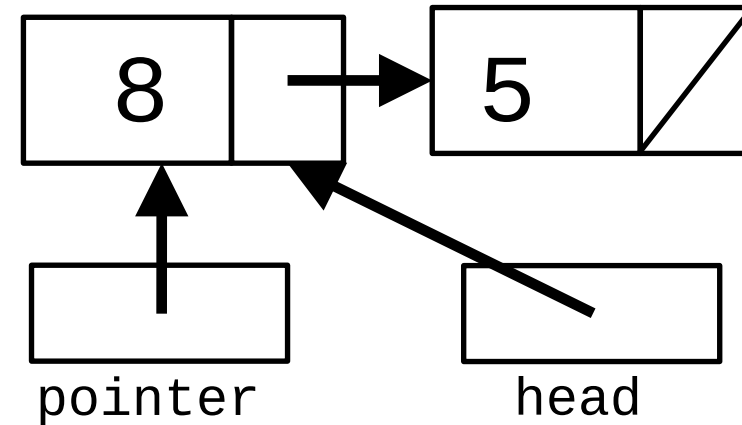
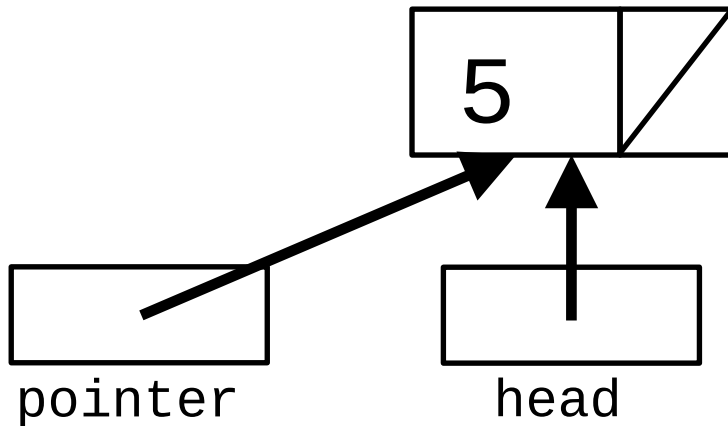
```
pointer      = new SLinkedList<int>::Node(8);  
pointer->next = head;  
head        = pointer;
```



to this, i.e. insert a node at the front

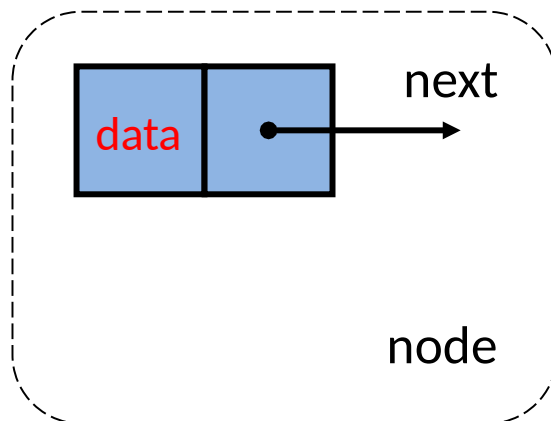
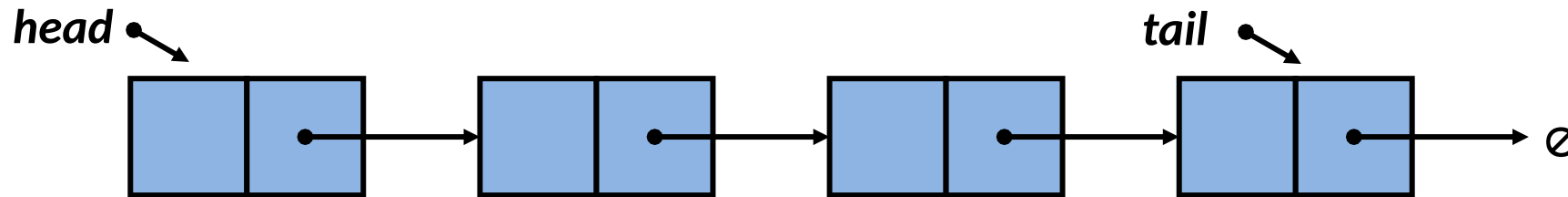
```
void  
prepend( const Data_t & element );
```

```
pointer      = new SLinkedList<int>::Node(8);  
pointer->next = head;  
head         = pointer;
```



# Singly Linked List and Doubly Linked List

What if want to access data in reverse order?

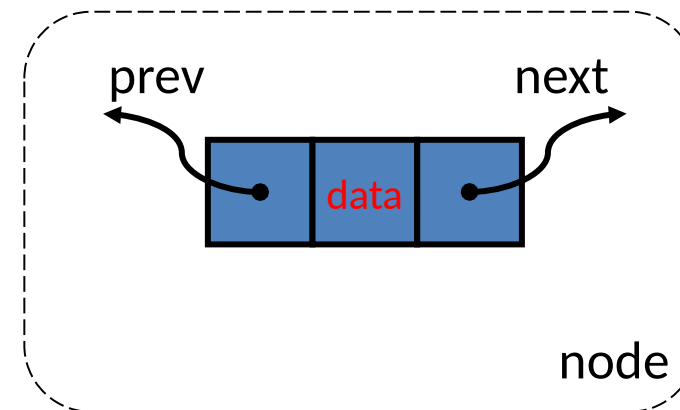


Node:

- element
- next pointer
- **previous pointer**

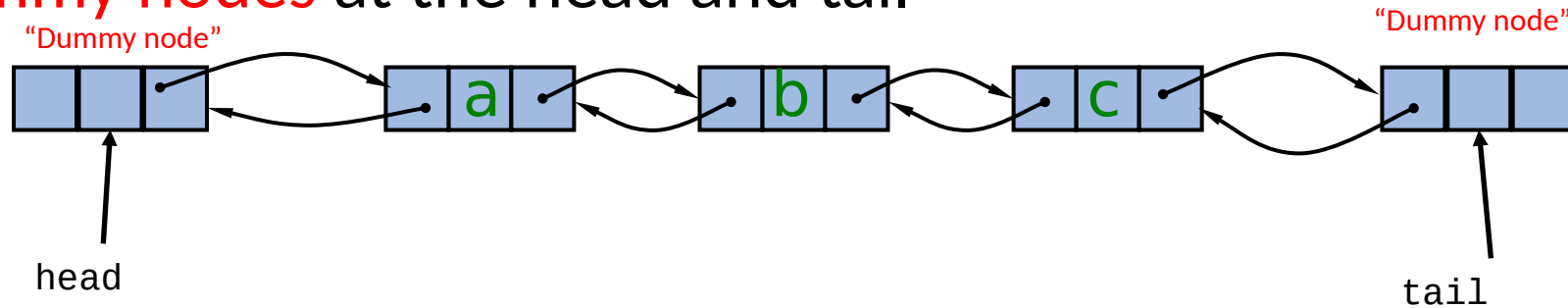
Linked List:

- length
- Head
- Tail

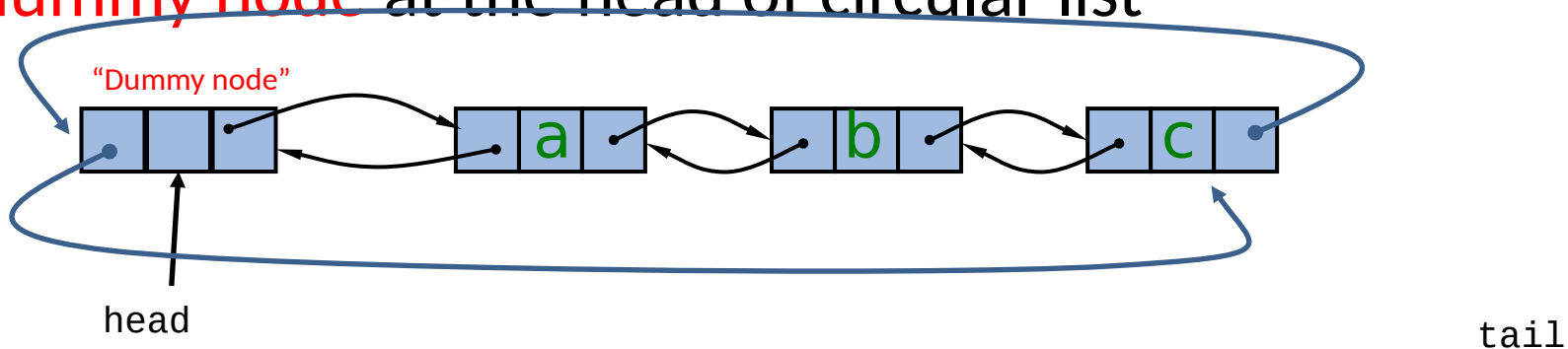


# Doubly linked lists

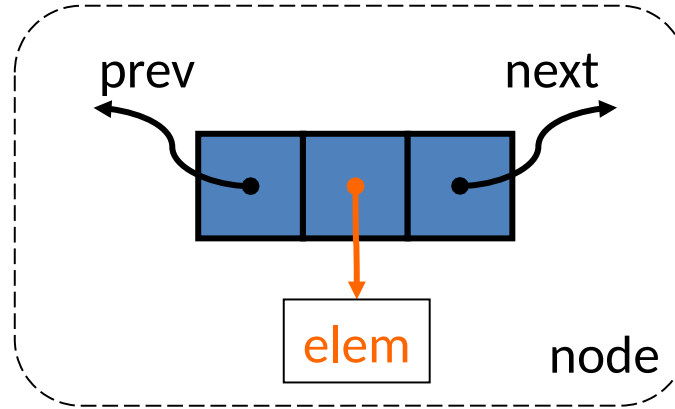
- Key ideas
  - Keep a **previous** pointer *in addition to a next pointer* at every node
  - Keep **dummy nodes** at the head and tail



- Keep a **dummy node** at the head of circular list



# Doubly Linked List Node



```
template <typename E>
struct DNode                                     // doubly linked list node
{
    DNode (const E & element) : elem(element);

    E          elem;                             // node element value
    DNode<E> *prev = nullptr;                     // previous node in the list
    DNode<E> *next = nullptr;                     // next node in the list
};
```



# Doubly Linked List constructor

- Create sentinels and interlink them

```
template <typename E>
DoublyLinkedList<E>:: DoublyLinkedList<E>() { // constructor
    head = new Node<E>; // create dummy nodes
    tail = new Node<E>;
    // have them point to each other
    head->next = tail;
    tail->prev = head;
}
```

- Empty linked list: only has two dummy nodes

