

Key switching:

$$e + \overbrace{pt}^{[ct_0 + ct_1 * S_{fr}]} \xrightarrow{S_{fr}} ct \in R_Q^2$$

#1

$$P \cdot S_{fr} \vec{g} + \vec{e} \xrightarrow{S_{to}} \vec{swr} \in (R_{QP}^2)^{DNUM}$$

$$\vec{g} \in \mathbb{Z}_{QP}^{DNUM}, \quad \vec{g}^{-1}(ct_1) \in R_{QP}^{DNUM}$$

$$s.t. \quad \vec{g} \circ \vec{g}^{-1}(ct_1) = ct_1 \in R_Q$$

$$\begin{array}{c}
 P S_{fr} \vec{g} * \vec{g}^{-1}(ct_1) \xrightarrow{S_{to}} \vec{g}^{-1}(ct_1) * \vec{swr} \in (R_{QP}^2)^{DNUM} \\
 + \vec{e} * \vec{g}^{-1}(ct_1) \\
 \downarrow \\
 P S_{fr} * ct_1 + \vec{e} \circ \vec{g}^{-1}(ct_1) \xrightarrow{S_{to}} \in R_{QP}^2 \\
 \downarrow \quad \quad \quad \oplus \downarrow \\
 S_{fr} * ct_1 + \text{small error} \xrightarrow{S_{to}} \tilde{ct} \in R_Q^2 \quad \quad \quad RS \\
 \downarrow \quad \quad \quad \downarrow \\
 pt + \text{small error} \xrightarrow{S_{to}} \boxed{[ct_0, 0] + \tilde{ct}} \in R_Q^2 \\
 \quad \quad \quad \text{out}
 \end{array}$$

template < int N, int L, int DNUM, int K >

```

void rs( const uint64_t q[L],
        const uint64_t p[K],
        const uint64_t swr[DNUM][2][full levelDNUM*K+K][N],
        const uint64_t c_hat[2][L][N],
        uint64_t out[2][L][N]) {

```

```

    uint64_t a[L][N];
    for(int i=0; i<L; i++)
        for(int j=0; j<N; j++) a[i][j] = c_hat[i][j][j];
    intt <N, L> (q, a);

```

$q = ct_1$

```
uint64_t ginv [DNUM][L+K][N];
```

```
gadget_ginva <N, L, DNUM, K> (g, p, a, ginv);
```

$$\vec{g}(\omega) \in R_{QP}^{DNUM}$$

```
uint64_t qp [L+K];
```

```
for (int i=0; i<L; i++) gp[i] = g[i];
```

```
for (int i=0; i<K; i++) gp[L+i] = p[i];
```

$$gp = [g_0 \dots g_{L-1} p_0 \dots p_{K-1}]$$

```
uint64_t sum [2][L+K][N];
```

```
for (int i=0; i<L+K; i++)
```

```
for (int j=0; j<N; j++) { sum[0][i][j] = 0;
```

```
sum[1][i][j] = 0; }
```

$$sum = \sum_d \vec{g}(\omega)[d] * \hat{sw}_k[d] \text{ in } R_{QP}$$

```
for (int d=0; d<DNUM; d++) {
```

```
ntt <N, L+K> (gp, ginv[d]);
```

```
for (int i=0; i<L; i++)
```

```
for (int j=0; j<N; j++) {
```

```
sum[0][i][j] = (sum[0][i][j] + mul_mod(ginv[d][i][j],
```

```
sw_k[d][0][i][j], q[i]) % q[i];
```

```
sum[1][i][j] = (sum[1][i][j] + mul_mod(ginv[d][i][j],
```

```
sw_k[d][1][i][j], q[i]) % q[i];
```

```
}
```

```
for (int i=0; i<K; i++)
```

```
for (int j=0; j<N; j++) {
```

```
sum[0][L+i][j] = (sum[0][L+i][j] + mul_mod(ginv[d][L+i][j],
```

```
sw_k[d][0][DNUM*K+i][j], p[i]) % p[i];
```

```
sum[1][L+i][j] = (sum[1][L+i][j] + mul_mod(ginv[d][L+i][j],
```

```
sw_k[d][1][DNUM*K+i][j], p[i]) % p[i];
```

```
}
```

```
}
```

```
RS_hat <N, L+K, L> (gp, sum, out);
```

```
for (int i=0; i<L; i++)
```

```
for (int j=0; j<N; j++) out[0][i][j] = (out[0][i][j] + sum[0][i][j]) % q[i];
```

modUp operation

$$QP = [q_0 \dots q_{L-1}] [p_0 \dots p_{K-1}]$$

$$a = [a]_Q \quad \text{extended}$$

$$a = \sum_{j=0}^{L-1} a[j] \frac{Q}{q[j]} \left(\frac{Q}{q[j]} \right)^{-1}_{q[j]}$$

```
template < int N, int LplusK >
```

```
void modUp( const uint64_t qp[LplusK], int L,
```

```
uint64_t a[LplusK][N]) {
```

```
const uint64_t* q = qp; int K = LplusK - L;
```

```
const uint64_t* p = qp + L;
```

```
uint64_t table1[LplusK][LplusK];
```

```
for(int j=0; j<L; j++)
```

```
for(int k=0; k<K; k++){
```

```
table1[j][k] = 1;
```

```
for(int i=0; i<L; i++)
```

```
if(i != j) table1[j][k] = mul_mod(table1[j][k],
```

```
q[i] % p[k], p[k]);
```

$$\text{table1}[j][k] = \text{mod}\left(\frac{Q}{q[j]}, p[k]\right)$$

```
uint64_t table2[LplusK];
```

```
for(int j=0; j<L; j++){
```

```
table2[j] = 1;
```

```
for(int i=0; i<L; i++)
```

```
if(i != j) table2[j] = mul_mod(table2[j],
```

```
inv_mod(q[i] % q[j], q[j]), q[j]);
```

$$\text{table2}[j] = \text{inv_mod}\left(\frac{Q}{q[j]}, q[j]\right)$$

```
uint64_t table3[LplusK];
```

```
for(int k=0; k<K; k++){
```

```
table3[k] = 1;
```

```
for(int j=0; j<L; j++)
```

```
table3[k] = mul_mod(table3[k],
```

```
q[j] % p[k], p[k]);
```

$$\text{table3}[k] = \text{mod}(Q, p[k]);$$

```
}
```

```
for(int i=0; i<N; i++){
```

```
uint64_t b[L+K]; int count=0;
for(int j=0; j<L; j++){
    b[j] = mul_mod(a[j][i],
                  table2[j], q[j]);
    if(2*b[j] >= q[j]) count++;
}
```

$$b[j] = \frac{a[j] \cdot \left(\frac{Q}{q[j]}\right)^{-1}}{\text{in } q[j]}$$

$$\text{count}++ \text{ if } b[j] \geq \frac{q[j]}{2}$$

```
for(int k=0; k<K; k++){
```

```
a[L+K][i] = 0;
for(int j=0; j<L; j++){
    a[L+K][i] = (a[L+K][i] +
                 mul_mod(b[j] % p[k], table1[j][k], p[k]) % p[k]);
    if(count > 0)
        a[L+K][i] = (a[L+K][i] + mul_mod(table3[k], p[k]-count, p[k]))
                     % p[k];
}
```

$$a[L+K]$$

$$= \text{mod}\left(\sum_j b[j] \frac{Q}{q[j]}, p[k], p[k]\right) - \text{count} \cdot Q$$

3

3

Switch Key is generated at the full level
and reused at every level

#5

Full level
($L = \text{DNUM} \times K$)
e.g. $L = 9$

$$Q_9 = \underbrace{q_0 q_1 q_2}_{D_0} \underbrace{q_3 q_4 q_5}_{D_1} \underbrace{q_6 q_7 q_8}_{D_2}, \quad P = \underbrace{P_0 P_1 P_2}_{2 \times \text{DNUM}}$$

$$S_{fr}, S_{to} \in \mathbb{Z}[x]/x^{N+1} \Rightarrow \text{swr} \in (R_{Q_9 P})^{2 \times \text{DNUM}}$$

At level 5,

modDown to $Q_5 P$

$$Q_5 = \underbrace{q_0 q_1 q_2}_{D_0} \underbrace{q_3 q_4}_{D_1} \underbrace{1}_{D_2}, \quad P = P_0 P_1 P_2$$

$$\vec{g} = [g_0, g_1, g_2] \in \mathbb{Z}_{Q_5 P}^{\text{DNUM}}, \quad P \vec{g} \in \mathbb{Z}_{Q_5 P}^{\text{DNUM}}$$

$$\begin{aligned} \mathbb{Z}_{Q_5 P} &\cong \mathbb{Z}_{q_0} \times \dots \times \mathbb{Z}_{q_4} \times \mathbb{Z}_{P_0} \times \mathbb{Z}_{P_1} \times \mathbb{Z}_{P_2} \\ P g_0 &= [[P]_{q_0}, [P]_{q_1}, [P]_{q_2}, 0, 0, 0, 0, 0] \\ P g_1 &= [0, 0, 0, [P]_{q_0}, [P]_{q_1}, 0, 0, 0] \\ P g_2 &= [0, 0, 0, 0, 0, 0, 0, 0] \end{aligned}$$

$$\vec{g}^{-1}: \mathbb{Z}_{Q_5} \rightarrow \mathbb{Z}_{Q_5 P}^{\text{DNUM}}$$

$$\vec{g}^{-1}(a)[0] = [[a]_{D_0} \mid \text{modUp extension}]$$

$$\vec{g}^{-1}(a)[1] = \{ \text{modUp extension} \mid [a]_{D_1} \mid \text{modUp extension} \}$$

$$\vec{g}^{-1}(a)[2] = [0 \quad 0 \quad 0]$$

template<int N, int L, int DNUM>

void swrgen(const int Sfr[N],
const int Sto[N],
const uint64_t q[L],
const uint64_t p[L/DNUM],
uint64_t swrk[DNUM][2][L+(L/DNUM)[N]]) {

const int K=L/DNUM;
uint64_t g[DNUM][L+K];
gadget-g<L, DNUM>(q, p, g);

$\vec{g} \in \mathbb{Z}_{pq}^{DNUM}$

assert(L==K*DNUM);

for(int n=0; n<DNUM; n++){
uint64_t pt[L+K][N];
for(int j=0; j<L; j++){

pt = P.Sfr g[n]

mod(P, q[j])

uint64_t P = 1;
for(int k=0; k<K; k++){
P = mul_mod(P, p[k]%q[j], q[j]);

pt = P.g[n].Sfr

uint64_t Pg = mul_mod(P, g[n][j], q[j]);
for(int i=0; i<N; i++){
pt[j][i] = mul_mod((q[j]+Sfr[i])%q[j], Pg, q[j]);

mod(pt, P)=0

for(int j=0; j<K; j++){
for(int i=0; i<N; i++){
pt[j+L][i] = 0;

$\vec{swrk} = \text{enc}(P, Sfr, \vec{g})_{Sto}$

uint64_t zp[L+K];
for(int i=0; i<L; i++) zp[i] = g[i];
for(int j=0; j<K; j++) zp[L+j] = p[j];
enc<N, L+K>(pt, Sto, zp, swrk[n]);

}

}

template< int L, int DNUM, int K> { L may not be the full level

#7

void gadget_g (const uint64_t g[L],
const uint64_t p[K],
uint64_t g[DNUM][L+K]) { $\vec{g} \in \mathbb{Z}_{QP}^{DNUM}$

for (int d=0; d<DNUM; d++)
for (int i=0; i<L+K; i++)
if ((d*K <= i) && (i < (d+1)*K) && (i < L)) g[d][i] = 1;
else g[d][i] = 0;

}

When L=5, K=3, DNUM=3,

see page #5

$\vec{g} = [11100000, 00011000, 00000000]$:

template< int N, int L, int DNUM, int K>

$a \in R_Q, \vec{g}(a) \in R_{QP}^{DNUM}$

void gadget_ginv (const uint64_t q[L],
const uint64_t p[K],
const uint64_t a[L][N],
uint64_t ginv[DNUM][L+K][N]) {

for (int d=0; d<DNUM; d++) {

uint64_t qp[L+K];

for (int i=0; i<L; i++) qp[i] = q[i];

for (int i=0; i<K; i++) qp[L+i] = p[i];

$QP = [q_0, q_1, q_2 | q_3, q_4 | p_0, p_1, p_2]$

int off1 = (d*K < L) ? (d*K) : L;

int off2 = ((d+1)*K < L) ? ((d+1)*K) : L;

uint64_t temp[K];

for (int i=off1; i<off2; i++) temp[i-off1] = qp[i];

for (int i=off1-1; i>=0; i--) qp[i+off2-off1] = qp[i];

for (int i=off1; i<off2; i++) qp[i-off1] = temp[i-off1];

for (int i=off1; i<off2; i++)
for (int j=0; j<N; j++)
ginv[d][i-off1][j] = a[i][j];

if (off2 > off1)
modUp < N, L+K> (qp, off2-off1, ginv[d]);

for (int j=0; j<N; j++) {

for (int i=0; i<off2-off1; i++) temp[i] = ginv[d][i][j];

for (int i=0; i<off1; i++) ginv[d][i][j] = ginv[d][i+off2-off1][j];

for (int i=0; i<off2-off1; i++) ginv[d][i+off1][j] = temp[i];

}

}

}

d=1: $q_0, q_1, q_2 | \boxed{q_3, q_4} | p_0, p_1, p_2$
 $\min(dK, L) \quad \min((d+1)K, L)$
d=2: $q_0, q_1, q_2, q_3, q_4 | p_0, p_1, p_2$
 $\min(dK, L) = \min((d+1)K, L)$

QP, a

reorder \uparrow off1 \uparrow off2

\uparrow off2-off1 \uparrow off2

$\vec{g}(a)[d]$

modUp reorder

$\vec{g}(a)[d]$