

PRACTICE MIDTERM EXAMINATION

Disclaimer: *this practice exam is my attempt to help you learn the material in this course. Things in here might not be in the real exam and vice versa. Don't use this as your main study. There might be some typos and/or mistakes.*

1. Implement a function `sumOdd(n)` in the x86-64 assembly language. This function computes and returns the sum of all odd numbers between 1 and n . For example, when $n = 6$, the program computes $1 + 3 + 5 = 9$ and returns 9 as the output.

```

                                sumOdd.s
.global sumOdd
.text
sumOdd:
    xor     %rax, %rax
    mov     $1, %rcx
loop:   cmp     %rdi, %rcx
        jg     done
        add     %rcx, %rax
        add     $2, %rcx
        jmp    loop
done:   ret

```

```

                                testsumOdd.c
#include<stdio.h>

long sumOdd(long n);

int main() {
    printf("sumOdd(%d) = %ld\n", 5, sumOdd(5));
    printf("sumOdd(%d) = %ld\n", 6, sumOdd(6));
    return 0;
}

```

2. Implement a function `manhattan(A, B, n)` in the x86-64 assembly language. This function computes the Manhattan distance between two vectors $A = [a_1, a_2, \dots, a_n]$ and $B = [b_1, b_2, \dots, b_n]$ of size n , where

$$\text{manhattan}(A, B) = |a_1 - b_1| + |a_2 - b_2| + \dots + |a_n - b_n|$$

```

                                manhattan.s
.global manhattan
.text
manhattan:
    xor     %rcx, %rcx
    xor     %rax, %rax
loop:   cmp     %rdx, %rcx
        jge     done
        mov     (%rdi), %r8
        sub     (%rsi), %r8
        test    %r8, %r8
        jge     g0
        neg     %r8
g0:     add     %r8, %rax
        inc     %rcx
        add     $8, %rdi
        add     $8, %rsi
        jmp    loop
done:   ret

```

testmanhattan.c

```
#include<stdio.h>

long manhattan(long A[], long B[], long n);

int main() {
    long X[] = {1,2,3,4};
    long Y[] = {2,2,2,2};
    long Z[] = {0,0,0,0};
    printf("Output = %ld\n", manhattan(X, Y, 4));
    printf("Output = %ld\n", manhattan(X, Z, 4));
    printf("Output = %ld\n", manhattan(Y, Z, 4));
    return 0;
}
```

3. Convert the following assembly function into C, Java or Python.

func.s

```
.global func
.text
func: xor    %rax, %rax
loop: cmp    %rsi, %rax
      jge    done
      mov    (%rdi), %r8
      cmp    %r8, %rdx
      jle    l1
      mov    %r8, (%rcx)
      add    $8, %rcx
l1:    add    $8, %rdi
      inc    %rax
      jmp    loop
done:  mov    $-999, %r8
      mov    %r8, (%rcx)
      ret
```

The function is used with the following C program.

testfunc.c

```
#include<stdio.h>

void func(long A[], long n, long x, long B[]);

int main() {
    long X[] = {1, 2, 3, 4, 5, 6, 7, 8};
    long Y[10];
    long i;

    func(X, 8, 4, Y);
    for(i=0; Y[i]!=-999; i++)
        printf("%ld\n", Y[i]);
    return 0;
}
```

func.c

```
void func(long A[], long n, long x, long B[]) {
    int i, j=0;
    for(i=0; i<n; i++) {
        if (x > A[i]) {
            B[j] = A[i];
            j++;
        }
    }
    B[j] = -999;
}
```

4. Write the control signals when the following instructions are executed in the single-cycle implementation of Y86.

```
0x2000: rmmov %rax, 0x100(%rcx)
        add  %rcx, %rdx
        push %edx
```

rmmov %rax, 0x100(%rcx)			
Control Signal	Value	Control Signal	Value
<i>needRegids</i>	1	<i>needValC</i>	1
<i>regA</i>	10	<i>regB</i>	10
<i>regE</i>	00	<i>regM</i>	00
<i>aluA</i>	10	<i>aluB</i>	1
<i>aluF</i>	00	<i>setCond</i>	0
<i>memRead</i>	0	<i>memWrite</i>	1
<i>memAddr</i>	0	<i>memData</i>	0
<i>newPC</i>	00		

add %rcx, %rdx			
Control Signal	Value	Control Signal	Value
<i>needRegids</i>	1	<i>needValC</i>	0
<i>regA</i>	10	<i>regB</i>	10
<i>regE</i>	10	<i>regM</i>	00
<i>aluA</i>	11	<i>aluB</i>	1
<i>aluF</i>	00	<i>setCond</i>	1
<i>memRead</i>	X	<i>memWrite</i>	0
<i>memAddr</i>	X	<i>memData</i>	X
<i>newPC</i>	00		

push %rdx			
Control Signal	Value	Control Signal	Value
<i>needRegids</i>	1	<i>needValC</i>	0
<i>regA</i>	10	<i>regB</i>	01
<i>regE</i>	10	<i>regM</i>	00
<i>aluA</i>	00	<i>aluB</i>	1
<i>aluF</i>	00	<i>setCond</i>	0
<i>memRead</i>	0	<i>memWrite</i>	1
<i>memAddr</i>	0	<i>memData</i>	0
<i>newPC</i>	00		

5. Suppose we want to add a new instruction **imrmov** that copies 6-byte value from a memory address specified by a constant to a register.

For example, **imrmov 0x8004, %rax** copies the 8-byte value at address **0x8004** to register **%rax**.

- (a) Design the format of this instruction type.

imrmov D, rA can be represented in machine code using a format similar to **irmov** instruction, i.e. **|icode|ifun|rA|F|D|** Here, we use **rA** instead of **rB** because of the link from the data memory to the register file.

- (b) Design the computation steps for this instruction type.

Fetch	<i>icode, ifun</i>	$icode:ifun \leftarrow M_1[PC]$
	<i>rA, rB</i>	$rA:rB \leftarrow M_1[PC + 1]$
	<i>valC</i>	$valC \leftarrow M_8[PC + 2]$
	<i>valP</i>	$valP \leftarrow PC + 10$
Decode	<i>valA, srcA</i>	
	<i>valB, srcB</i>	
Execute	<i>valE</i>	$valE \leftarrow 0 + valC$
	<i>cond</i>	
Memory	<i>valM</i>	$valM \leftarrow M_8[valE]$
Write back	<i>dstE</i>	
	<i>dstM</i>	$R[rA] \leftarrow valM$
PC update	<i>PC</i>	$PC \leftarrow valP$

- (c) Design the control signals for this instruction for the single-cycle implementation.

Control Signal	Value	Control Signal	Value
<i>needRegids</i>	1	<i>needValC</i>	1
<i>regA</i>	00	<i>regB</i>	00
<i>regE</i>	00	<i>regM</i>	10
<i>aluA</i>	10	<i>aluB</i>	0
<i>aluF</i>	00	<i>setCond</i>	0
<i>memRead</i>	1	<i>memWrite</i>	0
<i>memAddr</i>	0	<i>memData</i>	X
<i>newPC</i>	00		

6. Write the truth table for the **memData** control signal.

<i>instruction</i>	<i>icode</i>	<i>memData</i>
hlt	0000	X
nop	0001	X
rrmov	0010	X
irmov	0011	X
rmmov	0100	0
mrmov	0101	X
OP	0110	X
jXX	0111	X
call	1000	1
ret	1001	X
push	1010	0
pop	1011	X
	1100	X
	1101	X
	1110	X
	1111	X