

Lecture 17 : Rapidly Decreasing Singular Values

$X = n \times n$ real matrix

singular values $\sigma_1(x) \geq \sigma_2(x) \geq \dots \geq \sigma_n(x) \geq 0$

Fact: If $\sigma_{k+1}(x) = 0$, $\sigma_k(x) > 0$:

① $\text{rank}(x) = k$

② $x = u_1 v_1^T + \dots + u_k v_k^T$

$$\overbrace{\quad}^1 + \dots + \overbrace{\quad}^k$$

③ $\dim(\text{col. space}) = \dim(\text{rowspace}) = k$

Def x = low rank, if $2kn < n^2 \Rightarrow k < n/2$

Often, we demand $k \ll n/2$

What do low rank matrices look like?

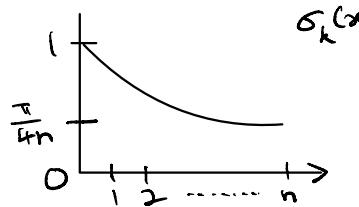
rank $\sim 1 \Rightarrow$ matrix highly aligned

Triangular flag

$$\begin{bmatrix} \text{triangular flag} \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ \vdots & \ddots \\ 0 & \dots & 1 \end{bmatrix} = x, x^{-1} = \begin{bmatrix} 1 & & & \\ \vdots & \ddots & & \\ 0 & \dots & 1 & 0 \\ 0 & 0 & \dots & 1 \end{bmatrix}, (x^T x)^{-1} = \begin{bmatrix} 1 & -1 & & \\ -1 & 2 & \dots & \\ & \ddots & \ddots & \\ & & \ddots & 2 \end{bmatrix}$$

$$\Rightarrow \sigma_k(x) = \left[2 \sin\left(\frac{\pi(2k-1)}{2(2n+1)}\right) \right]$$

$$\sigma_1(x) \approx \frac{2n}{\pi}, \sigma_n(x) \approx \frac{1}{2}$$



\Rightarrow Hence, triangular numbers are bad for low rank.

Japan Flag

$$\text{rk}(\begin{array}{|c|}\hline 0 \\ \hline \end{array})$$

$$\leq \text{rk}(\begin{array}{|c|}\hline 0 \\ \hline 0 \\ \hline \end{array}) + \text{rk}(\begin{array}{|c|}\hline 0 \\ \hline \end{array})$$

$$\leq \text{rk}(\begin{array}{|c|}\hline 0 \\ \hline 0 \\ \hline 0 \\ \hline \end{array}) + 1$$

$$\leq \text{rk}(\begin{array}{|c|}\hline 0 \\ \hline \end{array}) + \text{rk}(\begin{array}{|c|}\hline 0 \\ \hline \end{array}) + 1$$

$$\leq r_c(1 - \sqrt{2}/2)^T + r_c(1 - \sqrt{2}/2)^T + 1 = \frac{1}{2}r + 1$$

\Rightarrow low rank matrix

Numerical rank

for $0 < \varepsilon < 1$ (ε = tolerance)

$$\text{rank}_\varepsilon(x) = k \quad \sigma_{k+1}(x) \leq \varepsilon \sigma_1(x)$$

$$\sigma_k(x) > \varepsilon \sigma_1(x)$$

$$\text{rank}_0(x) = \text{rank}(x)$$

when $\varepsilon = 0$

Eckart - Young

$$\sigma_{k+1}(x) = \|x - x_k\|_2$$

x_k = best rank- k

Matrices of numerical low rank

- All low rank matrices

$$H_{jk} = \frac{1}{j+k-1}, \quad V = \begin{bmatrix} 1 & x_1 & \cdots & x_1^{n-1} \\ 1 & x_2 & \cdots & x_2^{n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & \cdots & x_n^{n-1} \end{bmatrix} \quad x_i? \text{ s real}$$

(Hilbert)

Vandermonde

often, we want V^{-1} , but difficult because V has low numerical rank.

Why so many low rank matrices in the world?

Reason 1: The world is smooth (Reade, 1983)

$$\text{Ex. } p(x,y) = 1 + xy$$

$$x_{jk} = p(j,k) = 1 + j + jk$$

$$x = \begin{bmatrix} 1 & 1 & \cdots & 1 \\ 1 & 2 & \cdots & 2 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & n & \cdots & n \end{bmatrix} + \begin{bmatrix} 0 & 0 & \cdots & 0 \\ 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 0 \end{bmatrix} + \begin{bmatrix} 1 & 2 & \cdots & n \\ 1 & 4 & \cdots & 2n \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 2n & \cdots & n^2 \end{bmatrix}$$

All rank-1 matrices

$$\text{rank}(x) \leq 3$$

In general, $p(x, y) = \sum_{s=0}^{m-1} \sum_{t=0}^{m-1} a_{st} x^s y^t$

$$x_{jk} = p(j,k), \text{ rank}(x) \leq m^2$$

If $x_{jk} = \frac{1}{j!k!}, f(x, y) = \frac{1}{x!y!}$ \Rightarrow smooth function like this can be well approximated by polynomials, therefore have low numerical rank.

$$\text{Find } p \quad |f(x, y) - p(x, y)| \leq \epsilon_n \|f\|_2$$

approx of f

$$Y_{jk} = p(j,k) = \text{finite rank}$$

$$\|x - Y\|_2 \leq \epsilon \|x\|_2$$

$$\text{rank}(H_{1000}) = 1000, \text{rank}_\epsilon(H_{1000}) = 28 \quad \epsilon = 10^{-15}$$

But, Reade's argument gets $\text{rank}_\epsilon(H_{1000}) \leq 719$

Another reason : The word is Sylvester, matrices satisfy $Ax - xB = C$
for some A, B, C

Thm : If x satisfies $Ax - xB = C$

(and A normal, B normal), then

$$G_{k+1}(x) \leq Z_k(E, F) G_k(x)$$

$$\text{rank}(C) = r \quad \text{Zolotarev no.}$$

$E = \text{set eigs of } A$

$F = \text{set eigs of } B$