#### A0201613Y – Computational Assignment Report

#### 1 Facility Location Problem

(a) We first note that there are m customers so  $i \in \{1, 2, ..., m\}$ . There are n facilities so  $j \in \{1, 2, ..., n\}$ . The linear inequalities listed under (FLP) are:

$$x_{i,j} \le y_j \text{ for all } i, j$$
 (1.1)

$$0 \le x_{i,j} \le 1 \tag{1.2}$$

For each facility j, there are m choices for customer i. With n facilities, there are mn possible choices of i, j. Hence, for (1.1), there are mn inequalities.

As for (1.2), this compound inequality consists of two separate inequalities,  $0 \le x_{i,j}$  and  $x_{i,j} \le 1$ , so there are 2mn inequalities.

Under (FLP), we have mn + 2mn = 3mn linear inequalities.

The linear inequalities listed under (AFL) are:

$$\sum_{i=1}^{m} x_{i,j} \le m y_j \text{ for all } j$$
 (1.3)

$$0 \le x_{i,j} \le 1 \tag{1.4}$$

For (1.3), the number of inequalities only depends on the possible choices of j. As such, there are only n inequalities for (1.3).

For (1.4), the inequality is the same as (1.2). For the same reasons, there are 2mn inequalities.

Under (AFL), we have n + 2mn = (2m + 1)n linear inequalities.

Thus, for  $m \ge 1$ , we have  $(2m+1)n \le 3mn$ . (AFL) has fewer inequalities for m > 1.

(b) For a solution to be feasible, all constraints must be satisfied. Comparing the two formulations, all the constraints are exactly the same except (1.1) and (1.3). I want to show that (1.1) and (1.3) are equivalent, i.e.

$$(1.1) \iff (1.3)$$

*Proof:* (1.1)  $\implies$  (1.3): Assuming  $x_{i,j} \leq y_j$  for all i, j, we sum up these inequalities over all i to get,

$$x_{1,j} + x_{2,j} + \dots + x_{m,j} \le y_j + y_j + \dots + y_j \text{ for all } j$$
$$\sum_{i=1}^m x_{i,j} \le my_j \text{ for all } j$$

(1.1) 
$$\iff$$
 (1.3): Assuming  $\sum_{i=1}^{m} x_{i,j} \leq my_j$  for all  $j$ ,
$$x_{1,j} + x_{2,j} + \dots + x_{m,j} \leq y_j + y_j + \dots + y_j \text{ for all } j$$
(1.5)

From (1.5), since  $y_j \in \{0, 1\}$  and  $0 \le x_{i,j} \le 1$ , if  $y_j = 0$ ,  $x_{i,j} = 0$  for all *i*. We get that at  $y_j = 0$ ,  $x_{i,j} = y_j$  for all *i*.

At  $y_j = 1$ , since  $0 \le x_{i,j} \le 1$ , we have  $x_{i,j} \le y_j$  for all i. Taking the union of these two possible cases of  $y_j \in \{0,1\}$ , we get

$$x_{i,j} \leq y_j$$
 for all  $i, j$ 

We have shown that constraints (1.1) and (1.3) are equivalent. Since all the constraints of (FLP) and (AFL) are equivalent, the set of feasible solutions of (FLP) and (AFL) are equal. With the same feasible region and also the same objective function, the two formulations are thus equivalent.

(c) , (d) , (e) After replacing the binary constraints with the box constraints for  $y_j$ , we arrive at the following relations:

$$(AFL-Val) = (FLP-Val)$$
 (1.6)

$$(FLP-LR-Val) \le (FLP-Val)$$
 (1.7)

$$(AFL-LR-Val) \le (AFL-Val) \tag{1.8}$$

For (1.6), since (FLP) and (AFL) are shown to be equivalent in part (b), this implies that their optimal values are equal.

Following that, we should note that their corresponding LRs may not necessarily have the same optimal values. When the binary constraint is dropped, (FLP-LR) and (AFL-LR) do not have the same set of feasible solutions and are thus not equivalent. In other words,

 $(1.1) \iff (1.3)$  may not hold when the binary constraint is dropped.

Proof with counter-example: (1.1)  $\iff$  (1.3): Suppose  $\sum_{i=1}^{m} x_{i,j} \leq my_j$  for all j,

$$x_{1,j} + x_{2,j} + \dots + x_{m,j} \le y_j + y_j + \dots + y_j \text{ for all } j$$
 (1.9)

The relation holds trivially for m=1. Hence, we find a counter-example at  $m \geq 2$ . For  $m \geq 2$ , we assume  $y_j = \gamma$ , where  $\gamma \in \{y_j : 0 \leq y_j \leq 1\} \setminus \{0,1\}$ . We assign  $x_{1,j} = \gamma + \varepsilon$  for  $0 < \varepsilon \leq \gamma$  while  $x_{i,j} = 0$  for all  $i \neq 1$ . Under (1.9), we have  $\gamma + \varepsilon \leq 2\gamma \leq m\gamma$  but  $x_{1,j} = \gamma + \varepsilon > \gamma = y_j$ . We have that  $x_{i,j} \leq y_j$  is not satisfied at i = 1.

For (1.7) and (1.8), the optimal values for the LRs are lower bounds of the corresponding optimal values of the MILPs. Since we are minimizing over a larger constraint set in the LRs, this means that the feasible region of the MILPs is contained in the feasible region of the corresponding LRs. Moreover, if the optimal solutions of the LRs are integral, these integral solutions are also feasible in the corresponding MILPs and their optimal values will be equal.

- (f) Since all the random variables are defined uniformly over a unit interval or unit square, np.random.rand() is used to generate them. With the coordinates of the customers and the facilities, we compute the distance matix using the Euclidean distance. To generate 100 sets of the random variables, for-loops are used. Here, we set a seed value to reference back to the same dataset to use as a numerical example for part (h).
- (g) For each of the four models, the for-loop is used when running the optimization, iterating through each set of random variables. At every iteration, we record the optimal value. As such, we will have 100 different optimal values recorded under each model.
- (h) Due to rounding error from computer arithmetic, to compare the equality of two optimal values for every set of random variables, we choose to check that their difference is approximately 0.

We first examine how often (FLP-Val) is equal to (FLP-LR-Val). Iterating through the 100 objective values of each model, the differences are all of magnitude  $10^{-15}$ , which is approximately 0. Hence, we conclude that (FLP-Val) is always equal to (FLP-LR-Val). This is an equality because we can check that (FLP-LR) always has an integral solution. This supports our answers in part (e), that this integral solution will also be feasible in (FLP), thus attaining the same optimal values.

For (AFL-Val) and (AFL-LR-Val), the smallest difference over the 100 pairs of objective values is approximately 0.668, which is significantly different. Hence, we conclude that (AFL-Val) is always different from (AFL-LR-Val), or specifically (AFL-Val) > (AFL-LR-Val). The reason for the strict inequality is that the solutions of (AFL-LR) are always non-integers. Hence, this non-integral solution will not be feasible for (AFL).

Following the same method, we can also check that (FLP-Val) and (AFL-Val) are always equal. On the other hand, their corresponding (FLP-LR-Val) and (AFL-LR-Val) always have different optimal values, as proven in part (e) that the feasible regions of the two LRs are not equivalent.

Overall, the findings support the relations explained in part (e).

### 2 Traveling Salesman Problem

- (a) To compute the total distance of a tour, a for-loop is used in the function to iterate through every consecutive pair of locations and sum up their distances. We then add the distance between the last and first location to indicate the return to the starting location.
- (b) For the function that outputs a perturbed tour, since we have to pick two unique indices i < j, we use random.sample() to pick two indices without replacement. For this first perturb rule, we use list slicing to inverse all the cities between indices i and j, inclusive.
- (c) We initialize a random tour using random.shuffle() to shuffle the indices of the cities. To compare results with the second perturb rule in part (f), we use .copy() to create the same initialized tour to use afterwards. The simulated annealing algorithm is set to run for 10000 loops, iterating through the sequence of steps given in the question.

$$\exp(-\frac{f_{\text{cand}} - f_{\text{curr}}}{T}) \ge u \tag{2.1}$$

Under step 3 of each iteration, if  $f_{\text{cand}} \leq f_{\text{curr}}$ , the condition in (2.1) always hold, i.e. the candidate tour of shorter distance gets accepted with 100% probability. However, when  $f_{\text{cand}} > f_{\text{curr}}$ , depending on the value of T and u, a bad candidate tour may be accepted even if it has a longer distance.

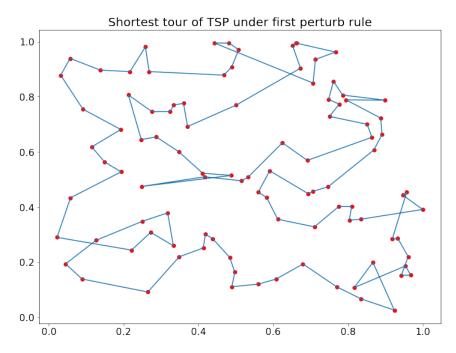
We infer that, at very small values of T and  $f_{\text{cand}} > f_{\text{curr}}$ , the probability of the bad candidate tour being accepted approaches 0. As such, if the initial value of T is set to be very small, we may be stuck with a local minimum solution.

This is because a large enough value of initial T allows the algorithm to search over the whole feasible region [1], even if a bad candidate tour is accepted in the process. This results in a higher probability of getting a global minimum solution. Likewise, the initial T should not be too large such that all the bad candidate tours are accepted throughout all iterations, making the algorithm inefficient.

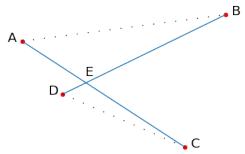
(d) Since the coordinates of the cities are defined uniformly over a unit square, we use np.random.rand() to generate a list of these coordinates. For reference, we name the cities by their indices in the list here. We can then compute the distance matrix using the Euclidean distance. As defined, the distance between any two cities in opposite directions will be the same. This problem can then be formulated as an undirected graph (as will be shown in the graph plot in part (e)).

We are now set to run the algorithm. After experimenting with a few choices of initial T (as mentioned in part (c)), the distance of the shortest tour generated is not very sensitive to the value of initial T. Hence, we just fix initial T = 100.

(e)

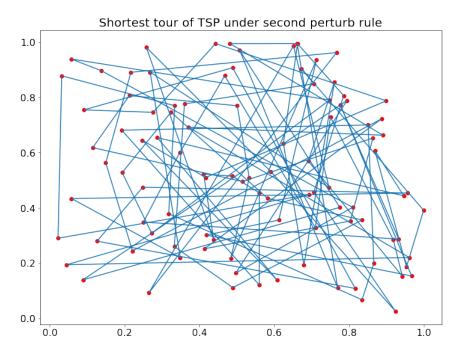


The above graph is one random instance of a solution. The red dots denote the location of the 100 cities. The blue line shows the shortest possible tour connecting every city. Evidently, the graph have edges that intersect, suggesting that the tour is not optimal.



We look at an intersection of edges formed by any four cities A, B, C and D as shown above, where E is the intersection point. By triangle inequality, |AB| < |AE| + |BE| and |CD| < |CE| + |DE|. By replacing the blue lines with the dotted lines, we get a shorter tour. Hence, the above solution generated is not optimal.

(f) Now, we create a new function that outputs a perturbed tour using a second perturb rule. This second perturb rule is now to swap only the city with index i and the city with index j. This can be done by reassigning the indices of the two cities in the list. We then run the algorithm again using the new function.



The above graph is one solution generated using the second perturb rule. Compared to the first graph, this graph has significantly more intersection of edges. This second perturb rule seems to be less efficient. To verify this, we compare their differences:

Distance of initialized tour = 50.5172279351468

	First perturb rule	Second perturb rule
Number of accepted candidate tours	889	5266
Distance of shortest tour generated	9.199323401302737	48.867230128219546



In general, the first perturb rule changes two edges, while the second perturb rule changes four edges for  $j-i \geq 3$ , as indicated by the orange edges in the above subgraph.

Since the first perturb rule changes less edges, it is more likely to pick solutions in its neighbourhood with similar total distance. It is thus able to approach a minimum steadily as bad candidate tours that get accepted do not increase the distance as much.

However, the second perturb rule changes more edges, so it may produce bad candidate tours that have relatively larger distance. The distance of the tour is more likely to increase than to decrease [1]. This results in a relatively large number of accepted candidate tours as the algorithm cannot approach a minimum efficiently. Hence, the distance of the shortest tour remains close to that of the initialized tour.

With the shortest tour generated by the first perturb rule being significantly better, we conclude that the first perturb rule is more efficient.

# References

[1] Wikipedia contributors. Simulated annealing — Wikipedia, The Free Encyclopedia. [Online; accessed 18-October-2021]. 2021. URL: https://en.wikipedia.org/wiki/Simulated\_annealing.

# Appendix

The codes for the report are appended below. Appendix 1 is used for Question 1 while Appendix 2 is used for Question 2. Some outputs have been cleared as they are too long to be displayed.

## 1. Facility Location Problem

```
In [1]:
         import gurobipy as gp
         from gurobipy import GRB
         from itertools import product
         import numpy as np
In [2]:
         # Set number of customers and facilities
         num\_customers = 20 # m = 20
         num_facilities = 15 # n = 15
         # Define distances
         def compute_distance(loc1, loc2):
             dx = loc1[0] - loc2[0]
             dy = loc1[1] - loc2[1]
             return np.sqrt(dx**2+dy**2)
         # Compute key parameters of model formulation
         cartesian_prod = list(product(range(num_customers), range(num_facilities)))
In [3]:
         # Generating 100 realizations of all the random variables
         customers = []
         facilities = []
         setup_cost = []
         distance = []
         np.random.seed(0) # Set the seed value to reference back to the same dataset
         for j in range(100):
             customers.append([(np.random.rand(),np.random.rand()) for i in range(num_customers)])
             facilities.append([(np.random.rand(),np.random.rand()) for i in range(num_facilities)])
             setup_cost.append([np.random.rand() for i in range(num_facilities)])
         for i in range(100):
             distance.append({(c,f) : compute_distance(customers[i][c], facilities[i][f]) for c, f in cartesian_prod})
```

### **FLP** model

## **AFL** model

### FLP\_LR model

#### AFL LR model

diff\_FLP\_AFL = []

```
In [ ]:
         # AFL_LR model formulation
         m_AFL_LR = gp.Model('facility_location_AFL_LR')
         # Replacing the binary constraint with the box constraint
         select_AFL_LR = m_AFL_LR.addVars(num_facilities, ub=1, vtype=GRB.CONTINUOUS, name='Select_AFL_LR')
         assign_AFL_LR = m_AFL_LR.addVars(cartesian_prod, ub=1, vtype=GRB.CONTINUOUS, name='Assign_AFL_LR')
         # Different constraint for the AFL_LR model
         m_AFL_LR.addConstrs((gp.quicksum(assign_AFL_LR[(c,f)] for c in range(num_customers)) <=</pre>
                              select_AFL_LR[f]*num_customers for c,f in cartesian_prod), name='Setup2ship_AFL_LR')
         m_AFL_LR.addConstrs((gp.quicksum(assign_AFL_LR[(c,f)] for f in range(num_facilities)) ==
                              1 for c in range(num_customers)), name='Demand_AFL_LR')
         Opt_AFL_LR_Val = [] # Record optimal values for AFL_LR model
         for i in range(100):
             m_AFL_LR.setObjective(select_AFL_LR.prod(setup_cost[i])+assign_AFL_LR.prod(distance[i]), GRB.MINIMIZE)
             m_AFL_LR.optimize()
             Opt_AFL_LR_Val.append(m_AFL_LR.objVal)
```

```
Comparisons of results from the four models
In [8]:
          # Comparing the last instance of the optimal values of the four models to get a general idea
          print(Opt_FLP_Val[-1])
          print(Opt_AFL_Val[-1])
          print(Opt_FLP_LR_Val[-1])
          print(Opt_AFL_LR_Val[-1])
         4.776366365335146
         4.776366365335146
         4.776366365335145
         3.4034990953486775
In [9]:
          # Check the difference between FLP_Val and FLP_LR_Val
          diff_FLP = []
          for i in range (len(Opt_FLP_Val)):
              diff_FLP.append(Opt_FLP_Val[i] - Opt_FLP_LR_Val[i])
          print(max(diff_FLP))
          print(min(diff_FLP))
         1.7763568394002505e-15
         -2.6645352591003757e-15
In [10]:
          # Check the difference between AFL_Val and AFL_LR_Val
          diff_AFL = []
          for i in range (len(Opt AFL Val)):
              diff_AFL.append(Opt_AFL_Val[i] - Opt_AFL_LR_Val[i])
          print(max(diff_AFL))
          print(min(diff_AFL))
         2.8961983924675803
         0.66761265347471
In [11]:
          # Check the difference between FLP_Val and AFL_Val
```

```
for i in range(len(Opt_FLP_Val)):
    diff_FLP_AFL.append(Opt_FLP_Val[i] - Opt_AFL_Val[i])

print(max(diff_FLP_AFL))

8.881784197001252e-16
0.0

In [12]:
# Check the difference between FLP_LR_Val and AFL_LR_Val

LR_diff_FLP_AFL = []
for i in range(len(Opt_FLP_LR_Val)):
    LR_diff_FLP_AFL.append(Opt_FLP_LR_Val[i] - Opt_AFL_LR_Val[i])

print(max(LR_diff_FLP_AFL))

print(max(LR_diff_FLP_AFL))

2.8961983924675803
0.66761265347471
```

## Checking solutions of the last instance of the four models

```
In [ ]:
         # Print solution for last instance of FLP model
         for v in m_FLP.getVars():
              print('%s %g' % (v.varName, v.x))
In [ ]:
         # Print solution for last instance of FLP_LR model
         for v in m_FLP_LR.getVars():
              print('%s %g' % (v.varName, v.x))
In [ ]:
         # Print solution for last instance of AFL model
         for v in m_AFL.getVars():
              print('%s %g' % (v.varName, v.x))
In [ ]:
         # Print solution for last instance of AFL_LR model
         for v in m_AFL_LR.getVars():
              print('%s %g' % (v.varName, v.x))
```

# 2. Traveling Salesman Problem

f\_cand = tour\_dist(cand\_tour)

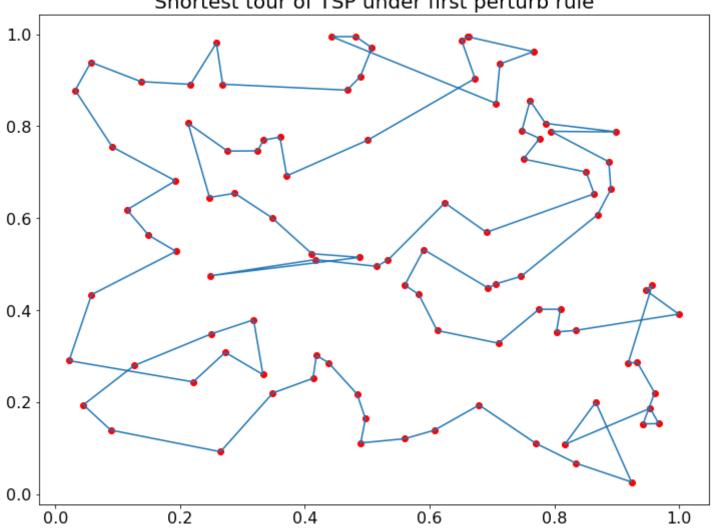
current\_tour = cand\_tour

if  $np.exp(-((f_cand - f_curr) / T)) >= u$ :

u = np.random.rand()

```
In [1]:
         import random
         import csv
         import numpy as np
         from itertools import product
         import matplotlib.pyplot as plt
In [2]:
         # Define distances
         def compute_distance(loc1, loc2):
             dx = loc1[0] - loc2[0]
             dy = loc1[1] - loc2[1]
             return np.sqrt(dx**2+dy**2)
         # Define function that takes in a tour and outputs the total distance
         def tour_dist(tour):
             total_dist = 0
             for i in range(len(tour)-1):
                 dist = distance[(tour[i], tour[i+1])]
                 total_dist += dist
             total_dist += distance[(tour[-1], tour[0])] # distance returning to starting Location
             return total_dist
         # Define function that takes in a tour and outputs a perturbed tour, with first perturb function
         def perturb(tour):
             indices = random.sample(range(0,len(tour)), 2) # random.sample ensures unique indices
             i = min(indices)
             j = max(indices)
             perturbed_tour = tour[:i] + tour[i:j+1][::-1] + tour[j+1:]
             return perturbed_tour
In [3]:
         # Setting parameters
         num_cities = 100 # n = 100
         cart_prod = list(product(range(num_cities), range(num_cities))) # Label the cities by their indices
In [4]:
         # Generate the coordinates of the cities
         cities = [(np.random.rand(),np.random.rand()) for i in range(num_cities)]
         # Compute distance matrix
         distance = ({(c1,c2) : compute_distance(cities[c1], cities[c2]) for c1, c2 in cart_prod})
       Simulated annealing algorithm with first perturb rule
In [5]:
         # Initialize same random tour for both perturb rules for comparison
         cities_index = [i for i in range(num_cities)]
         random.shuffle(cities_index)
         initial_tour = cities_index
In [6]:
         # Initial tour for first perturb rule
         current_tour = initial_tour.copy()
         # Distance of initial tour
         tour_dist(current_tour)
Out[6]: 50.5172279351468
In [7]:
         # Initialize temperature parameter
         T = 100
         T_{update} = 0.99 # eta = 0.99
         # Implement algorithm with first perturb rule
         accepted_count = 0
         for i in range(10000):
             f_curr = tour_dist(current_tour)
             cand_tour = perturb(current_tour)
```

```
accepted_count += 1
             T = T_update * T
        <ipython-input-7-d61a0cc8762c>:15: RuntimeWarning: overflow encountered in exp
          if np.exp(-((f_cand - f_curr) / T)) >= u:
In [8]:
         # Results for first perturb rule
         # Number of candidate tours accepted
         print(accepted_count)
         # Distance of shortest tour
         print(tour_dist(current_tour))
        889
        9.199323401302737
In [9]:
         # Plot the cities with the tour under first perturb rule
         fig, ax = plt.subplots(figsize=(12,9))
         # Plot the coordinates of the cities
         x_coord_of_all_cities = [coord[0] for coord in cities]
         y_coord_of_all_cities = [coord[1] for coord in cities]
         ax.scatter(x_coord_of_all_cities, y_coord_of_all_cities, color = 'red')
         # Plot the shortest tour connecting every city
         x_coord_of_tour = [cities[loc][0] for loc in current_tour]
         y_coord_of_tour = [cities[loc][1] for loc in current_tour]
         ax.plot(x coord of tour, y coord of tour, color = '#1f77b4')
         ax.plot([x_coord_of_tour[-1], x_coord_of_tour[0]], [y_coord_of_tour[-1], y_coord_of_tour[0]],
                 color = '#1f77b4')
         ax.set_title('Shortest tour of TSP under first perturb rule', fontsize = 20)
         plt.xticks(fontsize=16)
         plt.yticks(fontsize=16)
         plt.show()
                             Shortest tour of TSP under first perturb rule
```



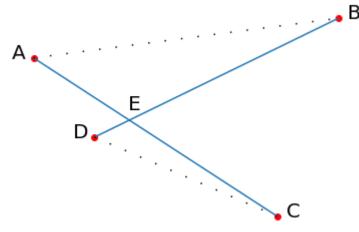
```
In [10]: # Plot the intersection of any two edges in general

fig, ax = plt.subplots()

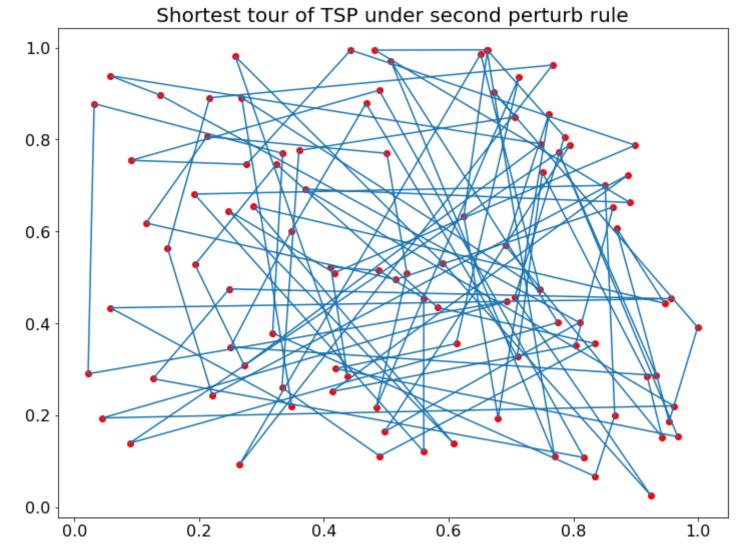
ax.scatter([0.12,0.2,0.1,0.18],[0.14,0.2,0.18,0.1], color = 'red')

ax.plot([0.12,0.2],[0.14,0.2], color = '#1f77b4')
 ax.plot([0.1,0.18],[0.18,0.1], color = '#1f77b4')
 ax.plot([0.1,0.2],[0.18,0.2], linestyle = (0, (1, 10)), color = 'black')
 ax.plot([0.12,0.18],[0.14,0.1], linestyle = (0, (1, 10)), color = 'black')

plt.text(0.093, 0.18, 'A', fontsize = 20)
 plt.text(0.203, 0.2, 'B', fontsize = 20)
 plt.text(0.113, 0.14, 'C', fontsize = 20)
 plt.text(0.113, 0.14, 'D', fontsize = 20)
 plt.text(0.131, 0.154, 'E', fontsize = 20)
 plt.axis('off')
 plt.show()
```



```
Simulated annealing algorithm with second perturb rule
In [11]:
          # Define new perturbed tour function, with second perturb rule
          def n perturb(tour):
              indices = random.sample(range(0,len(tour)), 2) # random.sample ensures unique indices
              i = min(indices)
              j = max(indices)
              tour[i], tour[j] = tour[j], tour[i]
              return tour
In [12]:
          # Initial tour for second perturb rule
          n_current_tour = initial_tour.copy()
          # Distance of initial tour
          tour_dist(n_current_tour)
Out[12]: 50.5172279351468
In [13]:
          # Initialize temperature parameter
          T = 100
          T update = 0.99 # eta = 0.99
          # Implement algorithm with second perturb rule, replace perturb with n_perturb
          n_accepted_count = 0
          for i in range(10000):
              f_curr = tour_dist(n_current_tour)
              cand_tour = n_perturb(n_current_tour)
              f_cand = tour_dist(cand_tour)
              u = np.random.rand()
              if np.exp(-((f_cand - f_curr) / T)) >= u:
                  n_current_tour = cand_tour
                  n_accepted_count += 1
              T = T_update * T
         <ipython-input-13-3cdb247b8186>:15: RuntimeWarning: overflow encountered in exp
           if np.exp(-((f_cand - f_curr) / T)) >= u:
In [14]:
          # Results for second perturb rule
          # Number of candidate tours accepted
          print(n_accepted_count)
          # Distance of shortest tour
          print(tour_dist(n_current_tour))
         5266
         48.867230128219546
In [15]:
          # Plot the cities with the tour under second perturb rule
          fig, ax = plt.subplots(figsize=(12,9))
          # Plot the coordinates of the cities
          x_coord_of_all_cities = [coord[0] for coord in cities]
          y_coord_of_all_cities = [coord[1] for coord in cities]
          ax.scatter(x_coord_of_all_cities, y_coord_of_all_cities, color = 'red')
          # Plot the shortest tour connecting every city
          x_coord_of_tour = [cities[loc][0] for loc in n_current_tour]
          y_coord_of_tour = [cities[loc][1] for loc in n_current_tour]
          ax.plot(x_coord_of_tour, y_coord_of_tour, color = '#1f77b4' )
          ax.plot([x_coord_of_tour[-1], x_coord_of_tour[0]], [y_coord_of_tour[-1], y_coord_of_tour[0]],
                  color = '#1f77b4')
          ax.set_title('Shortest tour of TSP under second perturb rule', fontsize = 20)
          plt.xticks(fontsize=16)
          plt.yticks(fontsize=16)
          plt.show()
```



```
In [16]:
           # Difference between first perturb rule and second perturb rule in general
           fig, ax = plt.subplots(1,2,figsize=(14,1))
           ax[0].scatter([0.05,0.1,0.15,0.2,0.25,0.3], [0.5,0.5,0.5,0.5,0.5,0.5], color = 'red')
           ax[0].plot([0.1,0.15,0.2,0.25], [0.5,0.5,0.5,0.5], color = '#1f77b4')
           ax[0].plot([0.05,0.1],[0.5,0.5], color = 'orange')
           ax[0].plot([0.25,0.3],[0.5,0.5], color = 'orange')
           ax[1].scatter([0.05,0.1,0.15,0.2,0.25,0.3], [0.5,0.5,0.5,0.5,0.5,0.5], color = 'red')
           ax[1].plot([0.15,0.2], [0.5,0.5], color = '#1f77b4')
           ax[1].plot([0.05,0.1,0.15],[0.5,0.5,0.5], color = 'orange')
           ax[1].plot([0.2,0.25,0.3],[0.5,0.5,0.5], color = 'orange')
           ax[0].text(0.0975, 0.475, 'i', fontsize = 20)
           ax[0].text(0.2475,0.475,'j',fontsize = 20)
           ax[0].text(0.04,0.475,'i-1',fontsize = 20)
           ax[0].text(0.29,0.475,'j+1',fontsize = 20)
           ax[1].text(0.0975,0.475,'i',fontsize = 20)
           ax[1].text(0.2475,0.475,'j',fontsize = 20)
ax[1].text(0.04,0.475,'i-1',fontsize = 20)
           ax[1].text(0.29,0.475,'j+1',fontsize = 20)
           ax[0].set_title('First perturb rule', fontsize = 20)
           ax[1].set_title('Second perturb rule', fontsize = 20)
           ax[0].axis('off')
           ax[1].axis('off')
           plt.show()
```

