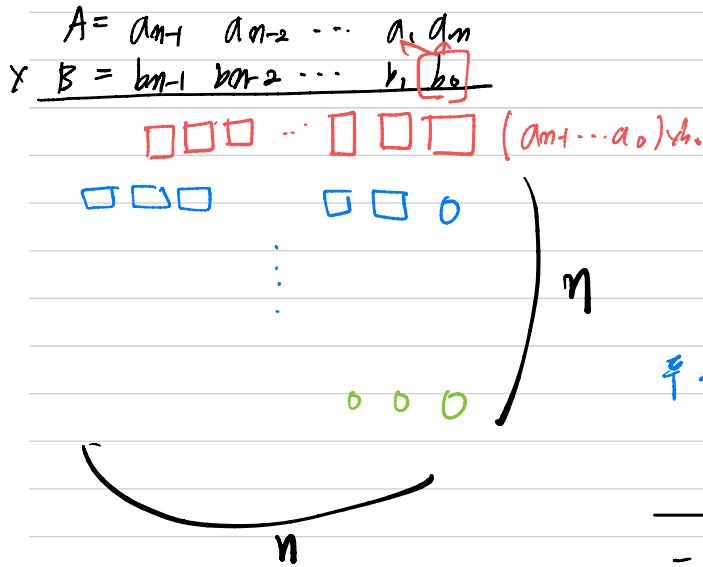


큰 수의 곱셈 : 분할 정복 알고리즘 - Karatsuba Algorithm

일반적인 $A \times B$ 의 연산 : $O(1)$ 시간 (RAM)
하지만, 무수히 큰 두 수의 곱은 $O(1)$ 이 불가능함.



$$\begin{aligned}
 &= n \times n = n^2 \text{의 기본 곱셈.} \\
 &\text{두 개의 자리수} = \text{자리수}(A \times B) = 2n \\
 &\quad \quad \quad n \quad m \text{ 자리수} \\
 &2n \times n = 2n^2 \text{의 기본 덧셈} \\
 &= 3n^2 \text{의 기본 덧셈과 기본 곱셈} \\
 &= O(n^2)
 \end{aligned}$$

- Karatsuba algorithm : u, v : n 의 자리의 두 수 (정수)

$u : a_{n-1} a_{n-2} \dots a_{\frac{n}{2}} a_{\frac{n}{2}+1} \dots a_1 a_0 = x \cdot 10^{\frac{n}{2}} + y$

address

$v : b_{n-1} b_{n-2} \dots b_{\frac{n}{2}} b_{\frac{n}{2}+1} \dots b_1 b_0 = w \cdot 10^{\frac{n}{2}} + z$

$$u \times v = (x \cdot 10^{\frac{n}{2}} + y)(w \cdot 10^{\frac{n}{2}} + z)$$

$u \times v = \underbrace{x \cdot w}_{\text{① } \frac{n}{2} \times \frac{n}{2}} \cdot 10^n + \underbrace{(x \cdot z + y \cdot w)}_{\text{② } \frac{n}{2} \times \frac{n}{2} + \frac{n}{2} \times \frac{n}{2}} \cdot 10^{\frac{n}{2}} + \underbrace{y \cdot z}_{\text{④ } \frac{n}{2} \times \frac{n}{2}}$

$$T(n) = 4T\left(\frac{n}{2}\right) + c \cdot n$$

$n \times n$ 자리수를
곱하기 위한
기본 덧셈

$$T(1) = c$$

$$n = 2^k = 4^k = (2^k)^2$$

$$T(n) = 4T\left(\frac{n}{2}\right) + cn$$

$$= 4^2 T\left(\frac{n}{2^2}\right) + 2cn + cn$$

$$= 4^k T\left(\frac{n}{2^k}\right) + cn(1 + 2 + \dots + 2^{k-1})$$

$$= O(4^k) = O(n^2)$$

$$U \times V = \underbrace{1 \cdot w}_{\substack{n \times n \\ \text{add}}} \cdot 10^n + \underbrace{(1z + yw)}_{\substack{n \times n \\ \text{add}}} \cdot 10^{\frac{n}{2}} + \underbrace{yz}_{\substack{n \times n \\ \text{add}}}$$

$n \text{ add} + n \text{ add} = O(n) \times 3n$

$$(1z + yw) = (1+y)(w+z) - 1w - yz$$

$$U \times V = \underbrace{1 \cdot w}_{\substack{n \times n \\ \text{add}}} \cdot 10^n + \underbrace{\left((1+y)(w+z) - 1w - yz \right)}_{\substack{\frac{n}{2} + \frac{n}{2} = n+1 \\ (2n)}} \cdot 10^{\frac{n}{2}} - \underbrace{yz}_{\substack{n \times n \\ \text{add}}}$$

$$n \text{ add} \times n \text{ add} = \left(\frac{n}{2} \text{ add} \times \frac{n}{2} \text{ add} \right) \times 2 \text{ add} + \left(\frac{n}{2} + 1 \right) \text{ add} \times \left(\frac{n}{2} + 1 \right) \text{ add} + CM$$

$1w \quad yz \quad (1+y)(z+w)$

$$T(n) = T\left(\frac{n}{2}\right) \times 2 + T\left(\frac{n}{2} + 1\right) + CM$$

$$= 3T\left(\frac{n}{2}\right) + CM = O(n^{\log_2 3}) = O(n^{1.7...})$$

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$$4T\left(\frac{n}{2}\right) + CM = O(n^2) = O(n^{\log_2 4})$$