Problem 1

(a) First, we compute P(A), P(B), $P(A \cap B)$

$$P(A) = 1 - p$$

$$P(B) = p\varepsilon_0 + (1 - p)\varepsilon_1$$

$$P(A \cap B) = (1 - p)\varepsilon_1$$

Next, we check when $P(A \cap B) = P(A) \cdot P(B)$

$$(1-p)\varepsilon_1 = (1-p)p\varepsilon_0 + (1-p)^2\varepsilon_1$$

$$\Rightarrow p(1-p)(\varepsilon_0 - \varepsilon_1) = 0$$

Therefore, A and B are independent if p = 0, p = 1, or $\varepsilon_0 = \varepsilon_1$.

(b) First, we compute P(A|C), P(B|C), $P(A \cap B|C)$

$$P(A|C) = \frac{P(A \cap C)}{P(C)} = \frac{(1-p)(\alpha_1 + \varepsilon_1)}{p(\alpha_0 + \varepsilon_0) + (1-p)(\alpha_1 + \varepsilon_1)}$$

$$P(B|C) = \frac{P(B \cap C)}{P(C)} = \frac{p\varepsilon_0 + (1-p)\varepsilon_1}{p(\alpha_0 + \varepsilon_0) + (1-p)(\alpha_1 + \varepsilon_1)}$$

$$P(A \cap B|C) = \frac{P((A \cap B) \cap C)}{P(C)} = \frac{(1-p)\varepsilon_1}{p(\alpha_0 + \varepsilon_0) + (1-p)(\alpha_1 + \varepsilon_1)}$$

Next, we check if $P(A \cap B|C) = P(A|C) \cdot P(B|C)$

$$P(A \cap B|C) = \frac{(1-p)\varepsilon_1}{p(\alpha_0 + \varepsilon_0) + (1-p)(\alpha_1 + \varepsilon_1)}$$

$$P(A|C) \cdot P(B|C) = \frac{(1-p)(\alpha_1 + \varepsilon_1)}{p(\alpha_0 + \varepsilon_0) + (1-p)(\alpha_1 + \varepsilon_1)} \cdot \frac{p\varepsilon_0 + (1-p)\varepsilon_1}{p(\alpha_0 + \varepsilon_0) + (1-p)(\alpha_1 + \varepsilon_1)}$$

We can simplify the check to

$$P((A \cap B) \cap C) \cdot P(C)$$

$$= (1 - p)\varepsilon_1 \cdot [p(\alpha_0 + \varepsilon_0) + (1 - p)(\alpha_1 + \varepsilon_1)]$$

$$= (1 - p)[p\varepsilon_1\alpha_0 + p\varepsilon_0\varepsilon_1 + (1 - p)\varepsilon_1\alpha_1 + (1 - p)\varepsilon_1^2]$$

$$P(A \cap C) \cdot P(B \cap C)$$

$$= (1 - p)[p\varepsilon_0\alpha_1 + p\varepsilon_0\varepsilon_1 + (1 - p)\varepsilon_1\alpha_1 + (1 - p)\varepsilon_1^2]$$

$$= (1 - p)(\alpha_1 + \varepsilon_1)[p\varepsilon_0 + (1 - p)\varepsilon_1]$$

Therefore, events A and B are conditionally independent given the event C if p=0, p=1, or $\varepsilon_0\alpha_1=\varepsilon_1\alpha_0$.

(c) By total probability theorem, we have

$$P(\text{Bit delivered correctly}) = p(a - \varepsilon_0 - \alpha_0) + (1 - p)(1 - \varepsilon_1 - \alpha_1)$$

Problem 2

(a) Let X be the number of "1"s transmitted in the time interval. Given Z = n, X follows a binomial distribution

$$X|(Z=n) \sim \text{Binomial}(n,p)$$

Then, the probability that X = k is

$$\begin{split} P(X=k) &= \sum_{n=k}^{\infty} P(X=k|Z=n) P(Z=n) \\ &= \sum_{n=k}^{\infty} \binom{n}{k} p^k (1-p)^{n-k} \frac{e^{-\lambda} \lambda^n}{n!} \\ &= \frac{e^{-\lambda} p^k \lambda^k}{k!} \sum_{n=k}^{\infty} \frac{(1-p)^{n-k} \lambda^{n-k}}{(n-k)!} \\ &= \frac{e^{-\lambda} p^k \lambda^k}{k!} \sum_{n=0}^{\infty} \frac{(1-p)^n \lambda^n}{n!} \\ &= \frac{e^{-\lambda} p^k \lambda^k}{k!} e^{\lambda (1-p)} \\ &= \frac{e^{-\lambda p} (\lambda p)^k}{k!} \end{split}$$

Therefore, $X \sim \text{Poisson}(\lambda p)$.

- (b) (i) This event occurs when all X_i are less than or equal to k. CDF of X is $F_X(k) = \prod_{i=1}^n P(X_i \le k) = \prod_{i=1}^n 1 P(X_i > k) = [1 (1-p)^k]^n$ PMF of X is $P(X = k) = F_X(k) F_X(k-1) = [1 (1-p)^k]^n [1 (1-p)^{k-1}]^n$
 - (ii) This event occurs when at least one X_i is less than or equal to k. CDF of Y is $F_Y(k) = 1 P(\min(X_1, \dots, X_n) > k) = 1 [(1-p)^k]^n = 1 (1-p)^{nk}$ PMF of Y is $P(Y = k) = F_Y(k) F_Y(k-1) = (1-p)^{n(k-1)} (1-p)^{nk}$
 - (iii) $P(Y=k)=(1-p)^{n(k-1)}-(1-p)^{nk}=(1-p)^{n(k-1)}[1-(1-p)^n]$ Let $p'=1-(1-p)^n$, then $P(Y=k)=(1-p')^{(k-1)}p'$, this is the PMF of a Geometric random variable with success probability p'.
- (c) According to the Binomial distribution, we have

$$P(X_{S_2} = k) = {123 \choose k} p_T^k (1 - p_T)^{123 - k}$$
 for $k = 0, 1, \dots, 123$

Problem 3

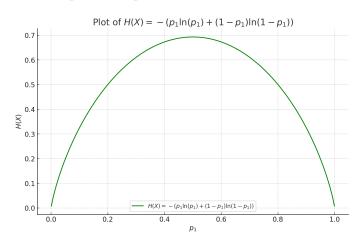
(a)

$$\begin{split} H(\frac{1}{2},\frac{1}{3},\frac{1}{6}) &= -(\frac{1}{2}\ln\frac{1}{2} + \frac{1}{3}\ln\frac{1}{3} + \frac{1}{6}\ln\frac{1}{6}) \\ &= (\frac{1}{2} + \frac{1}{6})\ln2 + (\frac{1}{3} + \frac{1}{6}\ln3) \\ &= \frac{2}{3}\ln2 + \frac{1}{2}\ln3 \\ H(\frac{1}{2},\frac{1}{2}) + \frac{1}{2}\cdot H(\frac{2}{3},\frac{1}{3}) = -(\frac{1}{2}\ln\frac{1}{2} + \frac{1}{2}\ln\frac{1}{2}) - \frac{1}{2}\cdot (\frac{2}{3}\ln\frac{2}{3} + \frac{1}{3}\ln\frac{1}{3}) \\ &= \ln2 + \frac{1}{2}(\frac{2}{3}(\ln3 - \ln2) + \frac{1}{3}\ln3) \\ &= \frac{2}{3}\ln2 + \frac{1}{2}\ln3 \end{split}$$

Therefore, $H(\frac{1}{2}, \frac{1}{3}, \frac{1}{6}) = H(\frac{1}{2}, \frac{1}{2}) + \frac{1}{2} \cdot H(\frac{2}{3}, \frac{1}{3}).$

(b) Maximum entropy: $H_{max} = \ln 2$ at $p_1 = 0.5$

Minimum entropy: $H_{min} = 0$ at $p_1 = 1$ or $p_1 = 0$



(c) Let $w_i = p_i$ and $x_i = \frac{1}{p_i}$.

$$\frac{p_1 \frac{1}{p_1} + \dots + p_n \frac{1}{p_n}}{1} = n \ge \prod_{i=1}^n (\frac{1}{p_i})^{p_i}$$

$$\Rightarrow \ln n \ge \ln(\prod_{i=1}^n (\frac{1}{p_i})^{p_i}) = -\sum_{i=1}^n p_i \ln p_i = H(X)$$

The equality holds when $p_1 = \cdots = p_n = \frac{1}{n}$. Therefore, $H(X)_{max} = \ln n$ when $\{p_i = \frac{1}{n}\}_{i=1}^n$.

(d) The entropy is minimized when the random variable X is deterministic. For a specific $i \in \{1, 2, \dots, n\}$

$$H(X) = -(1 \cdot \ln 1 + \sum_{j \neq i}^{n} 0 \cdot \ln 0) = 0$$

Therefore, $H(X)_{\min}=0$ is achieved by any PMF where exactly one $i\in\{1,2,\ldots,n\},\ p_i=1$ and $p_j=0$ for all $j\neq i$.

Problem 4

(a) (i)

$$E[X] = \sum_{k=1}^{\infty} k \cdot (1-p)^{k-1} p = p \sum_{k=1}^{\infty} k \cdot (1-p)^{k-1} = p \cdot \frac{1}{(1-(1-p))^2} = \frac{1}{p}$$

(ii)

$$E[e^{tX}] = \sum_{k=1}^{\infty} e^{tk} \cdot (1-p)^{k-1}p$$

$$= p \sum_{k=1}^{\infty} e^{tk} \cdot (1-p)^{k-1}$$

$$= pe^{t} \sum_{k=1}^{\infty} [e^{t}(1-p)]^{k-1}$$

$$= pe^{t} \cdot \frac{1}{1 - (1-p)e^{t}} = \frac{pe^{t}}{1 - (1-p)e^{t}}$$

Additionally, we should verify the convergence condition

$$|e^t(1-p)| < 1 \Rightarrow e^t < \frac{1}{1-p} \Rightarrow t < \ln \frac{1}{1-p} = -\ln(1-p)$$

(b) (i) To find $Var[Z] = E[Z^2] - (E[Z])^2$, we first compute E[Z] and $E[Z^2]$

$$E[Z^2] = \sum_{n=1}^{\infty} z_n^2 \cdot p_Z(z_n) = \sum_{n=1}^{\infty} n \cdot \frac{6}{(\pi n)^2} = \frac{6}{\pi^2} \sum_{n=1}^{\infty} \frac{1}{n} = \infty$$

Therefore, since $E[Z^2] = \infty$, the variance Var[Z] does not exist.

(ii)

$$\sum_{n=1}^{\infty} z_n^3 \cdot p_Z(z_n) = \sum_{n=1}^{\infty} (-1)^n (\sqrt{n})^3 \cdot \frac{6}{(\pi n)^2} = \frac{6}{\pi^2} \sum_{n=1}^{\infty} (-1)^n \frac{1}{\sqrt{n}} = -\frac{6}{\pi^2} \eta(\frac{1}{2}) \approx -\frac{6}{\pi^2} \cdot 0.6049$$

(iii)

$$E[Z^3] = \sum_{n=1}^{\infty} z_n^3 \cdot p_Z(z_n) = \sum_{n=1}^{\infty} n^{\frac{3}{2}} \cdot \frac{6}{(\pi n)^2} = \frac{6}{\pi^2} \sum_{n=1}^{\infty} \frac{1}{\sqrt{n}} = \infty$$

Therefore, ${\cal E}[Z^3]$ does not exist.

Problem 5

- (a) Partition: [0.7, 0.3] / Prior: Uniform
 - Test accuracy: 98.21%
 - Observations:
 - The classifier performs exceptionally well for the ham class, with a precision and recall close to 100%.
 - The spam detection is slightly weaker, with a recall of 89%, indicating that some spam messages were missed.
- (b) Partition: [0.7, 0.3] / Prior [0.5, 0.5]
 - Test accuracy: 97.67%
 - Observations:
 - * This prior slightly lowered the overall accuracy compared to the uniform prior.
 - * The spam detection was more balanced, achieving equal precision and recall, which implies the model's confidence in detecting both classes became more stable.
 - Partition: [0.7, 0.3] / Prior [0.9, 0.1]
 - Test accuracy: 98.27%
 - Observations:
 - * Setting a heavier prior towards ham improved the model's confidence in predicting spam but didn't significantly change the recall.
 - * The accuracy remains similar to the uniform prior but with slightly better precision for spam.
 - Partition: [0.8, 0.2] / Prior: Uniform
 - Test accuracy: 98.39%
 - Observations:
 - * Increasing the training data led to a slight improvement in accuracy.
 - * Spam detection precision increased to 99%, showing that more training data helps the model learn the nuances better.
 - Partition: [0.6, 0.4] / Prior: Uniform
 - Test accuracy: 98.34%
 - Observations:
 - * A larger test set still resulted in high accuracy.
 - * Spam detection improved slightly, demonstrating the model's robustness even with less training data.