Computer-Aided VLSI System Design

Homework 2: Simple MIPS CPU

Graduate Institute of Electronics Engineering, National Taiwan University



Goal

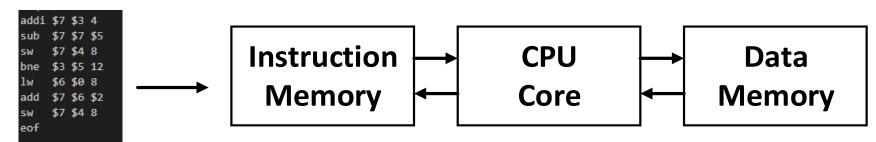
- In this homework, you will learn
 - How to write testbench
 - How to design FSM
 - How to use IP
 - Generate patterns for testing

Introduction



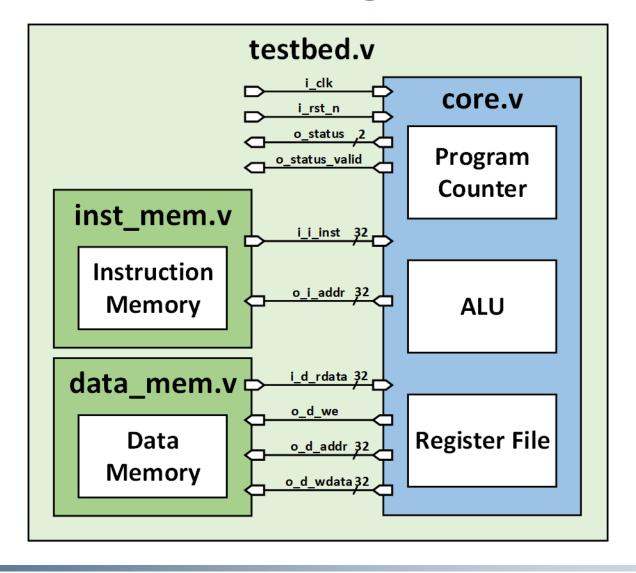
Central Processing Unit (CPU) is the important core in the computer system. In this homework, you are asked to design a simple MIPS CPU, which contains the basic module of program counter, ALU and register files. The instruction set of the simple CPU is similar to MIPS structure.

Instruction set



Block Diagram





Input/Output



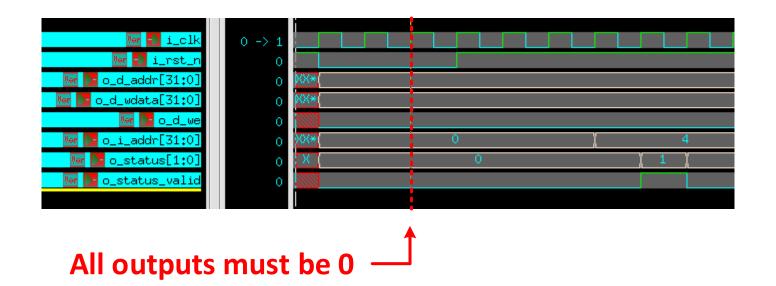
Signal Name	I/O	Width	Simple Description	
i_clk	I	1	Clock signal in the system.	
i_rst_n	I	1	Active low asynchronous reset.	
o_i_addr	0	32	Address from program counter (PC)	
i_i_inst	I	32	Instruction from instruction memory	
o_d_we	0	1	Write enable of data memory Set low for reading mode, and high for writing mode	
o_d_addr	0	32	Address for data memory	
o_d_wdata	0	32	Unsigned data input to data memory	
i_d_rdata	I	32	Unsigned data output from data memory	
o_status	0	2	Status of core processing to each instruction	
o_status_valid	0	1	Set high if ready to output status	



- All outputs should be synchronized at clock rising edge.
- Instruction memory and data memory are provided. All values in memory are reset to be zero.
- You should create 32 unsigned 32-bit registers in register file.
- Less than 1024 instructions are provided for each pattern.
- The whole processing time can't exceed 120000 cycles.



 You should set all your outputs and register file to be zero when i_rst_n is low. Active low asynchronous reset is used.



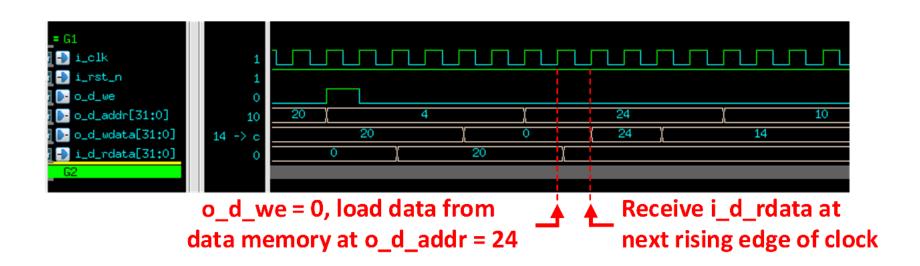


 After outputting o_i_addr to instruction memory, the core can receive the corresponding i_i_inst at the next rising edge of the clock.



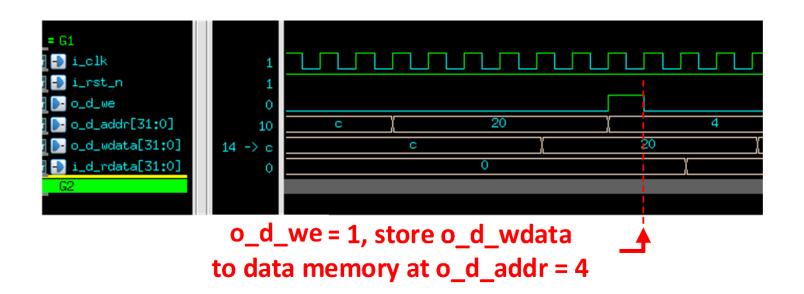


To load data from the data memory, set o_d_we to 0 and o_d_addr to relative address value. i_d_rdata can be received at the next rising edge of the clock.



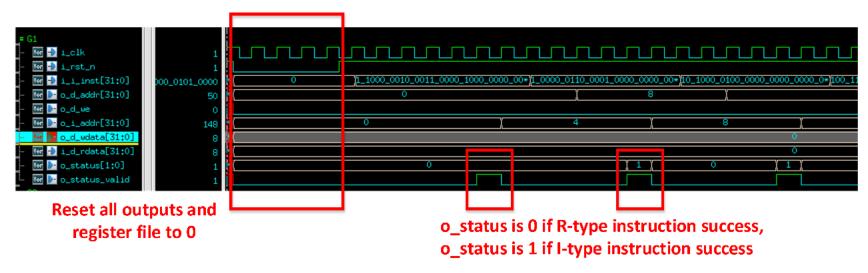


 To save data to the data memory, set o_d_we to 1, o_d_addr to relative address value, and o_d_wdata to the written data.



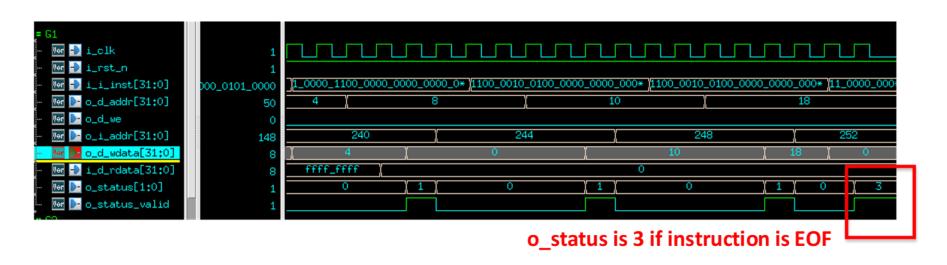


- Your o_status_valid should be turned to high for only one cycle for every o_status.
- The testbench will get your output at negative clock edge to check the o_status if your o_status_valid is high.



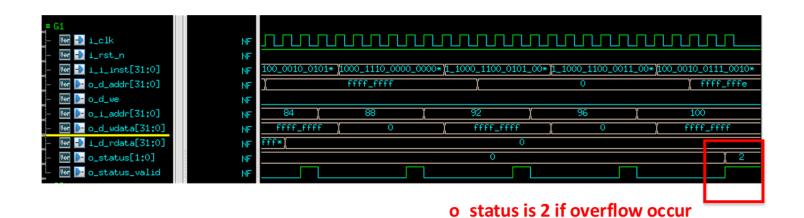


 When you set o_status_valid to high and o_status to 3, stop processing. The testbench will check your data memory value with golden data.





 If overflow happened, stop processing and raise o_status_valid to high and set o_status to 2. The testbench will check your data memory value with golden data.

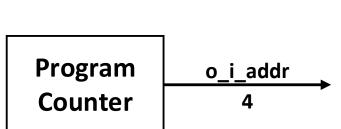


Program Counter



 Program counter is used to control the address of instruction memory.

\$pc = \$pc + 4 for every instruction (except beq, bne)



Instruction			
	Memory		
Addr.	Instruction		
0	addi \$1 \$0 20		
4	addi \$2 \$0 12		
•	:		

Instruction mapping



R-type

[31:26]	[25:21]	[20:16]	[15:11]	[10:0]
opcode	\$s2	\$s3	\$s1	Not used

31

I-type

	[31:26]	[25:21]	[20:16]	[15:0]
	opcode	\$s2	\$s1	im
31	_			Ó

EOF

	[31:26]	[25:0]
	opcode	Not used
31		0

Instruction



Operation	Assemble	Opcode	Туре	Meaning	Note
Add (integer)	add	6'd1	R	\$s1 = \$s2 + \$s3	Signed Operation
Subtract (integer)	sub	6'd2	R	\$s1 = \$s2 - \$s3	Signed Operation
Multiply (integer)	mul	6'd3	R	\$s1 = \$s2 * \$s3	Signed Operation
Add (floating point)	fp_add	6'd13	R	\$s1 = \$s2 + \$s3	Floating Point Operation
Subtract (floating point)	fp_sub	6'd14	R	\$s1 = \$s2 - \$s3	Floating Point Operation
Multiply (floating point)	fp_mul	6'd15	R	\$s1 = \$s2 * \$s3	Floating Point Operation
Add immediate	addi	6'd4	I	\$s1 = \$s2 + im	Signed Operation
Load word	lw	6'd5	I	\$s1 = Mem[\$s2 + im]	Signed Operation
Store word	SW	6'd6	I	Mem[\$s2 + im] = \$s1	Signed Operation

Instruction (cont'd)



Operation	Assemble	Opcode	Туре	Meaning	Note
AND	and	6'd7	R	\$s1 = \$s2 & \$s3	Bit-wise
OR	or	6'd8	R	\$s1 = \$s2 \$s3	Bit-wise
NOR	nor	6'd9	R	\$s1 = ~(\$s2 \$s3)	Bit-wise
Branch on equal	beq	6'd10	I	if(\$s1==\$s2), \$pc = \$pc + im; else, \$pc = \$pc + 4	PC-relative Unsigned Operation
Branch on not equal	bne	6'd11	ı	if(\$s1!=\$s2), \$pc = \$pc + im; else, \$pc = \$pc + 4	PC-relative Unsigned Operation
Set on less than	slt	6'd12	R	if(\$s2<\$s3), \$s1 = 1; else, \$s1 = 0	Signed Operation
Shift left logical	sll	6'd16	R	\$s1 = \$s2 << \$s3	Unsigned Operation
Shift right logical	srl	6'd17	R	\$s1 = \$s2 >> \$s3	Unsigned Operation
End of File	eof	6'd18	EOF	Stop processing	Last instruction in the pattern

Note: The notation of **im** in I-type instruction is **2's complement.**

Note: Signed operations indicates that the data in register file are expressed in 2's complement.

Floating Point



- For instructions fp_add, fp_sub, fp_mul, you will have to implement operations with floating point format
- Only normal numbers will be provided in test patterns
 - You do not need to consider infinity, denormalized numbers, and NaN
- IEEE 754 single precision format [1]
 - Round to nearest even [2]

[31]	[30:23]	[22:0]
sign	exponent	mantissa

31

Floating Point



Rounding

1.BBGRXXX

Guard bit: LSB of result

Sticky bit: OR of remaining bits

Round bit: 1st bit removed

Round up conditions

Round = 1, Sticky = 1 → > 0.5

Guard = 1, Round = 1, Sticky = 0 → Round to even

Value	Fraction	GRS	Incr?	Rounded
128	1.0000000	000	N	1.000
15	1.1010000	100	N	1.101
17	1.0001000	010	N	1.000
19	1.0011000	110	Y	1.010
138	1.0001010	011	Y	1.001
63	1.1111100	111	Y	10.000

Memory IP



- Instruction memory
 - Size: 1024 × 32 bit
 - i_addr[11:2] for address mapping in instruction memory
- Data memory
 - Size: 64×32 bit
 - i addr[7:2] for address mapping in data memory

Status



4 statuses of o_status

o_status[1:0]	Definition
2'd0	R_TYPE_SUCCESS
2'd1	I_TYPE_SUCCESS
2'd2	MIPS_OVERFLOW
2'd3	MIPS_END

Overflow



- Overflow may be happened.
 - Situation1: Overflow happened at arithmetic instructions (add, sub, mul, addi)
 - Situation2: If output address are mapped to unknown address in data/instruction memory. (Do not consider the case if instruction address is beyond eof, but the address mapping is in the size of instruction memory)

rtl.f



Filelist

```
// Simulation: HW2 simple mips CPU
// -----
// define files
// -----
../00 TESTBED/define.v
// testbench
 ______
../00_TESTBED/testbed.v
../00 TESTBED/inst mem.vp
../00_TESTBED/data_mem.vp
// design files
./core.v
```

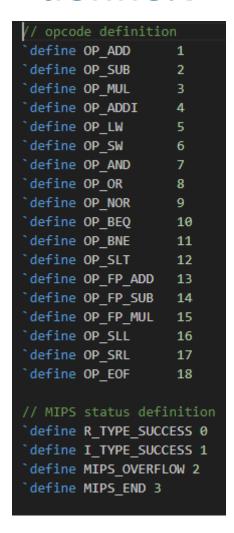
core.v



Do not modify interface

```
nodule core #( // DO NOT MODIFY!!!
   parameter ADDR_WIDTH = 32,
   parameter INST WIDTH = 32,
   parameter DATA WIDTH = 32
   input
                             i clk,
   input
                             i rst n,
   output [ADDR_WIDTH-1:0] o_i_addr,
   input
           [ADDR_WIDTH-1:0] i_i_inst,
   output
                             o_d_we,
           [ADDR_WIDTH-1:0] o_d_addr,
   output
   output [DATA_WIDTH-1:0] o_d_wdata,
   input
           [DATA_WIDTH-1:0] i_d_rdata,
   output
                       1:0] o_status,
                             o_status_valid
   output
endmodule
```

define.v





testbed_temp.v

- Things to add in your testbench
 - Clock
 - Reset
 - Waveform file
 - Function test

— ...

```
core u core (
    .i_clk(),
    .i_rst_n(),
    .o i addr(),
    .i_i_inst(),
    .o d we(),
    .o d addr(),
    .o d wdata(),
    .i d rdata(),
    .o status(),
    .o status valid()
    .i clk(),
    .i rst n(),
    .i addr(),
    .o_inst()
);
data_mem u_data_mem (
    .i_clk(),
    .i_rst_n(),
    .i_we(),
    .i_addr(),
    .i wdata(),
    .o_rdata()
```

Protected Files

- The following files are protected
 - inst_mem.vp
 - data_mem.vp

```
module inst mem (
    input
                      i_clk,
    input
                      i rst n,
    input [ 31 : 0 ] i addr,
    output [ 31 : 0 ] o inst
 protected
Ndi5kSQH5DT^<D9i:i7T7ceFn3@o:C2]Ke:L;dfq^QGQOG?3K:ogIe8]1ge<gcg3
1CH3E]ekmLN<RVkKa1o39E7E21a; hJRSFMUb2pAgL?TeZdH>]^RK;KWYU@>G2G6
H[IMYG;D<[Z>[;0]] NbPoEAQM<_ZfDbp1HN@HmqSOQ<5[53C:9UD4^:Y44]9a^e
PDH[cdHb;HPi\R4k7mAlPdY8ZpI=4?nNZgQ2I>QUg[agM4j@cTl]hnMoC<i1F9DR
[kf;]ULlecpF`H;9L2DeZa>@LdfLgfB8l4bWgT:_P3?ENhifQW@_Ne;gMZE9@f0A
OERY:F4d68KqAIn]N1dj4LN7 8:Uigk?9UJ9JYQM4l=Lq\TEXDQO1>Zo^SJq=Cge
?kp68am:9p81Q1[<jSXm?;GhoPHHYKp\Q][2epXn_18k8LA5g=N7=D?=VOX<Ham8
[A:Qc;RlpO38>d9 Qk9cfk?:5hXP>LT3n=DP08A ]WPa6nA3cYZjGl32qB9]I4kp
>=:4m9P`dCB8@?ip`@VR7AahIggjNR:M1: \KXElBFOm<Bb@ZS[^W7EheJ18mX8;
?7F`Pg\CCA8igfFUoWY@k>Yq=U3 4>E50 nJ\`aUGcfWD 89dab]cUQfF<?2P?OG
qWglWC[\iqnjC<OipHHnb<T4Sg<:UORVSVocI g?<a@o <PQ493cZIE;7^Sp1AQ
G<cl7[]R\>VT]]LA\7?Uk=]\bG19MT9N;K<Y92[iKOged92EIkQZliW>qlG]QI?5
ST06RFN<KJl@VM1EWKSmB1B5U:BaX`E7of7mqOJBgO`9k$
 endprotected
endmodule
```

Command



- 01_run
 - Usage: ./01_run p0

```
vcs -f rtl.f -full64 -R -debug_access+all
+define+$1 +v2k
```

99_clean_up

```
rm -rf *.history *.key *.log
rm -rf novas.rc novas.fsdb novas.conf
rm -rf INCA_libs nWaveLog BSSLib.lib++
```

PATTERN

Files in PATTERN are for your references

inst_assemble.dat

R-type	\$52	\$ s3	\$ 51
I-type	\$52	\$ s1	im
and	\$1	\$3	\$1
lw	\$ 3	\$1	8
bne	\$2	\$0	8
add	\$7	\$1	\$4
slt	\$6	\$5	\$4
slt	\$4	\$1	\$1
lw	\$1	\$3	12
lw	\$7	\$7	4
bne	\$6	\$7	8
lw	\$6	\$5	8
lw	\$ 5	\$2	8

Grading Policy



TA will run your code with following command

- Pass the patterns to get full score
 - Provided pattern: 70%
 - 30% for each test (data from data memory: 15%, status check: 15%)
 - **10%** for spyglass check
 - Hidden pattern: 30% (20 patterns in total)
 - 1.5% for each test (data & status both correct)
- No delay submission is allowed
- Lose 5 point for any wrong naming rule or format for submission

Submission



 Create a folder named studentID_hw2, and put all below files into the folder

```
r11943133_hw2/

-- 01_RTL
-- core.v
-- rtl.f
-- (other design files)
```

- Compress the folder studentID_hw2 in a tar file named studentID_hw2_vk.tar (k is the number of version, k =1,2,...)
 - Use lower case for student ID. (Ex. r11943133_hw2_v1.tar)
- Submit to NTU Cool

Hint

- Design your FSM with following states
 - 1. Idle
 - 2. Instruction Fetching
 - 3. Instruction decoding
 - 4. ALU computing/Load data
 - 5. Data write-back
 - 6. Next PC generation
 - 7. Process end

Reference



- [1] IEEE 754 Single Precision Format
 - https://zh.wikipedia.org/zh-tw/IEEE 754
- [2] Round to Nearest Even
 - https://www.cs.cmu.edu/afs/cs/academic/class/15213s16/www/lectures/04-float.pdf