# Computer-Aided VLSI System Design Homework 2: Simple MIPS CPU

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### **Data Preparation**

1. Decompress 1112 hw2.tar with following command

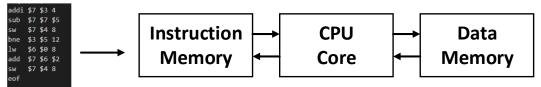
tar -xvf 1112\_hw2.tar

Folder	File	Description		
	inst_mem.vp	Module of instruction memory (protected)		
00 TESTDED	data_mem.vp	Module of data memory (protected)		
00_TESTBED	define.v	File of definition		
	testbed_temp.v	Testbench template		
	inst.dat	Pattern of instruction in binary format		
00 TESTDED/	inst assemble.dat	Corresponding assembly code of the		
00_TESTBED/ PATTERN/p*	mst_assemble.dat	instruction pattern		
PALLEKIN/p	data.dat	Pattern of final data in data memory		
	status.dat	Pattern of corresponding status		
	core.v	Your design		
	rtl.f	File list		
01_RTL	flist.v	File list (for spyglass check)		
	lint.tcl	Spyglass script		
	01_run	VCS command		
	02_lint	Spyglass command		
	99_clean	Command to clean temporary data		

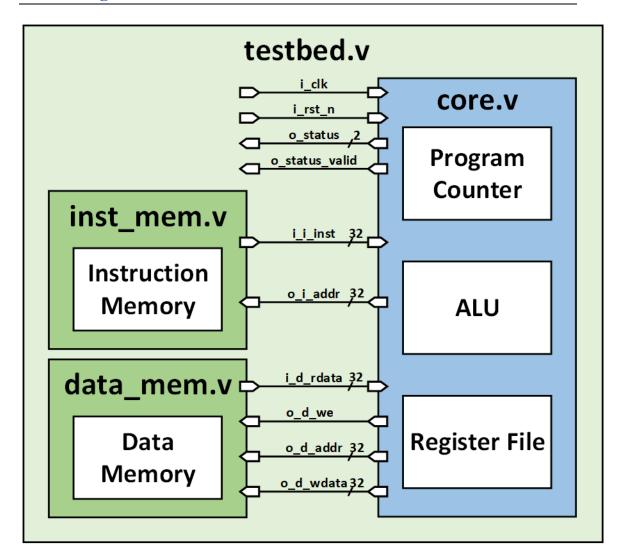
#### Introduction

Central Processing Unit (CPU) is the important core in the computer system. In this homework, you are asked to design a simple MIPS CPU, which contains the basic module of program counter, ALU and register files. The instruction set of the simple CPU is similar to MIPS structure. Since the files of testbench (testbed.v, inst\_mem.v, data mem.v) are protected, you also need to design the testbench to test your design.

#### Instruction set



### **Block Diagram**



# **Specifications**

- 1. Top module name: core
- 2. Input/output description:

Signal Name	I/O	Width	Simple Description	
i_clk	I	1	Clock signal in the system.	
i_rst_n	Ι	1	Active <b>low</b> asynchronous reset.	
o_i_addr	О	32	Address from program counter (PC)	

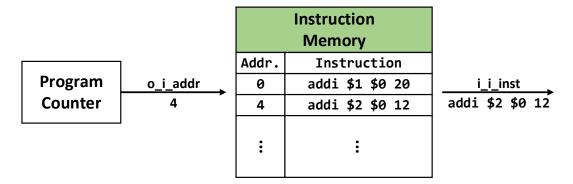
i_i_inst	I	32	Instruction from instruction memory
1 0		1	Write enable of data memory
o_d_we	0	1	Set <b>low</b> for reading mode, and <b>high</b> for writing mode
o_d_addr	О	32	Address for data memory
o_d_wdata	О	32	Unsigned data input to data memory
i_d_rdata	I	32	Unsigned data output from data memory
o_status	О	2	Status of core processing to each instruction
o_status_valid	О	1	Set <b>high</b> if ready to output status

- 3. All outputs should be synchronized at clock **rising** edge.
- 4. You should set all your outputs and register file to be zero when i\_rst\_n is **low**. Active low asynchronous reset is used.
- 5. Instruction memory and data memory are provided. All values in memory are reset to be zero.
- 6. You should create 32 unsigned 32-bit registers in register file.
- 7. After outputting o\_i\_addr to instruction memory, the core can receive the corresponding i\_i\_inst at the next rising edge of the clock.
- 8. To load data from the data memory, set o\_d\_we to **0** and o\_d\_addr to relative address value. i d rdata can be received at the next rising edge of the clock.
- 9. To save data to the data memory, set o\_d\_we to 1, o\_d\_addr to relative address value, and o d wdata to the written data.
- 10. Your o status valid should be turned to **high** for only **one cycle** for every o status.
- 11. The testbench will get your output at negative clock edge to check the o\_status if your o status valid is **high**.
- 12. When you set o\_status\_valid to **high** and o\_status to **3**, stop processing. The testbench will check your data memory value with golden data.
- 13. If overflow happened, stop processing and raise o\_status\_valid to **high** and set o status to **2**. The testbench will check your data memory value with golden data.
- 14. Less than 1024 instructions are provided for each pattern.
- 15. The whole processing time can't exceed 120000 cycles.

#### **Design Description**

1. Program counter is used to control the address of instruction memory.

**\$pc** = **\$pc** + **4** for every instruction (except **beq**, **bne**)



- 2. Register file contains 32 unsigned 32-bit registers for operation.
- 3. Instruction mapping
  - a. R-type

	[31:26]	[25:21]	[20:16]	[15:11]	[10:0]	
	opcode	\$s2	\$s3	\$s1	Not used	
32	L					0

### b. I-type

	[31:26]	[25:21]	[20:16]	[15:0]	
	opcode	\$s2	\$s1	im	
31	L	-	-	0	)

#### c. EOF

	[31:26]	[25:0]
	opcode	Not used
31		O

4. The followings are the instructions you need to design for this homework:

Operation	Assemble	Opcode	Type	Meaning	Note
Add	add	6'd1	R	\$s1 = \$s2 + \$s3	Signed Operation
Subtract	sub	6'd2	R	\$s1 = \$s2 - \$s3	Signed Operation
Multiply	mul	6'd3	R	\$s1 = \$s2 * \$s3	Signed Operation
Add	fn add	6'd13	R	\$s1 = \$s2 + \$s3	Floating Point
(floating point)	fp_add	0 013	K	\$51 = \$52 + \$55	Operation
Subtract	£n cub	6'd14	R	\$s1 = \$s2 - \$s3	Floating Point
(floating point)	fp_sub	p_sub   6 d14	K	\$51 = \$52 - \$53	Operation
Multiply	fp_mul	(2.11.5	D	d-4	Floating Point
(floating point)		6'd15	R	\$s1 = \$s2 * \$s3	Operation

Add immediate	addi	6'd4	I	\$s1 = \$s2 + im	Signed Operation
Load word	1w	6'd5	I	\$s1 = Mem[\$s2 + im]	Signed Operation
Store word	SW	6'd6	I	Mem[\$s2 + im] = \$s1	Signed Operation
AND	and	6'd7	R	\$s1 = \$s2 & \$s3	Bit-wise
OR	or	6'd8	R	\$s1 = \$s2   \$s3	Bit-wise
NOR	nor	6'd9	R	\$s1 = ~(\$s2   \$s3)	Bit-wise
Branch on equal	beq	6'd10	I	<pre>if(\$s1==\$s2), \$pc = \$pc +</pre>	PC-relative Unsigned Operation
Branch on not equal	bne	6'd11	Ι	<pre>if(\$s1!=\$s2), \$pc = \$pc +</pre>	PC-relative Unsigned Operation
Set on less than	slt	6'd12	R	if(\$s2<\$s3), \$s1 = 1; else, \$s1 = 0	Signed Operation
Shift left logical	sll	6'd16	R	\$s1 = \$s2 << \$s3	Unsigned Operation
Shift right logical	srl	6'd17	R	\$s1 = \$s2 >> \$s3	Unsigned Operation
End of File	eof	6'd18	EOF	Stop processing	Last instruction in the pattern

Note: The notation of im in I-type instruction is 2's complement.

Note: Signed operations indicates that the data in register file are expressed in 2's complement.

- 5. Interface of instruction memory (size: 1024×32 bit)
  - i addr[11:2] for address mapping in instruction memory

- 6. Interface of data memory (size: 64×32 bit)
  - i addr[7:2] for address mapping in data memory
  - To fetch data of data memory in your testbench, use following instance name

```
u_data_mem.mem_r[i]
```

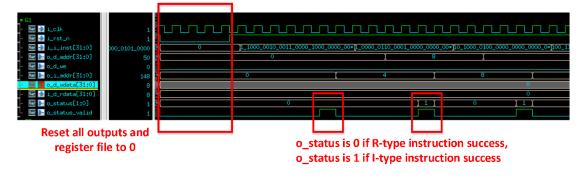
- 7. Overflow may be happened.
  - <u>Situation1</u>: Overflow happened at arithmetic instructions (add, sub, addi)
  - <u>Situation2</u>: If output address is mapped to unknown address in data/instruction memory. (Do not consider the case if instruction address is beyond eof, but the address mapping is in the size of instruction memory)
- 8. 4 statuses of o status

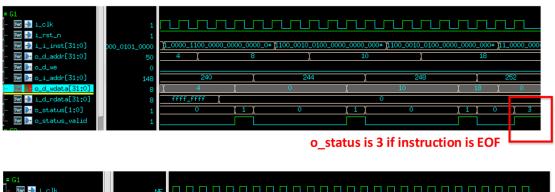
o_status[1:0]	Definition
2'd0	R_TYPE_SUCCESS
2'd1	I_TYPE_SUCCESS
2'd2	MIPS_OVERFLOW
2'd3	MIPS_END

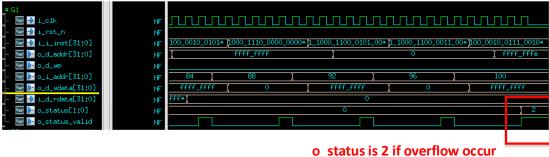
- 9. Last instruction would be **eof** for every pattern.
- 10. There is no unknown opcode in the pattern.

# **Sample Waveform**

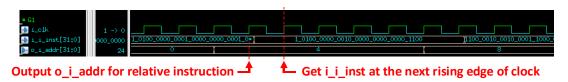
1. Status check



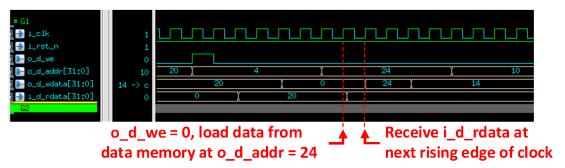




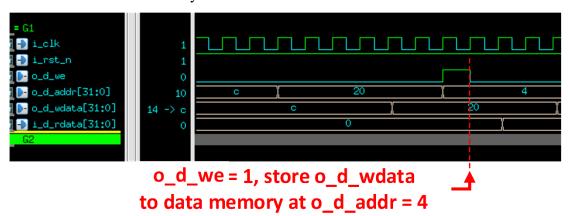
2. Read instruction from instruction memory



3. Load data from data memory



4. Save data to data memory



### **Testbed**

- 1. Things to add in your testbench
  - Clock
  - Reset
  - Waveform file (.fsdb)
  - Function test
  - ...

#### **Submission**

- 2. Create a folder named **studentID\_hw2**, and put all below files into the folder
  - **rtl.f** (your file list)
  - core.v
  - all other design files in your file list (optional)

Note: Use lower case for the letter in your student ID. (Ex. r11943133 hw2)

3. Compress the folder studentID\_hw2 in a tar file named studentID\_hw2\_vk.tar (k is the number of version, k=1,2,...)

```
tar -cvf studentID_hw2_vk.tar studentID_hw2
```

TA will only check the last version of your homework.

Note: Use **lower case** for the letter in your student ID. (Ex. r11943133\_hw2\_v1)

4. Submit to NTU Cool

### **Grading Policy**

1. TA will run your code with following format of command. Make sure to run this command with no error message.

```
vcs -f rtl.f -full64 -R -debug_access+all +define+p0 +v2k
```

- 2. Pass the patterns to get full score.
  - Provided pattern: **70%** (patterns: p0, p1)
    - 30% for each pattern (data in data memory: 15%, status check: 15%)
    - 10% for spyglass check
    - Don't implement the answers in your design directly!
  - Hidden pattern: **30%** (20 patterns in total)
    - 1.5% for each pattern (data & status both correct)
- 3. Delay submission
  - No delay submission is allowed
  - Lose **5 point** for any wrong naming rule. Don't compress all homework folder.

## Hint

- 1. Design your FSM with following states
  - Idle
  - Instruction Fetching
  - Instruction decoding
  - ALU computing/ Load data
  - Data write-back
  - Next PC generation
  - Process end