

Balanced Search Tree

AVL Tree

Outline

- 2-3 tree

- 2-3-4 tree

- **AVL tree**

 - [Adelson-Velskii & Landis, 1962]

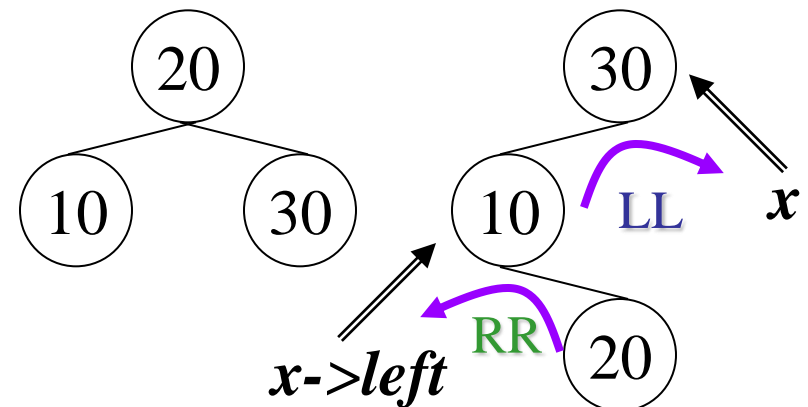
- **Red-black tree**

 - [Rudolf Bayer, 1972]... **B-tree**

AVL Tree: *Double Rotations*

- Let x be the node at which $x \rightarrow \text{left}$ and $x \rightarrow \text{right}$ differ by more than 1; Assume that the height of x is 3
 - Height of $x \rightarrow \text{left}$ is 2 (i.e. height of $x \rightarrow \text{right}$ is 0)
 - 3. Height of $x \rightarrow \text{left} \rightarrow \text{right}$: 1 \Rightarrow double rotation with the left child (LR)
 - $\text{BF}(x) = +2$
 - $\text{BF}(x \rightarrow \text{left}) = -1$

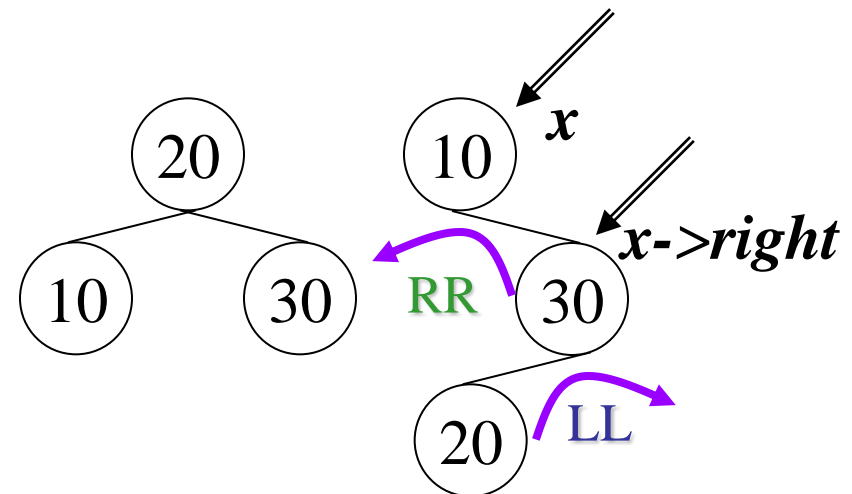
Tree height: shorter



AVL Tree: *Double Rotations*

- Let x be the node at which $x \rightarrow \text{left}$ and $x \rightarrow \text{right}$ differ by more than 1; Assume that the height of x is 3
 - Height of $x \rightarrow \text{right}$ is 2 (i.e. height of $x \rightarrow \text{left}$ is 0)
 - 4. Height of $x \rightarrow \text{right} \rightarrow \text{left}$: 1 \Rightarrow double rotation with the right child (RL)
 - $\text{BF}(x) = -2$
 - $\text{BF}(x \rightarrow \text{right}) = +1$

Tree height: shorter

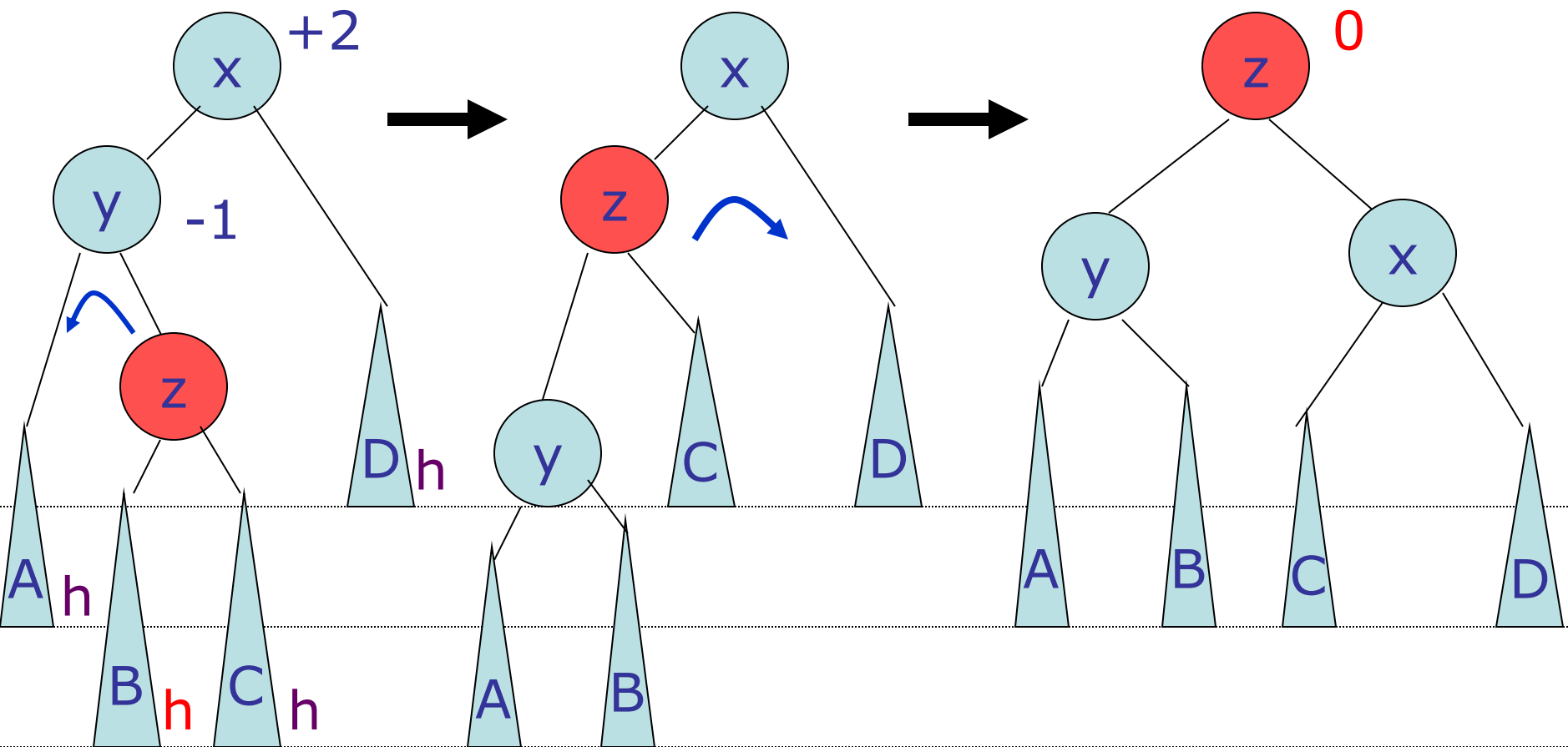


AVL Tree: *Double Rotations*

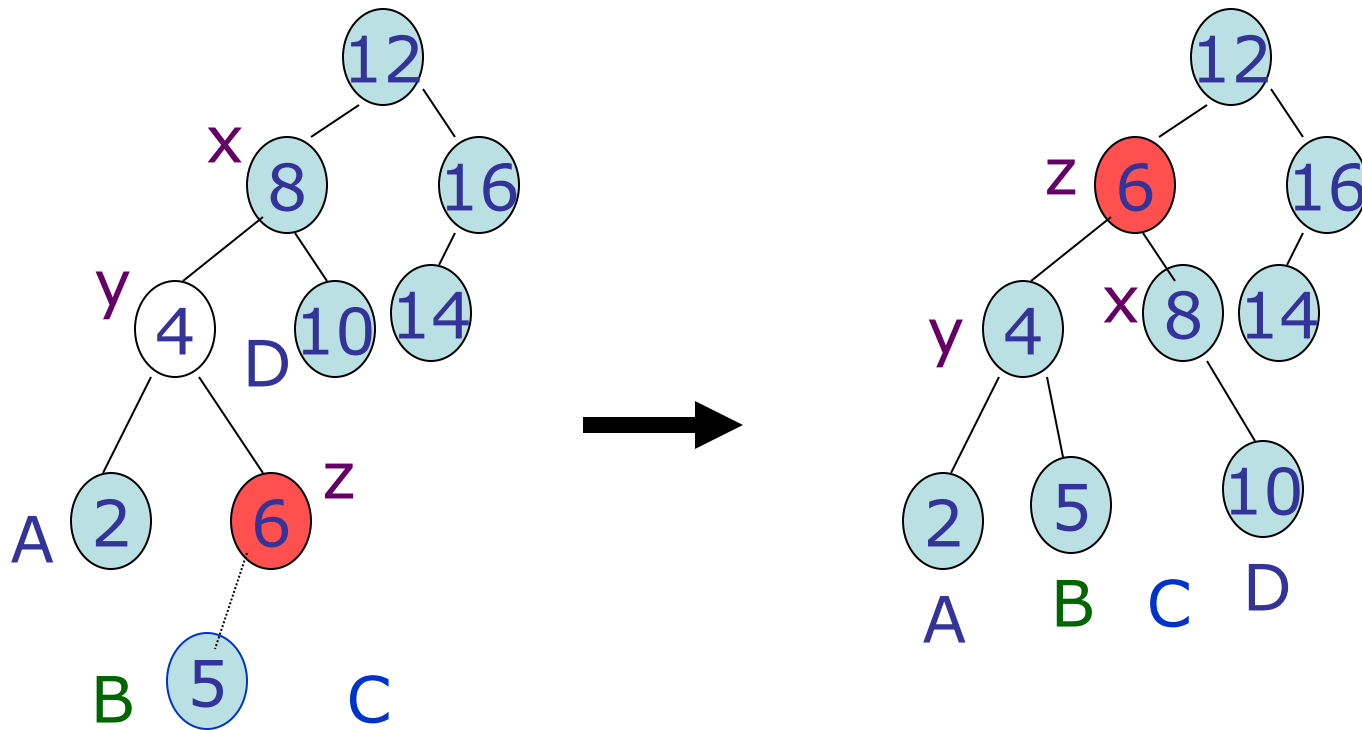
- Let x be the node at which $x \rightarrow \text{left}$ and $x \rightarrow \text{right}$ differ by more than 1; Assume that the height of x is $h+3$
 - Heights of two subtrees: $h+2$, h
 3. Height of $x \rightarrow \text{left} \rightarrow \text{right}$: $h+1 \Rightarrow$ double rotation with the left child (LR): $RR \rightarrow LL$
 - $BF(x) = +2$, $BF(x \rightarrow \text{left}) = -1$
 4. Height of $x \rightarrow \text{right} \rightarrow \text{left}$: $h+1 \Rightarrow$ double rotation with the right child (RL): $LL \rightarrow RR$
 - $BF(x) = -2$, $BF(x \rightarrow \text{right}) = +1$

Q&A: What are the conditions for double rotations?

Double Rotations: *LR*



Double Rotations: *LR*

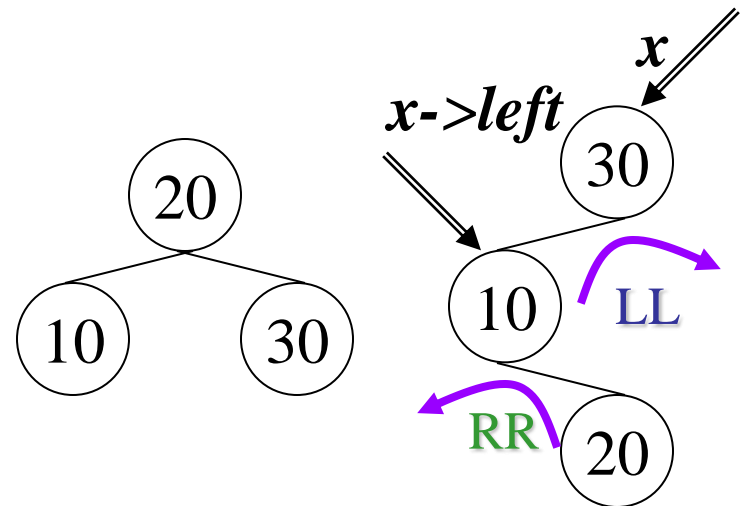


Insert "5"

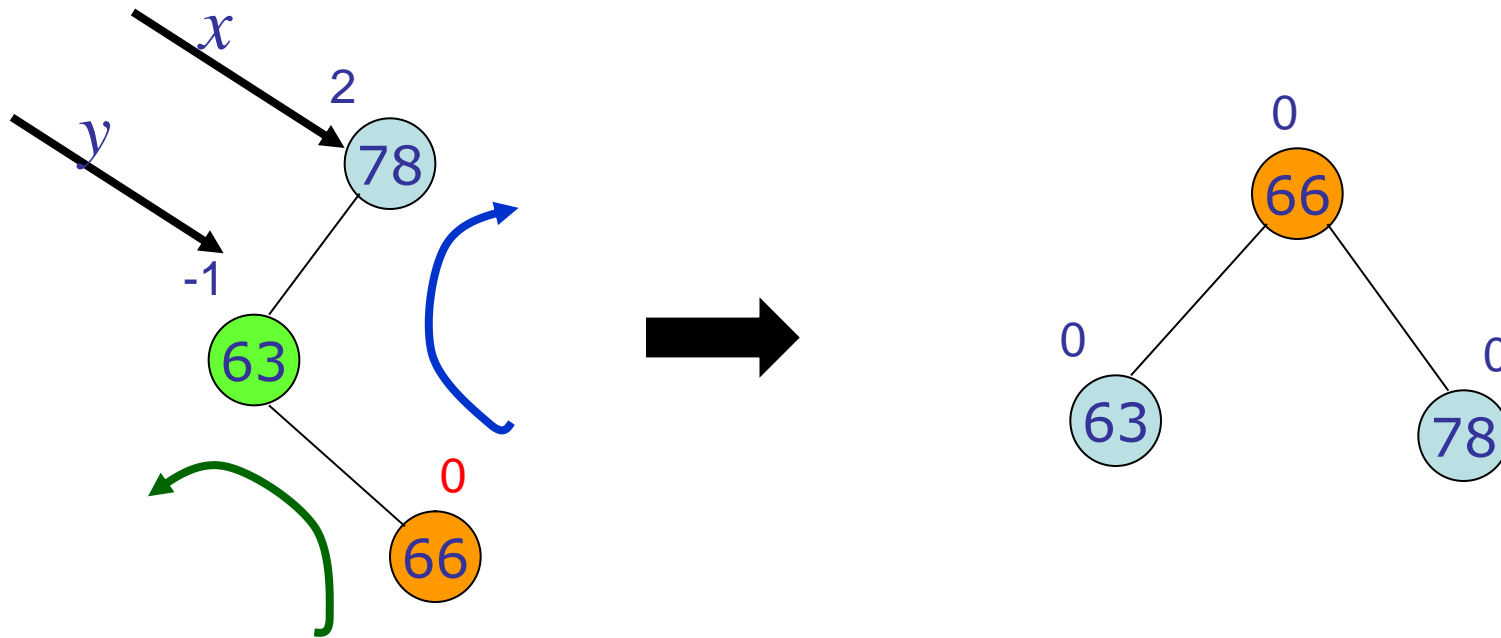
LR Rotation: *Pseudocode*

// first rotate left child with its right child; RR
// then, rotate node x with its new left child. LL

```
nodeType rotateLR(nodeType x)
{
    x->left = rotateRR (x->left);
    return rotateLL(x);
}
```

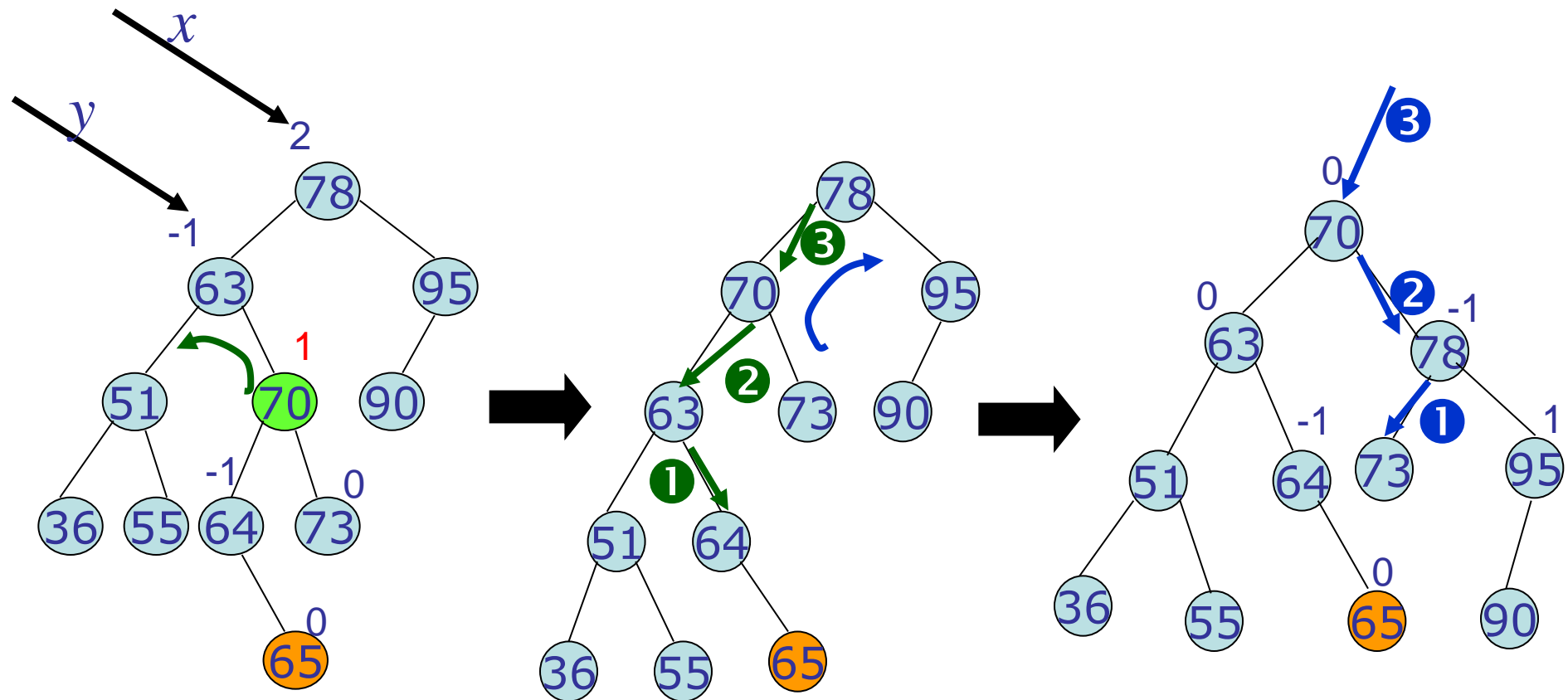


LR Rotation: *Examples*



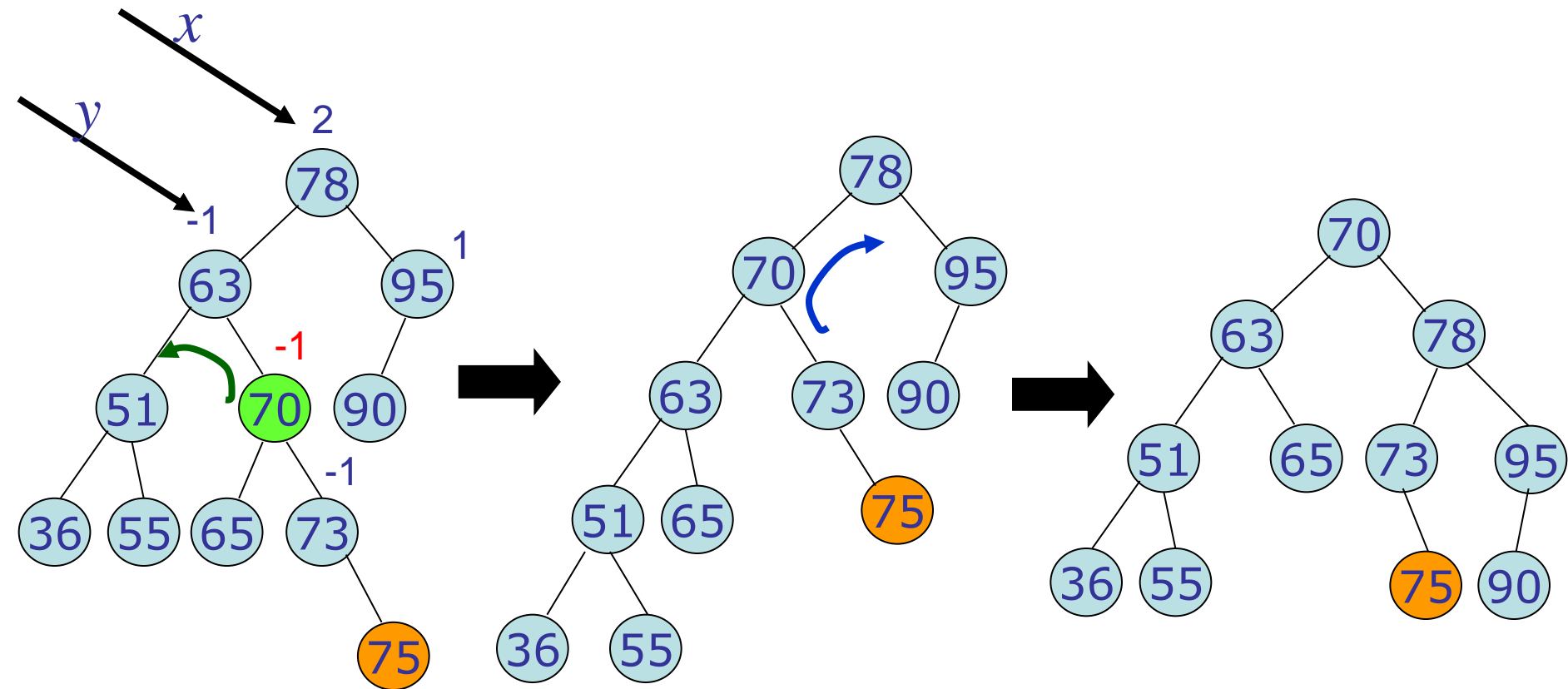
LR: case 1

LR Rotation: *Examples*



LR: case 2

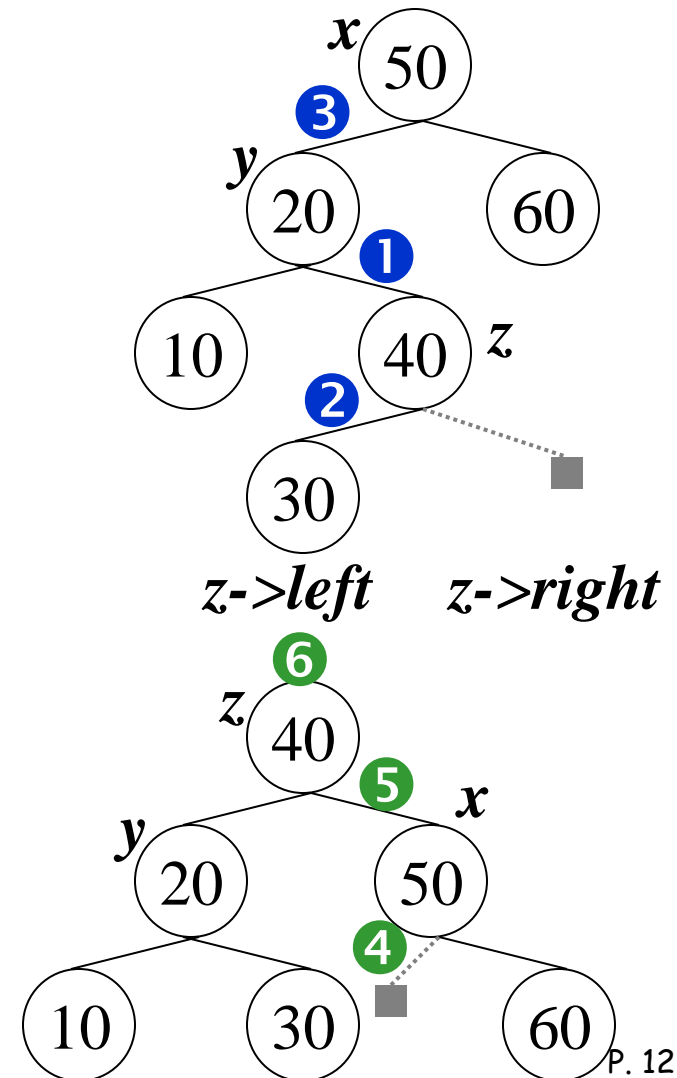
LR Rotation: *Examples*



LR: case 3

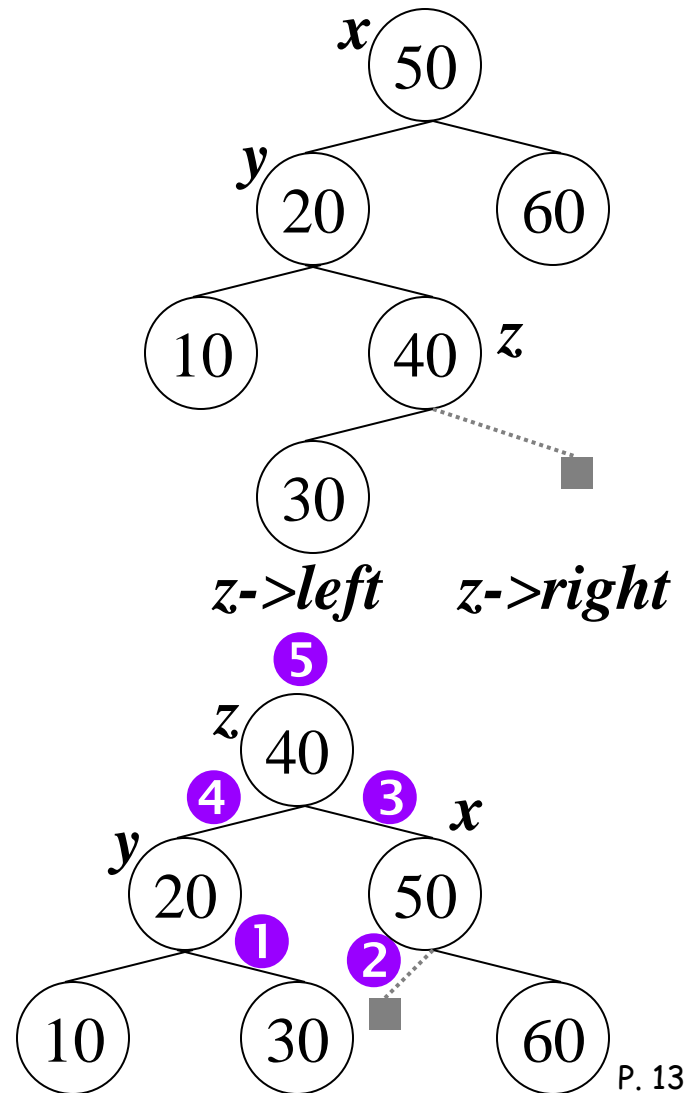
LR Rotation: *Version I.*

```
nodeType rotateLR1(nodeType x)
{
  nodeType y = x->left;
  nodeType z = y->right;
  // RR rotation on y
  y->right = z->left;      ①
  z->left = y;             ②
  x->left = z;             ③
  // LL rotation on x
  x->left = z->right;      ④
  z->right = x;           ⑤
  return z;              ⑥
}
```

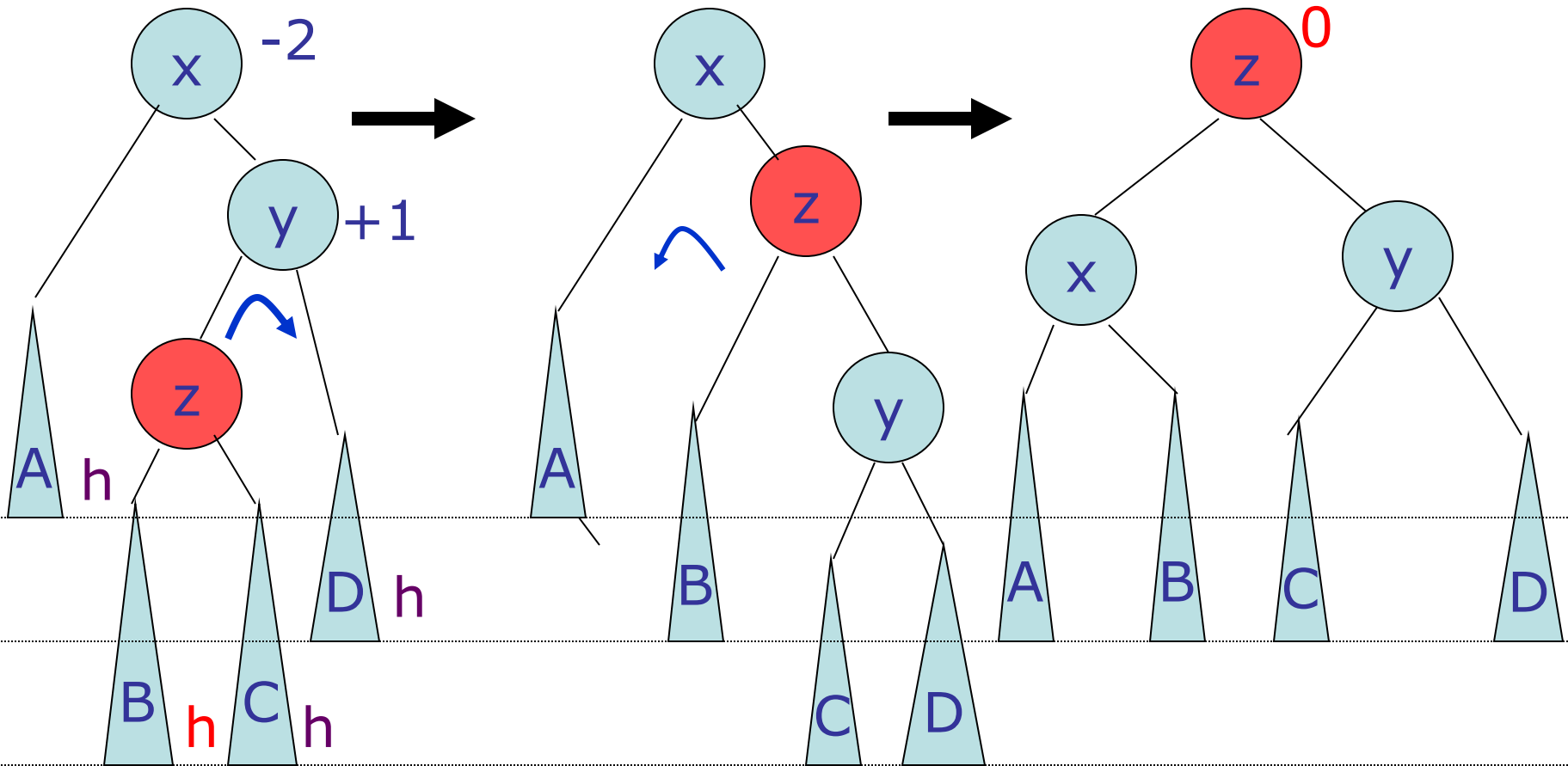


LR Rotation: *Version II.*

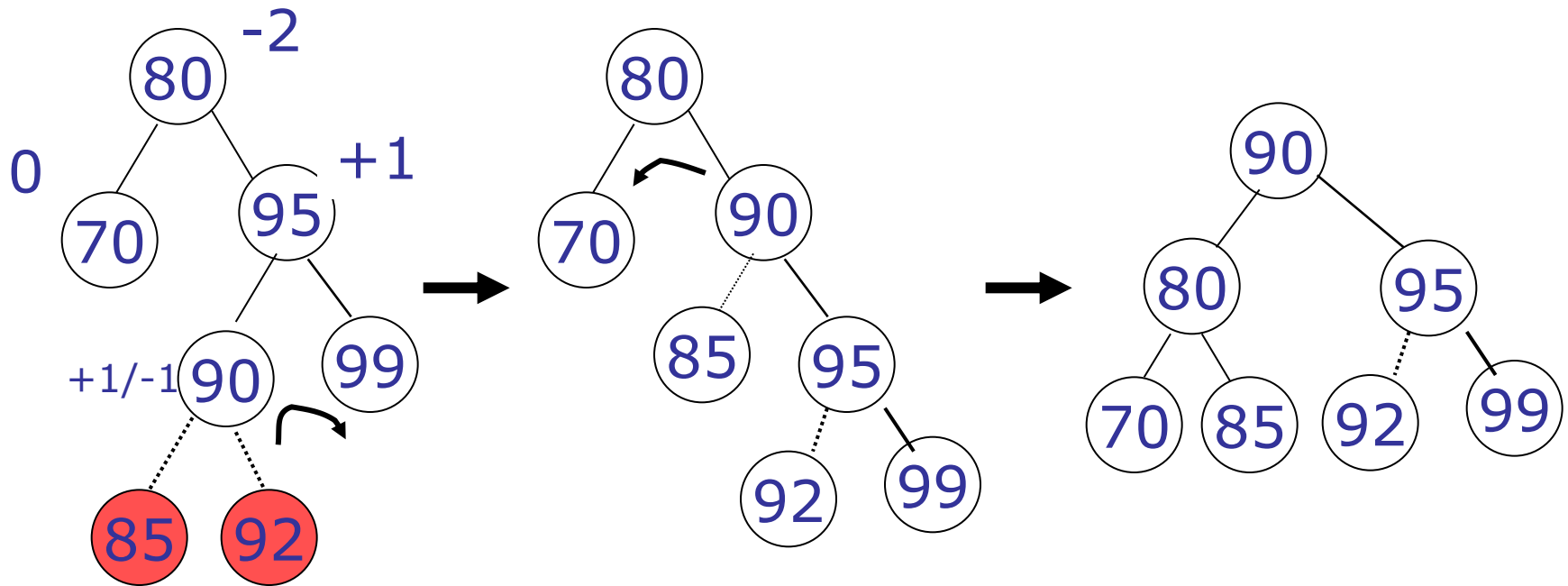
```
nodeType rotateLR2(nodeType x)
{
  nodeType y = x->left;
  nodeType z = y->right;
  y->right = z->left;           ①
  x->left = z->right;           ②
  z->right = x;                 ③
  z->left = y;                  ④
  return z;                     ⑤
}
```



Double Rotations: *RL*



Double Rotations: *RL*



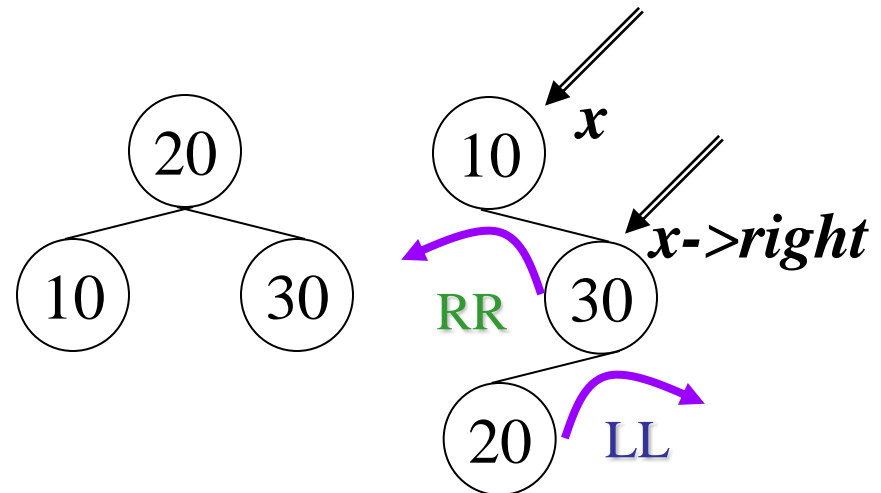
Insert 85 or 92

RL Rotation: *Pseudocode*

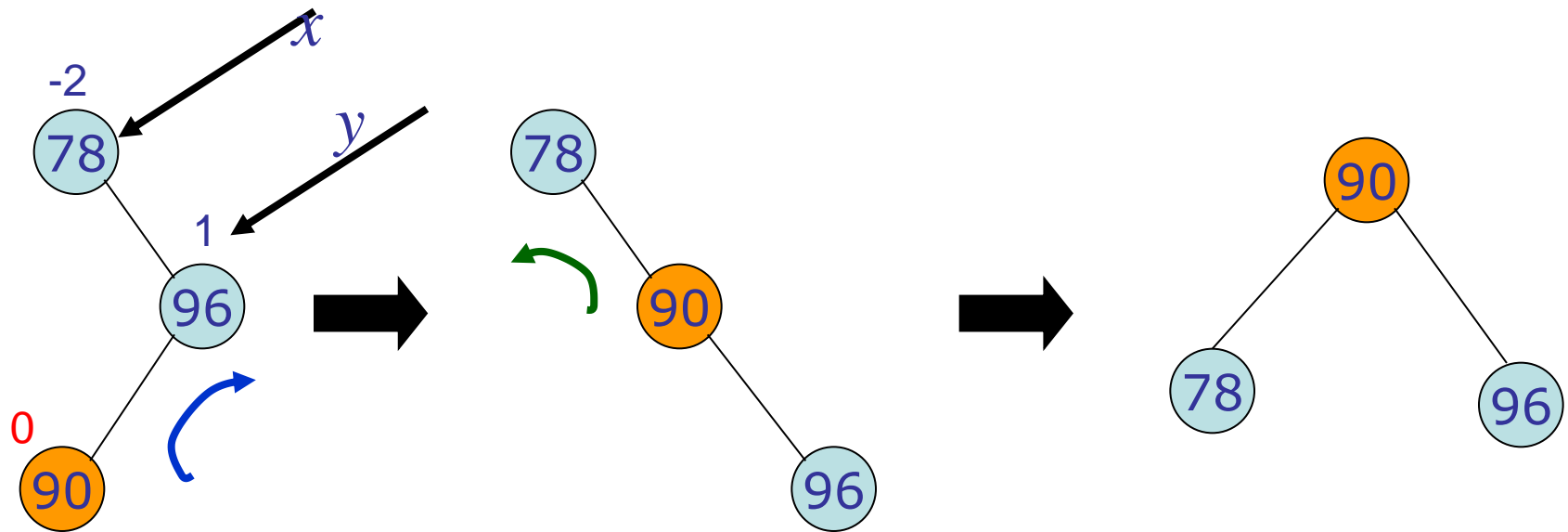
// first rotate right child with its left child; LL
// then, rotate node x with its new right child. RR

nodeType *rotateRL*(nodeType **x**)

```
{  
    x->right = rotateLL (x->right);  
    return rotateRR(x);  
}
```

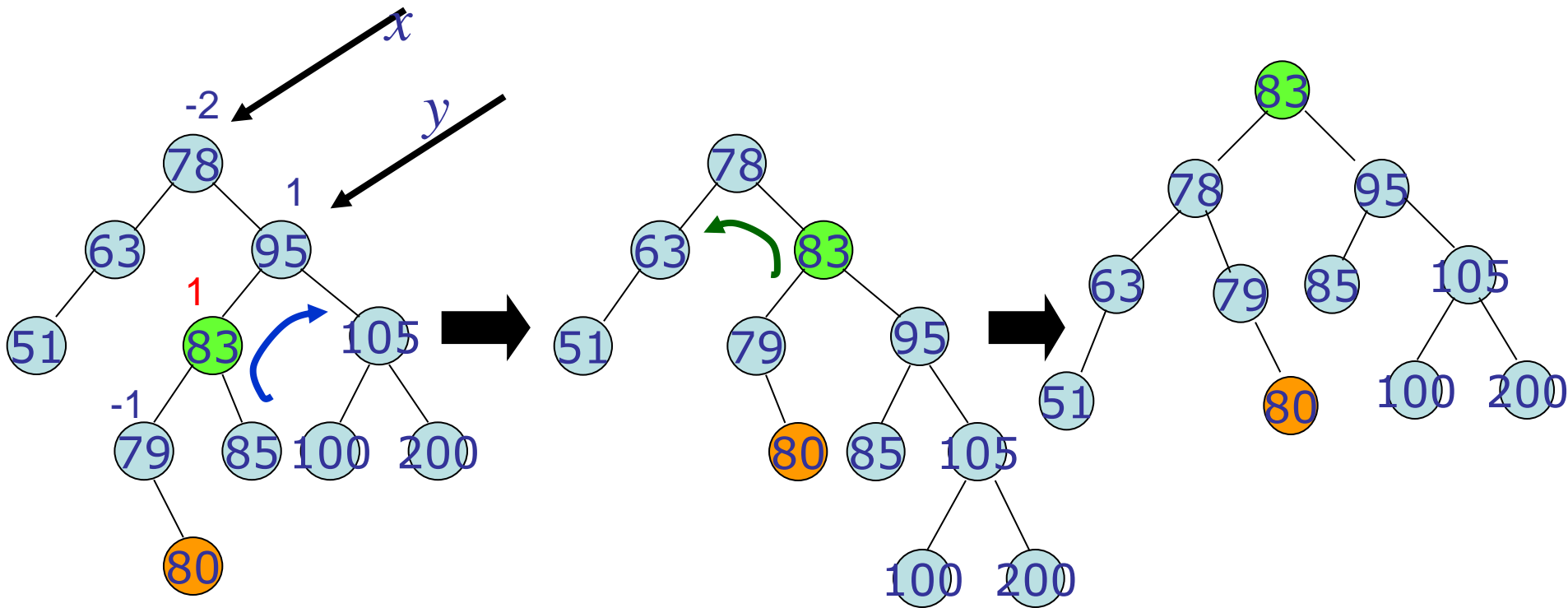


RL Rotation: *Examples*



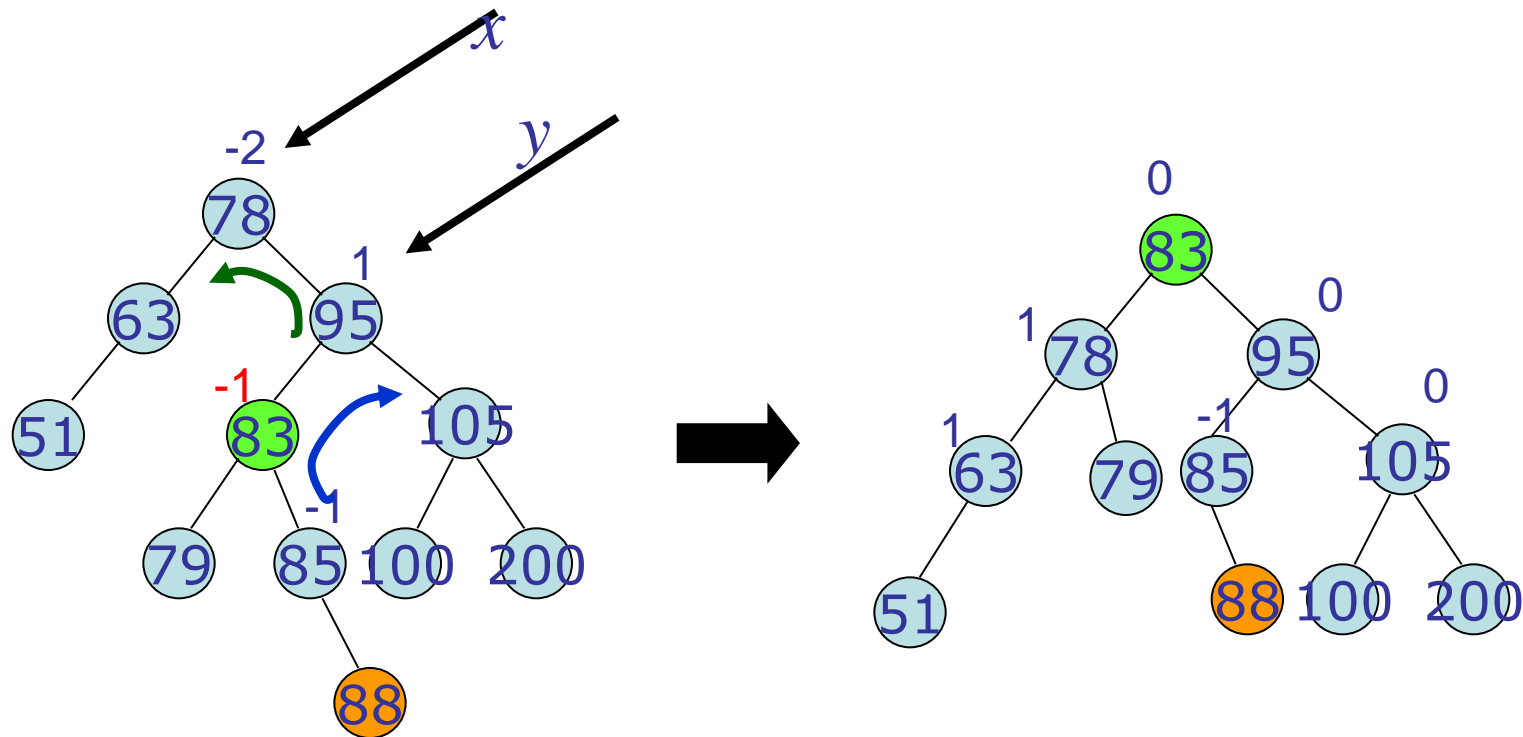
RL: case 1

RL Rotation: *Examples*



RL: case 2

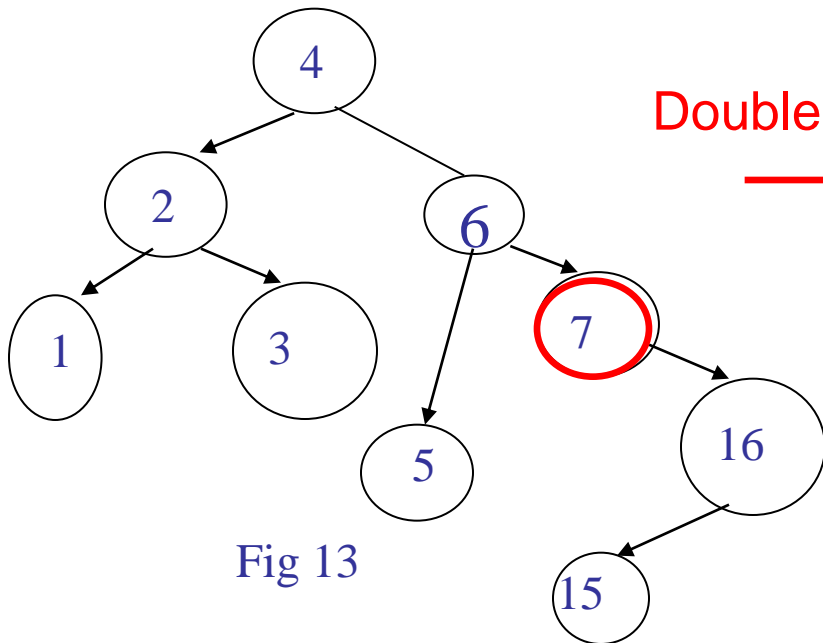
RL Rotation: *Examples*



RL: case 3

Try it!

Insert *15, 14*



Double rotation

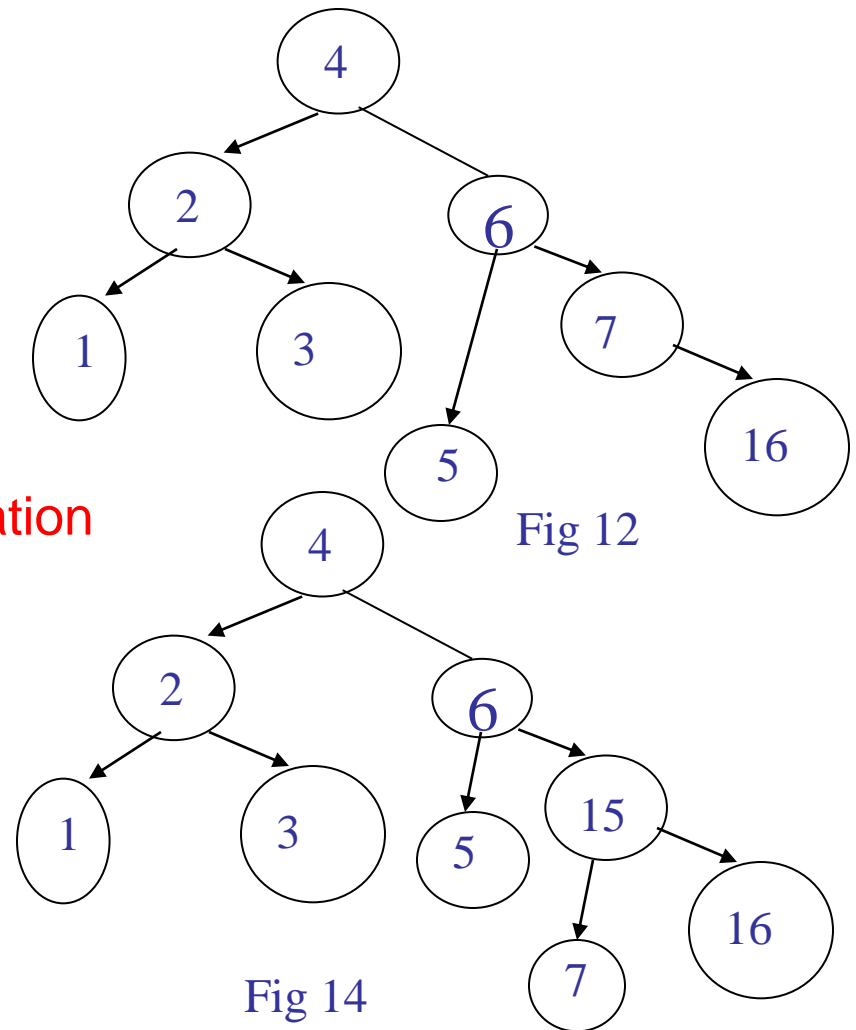
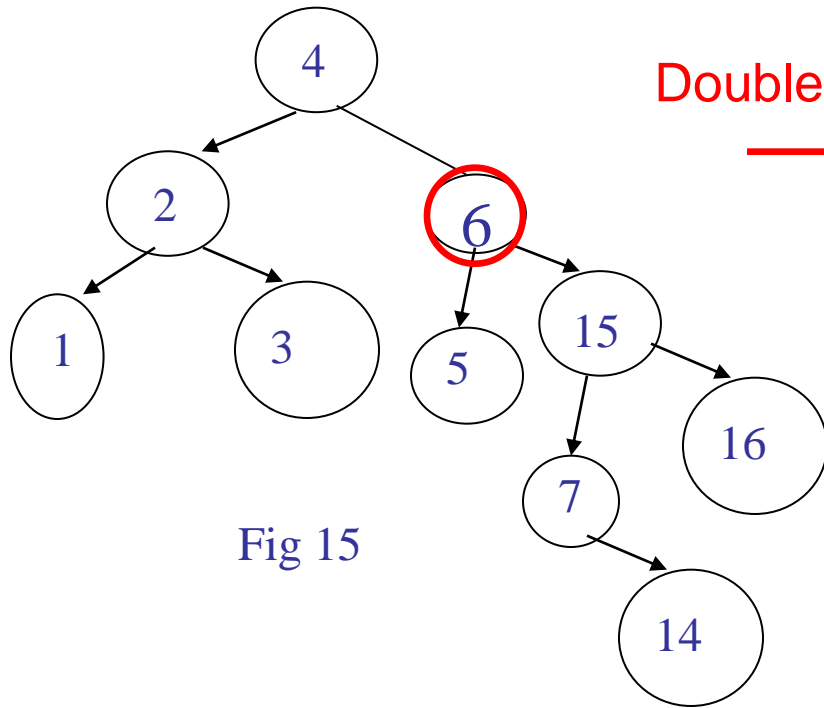


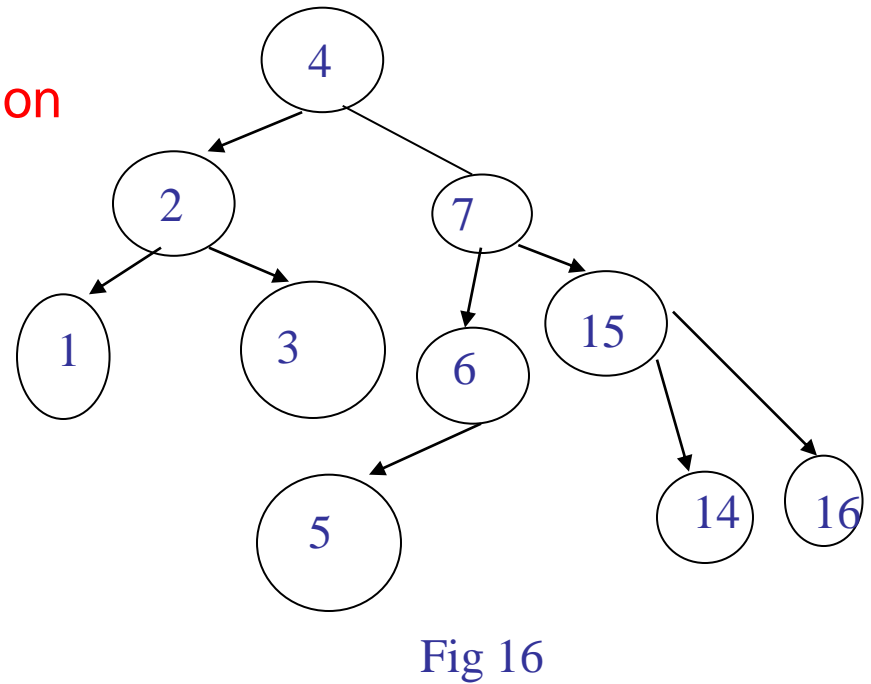
Fig 12

Try it!

Insert *15*, *14*



Double rotation



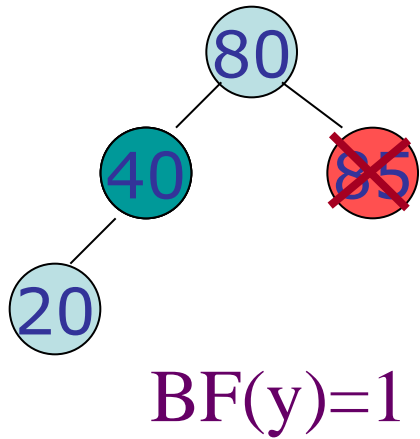
Practice 1: *Inserting* into AVL Tree

□ Input order: 10 12 30 8 60 40 70

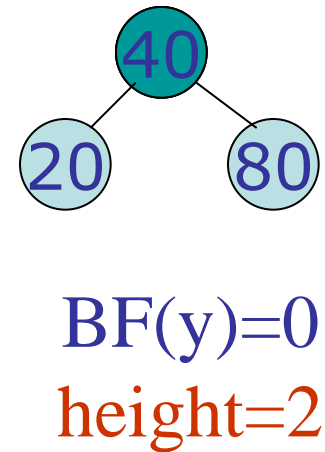
AVL Tree: *Deletion*

1. Delete a node x as in a *binary search tree* and the last node deleted is a **leaf**.
2. Trace the path *from the new leaf towards the root*.
 - For each node x encountered, check if the heights of $left(x)$ and $right(x)$ differ by **at most 1**.
 - If **NOT**, restructure by either a **single rotation** or a **double rotation**
3. After we perform a rotation at x , we may have to perform a rotation at some ancestor of x .
 - We must **continue** to trace the path until the root.

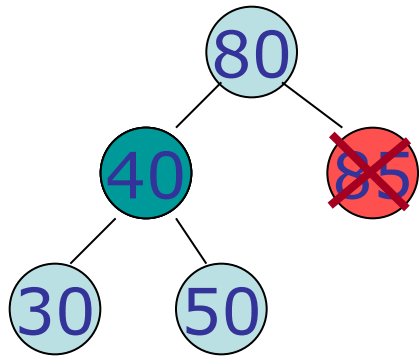
After delete 85,...



Single LL
rotations

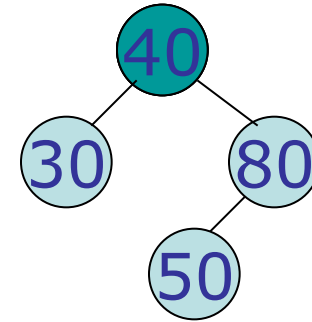


After delete 85,...



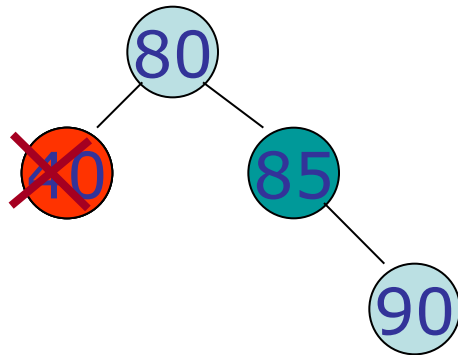
$BF(y)=0$

Single LL
rotations



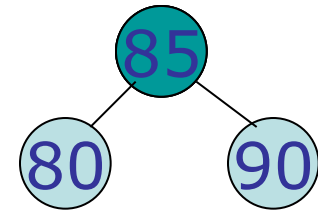
$BF(y)= -1$
height=3

After delete 40,...



$BF(y) = -1$

Single RR
rotations

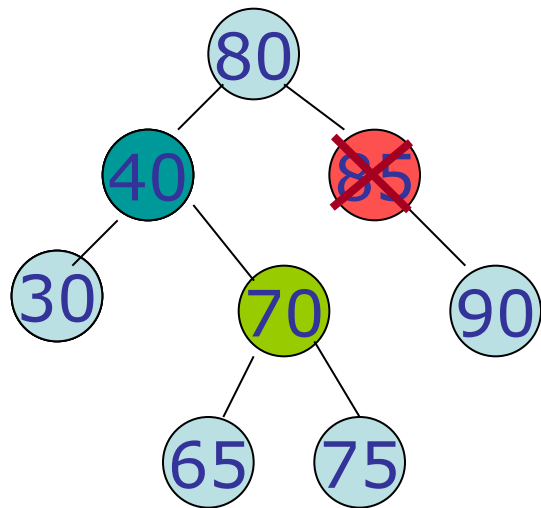


$BF(y) = 0$
height=2

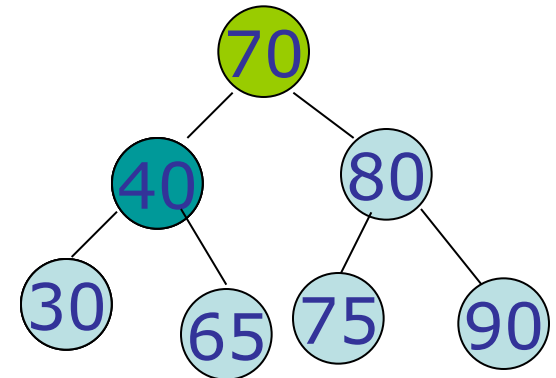
Rotations in Deletion

- There are *4 cases* for **single rotations**, but we do not need to distinguish among them.
- There are exactly *two cases* for **double rotations** (*as in the cases of insertion*)
- Therefore, we can **reuse** exactly the same procedure for insertion to determine which rotation to perform!

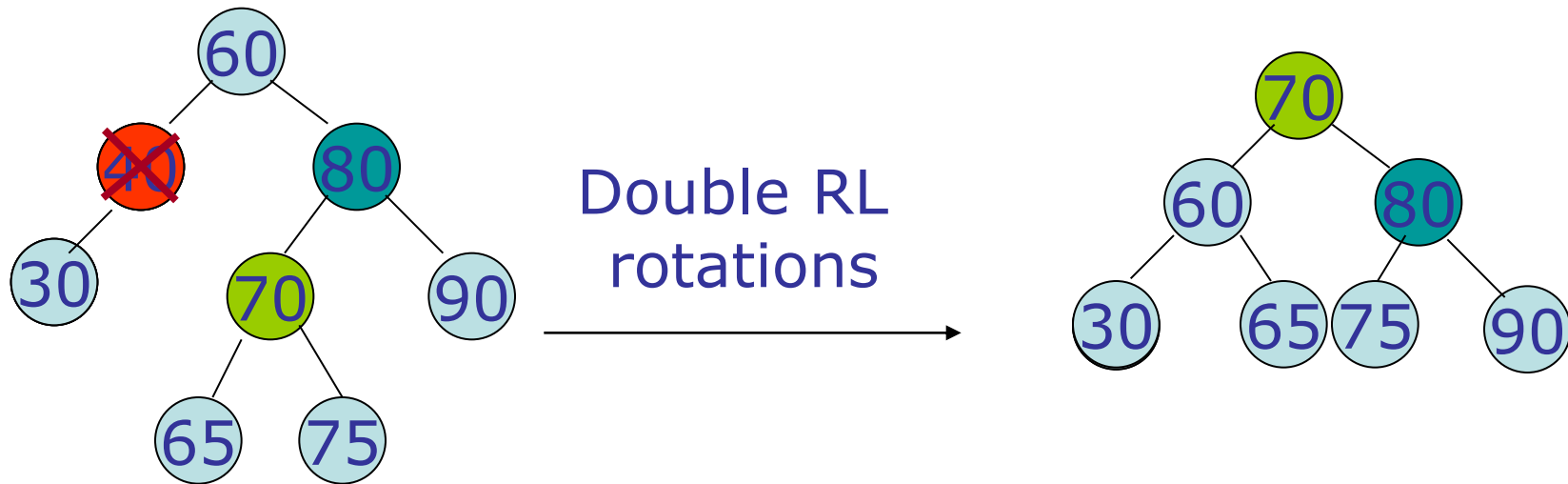
After delete 85,...



Double LR
rotations



After delete 40,...



Practice 2: *Deleting* from AVL Tree

□ After the deletions of 30, 10, 60

