

Discrete Mathematics with Applications I

COMPSCI&SFWRENG 2DM3

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Raymond Smullyan posed many puzzles about an island that has two kinds of inhabitants:

- **knights**, who always tell the truth, and
- **knaves**, who always lie.

You encounter two people A and B .
What are A and B if

- 1 A says “We are both knaves.”?
- 2 A says “At least one of us is a knave.”?
- 3 A says “If I am a knight, then so is B .”?
- 4 A says “We are of the same type.”?
- 5 A says “ B is a knight” and
 B says “The two of us are opposite types.”?

Plan for Today

- A Theorem of the Integers — brief outlook using LADM Chapter 15
- Boolean Expressions — picking up pieces of LADM Chapter 2

Proving Zero of Multiplication

- (1.2) **Axiom, Reflexivity of =:** $a = a$
 (15.3) **Axiom, Additive identity:** $a + 0 = a$
 (15.5) **Axiom, Distributivity:** $a \cdot (b + c) = a \cdot b + a \cdot c$
 (15.8) **Cancellation of +:** $a + b = a + c \quad \equiv \quad b = c$

Proving (15.9) $a \cdot 0 = 0$:

$$\begin{aligned}
 & a \cdot 0 = 0 \\
 \equiv & \langle \text{Cancellation of + (15.8), with } a, b, c := a \cdot d, a \cdot 0, 0 \rangle \\
 & a \cdot d + a \cdot 0 = a \cdot d + 0 \\
 \equiv & \langle \text{Distributivity of } \cdot \text{ over + (15.5)} \rangle \\
 & a \cdot (d + 0) = a \cdot d + 0 \\
 \equiv & \langle \text{Identity of + (15.3), twice} \rangle \\
 & a \cdot d = a \cdot d \quad \text{— This is Reflexivity of = (1.2)}
 \end{aligned}$$

Proving Zero of Multiplication

- (1.2) **Axiom, Reflexivity of =:** $a = a$
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 (15.5) **Axiom, Distributivity:** $a \cdot (b + c) = a \cdot b + a \cdot c$
 (15.8) **Cancellation of +:** $a + b = a + c \quad \equiv \quad b = c$

Proving (15.9) $a \cdot 0 = 0$:

$$\begin{aligned}
 & a \cdot 0 = 0 \\
 \equiv & \langle \text{Cancellation of + (15.8), with } a, b, c := a \cdot 42, a \cdot 0, 0 \rangle \\
 & a \cdot 42 + a \cdot 0 = a \cdot 42 + 0 \\
 \equiv & \langle \text{Distributivity of } \cdot \text{ over + (15.5)} \rangle \\
 & a \cdot (42 + 0) = a \cdot 42 + 0 \\
 \equiv & \langle \text{Identity of + (15.3), twice} \rangle \\
 & a \cdot 42 = a \cdot 42 \quad \text{— This is "Reflexivity of =" (1.2)}
 \end{aligned}$$

LADM Theory of Integers — (15.20)

(15.5) Distributivity $a \cdot (b + c) = a \cdot b + a \cdot c$	(15.3) Identity of + $0 + a = a$	(15.4) Identity of · $1 \cdot a = a$
(15.8) Cancellation of + $a + b = a + c \quad \equiv \quad b = c$	(15.13) Unary minus $a + (-a) = 0$	(15.14) Subtraction $a - b = a + (-b)$

(15.20) $-a = -1 \cdot a$

— Prove this here!

Truth Values and Equivalence — Remember

Boolean constants/values: *false, true*

The set/type of Boolean values: \mathbb{B}

Equality of boolean values is also called **equivalence** and written \equiv

$p \equiv q$ can be read as: p is equivalent to q

or: p exactly when q

or: p if-and-only-if q

or: p iff q

In many current notebooks, the following is added as an axiom:

(15.8) **Cancellation of +:** $a + b = a + c \equiv b = c$

Remember: Equivalence is just equality of truth values!

— but written with a different symbol

— with different notational conventions

Existential Quantification Examples

$(\exists k : \mathbb{N} \bullet k > 9999)$

- “There exists a natural number k such that $k > 9999$ (holds)”
- “For some natural number k , we have $k > 9999$ ”
- “Some natural number is greater than 9999”

$(\exists x : \mathbb{R} \mid x > 0 \bullet x \cdot x = x + 1)$

- “There exists a real number x with $x > 0$ such that $x \cdot x = x + 1$ (holds)”
- “For some positive real number x , we have $x \cdot x = x + 1$ ”

$(\exists r, s : \mathbb{Q} \mid r < s < r + 1/1000 \bullet r < \pi < s)$

- “There exist rational numbers r and s with $r < s < r + 1/1000$ such that $r < \pi < s$ (holds)”
- “For some rational numbers r and s with $r < s$ and $s - r < 1/1000$, we have $r < \pi < s$ ”
- “ π can be enclosed within rational bounds that are less than 1/1000 apart”

Universal Quantification Examples

$(\forall k : \mathbb{N} \bullet 2 \cdot k \geq k)$

- “For all natural numbers k , we have $2 \cdot k \geq k$ ”

$(\forall x, y : \mathbb{R} \bullet x \cdot y = y \cdot x)$

- “For all real numbers x and y , we have $x \cdot y = y \cdot x$ ”
- “Multiplication of real numbers is symmetric (commutative)”

$(\forall x : \mathbb{R} \mid x > 5 \bullet x \cdot x > 10)$

- “For all real numbers x with $x > 5$, we have $x \cdot x > 10$ ”
- “The square of a real number greater than 5 is greater than 10.”

$(\forall m, n : \mathbb{N} \mid m \neq n \bullet m \cdot m \neq n \cdot n)$

- “For all natural numbers m and n with $m \neq n$, we have $m \cdot m \neq n \cdot n$ ”
- “Different natural numbers have different squares.”

Combined Quantification Examples

- “Every integer has an additive inverse.”
- “For every integer k , there exists an integer n such that $k + n = 0$ (holds).”

$$(\forall k : \mathbb{Z} \bullet (\exists n : \mathbb{Z} \bullet k + n = 0))$$

- “There is a least natural number.”
- “There exists a natural number b such that every natural number n is at least b ”.
- “There exists a natural number b such that for every natural number n , we have $b \leq n$ ”.

$$(\exists b : \mathbb{N} \bullet (\forall n : \mathbb{N} \bullet b \leq n))$$

Combined Quantification Examples (ctd.)

- “There is a least integer.”
- “There exists an integer b such that every integer n is at least b ”.
- “There exists an integer b such that for every integer n , we have $b \leq n$ ”.

$$(\exists b : \mathbb{Z} \bullet (\forall n : \mathbb{Z} \bullet b \leq n))$$

- “ π can be enclosed within rational bounds that are less than any ε apart”
- “For every positive real number ε , there are rational numbers r and s with $r < s < r + \varepsilon$, such that $r < \pi < s$ ”

$$(\forall \varepsilon : \mathbb{R} \mid 0 < \varepsilon \bullet (\exists r, s : \mathbb{Q} \mid r < s < r + \varepsilon \bullet r < \pi < s))$$

Truth Values and Equivalence — Remember

Boolean constants/values: *false, true*

The set/type of Boolean values: \mathbb{B}

Equality of boolean values is also called **equivalence** and written \equiv

$p \equiv q$ can be read as: p is **equivalent** to q
or: p **exactly when** q
or: p **if-and-only-if** q
or: p **iff** q

For now, treat the following as an axiom:

$$(15.8) \quad \text{Cancellation of } +: \quad a + b = a + c \quad \equiv \quad b = c$$

Remember: Equivalence is just equality of truth values!

— but written with a different symbol

— with different notational conventions

Boolean Expressions

- Boolean constants: *false*, *true*
- **Proposition symbols** p, q — **variables of type** \mathbb{B}
- Applications of Boolean-valued operators to expressions of their argument types:
Number types \mathbb{N}, \mathbb{Z} :
 - $1 = 2$
 - $42 \leq 56$
 - $a + b = a + c$**String:** "Hello" \leq "Hello World!"
Set: $\{1, 2, 3\} \cap \{2, 3, 4\} = \{2, 3\}$
 \mathbb{B} :
 - $(a + b = a + c) \equiv (b = c)$
 - $(a \leq b) \Rightarrow (c - b \leq c - a)$
 - The inscription on the gold casket is true \neq The portrait is in the gold casket
 - $G \neq gc$
 - $(p \wedge q) \Rightarrow p$

Truth Values & Unary Boolean Operators

Boolean constants/values: *false*, *true*

Unary Boolean operators:

Arg.	Result (one column per operator)				
	<i>id</i>	\neg			
<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>true</i>	$\neg \text{false} = \text{true}$
<i>true</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>true</i>	$\neg \text{true} = \text{false}$

This table shows all four possible functions that map one Boolean argument to a Boolean result.

Unary Boolean Operators — f_1

Unary Boolean operators:

Arg.	Result (one column per operator)				
	f_1	<i>id</i>	\neg	f_4	
<i>false</i>	<i>false</i>	<i>false</i>	<i>true</i>	<i>true</i>	$f_1(\text{false}) = \text{false}$
<i>true</i>	<i>false</i>	<i>true</i>	<i>false</i>	<i>true</i>	$f_1(\text{true}) = \text{false}$

Unary Boolean Operators — *id*

Unary Boolean operators:

Arg.	Result (one column per operator)				
	f_1	id	\neg	f_4	
<i>false</i>	false	false	true	true	$id(\text{false}) = \text{false}$
<i>true</i>	false	true	false	true	$id(\text{true}) = \text{true}$

Arg.	Result (one column per operator)				
	f_1	id	\neg	f_4	
<i>false</i>	false	false	true	true	$id(false) = false$
<i>true</i>	false	true	false	true	$id(true) = true$

Unary Boolean Operators — f_4

Unary Boolean operators:

Arg.	Result (one column per operator)				
	f_1	id	\neg	f_4	
<i>false</i>	false	false	true	true	$f_4(\text{false}) = \text{true}$
<i>true</i>	false	true	false	true	$f_4(\text{true}) = \text{true}$

Arg.	Result (one column per operator)				
	f_1	id	\neg	f_4	
false	false	false	true	true	$f_4(\text{false}) = \text{true}$
true	false	true	false	true	$f_4(\text{true}) = \text{true}$

How Many Binary Boolean Operators Are There?

- How many arguments does a binary Boolean operator take?
- How many different argument value (\mathbb{B}) combinations are there?
- How many ways are there to map all argument value combinations to \mathbb{B} ?

Args.		Result (one column per operator)														
		\wedge	\nrightarrow	\Leftarrow	\neq	\vee	nor	\equiv	$=$	\Leftarrow	\Rightarrow	nand				
F	F	F	F	F	F	F	F	T	T	T	T	T	T	T	T	T
F	T	F	F	T	T	T	F	F	F	F	T	T	T	T	T	T
T	F	F	T	F	F	T	F	F	T	T	F	F	T	T	T	T
T	T	F	T	F	T	F	F	T	F	T	F	T	F	T	F	T

- How many arguments does a binary Boolean operator take?
- How many different argument value (\mathbb{B}) combinations are there?
- How many ways are there to map all argument value combinations to \mathbb{B} ?

Args.		Result (one column per operator)									
		\wedge	\nrightarrow	$\not\leftarrow$	\neq	\vee	nor	$=$	\Leftarrow	\Rightarrow	nand
F	F	F	F	F	F	F	T	T	T	T	T
F	T	F	F	T	T	T	F	F	F	T	T
T	F	F	T	T	F	T	F	F	T	F	T
T	T	F	T	F	T	F	T	F	T	F	T

Binary Boolean Operators

Args.		Result (one column per operator)													
		\wedge	\nrightarrow	\nleftarrow	\neq	\vee	nor	\equiv	\Leftarrow	\Rightarrow	nand				
F	F	F	F	F	F	F	F	T	T	T	T	T	T	T	T
F	T	F	F	F	T	T	T	F	F	F	F	T	T	T	T
T	F	F	T	T	F	F	T	F	F	T	T	F	F	T	T
T	T	F	T	F	T	F	T	F	T	F	T	F	T	F	T

For example:

$$\begin{aligned} \text{true} \wedge \text{false} &= \text{false} \\ \text{true} \neq \text{false} &= \text{true} \\ \text{false} \Leftarrow \text{true} &= \text{false} \end{aligned}$$

Names of Binary Boolean Operators & their Arguments

		\wedge				\neq	\vee	nor	\equiv		\Leftarrow	\Rightarrow	nand		
F	F	F	F	F	F	F	F	T	T	T	T	T	T	T	T
F	T	F	F	F	T	T	T	F	F	F	F	T	T	T	T
T	F	F	F	T	F	F	T	F	F	T	T	F	F	T	T
T	T	F	T	F	T	F	T	F	T	F	T	F	T	F	T

- $b = c$ reads “ b equals c ” — conventional **equality**
- $b \neq c$ reads “ b differs from c ” — conventional **inequality**
- $b \vee c$ reads “ b or c ” — **disjunction**; b and c are **disjuncts**
- $b \wedge c$ reads “ b and c ” — **conjunction**; b and c are **conjuncts**
- $b \Rightarrow c$ reads “ b implies c ” or “if b then c ” — **implication**;
 b is the **antecedent**; c is the **consequent**
- $b \Leftarrow c$ reads “ b follows from c ” or “ b if c ” — **consequence**;
 b is the **consequent**; c is the **antecedent**

Table of Precedences

- $[x := e]$ (textual substitution) (highest precedence)
- $.$ (function application)
- unary prefix operators $+$, $-$, \neg , $\#$, \sim , \mathcal{P}
- $**$
- \cdot $/$ \div **mod** **gcd**
- $+$ $-$ \cup \cap \times \circ \bullet
- \downarrow \uparrow
- $\#$
- \triangleleft \triangleright \wedge
- $=$ \neq $<$ $>$ \in \subset \subseteq \supset \supseteq $|$ (conjunctive)
- \vee \wedge
- \Rightarrow \nrightarrow \Leftarrow \nleftarrow
- \equiv \neq (lowest precedence)

All non-associative binary infix operators associate to the left, except $**$, \triangleleft , \Rightarrow , \rightarrow , which associate to the right.

Modeling English Propositions — Recipe

- Transform into shape with clear subpropositions
- Introduce Boolean variables to denote subpropositions
- Replace these subpropositions by their corresponding Boolean variables
- Translate the result into a Boolean expression, using (no perfect translation rules are possible!) **for example**:

and, but	becomes	\wedge
or	becomes	\vee
not	becomes	\neg
it is not the case that	becomes	\neg
if p then q	becomes	$p \Rightarrow q$

Binary Boolean Operators: “but”

Args.			
		\wedge	
F	F	F	The moon is green, but $2 + 2 = 7$.
F	T	F	The moon is green, but $1 + 1 = 2$.
T	F	F	$1 + 1 = 2$, but the moon is green.
T	T	T	$1 + 1 = 2$, but the sun is a star.

Binary Boolean Operators: “if”

Args.			
		\Leftarrow	
F	F	T	The moon is green if $2 + 2 = 7$.
F	T	F	The moon is green if $1 + 1 = 2$.
T	F	T	$1 + 1 = 2$ if the moon is green.
T	T	T	$1 + 1 = 2$ if the sun is a star.

See also textbook p. 36:

$$\begin{array}{lcl}
 \boxed{\text{To stay dry, it's \textbf{sufficient} to wear a raincoat.}} & = & \boxed{\text{You will stay dry if you wear a raincoat.}} \\
 & = & \left(\boxed{\text{You will stay dry.}} \Leftarrow \boxed{\text{You wear a raincoat.}} \right) \\
 \text{“}p \text{ if } q\text{.”} & = & \text{“If } q \text{ then } p\text{.”}
 \end{array}$$

Binary Boolean Operators: "only if"

Args.			
		\Rightarrow	
F	F	T	The moon is green only if $2 + 2 = 7$.
F	T	T	The moon is green only if the sun is a star.
T	F	F	The sun is a star only if the moon is green.
T	T	T	$1 + 1 = 2$ only if the sun is a star.

See also textbook p. 36:

$$\begin{aligned}
 &\boxed{\text{To stay dry, it's necessary to wear a raincoat.}} = \boxed{\text{You will stay dry only if you wear a raincoat.}} \\
 &= \left(\boxed{\text{You will stay dry.}} \Rightarrow \boxed{\text{You wear a raincoat.}} \right) \\
 &\text{"}p \text{ only if } q\text{"} = \text{"If } p \text{ then } q\text{"}
 \end{aligned}$$

Necessary and Sufficient Conditions

(Textbook p. 36)

$$\begin{aligned}
 &\boxed{\text{To stay dry, it's necessary to wear a raincoat.}} = \boxed{\text{You will stay dry only if you wear a raincoat.}} \\
 &= \left(\boxed{\text{You will stay dry.}} \Rightarrow \boxed{\text{You wear a raincoat.}} \right) \\
 &\boxed{\text{To stay dry, it's sufficient to wear a raincoat.}} = \boxed{\text{You will stay dry if you wear a raincoat.}} \\
 &= \left(\boxed{\text{You will stay dry.}} \Leftarrow \boxed{\text{You wear a raincoat.}} \right)
 \end{aligned}$$

Binary Boolean Operators: "even if"

Args.			
p	q	p	
F	F	F	The moon is green, even if $2 + 2 = 7$.
F	T	F	The moon is green, even if $1 + 1 = 2$.
T	F	T	$1 + 1 = 2$, even if the moon is green.
T	T	T	$1 + 1 = 2$, even if the sun is a star.

Args.		Transforming “even if”	
p	q	p	
F	F	F	The moon is green, even if $2 + 2 = 7$.
F	T	F	The moon is green, even if $1 + 1 = 2$.
T	F	T	$1 + 1 = 2$, even if the moon is green.
T	T	T	$1 + 1 = 2$, even if the sun is a star.

$1 + 1 = 2$, and, if the sun is a star, we still have $1 + 1 = 2$.

Declarations:

$t \quad := \quad 1 + 1 = 2$

$s \quad := \quad \text{The sun is a star}$

Formalisation:

$t \wedge (s \Rightarrow t)$