### **COMPSCI 2GA3 Tutorial 4 Note**

#### Note:

This note does NOT cover all the materials in Chapter 2 -- Only the ones rated to sample questions of this tutorial are included.

For any questions about the tutorials and courses, feel free to contact me. (Email: wangm235@mcmaster.ca)

GLHF: )
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NO ADDITIONAL MATERIAL IN THIS WEEK'S NOTE --- HAPPY READING WEEK:)

### What is the "Caller"? What is the "Callee"?

```
public class Main {
   public static void main(String[] args) {
      System.out.println("Hello World");
   }
}
```

In this simple Java code:

```
main is the caller,

System.out.println is the callee.

Quite easy, right?:)
```

## Leaf Procedure Example

• C code:

```
long long int leaf_example (
long long int g, long long int h,
long long int i, long long int j) {
  long long int f;
  f = (g + h) - (i + j);
  return f;
}

By convention, we only need to store x20 (callee).
  Saving x5, x6 is just for the purpose of example.
```

- Arguments *g*, *h*, *i*, *j* in x10, ..., x13, *f* in x20
- Temporaries x5, x6
- Need to save x5, x6, x20 on stack

# Leaf Procedure Example

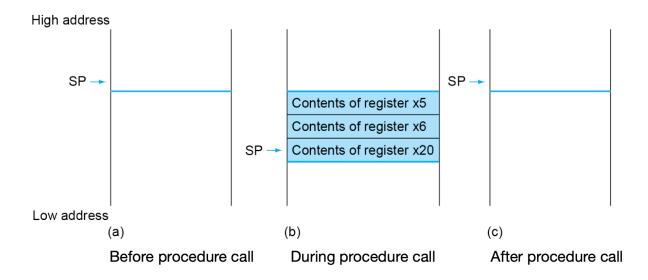
Assume 64-bit (8 bytes) memory.

RISC-V code:

```
leaf example:
                       // label of the procedure
  addi sp,sp,-24
                       // adjust stack to make room for 3 items
  sd
       x5,16(sp)
                      // save register x5
       x6,8(sp)
                       // save register x6
                       // save register x20
  x20,0(sp)
  add x5,x10,x11
                      // register x5 contains g + h
  add x6,x12,x13
                       // register x6 contains i + j
                       // f = x5 - x6
  x20,x5,x6
                       // \text{ returns f } (x10 = x20 + 0)
  addi x10,x20,0
  1d \times 20,0(sp)
                       // restore register x20 for caller
                      // restore register x6 for caller
  1d \times 6.8(sp)
  ld x5,16(sp)
                      // restore register x5 for caller
  addi sp,sp,24
                       // adjust stack to delete 3 items
                      // branch back to calling routine
  jalr x0,0(x1)
```

(Note: Leaf procedure --- procedure that doesn't call another precedure in its body. Non-leaf procedure -- procedure that calls another precedure in its body.)

## **Local Data on the Stack**



The values of the stack pointer and the stack

(More example: Textbook P98 - P102)

### **Procedure Call Instructions**

another procedure

• Procedure call jump and link

jal x1, ProcedureLabel

- Address of following instruction put in x1
- Jumps to target address

to your caller

Procedure return jump and link register

jalr x0, 0(x1)

- Like jal, but jumps to 0 + address in x1
- Use x0 as rd (x0 cannot be changed)
- Can also be used for computed jumps e.g., for case/ switch statements

One word: when you call another procedure, use jal; when you want to return to your caller (the procedure which calls you), use jalr.

### **Calling Convention**

RISC-V Calling Convention			
Register	ABI Name	Saver	Description
x0	zero		Hard-wired zero
x1	ra	Caller	Return address
x2	sp	Callee	Stack pointer
<b>x</b> 3	gp		Global pointer
x4	tp		Thread pointer
x5-7	t0-2	Caller	Temporaries
x8	s0/fp	Callee	Saved register/frame pointer
x9	s1	Callee	Saved register
x10-11	a0-1	Caller	Function arguments/return values
x12-17	a2-7	Caller	Function arguments
x18-27	s2-11	Callee	Saved registers
x28-31	t3-t6	Caller	Temporaries

In another word, when you are writing RISC-V code:

If that register is callee saved by convention, we should make sure that register's value not changed after your procedure. (Solution: use a stack to keep its original value at the beginning of our procedure and restore the original value at the end of our code.)

If that register is caller saved by convention, NO need to do anything. Because it's NOT our (callee) responsibility to keep that's register's value not changed. However, that also means when you try to call another procedure and your data is stored on a caller saved register (this time, you are the caller), your original data could be lost after calling, because it's NOT callee's responsibility to maintain this register's value. (Solution: store your data in callee saved registers / use a stack to store your data before calling another procedure).

e.g.

addi x10, x6, 1 // x10 = x6 + 1 jal x1, g // call function g what's in x10? God knows!!!