
Question 1) 2.27

Translate the following loop into C. Assume that the C-level integer \pm is held in register \pm 6, \pm 5 holds the C-level integer called result, and \pm 10 holds the base address of the integer

```
.data
                         // array of 100 words
MemArray:
                   0
          .word
          .word 1
// ...
          .word 99
.text
     addi x6, x0, 0
     addi x5, x0, 0
     addi x29, x0, 100
LOOP: lw x7, 0(x10)
     add x5, x5, x7
     addi x10, x10, 4
     addi x6, x6, 1
     blt x6, x29, LOOP
Solution:
int i;
int result = 0;
for (i = 0; i < 100; i++) {
     result += *MemArray;
     MemArray++;
}
return result;
// Many people would write the code this way
int i;
int result = 0;
for (i = 0; i < 100; i++) {
     result += MemArray[i];
return result;
```

Question 2) 2.31

jalr x0, 0(x1)

Translate function f into RISC-V assembly language. Assume the function declaration for $\ensuremath{\sigma}$ is

```
int g(int a, int b).
The code for function f is as follows:
int f(int a, int b, int c, int d){
                            return g(g(a,b), c+d);
 }
Solution:
f:
                            addi x2, x2, -8 // Allocate stack space for 2 words
                            sw x1, 0(x2)
                                                                                                                                        // Save return address
                            add x5, x12, x13 // x5 = c+d
                                                                                                                         // Save c+d on the stack
                            sw x5, 4(x2)
                                                                                                                                          // Call x10 = g(a,b)
                            jal x1, g
                            lw x11, 4(x2) // Reload x11= c+d from the stack \frac{1}{2} \frac{1}{2
                            jal x1, q
                                                                                                                                         // Call x10 = g(g(a,b), c+d)
                                                                                                                           // Restore return address
                            lw
                                                        x1, 0(x2)
                            addi x2, x2, 8
                                                                                                                                        // Restore stack pointer
```

Question 3) 2.37

Write the RISC-V assembly code to implement the following C code as an atomic "set max" operation using the lr.d and sc.d instructions. Here, the argument shvar contains the address of a shared variable which should be replaced by y if y is greater than the value it points to. Assume x10 is address of integer pointed by shvar and value of y is in x11.

```
void setmax(int* shvar, int y) {
    // Begin critical section
    if (y > *shvar)
        *shvar = y;
    // End critical section}
}
```

Solution:

```
setmax:
    try:
    lr.d x5, (x10)  # Load-reserve *shvar
    bge x5, x11, release # Skip update if *shvar > y
    addi x5, x11, 0

release:
    sc.d x7, x5, (x10)
    bne x7, x0, try # If store-conditional failed, try again
    jalr x0, 0(x1)
```