

APPENDIX A

OVERVIEW ON THE INDUCTION PHASE

Notice that $\text{suf}(i) < \text{suf}(j)$ if (1) $x[i] < x[j]$ or (2) $x[i] = x[j]$ and $\text{suf}(i+1) < \text{suf}(j+1)$; otherwise, $\text{suf}(i) > \text{suf}(j)$. This observation is utilized by the IS algorithms to sort suffixes as follows:

- S1 Clear S-type sub-buckets in sa . Scan sa^* leftward and insert each element into the current rightmost empty position in the corresponding S-type sub-bucket.
- S2 Clear L-type sub-buckets in sa and insert $n-1$ into the leftmost position in $sa_bkt_L(x[n-1])$. Scan sa rightward with i increasing from 0 to $n-1$. For each scanned non-empty $sa[i]$ with $t[sa[i]-1] = 0$, insert $sa[i]-1$ into the current leftmost empty position in $sa_bkt_L(x[sa[i]-1])$.
- S3 Clear S-type sub-buckets in sa . Scan sa leftward with i decreasing from $n-1$ to 0. For each scanned non-empty $sa[i]$ with $t[sa[i]-1] = 1$, insert $sa[i]-1$ into the current rightmost empty position in $sa_bkt_S(x[sa[i]-1])$.

In brief, given sa^* , S1 inserts all the S*-type suffixes into sa in their lexical order. Then, S2-S3 induce the order of L- and S-type suffixes from those already sorted in sa , respectively, where the relative order of two suffixes induced into the same sub-bucket matches their insertion order according to the rule stated above. To be more specific, we show in Fig. 4 a running example of the induction phase.

As depicted, the input string x contains 6 S*-type suffixes sorted in line 3. When finished inserting the S*-type suffixes in lines 5-6, we first find the head of each L-type sub-bucket (marked by the symbol \wedge) and insert $\text{suf}(13)$ into sa . Notice that $\text{suf}(13)$ consists of only one character; it must be the smallest L-type suffixes starting with 1. Thus, we put $\text{suf}(13)$ into the leftmost position in $sa_bkt_L(1)$ in line 8. Then, we scan sa from left to right for inducing the order of all the L-type suffixes. In lines 10-11, when visiting $sa[0] = 13$ (marked by the symbol $@$), we check the type array t to find $x[12] = 2$ is L-type and hence insert $\text{suf}(12)$ into the current leftmost empty position in $sa_bkt_L(2)$. Similarly, in lines 12-13, we visit the next scanned item $sa[1] = 11$ and see that $t[10] = 0$, thus we place $\text{suf}(10)$ into the current head of $sa_bkt_L(3)$. Following this way, we get all the L-type suffixes sorted in sa . After that, we first find the end of each S-type sub-bucket in lines 25-26 and scan sa leftward for inducing the order of all the S-type suffixes in lines 27-40. When visiting $sa[13] = 2$, we see $x[1]$ is S-type and thus put $\text{suf}(1)$ into the current rightmost empty position in $sa_bkt_S(1)$. Then, at $sa[12] = 8$, we see $x[7] = 1$ is S-type and thus put $\text{suf}(7)$ into the current rightmost empty position in $sa_bkt_S(1)$. To repeat scanning sa in this way, we get all the S-type suffixes sorted in sa .

The work in [14] describes how to compute the LCP array during the execution of S2-S3. Given two suffixes placed at the neighboring positions in sa , their LCP-value can be computed according to one of the following two cases in respect to whether or not they are inserted into the same sub-bucket: if yes, then their LCP-value is one greater than that of the two suffixes from which inducing them; otherwise, their LCP-value equals to zero. In this way, we can determine $\text{lcp}[i]$ immediately after the computation of $sa[i]$. The problem here is how to obtain the LCP-values of these inducing suffixes starting at the next positions in x ,

| | | | | | | | | | | | | | | | |
|----|--|------------|----|----|----|----|----|-----|------------|----|-----|----------|----|------------|-----|
| 00 | p : | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 |
| 01 | $x[p]$: | 2 | 1 | 3 | 1 | 3 | 1 | 2 | 1 | 3 | 1 | 3 | 1 | 2 | 1 |
| 02 | $t[p]$: | L | S* | L | S* | L | S* | L | S* | L | S* | L | S* | L | L |
| 03 | $sa^*[p]$: | 11 | 5 | 9 | 3 | 7 | 1 | | | | | | | | |
| 04 | Insert the sorted S*-type suffixes into sa^* : | | | | | | | | | | | | | | |
| 05 | bucket: | | | 1 | | | | | | 2 | | | 3 | | |
| 06 | $sa^*[p]$: | {-1 | 11 | 5 | 9 | 3 | 7 | 1} | {-1 | -1 | -1} | {-1 | -1 | -1 | -1} |
| 07 | Sort L-type suffixes: | | | | | | | | | | | | | | |
| 08 | $sa^*[p]$: | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {-1 | -1 | -1} | {-1 | -1 | -1 | -1} |
| 09 | | \wedge | | | | | | | \wedge | | | \wedge | | | |
| 10 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | -1 | -1} | {-1 | -1 | -1 | -1} |
| 11 | | @ \wedge | | | | | | | \wedge | | | \wedge | | | |
| 12 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | -1 | -1} | {10 | -1 | -1 | -1} |
| 13 | | \wedge | @ | | | | | | \wedge | | | \wedge | | | |
| 14 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | -1 | -1} | {10 | 4 | -1 | -1} |
| 15 | | \wedge | @ | | | | | | \wedge | | | \wedge | | | |
| 16 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | -1 | -1} | {10 | 4 | 8 | -1} |
| 17 | | \wedge | @ | | | | | | \wedge | | | \wedge | | | |
| 18 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | -1 | -1} | {10 | 4 | 8 | 2} |
| 19 | | \wedge | @ | | | | | | \wedge | | | \wedge | | | |
| 20 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | 6 | -1} | {10 | 4 | 8 | 2} |
| 21 | | \wedge | @ | | | | | | \wedge | | | \wedge | | | |
| 22 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 23 | | \wedge | @ | | | | | | \wedge | | | \wedge | | | |
| 24 | Sort S-type Suffixes: | | | | | | | | | | | | | | |
| 25 | | {13 | -1 | -1 | -1 | -1 | -1 | -1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 26 | | \wedge | | | | | | | \wedge | | | \wedge | | | |
| 27 | | {13 | -1 | -1 | -1 | -1 | -1 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 28 | | \wedge | | | | | | | \wedge | | | \wedge | | @ \wedge | |
| 29 | | {13 | -1 | -1 | -1 | -1 | 7 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 30 | | \wedge | | | | | | | \wedge | | | \wedge | | @ \wedge | |
| 31 | | {13 | -1 | -1 | -1 | 3 | 7 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 32 | | \wedge | | | | | | | \wedge | | | \wedge | | @ \wedge | |
| 33 | | {13 | -1 | -1 | 9 | 3 | 7 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 34 | | \wedge | | | | | | | \wedge | | | \wedge | | @ \wedge | |
| 35 | | {13 | -1 | -1 | 9 | 3 | 7 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 36 | | \wedge | | | | | | | \wedge | | | \wedge | | @ \wedge | |
| 37 | | {13 | -1 | 5 | 9 | 3 | 7 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 38 | | \wedge | | | | | | | @ \wedge | | | \wedge | | @ \wedge | |
| 39 | | {13 | 11 | 5 | 9 | 3 | 7 | 1} | {12 | 6 | 0} | {10 | 4 | 8 | 2} |
| 40 | | \wedge | | | | | | | @ \wedge | | | \wedge | | @ \wedge | |

Fig. 4. An Example for inducing the suffix and LCP arrays.

which is modeled as a range minimum query in [14] and can be answered within amortized $\mathcal{O}(1)$ time. For example, when scanning $sa[0]$ and $sa[5]$ in lines 10-11 and 20-21 of Fig. 4, $\text{suf}(12)$ and $\text{suf}(6)$ are sequentially induced into the neighboring positions in $sa_bkt_L(2)$. In the meantime, if we keep recording the minimum over $\text{lcp}(0, 5]$, then we can obtain the LCP-value of the inducing suffixes $\text{suf}(13)$ and $\text{suf}(7)$ when putting the induced suffix $\text{suf}(6)$ into sa .