

Quantized Tensor Train

Entanglement analysis of encoding functions in quantum states

UC Berkeley / Lawrence Berkeley Laboratory Applied Mathematics Seminar

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Chris (Jielun) Chen

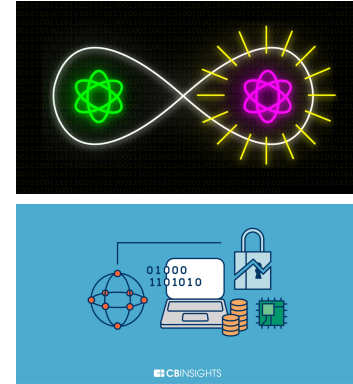
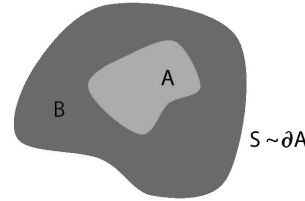


Caltech

Entanglement is one of the most important feature in quantum physics. (2022 Nobel Prize in Physics)

Often studied in the context of:

- Many-body physics
- Quantum cryptography
- Information theory
- ...



These topics often consider physical (non-)locality:

$$|...010...\rangle \quad |...111...\rangle$$

However, in quantum algorithms, physical qubits store information “digitally”:

$$\sum f(x)|x\rangle$$

No notion of
physical locality!

$$|x\rangle = |000\dots 00\rangle, \dots |011\dots 11\rangle, |100\dots 00\rangle, \dots$$

What happens to the entanglement there?

i.e. what properties of f affect the state's entanglement?

$\sum f(x)|x\rangle$ What properties of f affect the state's entanglement?

Answering this question will help:

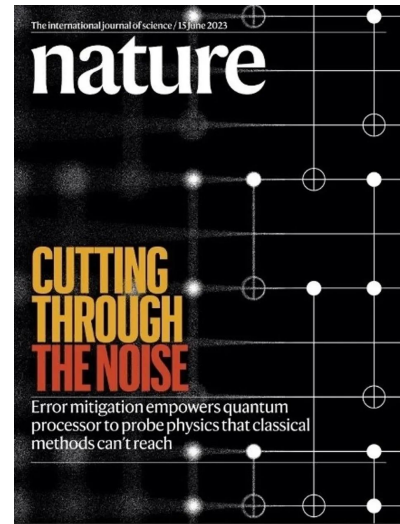
1. Understand when might quantum algorithms beat classical (tensor network) algorithms.
2. Design quantum-inspired classical algorithms.

Good classical
algorithm by just
“simulation”

Classically
doable

Quantum
Advantage?

Entanglement



**Quantized Tensor
Train (QTT)**

Quantized Tensor Train

Outline

1. Notations
2. Introduction to QTT
3. Examples of efficient QTT
4. Applications of QTT
5. Summary & Discussion

Focus on QTT representations
of functions & operators,
rather than QTT algorithms

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Notations: bra-ket & entanglement

$$\vec{v} = \begin{pmatrix} v_0 \\ v_1 \\ \vdots \\ v_{N-1} \end{pmatrix} = |v\rangle \quad \langle v| = (|v\rangle)^\dagger = (v_0^* \ v_1^* \ \dots \ v_{N-1}^*)$$

$$|ab\rangle = |a\rangle|b\rangle = |a\rangle \otimes |b\rangle$$



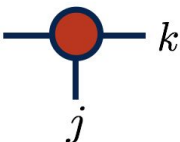
What I meant by “entanglement”:

Schmidt decomposition
= SVD

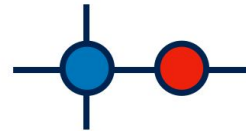
$$|\psi\rangle = \sum_{i=1}^{\chi} \sigma_i |L_i\rangle |R_i\rangle$$

Schmidt rank \approx entanglement
or some measure on how fast singular values decay

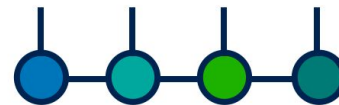
Notations: tensor network

vector	v_j	
matrix	M_{ij}	
3-index tensor	T_{ijk}	

Tensor Contraction



$$= \sum_k T_{ijkl} V_{km}$$



$$= \sum_{\alpha_1, \alpha_2, \alpha_3} A_{\alpha_1}^{s_1} B_{\alpha_1 \alpha_2}^{s_2} C_{\alpha_2 \alpha_3}^{s_3} D_{\alpha_3}^{s_4}$$

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Quantized Tensor Train: **(very brief) history**

Condensed matter physics

1990s: Density Matrix Renormalization Group, Matrix Product State

2000s: Tensor Networks

Independently developed from two communities.

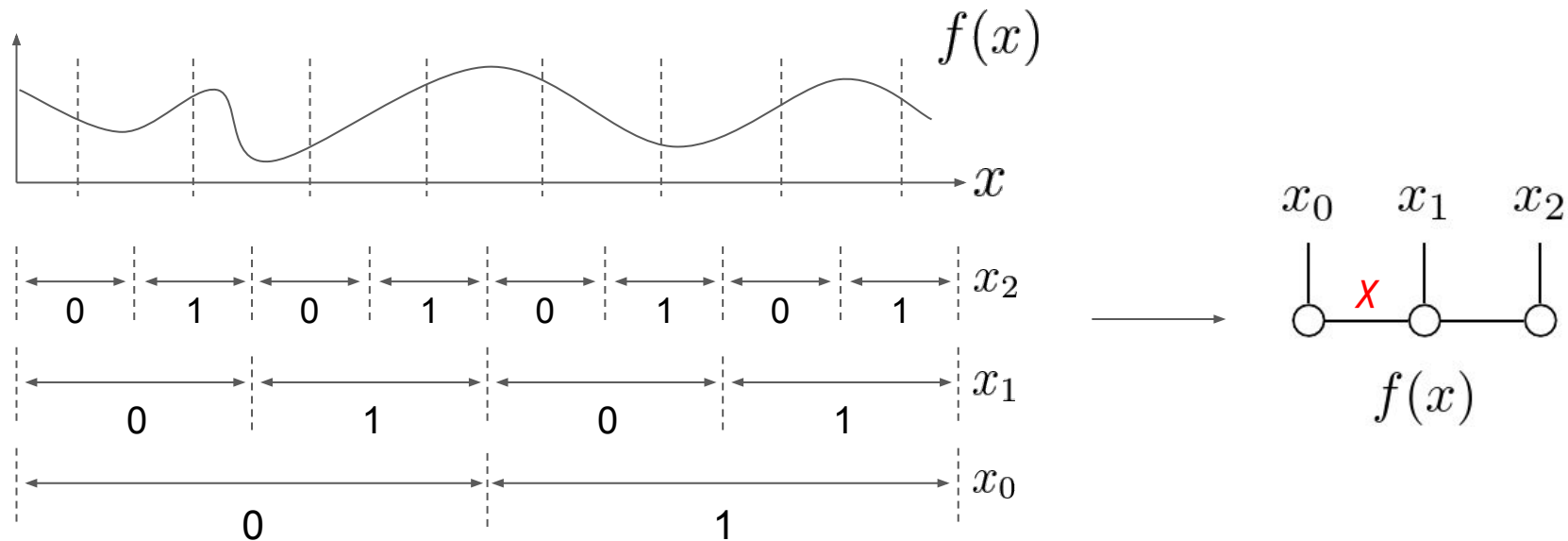
Many things about TT/QTT already known by physicists, but also many new things from math perspective!

Applied Math

2009: Tensor Train

2010s: Quantized Tensor Train

Quantized Tensor Train for **functions**



Efficient criteria: bond-dimension $\chi = \text{poly}(n)$ or even $\chi = O(1)$

Space savings: $2^n \rightarrow O(n\chi^2)$

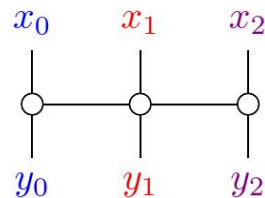
Quantized Tensor Train for **operators & 2D functions**

		y_2	y_1	y_0				
x_2	a_{11}	a_{12}	a_{13}	a_{14}	a_{15}	a_{16}	a_{17}	a_{18}
x_1	a_{21}	a_{22}	a_{23}	a_{24}	a_{25}	a_{26}	a_{27}	a_{28}
x_0	a_{31}	a_{32}	a_{33}	a_{34}	a_{35}	a_{36}	a_{37}	a_{38}
	a_{41}	a_{42}	a_{43}	a_{44}	a_{45}	a_{46}	a_{47}	a_{48}
	a_{51}	a_{52}	a_{53}	a_{54}	a_{55}	a_{56}	a_{57}	a_{58}
	a_{61}	a_{62}	a_{63}	a_{64}	a_{65}	a_{66}	a_{67}	a_{68}
	a_{71}	a_{72}	a_{73}	a_{74}	a_{75}	a_{76}	a_{77}	a_{78}
	a_{81}	a_{82}	a_{83}	a_{84}	a_{85}	a_{86}	a_{87}	a_{88}

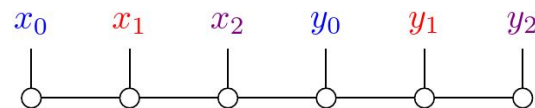
operator



2D function



subtlety with
ordering...



Space savings: $2^{2n} \rightarrow O(n\chi^2)$

Similar scheme for multilinear map & high-dimensional function

Quantized Tensor Train

Outline

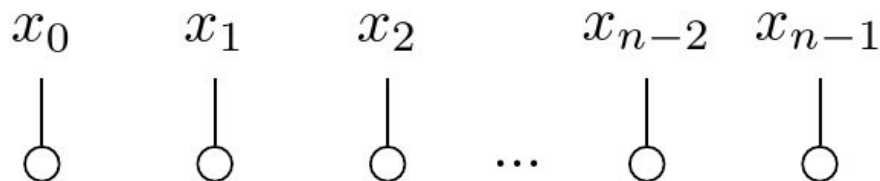
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Efficient QTT functions: $\exp(x)$

$\exp(x)$ is a product of exponentials of individual bits:

$$\begin{aligned} e^x &= e^{x_0 + 2x_1 + 2^2x_2 + \dots} \\ &= \prod e^{2^i x_i} \end{aligned} \longrightarrow \begin{pmatrix} 1 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ e^2 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ e^{2^2} \end{pmatrix} \otimes \dots$$

which corresponds to a $\chi = 1$ QTT:

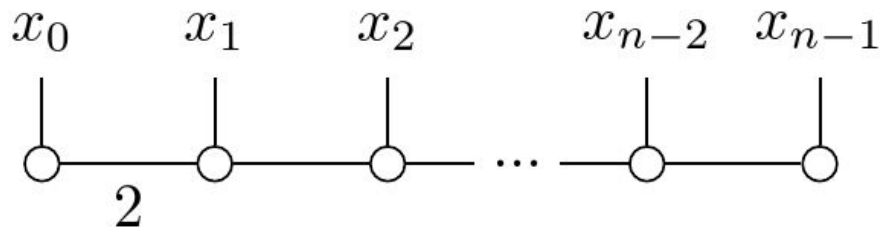


Efficient QTT functions: $\cos(x)$ & $\sin(x)$

$\cos(x)$ & $\sin(x)$ = a sum of two exponentials:

$$\cos(x) = \frac{e^{ix} + e^{-ix}}{2} \quad \sin(x) = \frac{e^{ix} - e^{-ix}}{2i}$$

which corresponds to a $\chi = 2$ QTT: (canonical polyadic decomposition)



Generally:
 $\chi \leq \text{\#Fourier coefficients}$

Efficient QTT functions: **polynomials (1st order)**

First-order polynomial: $x = x_0 + 2x_1 + 2^2x_2 \dots$

Want to construct QTT state $|x^1\rangle$ s.t. $\langle x_0x_1x_2\dots|x^1\rangle = x_0 + 2x_1 + 2^2x_2 \dots$

Solution:

$$|x^1\rangle = \sum_{j=0}^{n-1} |+\rangle^{\otimes j} |v_j^1\rangle |+\rangle^{\otimes n-j-1}$$

$$|v_j^k\rangle = \begin{pmatrix} 0 \\ 2^{jk} \end{pmatrix} \quad |+\rangle = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$\begin{aligned} \langle x_j | v_j^1 \rangle &= 2^j x_j \\ \langle x_j | + \rangle &= 1 \end{aligned}$$

$\chi = 2$ QTT format:

$$\begin{array}{c} | \\ (|+\rangle \quad |v_0^1\rangle) \end{array} \text{ --- } \begin{array}{c} | \\ \left(\begin{array}{c} |+\rangle \\ 0 \end{array} \quad \begin{array}{c} |v_1^1\rangle \\ |+\rangle \end{array} \right) \end{array} \text{ --- } \begin{array}{c} | \\ \left(\begin{array}{c} |+\rangle \\ 0 \end{array} \quad \begin{array}{c} |v_2^1\rangle \\ |+\rangle \end{array} \right) \end{array} \text{ --- } \dots \text{ --- } \begin{array}{c} | \\ \left(\begin{array}{c} |v_{n-1}^1\rangle \\ |+\rangle \end{array} \right) \end{array}$$

Efficient QTT functions: **polynomials (2nd order)**

Second-order polynomial: $x^2 = (x_0 + 2x_1 + 2^2x_2\dots)^2 = \sum x_j x_k 2^{j+k}$

In the first-order example we defined $|v_j^1\rangle$ s.t. $\langle x_j | v_j^1 \rangle = x_j 2^j$

Extending to second order:

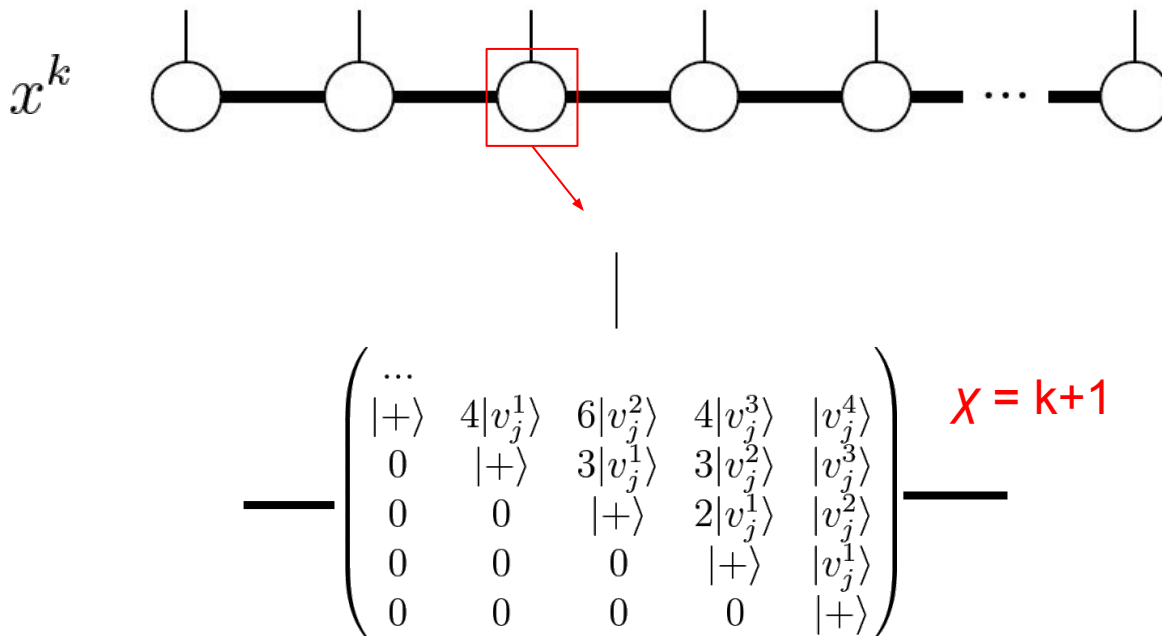
$$|x^2\rangle = 2 \sum_{j < k} |+\rangle^{\otimes j} |v_j^1\rangle |+\rangle^{\otimes k-j-1} |v_k^1\rangle |+\rangle^{\otimes n-k-1} + \sum_j |+\rangle^{\otimes j} |v_j^2\rangle |+\rangle^{\otimes n-j-1}$$

Corresponding to a **$\chi = 3$** QTT format:

$$\begin{array}{c} | \\ (|+\rangle \quad 2|v_0^1\rangle \quad |v_0^2\rangle) \end{array} \text{ --- } \begin{array}{c} | \\ \left(\begin{array}{ccc} |+\rangle & 2|v_1^1\rangle & |v_1^2\rangle \\ 0 & |+\rangle & |v_1^1\rangle \\ 0 & 0 & |+\rangle \end{array} \right) \end{array} \text{ --- } \begin{array}{c} | \\ \left(\begin{array}{ccc} |+\rangle & 2|v_2^1\rangle & |v_2^2\rangle \\ 0 & |+\rangle & |v_2^1\rangle \\ 0 & 0 & |+\rangle \end{array} \right) \end{array} \text{ --- } \dots \text{ --- } \begin{array}{c} | \\ \left(\begin{array}{c} |v_{n-1}^2\rangle \\ |v_{n-1}^1\rangle \\ |+\rangle \end{array} \right) \end{array}$$

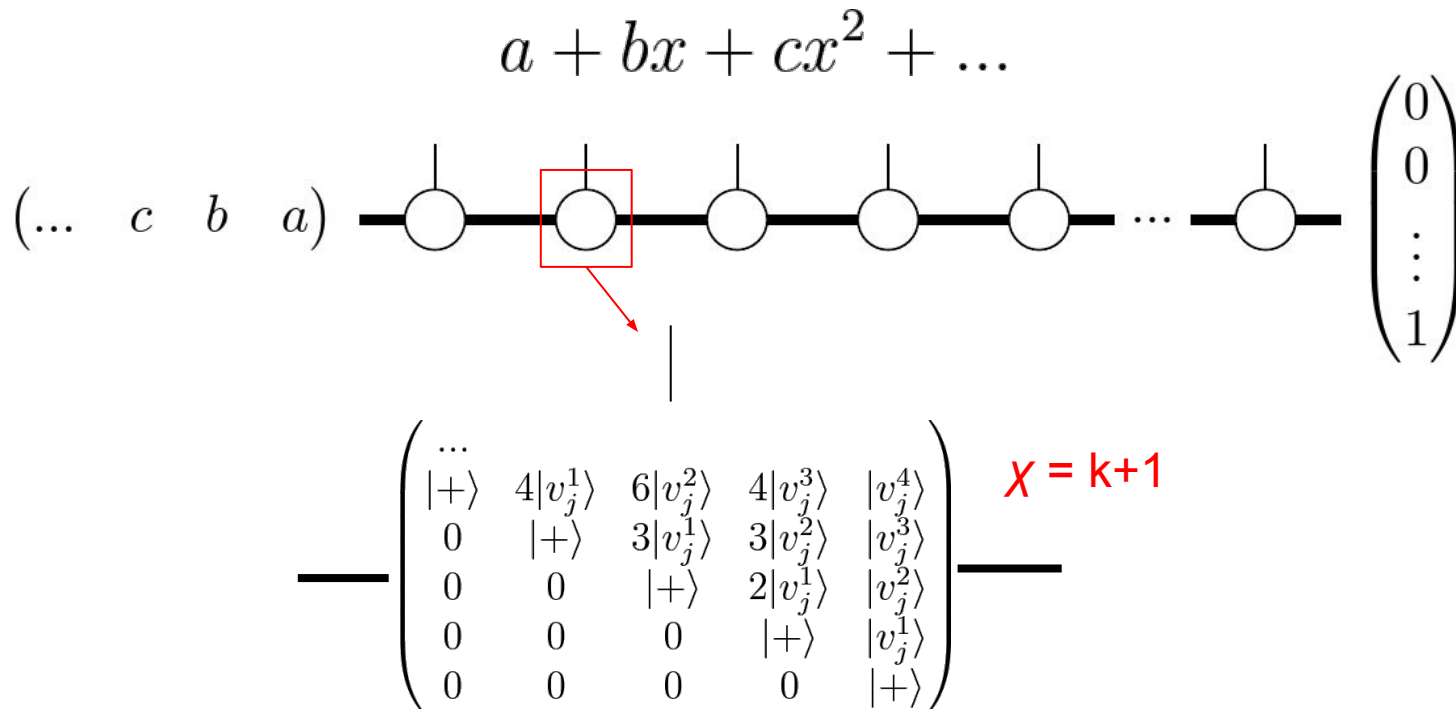
Efficient QTT functions: **polynomials (higher order)**

Extending to higher order:



Efficient QTT functions: **polynomials (general coefficients)**

Boundary tensor determines polynomial coefficients:



Efficient QTT functions: **polynomials (finite state machine)**

Can be viewed as a finite state machine.

Related: Hamiltonian MPO as FSM

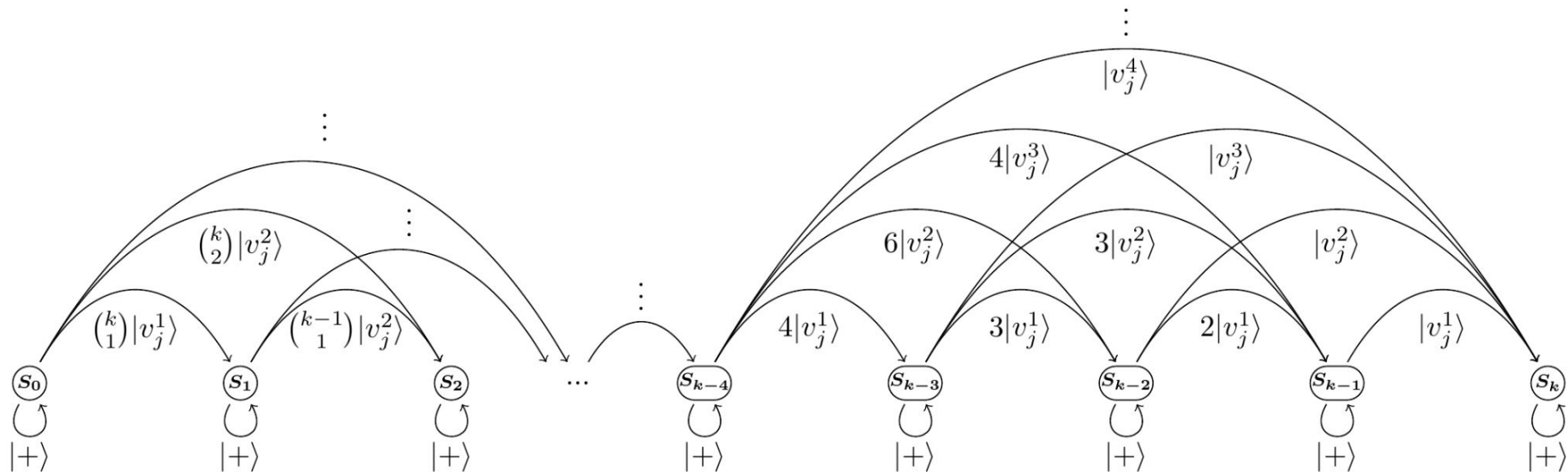
Crosswhite, Bacon, Phys. Rev. A 78, 012356 (2008)

Crosswhite, Doherty, Vidal, Phys. Rev. B 78, 035116 (2008)

Motruk, Zaletel, Mong, Pollmann, Phys. Rev. B 93, 155139 (2016)

...

$$\begin{pmatrix} \dots & & & & \\ |+\rangle & 4|v_j^1\rangle & 6|v_j^2\rangle & 4|v_j^3\rangle & |v_j^4\rangle \\ 0 & |+\rangle & 3|v_j^1\rangle & 3|v_j^2\rangle & |v_j^3\rangle \\ 0 & 0 & |+\rangle & 2|v_j^1\rangle & |v_j^2\rangle \\ 0 & 0 & 0 & |+\rangle & |v_j^1\rangle \\ 0 & 0 & 0 & 0 & |+\rangle \end{pmatrix}$$

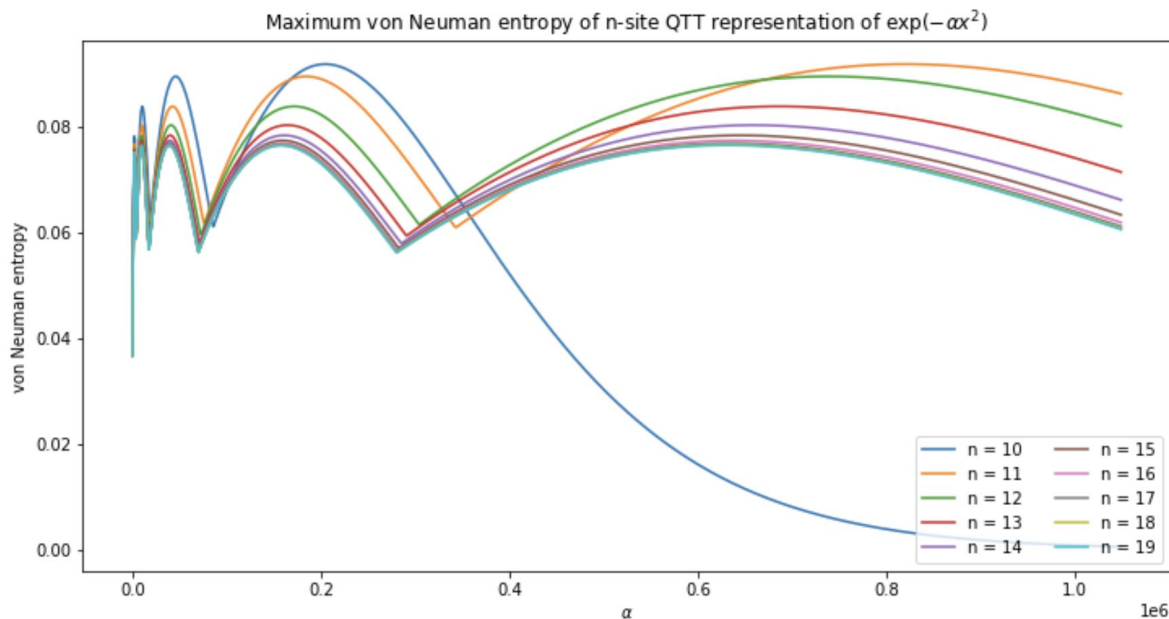


Efficient QTT functions: **Gaussian**

QTT for $e^{-\alpha x^2}$ has error upper-bounded by $\sim O(\chi e^{-\chi^2/\alpha})$

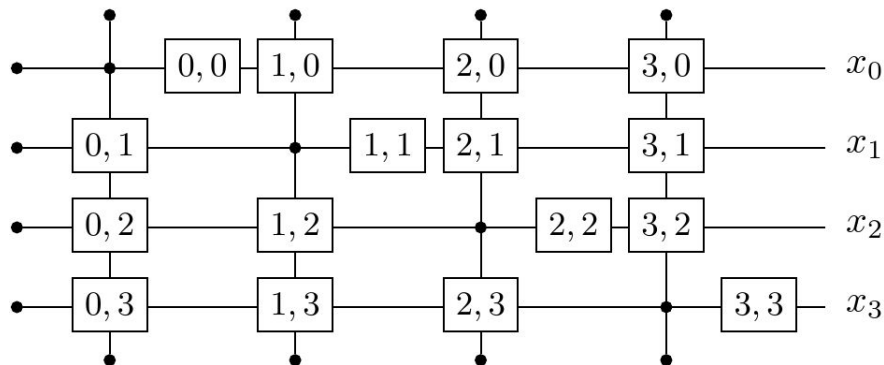
Dolgov, Khoromskij, Oseledets,
SIAM (2012), 34, 6

Numerical experiments showed for almost all α , $\chi = O(1)$



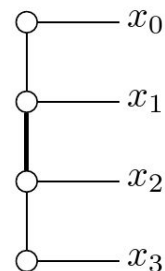
Efficient QTT functions: **Gaussian**

Hard to write out each QTT site, but can contract an $n \times n$ tensor network



$$e^{(x_0 + 2x_1 + 2^2x_2 + 2^3x_3)^2}$$

contract
 \Rightarrow



$$d - \begin{array}{c} a \\ | \\ \boxed{i, j} \\ | \\ c \end{array} - b = \delta_{a,c} \delta_{b,d} e^{ab2^{i+j}}$$

$$a - \boxed{i, j} - b = \delta_{a,b} e^{a2^{i+j}}$$

$$\begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ \cdot \\ \diagup \quad \diagdown \\ c \quad d \\ \vdots \\ e \end{array} = \delta_{a,b,c,d,e,\dots}$$

Quantized Tensor Train

Outline

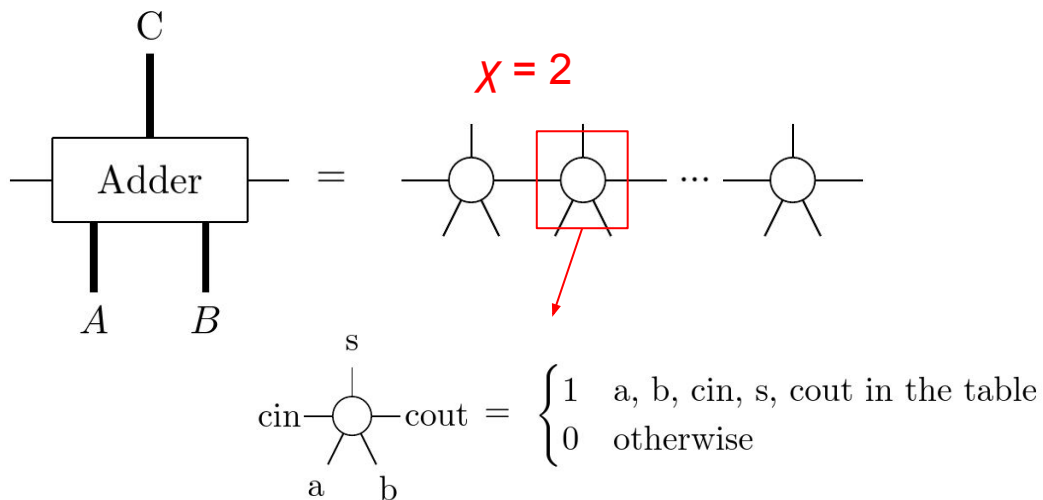
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Efficient QTT operators: **addition**

Addition is defined to be the following linear map:

$$|A\rangle|B\rangle \rightarrow |A + B\rangle$$

The QTT can be obtained directly from a **Ripple-carry adder** circuit:

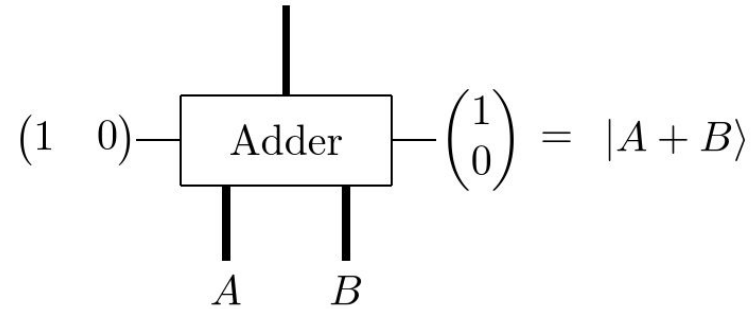


Full adder truth table

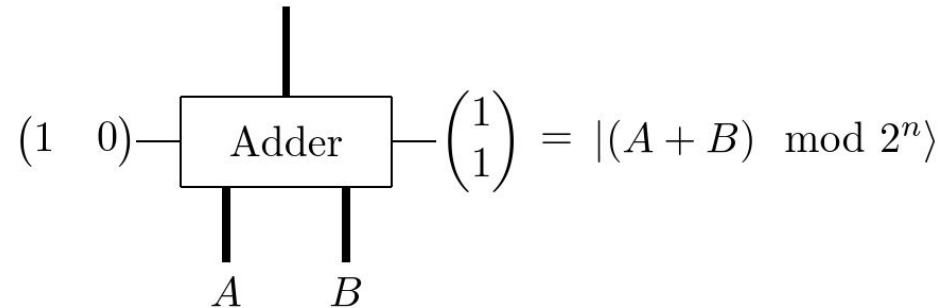
a	b	cin	s	cout
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

Efficient QTT operators: **addition**

Boundary tensors determines modulation:


$$\begin{pmatrix} 1 & 0 \end{pmatrix} \text{---} \text{Adder} \text{---} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = |A + B\rangle$$

Building block for many operators!

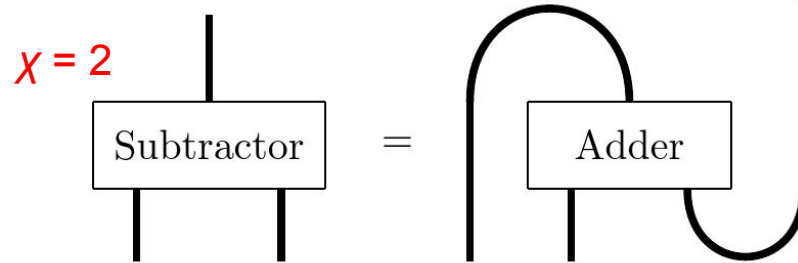

$$\begin{pmatrix} 1 & 0 \end{pmatrix} \text{---} \text{Adder} \text{---} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = |(A + B) \bmod 2^n\rangle$$

Efficient QTT operators: **subtraction**

Subtraction is defined to be the following linear map:

$$|A\rangle|B\rangle \rightarrow |A - B\rangle$$

Subtractor in QTT = **reshaped** adder in QTT:

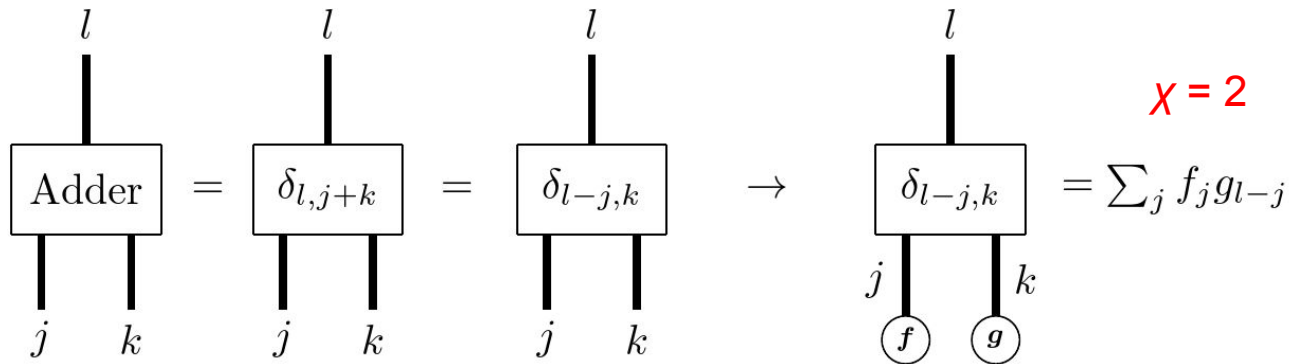


Efficient QTT operators: **convolution**

Convolution is defined as the linear map:

$$\sum_j f_j |j\rangle \otimes \sum_k g_k |k\rangle \rightarrow \sum_l \left(\sum_j f_j g_{l-j} \right) |l\rangle$$

It turns out convolution in QTT = addition in QTT:



circular convolution = **modulo** adder

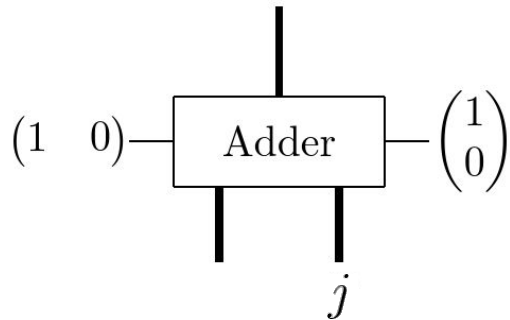
Efficient QTT operators: **shift matrix**

A (non-)circular shift matrix is defined as:

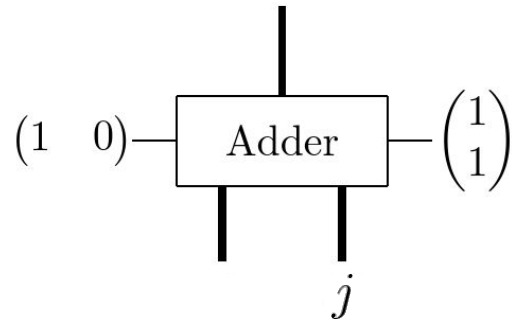
$$j \left\{ \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 \\ 1 & 0 & 0 & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & 0 & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 1 & 0 & \dots & 0 \end{bmatrix} \right.$$

$$j \left\{ \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 0 & 0 & \dots & 1 \\ 1 & 0 & 0 & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & 0 & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 1 & 0 & \dots & 0 \end{bmatrix} \right.$$

Corresponding to adding index by j :



$\chi = 2$



Efficient QTT operators: **Toeplitz matrix**

A Toeplitz matrix has the form:

$$\begin{bmatrix} a_0 & a_{-1} & a_{-2} & \cdots & \cdots & a_{-(n-1)} \\ a_1 & a_0 & a_{-1} & \ddots & & \vdots \\ a_2 & a_1 & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & a_{-1} & a_{-2} \\ \vdots & & \ddots & a_1 & a_0 & a_{-1} \\ a_{n-1} & \cdots & \cdots & a_2 & a_1 & a_0 \end{bmatrix}$$

Appear frequently in signal processing, numerical analysis, differential equations...

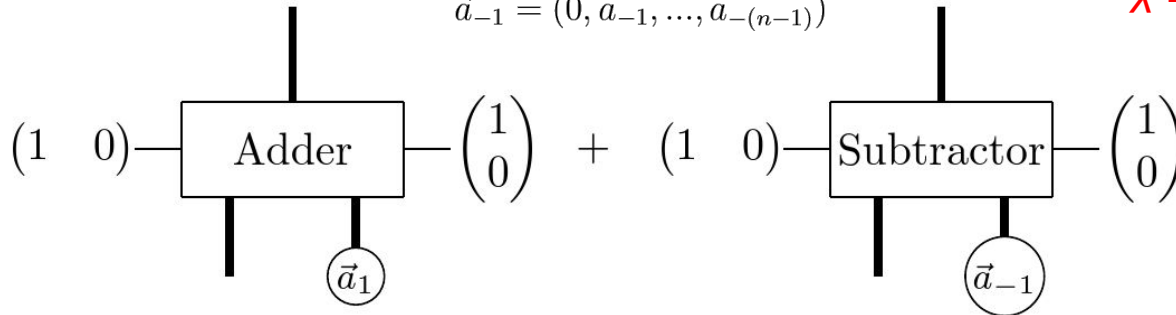
Corresponding to the sum:

$$\vec{a}_1 = (a_0, \dots, a_{n-1})$$

$$\vec{a}_{-1} = (0, a_{-1}, \dots, a_{-(n-1)})$$

$$\chi \leq 2\chi(\vec{a}_1) + 2\chi(\vec{a}_{-1})$$

Small χ for e.g. banded Toeplitz



Efficient QTT operators: **circulant matrix**

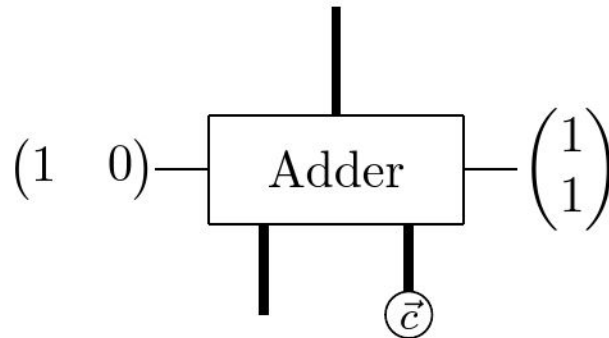
A circulant matrix has the form:

$$\begin{bmatrix} c_0 & c_{n-1} & \cdots & c_2 & c_1 \\ c_1 & c_0 & c_{n-1} & & c_2 \\ \vdots & c_1 & c_0 & \ddots & \vdots \\ c_{n-2} & & \ddots & \ddots & c_{n-1} \\ c_{n-1} & c_{n-2} & \cdots & c_1 & c_0 \end{bmatrix}$$

Special case of Toeplitz

Diagonalizable by discrete
Fourier transform

Corresponding to circular convolution with vector $\vec{c} = (c_0, c_1, \dots, c_{n-1})$:



$$\chi \leq 2\chi(\vec{c})$$

Efficient QTT operators: **discrete Fourier transform**

Discrete Fourier transform (DFT):

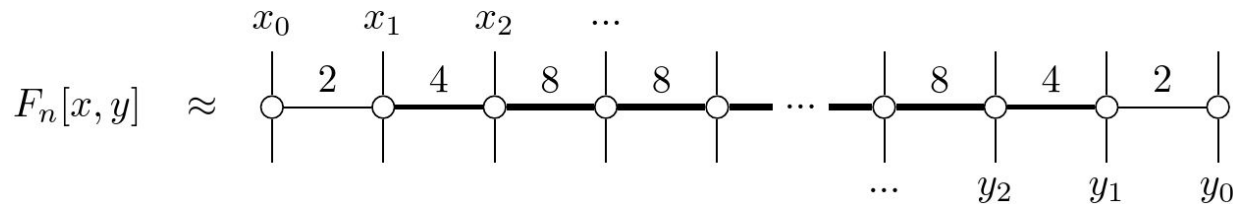
$$F_n = \frac{1}{\sqrt{2^n}} \begin{pmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \omega & \omega^2 & \dots & \omega^{2^n-1} \\ 1 & \omega^2 & \omega^4 & \dots & \omega^{2(2^n-1)} \\ \dots & \dots & \dots & \dots & \dots \\ 1 & \omega^{2^n-1} & \omega^{2(2^n-1)} & \dots & \omega^{(2^n-1)(2^n-1)} \end{pmatrix} \quad \omega = \exp(i2\pi/2^n)$$

DFT is well-approximated by a QTT with error $O(ne^{-\chi \log(\chi/3)})$.

i.e. χ grows **sub-logarithmically** to maintain a constant global error.

Numerics suggest $\chi = 8$ gives error below 10^{-15} .

JC, Stoudenmire, White,
arXiv:2210.08468
(accepted to PRX
quantum)



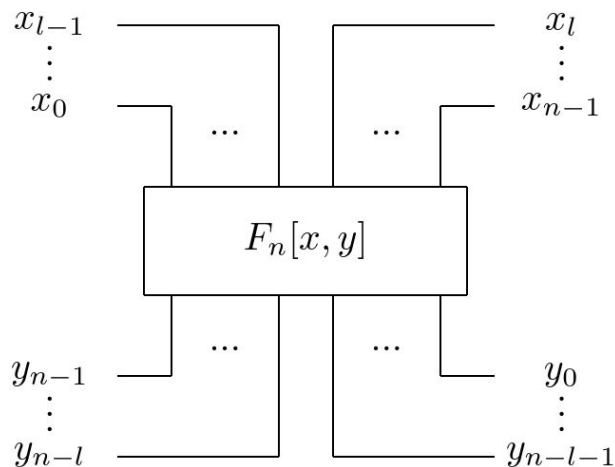
**reversed
ordering is
important!**

Efficient QTT operators: **discrete Fourier transform**

Why is DFT compressible in QTT:

$$F_n^S = 2^l \begin{pmatrix} \overset{2^{n-l}}{\boxed{\begin{matrix} 1 & 1 & 1 & \dots \\ 1 & \omega & \omega^2 & \dots \\ 1 & \omega^2 & \omega^4 & \dots \\ \dots & \dots & \dots & \dots \end{matrix}}} & \begin{matrix} 1 \\ \omega^{2^{n-1}} \\ \omega^{2(2^n-1)} \\ \dots \end{matrix} \\ \begin{matrix} 1 \\ \omega^{2^{n-1}} \\ \omega^{2(2^n-1)} \\ \dots \end{matrix} & \begin{matrix} \omega^{(2^n-1)(2^n-1)} \\ \dots \end{matrix} \end{pmatrix}$$

DFT's QTT rank

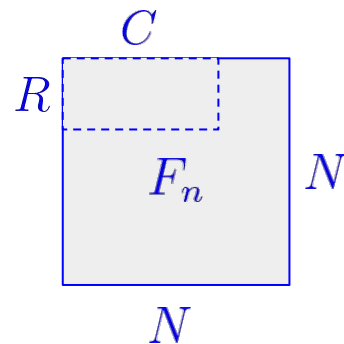


DFT's **submatrix** rank

$$= \underbrace{\begin{bmatrix} V_1 \end{bmatrix}}_{2^{2l} \times 2^l} \underbrace{\begin{bmatrix} F_n^S \end{bmatrix}}_{2^l \times 2^{n-l}} \underbrace{\begin{bmatrix} V_2 \end{bmatrix}}_{2^{n-l} \times 2^{2(n-l)}}$$

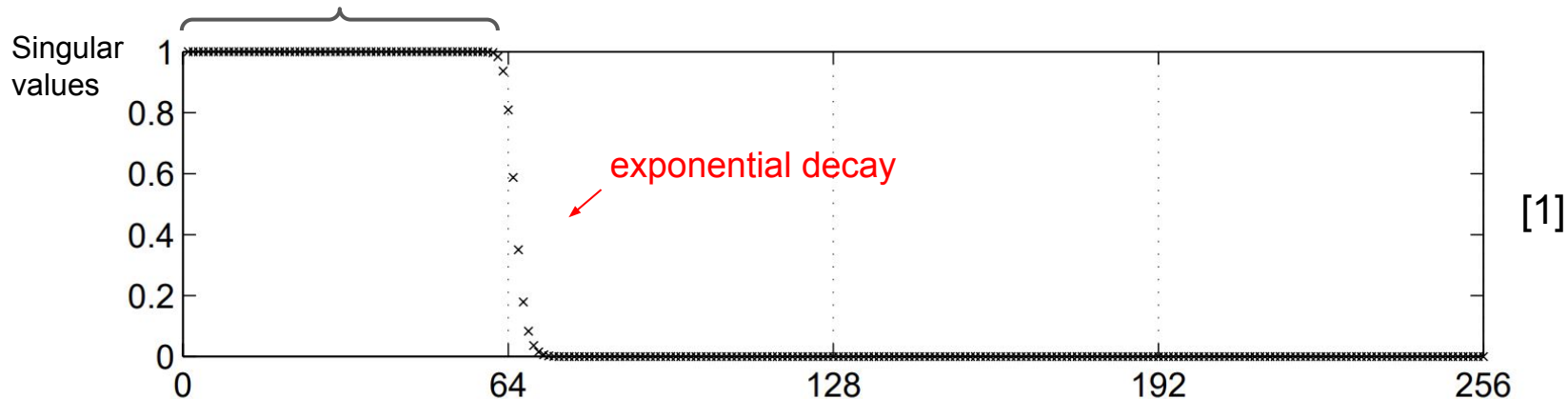
Efficient QTT operators: **discrete Fourier transform**

For an $R \times C$ submatrix of the $N \times N$ DFT, its effective rank is very small.



$$2^l \times 2^{n-l} / 2^n = 1$$

$$\sim RC/N$$



Efficient QTT operators: **derivatives**

Option 1: finite difference method $\chi \leq 2(\text{FDM order} + \text{derivative order})$

$$\frac{\partial^2}{\partial x^2} \sim \begin{pmatrix} -2 & 1 & 0 & \dots & 0 & 1 \\ 1 & -2 & 1 & \dots & 0 & 0 \\ 0 & 1 & -2 & \dots & 0 & 0 \\ \dots & & & & & \\ 1 & 0 & 0 & \dots & 1 & -2 \end{pmatrix} = -2I + \begin{pmatrix} 0 & 0 & 0 & \dots & 0 & 1 \\ 1 & 0 & 0 & \dots & 0 & 0 \\ 0 & 1 & 0 & \dots & 0 & 0 \\ \dots & & & & & \\ 0 & 0 & 0 & \dots & 1 & 0 \end{pmatrix} + \begin{pmatrix} 0 & 1 & 0 & \dots & 0 & 0 \\ 0 & 0 & 1 & \dots & 0 & 0 \\ 0 & 0 & 0 & \dots & 0 & 0 \\ \dots & & & & & \\ 1 & 0 & 0 & \dots & 0 & 0 \end{pmatrix}$$

Option 2: diagonalization by DFT $\chi \leq \chi(\text{DFT})^2 * (\text{derivative order} + 1)$

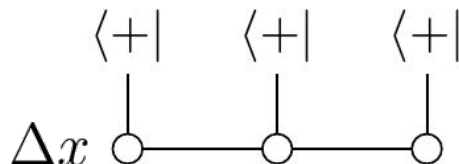
$$\frac{\partial^2}{\partial x^2} \sim \text{DFT}^{-1} \cdot \text{diag}(x^2) \cdot \text{DFT}$$

Efficient QTT operators: **integral**

First order approximation to the integral:

$$\int_{x_0}^{x_N} f(x) dx \approx \sum_{j=0}^{N-1} f(x_j) \Delta x = \Delta_x \langle + |^{\otimes n} | f \rangle$$

Corresponding to inner product with **$\chi = 1$** QTT:

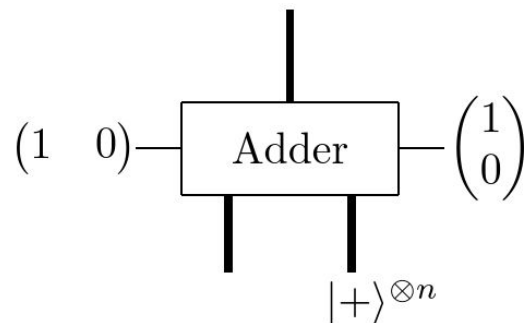


Efficient QTT operators: **integral with variable range**

Integral with variable range:

$$g(x) = \int_0^x f(x') dx \approx \Delta x \begin{bmatrix} 1 & 0 & \dots & 0 & 0 \\ 1 & 1 & \dots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & 1 & \dots & 1 & 0 \\ 1 & 1 & \dots & 1 & 1 \end{bmatrix} \begin{bmatrix} f(x'_0) \\ f(x'_1) \\ \vdots \\ f(x'_{N-2}) \\ f(x'_{N-1}) \end{bmatrix}$$

The matrix corresponds to a **$\chi = 2$** QTT:



Efficient QTT for e.g.

$$\operatorname{erf} z = \frac{2}{\sqrt{\pi}} \int_0^z e^{-t^2} dt.$$

Other efficient QTT in literature

Wavelets as QTT

Oseledets Tyrtysnikov, Algebraic
Wavelet Transform via Quantics
Tensor Train Decomposition

Image Compression with QTT

Latorre, Image compression and entanglement, arXiv:quant-ph/0510031, 2005

Green's functions of quantum many-body systems as QTT

Shinaoka, Wallerberger, Murakami, Nogaki, Sakurai, Werner, and
Kauch, Multiscale Space-Time Ansatz for Correlation Functions
of Quantum Systems Based on Quantics Tensor Trains

...

Major open question: **when is QTT efficient in general?**

smoothness?

QTT can embed both very **smooth** or very **sharp** functions

Uniform distribution Lipschitz-continuous functions

Exponential cusps

Delta-function

Small rank

Small rank

#Fourier coeffs

Smoothness

Outlier: Gaussian?

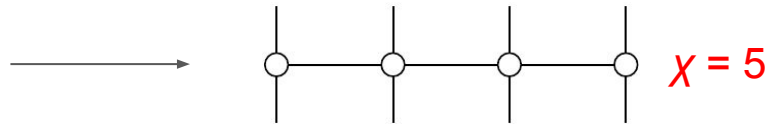
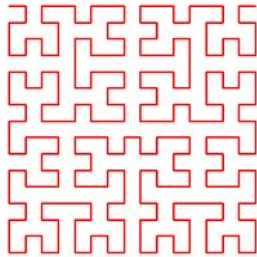
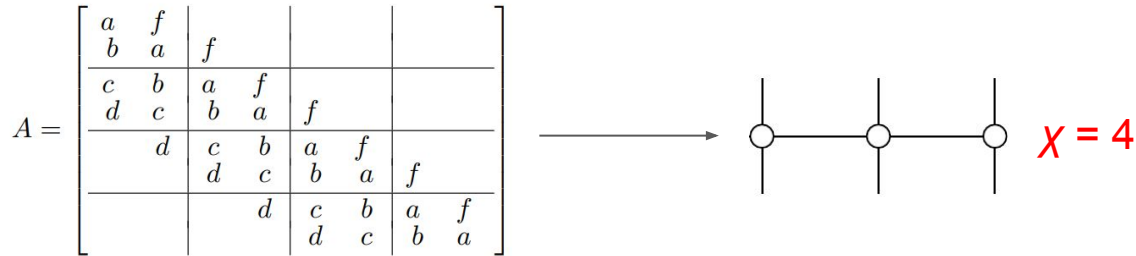
Some rigorous results, e.g. **classical Besov smoothness implies QTT**

Ali & Nouy, Constructive Approximation volume 58, pages 463–544 (2023)

Major open question: when is QTT efficient in general?

Recursion & Fractal structure?

Recursive construction \rightarrow QTT



How to formalize?

Entropy of fractal systems

Zmeskal, Dzik, Vesely, Computers
& Mathematics with Applications,
Volume 66, Issue 2, 2013,

Quantized Tensor Train

Outline

1. Notations
2. Introduction to QTT
3. Examples of efficient QTT
4. Applications of QTT
5. Summary & Discussion

Applications: **plasma physics**

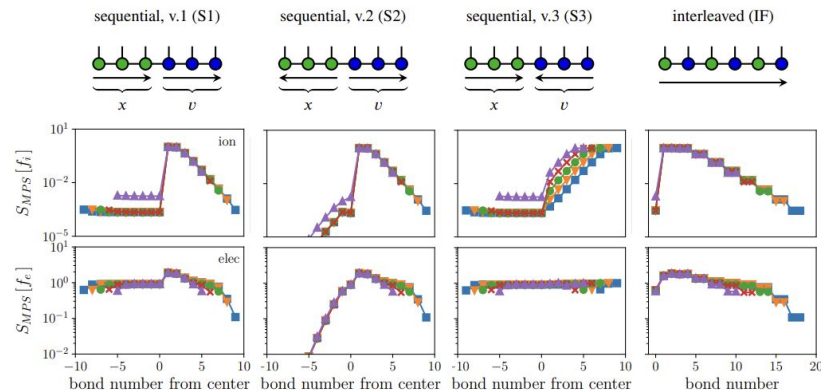
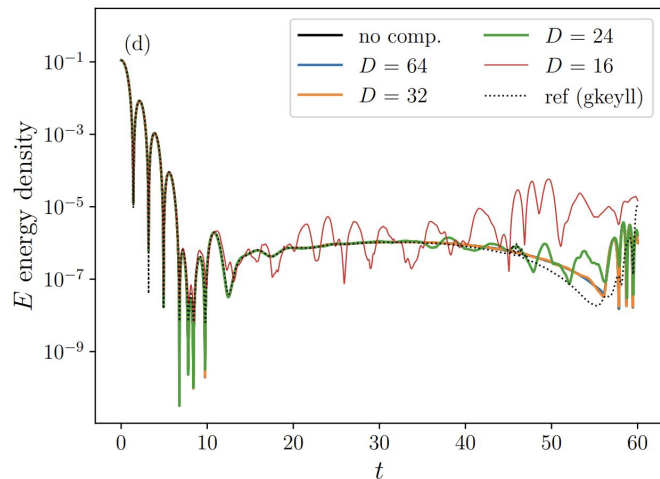
Ye, Loureiro, Phys. Rev. E 106, 035208 (2022)

Solving the Vlasov-Poisson equation by time evolution in QTT:

$$\frac{\partial f_s}{\partial t} + \mathbf{v}_s \cdot \nabla_{\mathbf{r}} f_s + \frac{q_s}{m_s} \mathbf{E} \cdot \nabla_{\mathbf{v},s} f_s = \mathcal{C}[f_s]$$

$$\frac{\partial f}{\partial t} = F(f) = \underbrace{\begin{array}{c} \text{---} \square \text{---} \square \text{---} \square \text{---} \square \text{---} \square \text{---} \\ | \quad | \quad | \quad | \quad | \\ \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \end{array}} + \begin{array}{c} \text{---} \square \text{---} \square \text{---} \square \text{---} \square \text{---} \square \text{---} \\ | \quad | \quad | \quad | \quad | \\ \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \end{array} + \dots$$

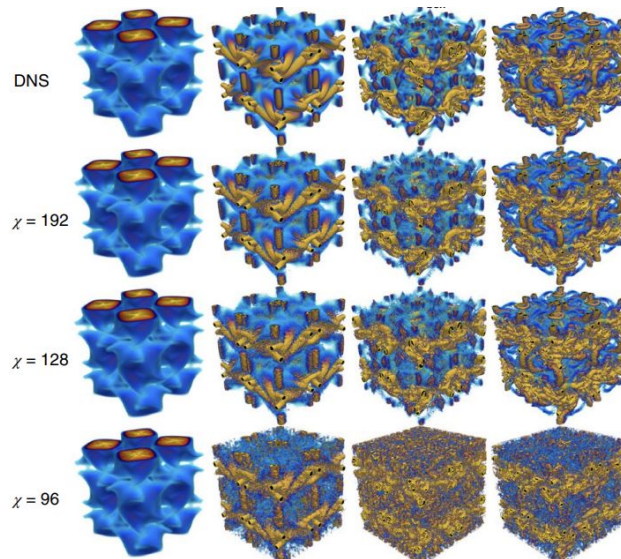
$$f(t + \Delta t) \approx f(t) + \frac{\partial f}{\partial t} \Delta t$$



Solving the incompressible Navier–Stokes equations iteratively in QTT :

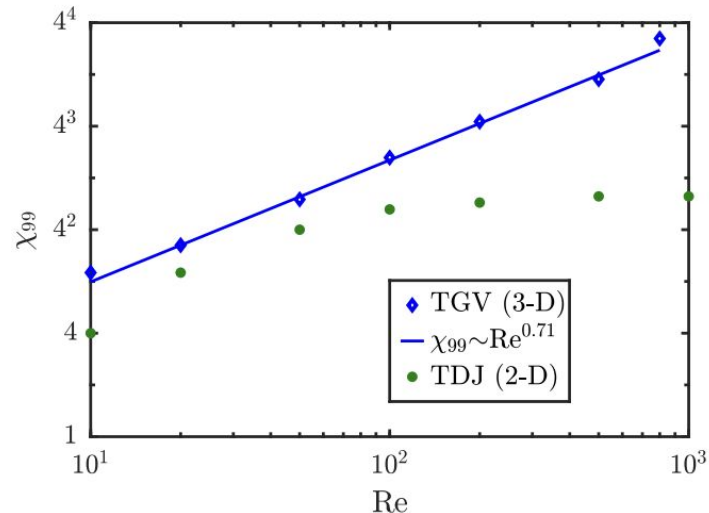
$$\nabla \cdot V = 0$$

$$\frac{\partial V}{\partial t} + (V \cdot \nabla) V = -\nabla p + \nu \nabla^2 V,$$



2D: QTT-rank saturates for Reynold number ≥ 200

3D: QTT-rank increases according to a power law



Applications: **quantum chemistry**

Physical orbitals tend to be smooth \rightarrow efficient QTT

Jolly, Fernández, Waintal, arXiv:2308.03508

Solve Hartree-fock in QTT iteratively using DMRG (minimization):

$$\left(-\frac{1}{2}\nabla^2 + V_{ion} + J[\rho] + K[\{\phi_j\}] \right) \phi_i = \epsilon_i \phi_i$$

$$V_{ion} = \sum_A \frac{Z_A}{|R_A - r|}$$

$$J[\rho] = \int \frac{\rho(r')}{|r - r'|} dr'$$

$$K[\{\phi_j\}]\phi_i = \sum_j \phi_j(r) \int \frac{\phi_j^*(r')\phi_i(r')}{|r - r'|} dr'$$

Work in progress with
Sandeep Sharma &
Garnet Chan

$$\frac{1}{r} \approx \sum_i c_i e^{-\alpha_i r^2}$$

Applications: “superfast” Fourier transform

Assume an input vector v has length $N = 2^n$; want to compute $\text{DFT}(v)$.

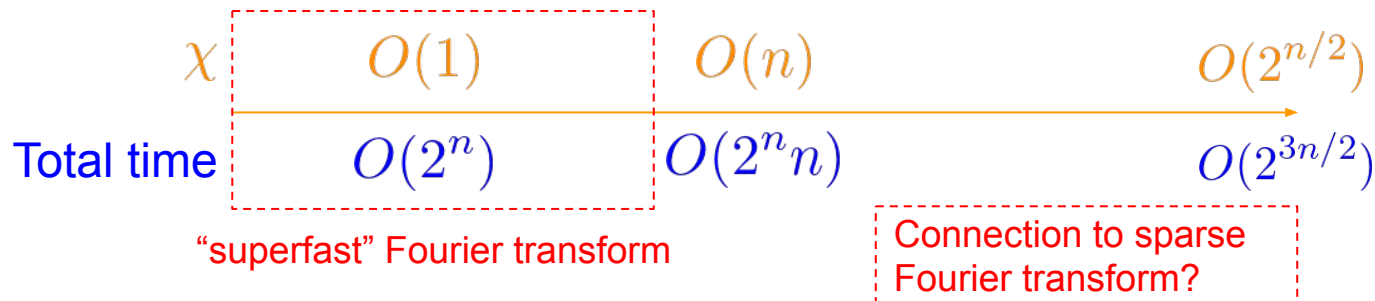
- Time complexity for the fast Fourier transform:

$$O(N \log N) = O(2^n n)$$

dominates time complexity

- Total time complexity for converting v to QTT with rSVD + DFT QTT:

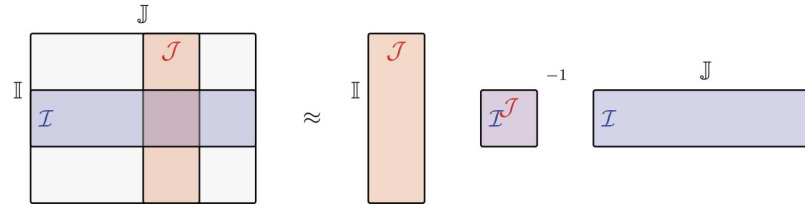
$O(2^n \chi)$ if data can be compressed into an QTT with rank χ



Side note: convert vector into QTT

Converting an exponentially-long vector to QTT takes exponential time with SVD, even when QTT is efficient. What are some other methods?

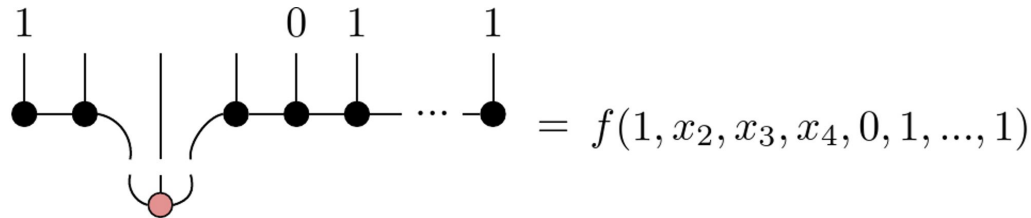
Cross-interpolation (iterate through all cuts):



Dolgov, Savostyanov, Computer Physics Communications, Volume 246, 2020

DMRG-like method:

Initial guess \rightarrow sampling environment \rightarrow solve local LSE \rightarrow sweep



Improve &
Rigorous
guarantee?

Summary & Discussion

- Efficient QTT construction for many important functions & operators
- Formalize efficient criteria for QTT?
- Directly connect to entanglement in quantum algorithms.
- Already been applied to many real-world differential equations.

Acknowledgement

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