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CS331 HW1

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1. The message displays *Fifty fearless firefighters found foliage for friendly florists.*
2. (a) C++ type checking is primarily static.
(b) The difference between the two is that static type checking is done before runtime, whereas dynamic type checking occurs at runtime.
3. String 5 is generated.
4. An equal number of 0 or more x's on either side of 0 or more pairs of concatenated y's.
5. The strings that are matched by the regular expression are strings 3, 4, 5, and 7.
6. $a(a|b)^*$

7. (a) Leftmost Derivation

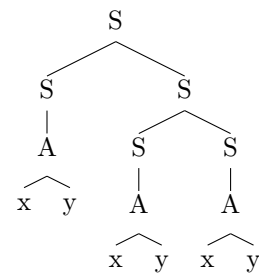
S
SS
AS
xyS
xyA
xyxy

- (b) Rightmost Derivation

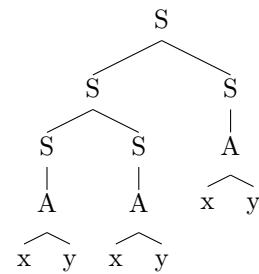
S
SS
SA
Sxy
Axy
xyxy

(c) xyxyxy

i. parse tree 1

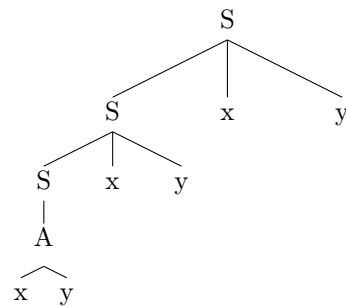


ii. parse tree 2

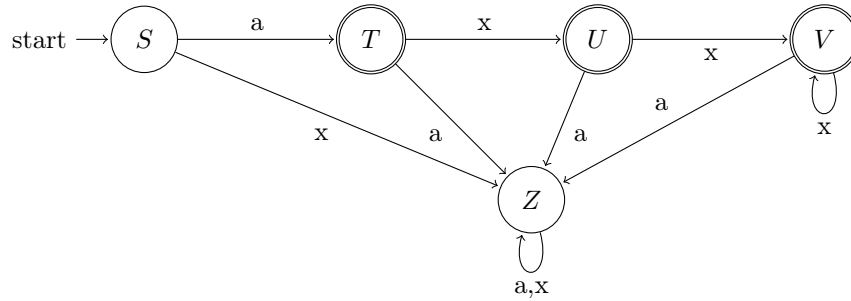


(d) Non-ambiguous grammar

$S \rightarrow Sxy | A$
 $A \rightarrow xy$

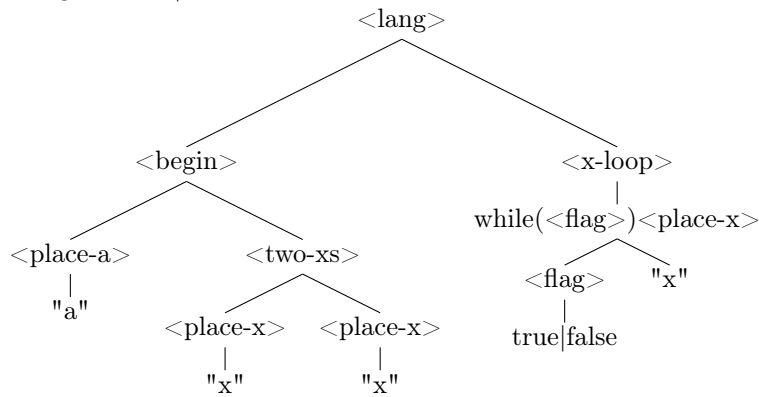


8. (a) $axxx^*$



(b) DFA

9. $\langle \text{lang} \rangle ::= \langle \text{begin} \rangle \langle \text{x-loop} \rangle$
 $\langle \text{begin} \rangle ::= \langle \text{place-a} \rangle \langle \text{two-xs} \rangle$
 $\langle \text{x-loop} \rangle ::= \text{while}(\langle \text{flag} \rangle) \langle \text{place-x} \rangle$
 $\langle \text{two-xs} \rangle ::= \langle \text{place-x} \rangle \langle \text{place-x} \rangle$
 $\langle \text{place-x} \rangle ::= \text{"x"}$
 $\langle \text{place-a} \rangle ::= \text{"a"}$
 $\langle \text{flag} \rangle ::= \text{true} | \text{false}$



10. My grammar is ambiguous; traversing through the grammar yields a single tree that generates the concatenation of one 'a' followed by at least two 'x's, or $axxx^*$.