

# *TOTAL CHORD:* THE COMPLETE SEARCH OVER CHORD SOLUTION

CSCI-GA.2620: NETWORKS AND MOBILE SYSTEMS

# Overview

- ❑ In a Chord peer-to-peer lookup service a user can only search via specific predetermined indexes (keys) omitting information embedded within a file and/or other content
- ❑ Creating an unrestricted content search option over Chord service allows users freedom to:
  - ❑ Review information embedded within any file
  - ❑ Search full P2P content using any keyword
  - ❑ Supply keywords without imposed constraints

# Defining the Problem

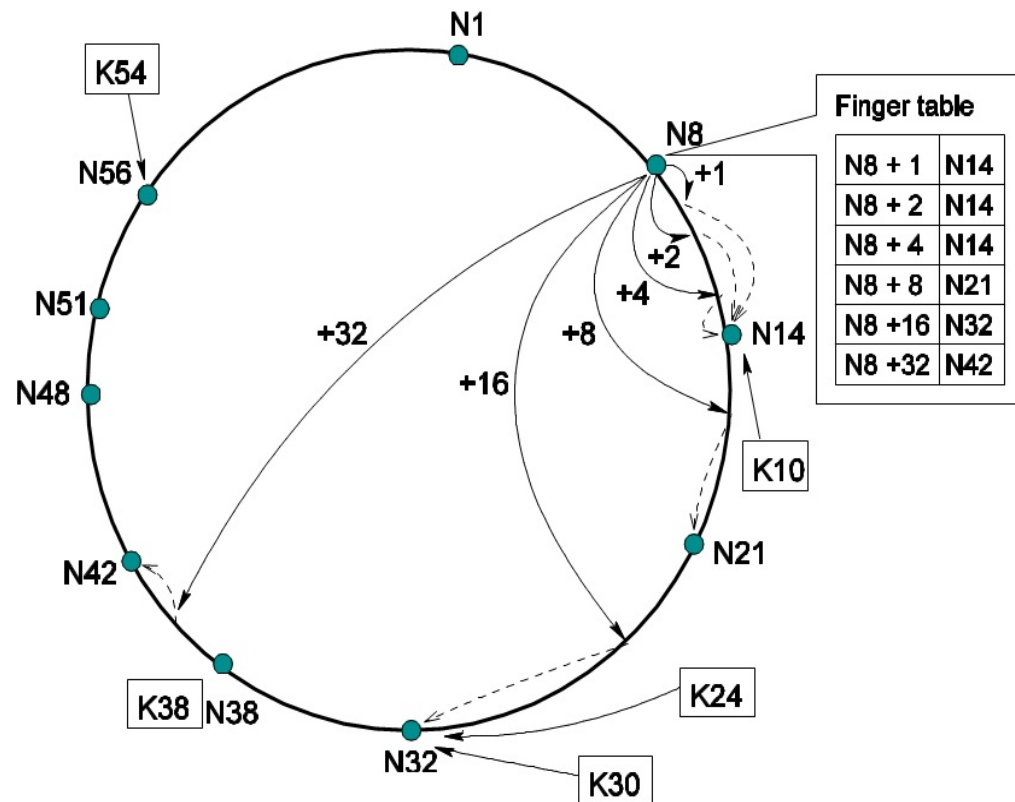
- A Chord P2P service limits a users ability to search for information within files
- Nodes only store (key, value) pairs
  - ▣ Key generated from filename, contents, or any arbitrary index.
  - ▣ The value can be file contents, pointers, simple strings, etc.
  - ▣ In current form Chord service does not provide a method to search the values' content

# The Solution

- Solution is a distributed and parallel search
  - ▣ The search occurs across every node
  - ▣ Each node initiates additional remote searches
  - ▣ Results are cached
  - ▣ Chord strengths (simplicity, scalability, flexibility, availability, load balance) are leveraged for caching and searching
- Implementing this solution will:
  - ▣ Reduce search time for users
  - ▣ Balance load across the network
  - ▣ Allows users to perform any search(es)

# Grasping Chord

- Chord is a distributed lookup service
  - ▣ Supports one operation: given key, it maps the key onto a node
  - ▣ Maps both keys and nodes (node IPs) to the same ID space
  - ▣ Uses consistent hashing:
    - Spreading keys evenly over nodes
    - Node responsible for K/N keys
  - ▣ Lookups needs  $O(\log N)$  messages (finger table)
  - ▣ Allows flexibility of how to map names to Chord keys
  - ▣ Fully distributed



\* Modified Chord figure from "Chord: A Scalable Peer-to-peer Lookup Protocol for Internet Applications", Ion Stoica et al., IEEE/ACM Transactions on Networking, Feb. 2003.

# Implementation

## □ Build on top of Chord solution

```
// ask node n to find the successor of id
n.find_successor(id)
  if (id ∈ (n, successor])
    return successor;
  else
    n' = closest_preceding_node(id);
    return n'.find_successor(id);
```

```
// search the local table for the highest predecessor of id
n.closest_preceding_node(id)
  for i = m downto 1
    if (finger[i] ∈ (n, id))
      return finger[i];
  return n;
```

```
// create a new Chord ring.
n.create()
  predecessor = nil;
  successor = n;
```

```
// join a Chord ring containing node n'.
n.join(n')
  predecessor = nil;
  successor = n'.find_successor(n);
```

```
// called periodically. verifies n's immediate
// successor, and tells the successor about n.
n.stabilize()
  x = successor.predecessor;
  if (x ∈ (n, successor))
    successor = x;
  successor.notify(n);
```

```
// n' thinks it might be our predecessor.
n.notify(n')
  if (predecessor is nil or n' ∈ (predecessor, n))
    predecessor = n';
```

```
// called periodically. refreshes finger table entries.
// next stores the index of the next finger to fix.
```

```
n.fix_fingers()
  next = next + 1;
  if (next > m)
    next = 1;
  finger[next] = find_successor(n + 2next-1);
```

```
// called periodically. checks whether predecessor has failed.
n.check_predecessor()
  if (predecessor has failed)
    predecessor = nil;
```

```
// called to perform search using the finger table over the entire ring
n.search(wallID, keyword)
  for i = m downto 1
    if (finger[i] ∈ (n, wallID))
      finger[i].search(wallID, keyword)
      wallID = finger[i]
  // perform search over data in n and return results
```

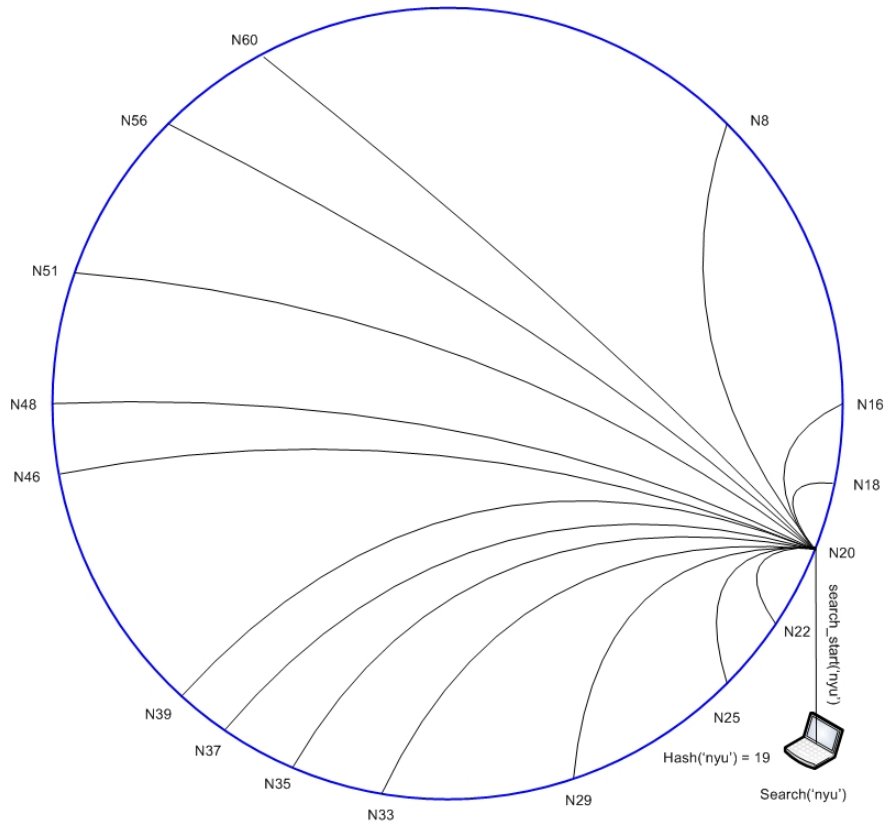
```
// called to initiate search
n.search_start(keyword)
  n.search(n, keyword)
```

```
// called to save keys for particular keyword search
n.cache_put_keys(keyword, keys)
  // hash keyword and save the keys at that location
```

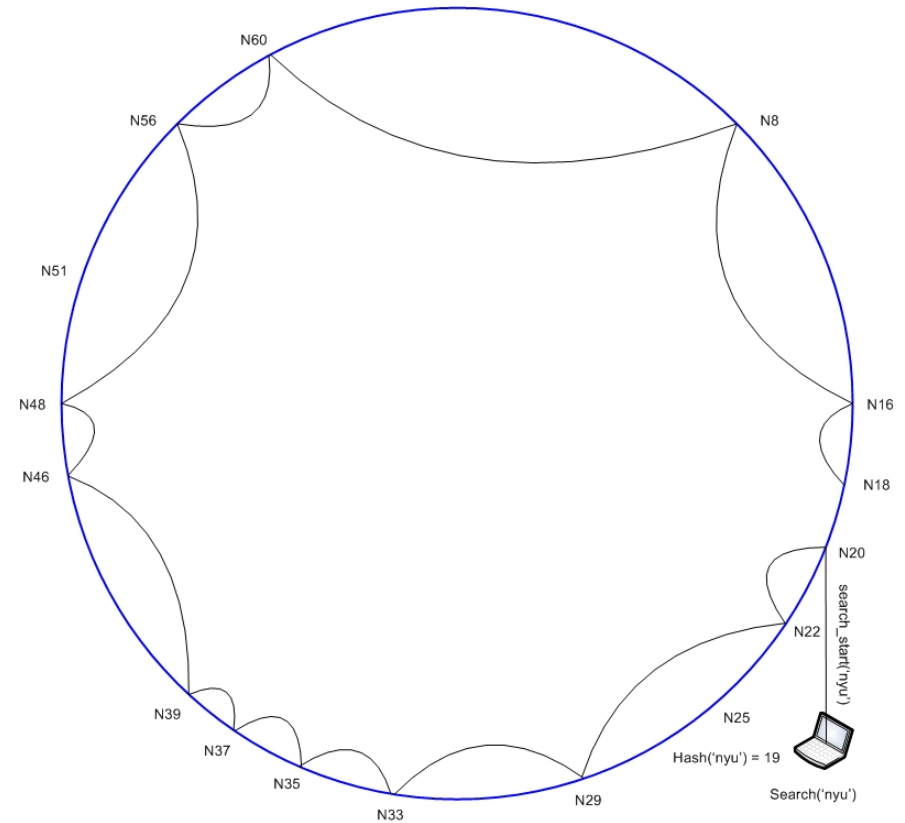
```
// called to retrieve keys for particular keyword
n.cache_get_keys(keyword)
  // hash keyword and return keys at that location for that keyword
```

# Implementation

## Eliminating Alternate Solutions



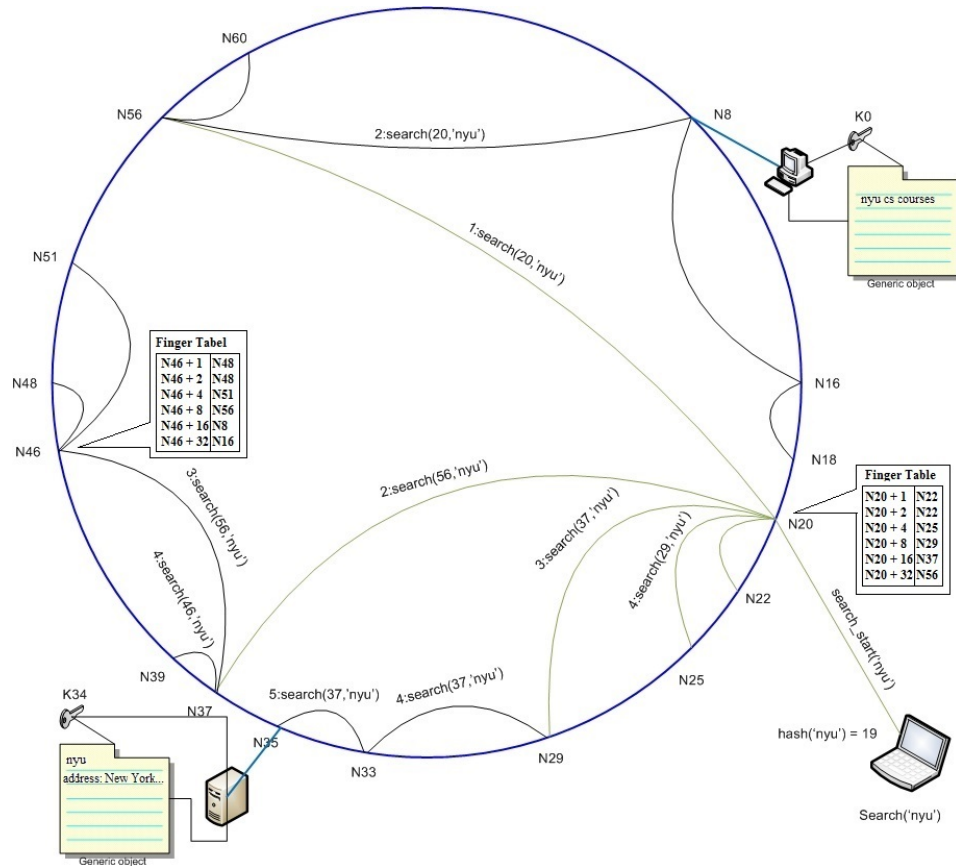
- ❑ **Impractical.** This approach does not scale well with large numbers of nodes.



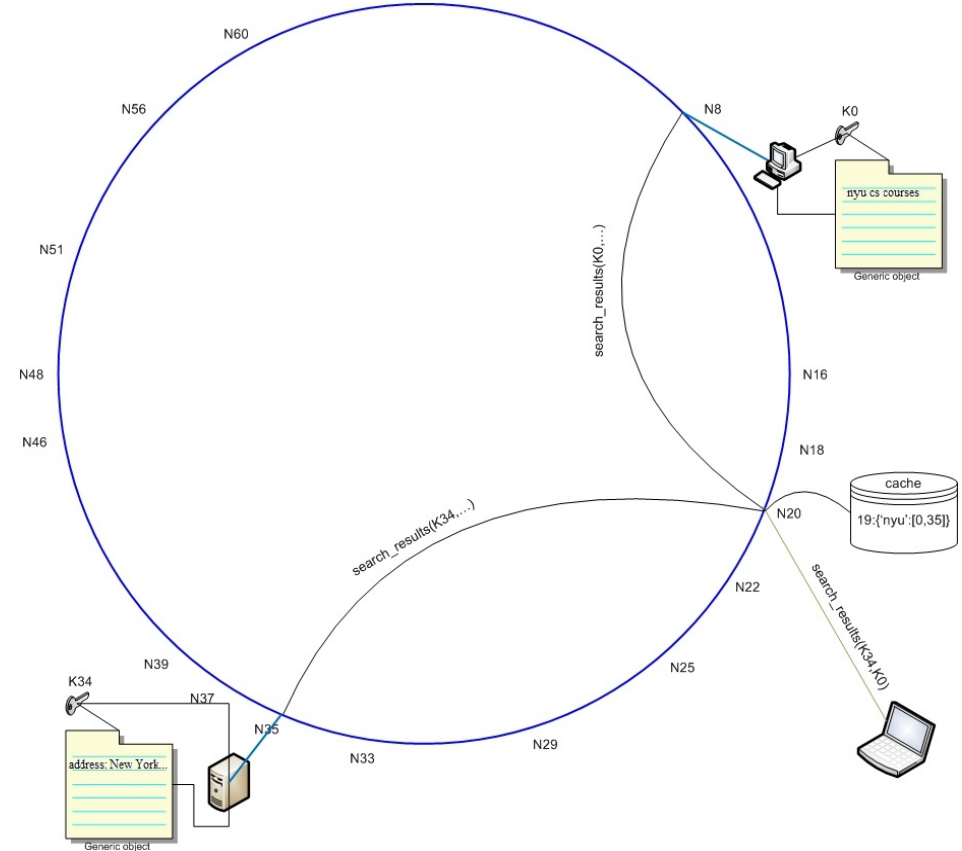
- ❑ **Slow.** Search happens one node at a time.

# Implementation

## Final Proposed Solution



□ 1: Search Start.



□ 2: Search Results.



# Summary

- Implemented Total Chord project using python with 2000 lines of code
- Tested on local machine and over Seattle peer-to-peer computing testbed
- Conclusion is that program works!
  - ▣ Initial top line test results indicated that Total Chord solves the problem of limited user search capabilities
  - ▣ Total Chord delivers on outlined solution providing flexibility of searching and retrieving more content
  - ▣ Next steps include further more-in-depth testing with increased number of machines required to understand full potential and measure reliable performance over continued usage