

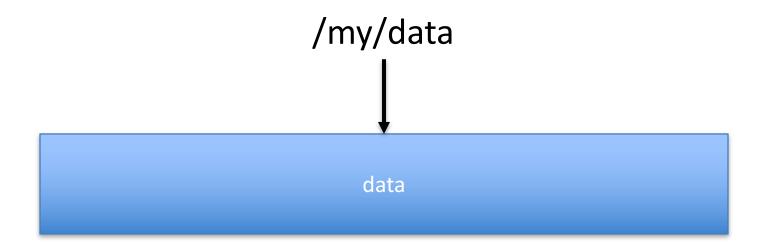
(The Futility of) Data Privacy in Content-Centric Networking

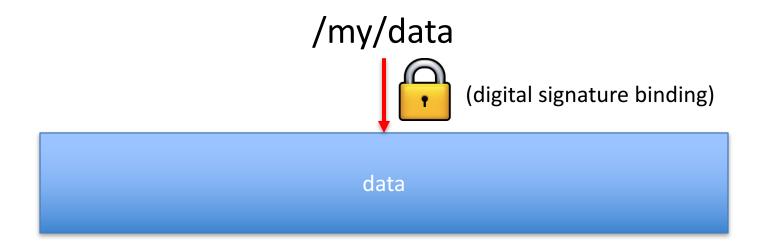
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ACM WPES 2016

Outline

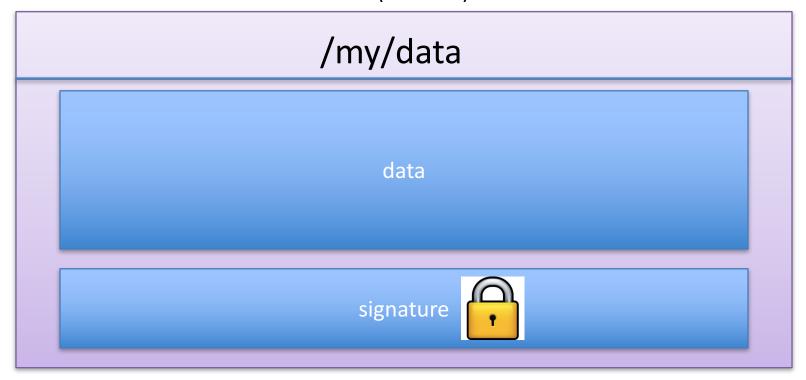
- CCN overview
- Privacy in IP vs CCN
- Privacy attacks
- Privacy requirements
- Looking ahead

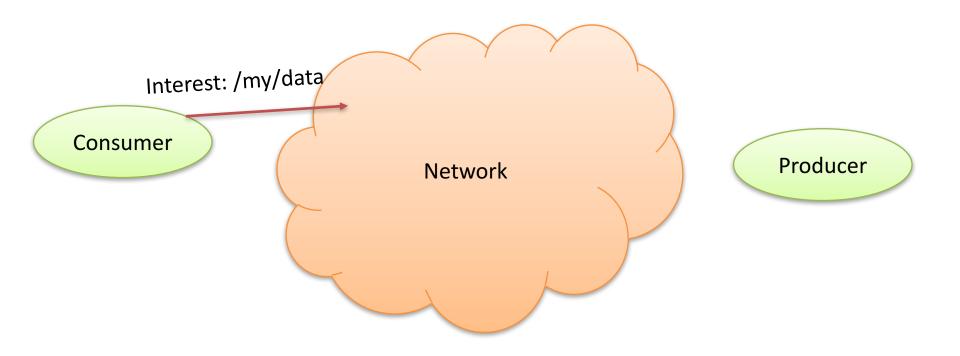
data

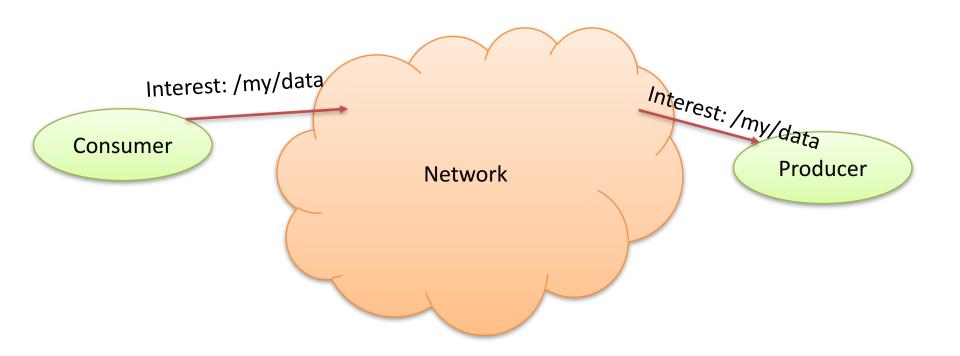


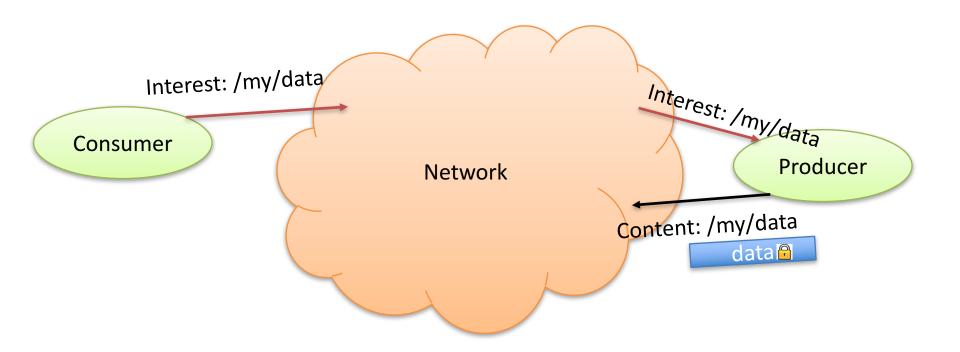


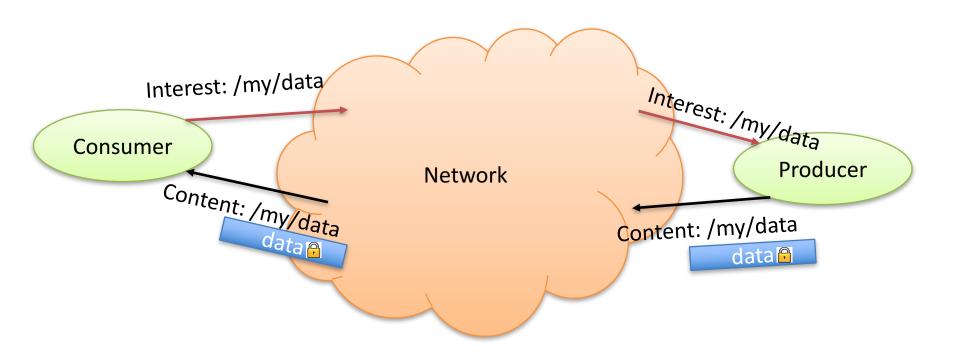
Named Data (Content) Packet

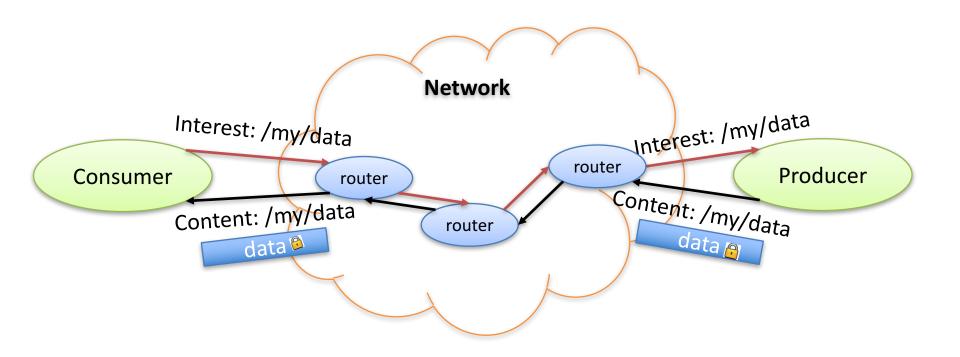


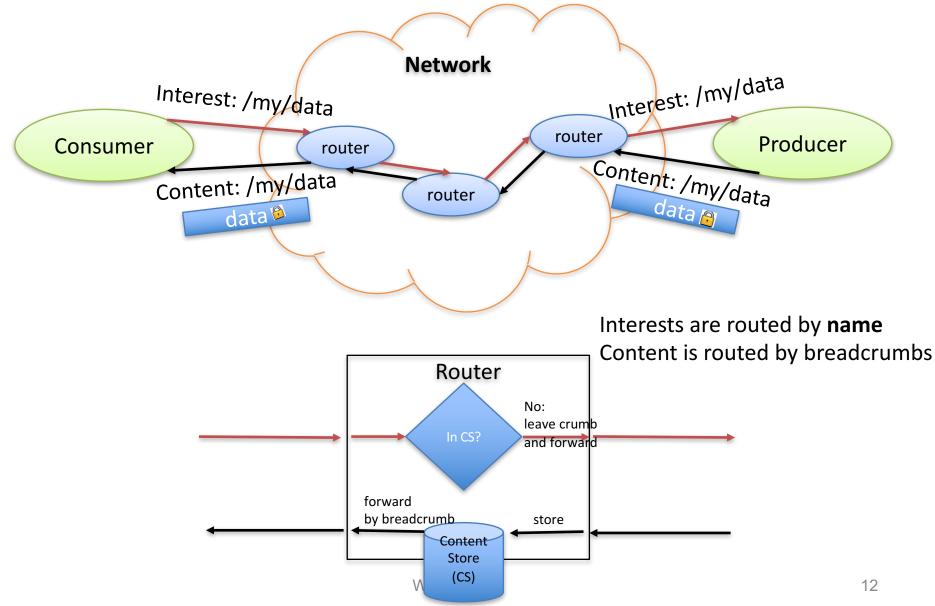


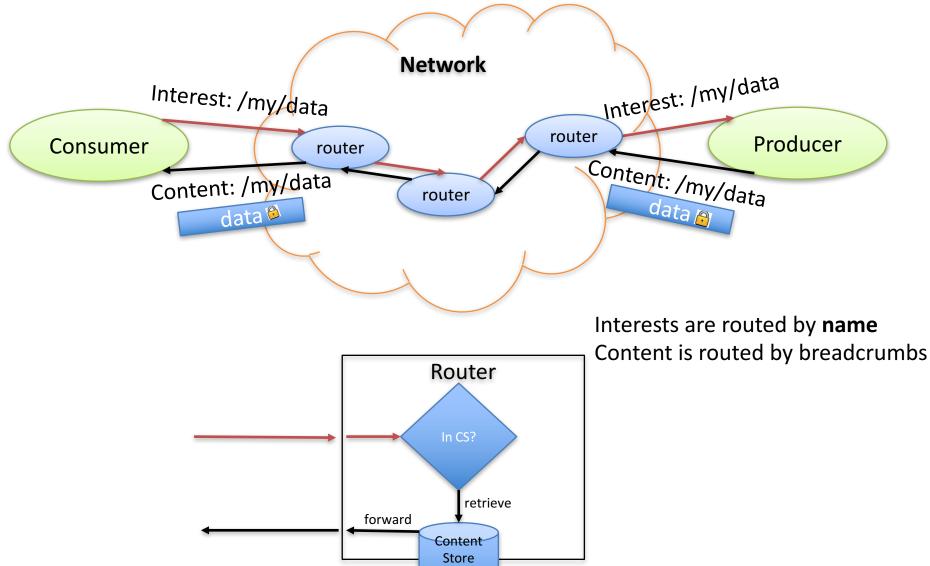












(CS)

Onto Privacy

IP Privacy

Turns this...



IP Privacy

Into this... (with IPSec or TLS)



What's revealed?

- Source and destination addresses and port #
- Timing
- Packet sizes

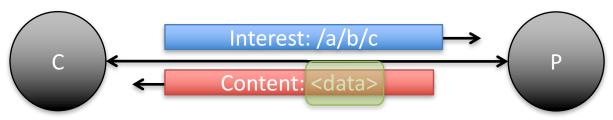
CCN Privacy

Turns this...



CCN Privacy

Into this...



encrypted content?

What's revealed?

- Consumer and producer locations
- Timing
- Packet sizes
- Interest name
- Producer identity

Properties of the (application) data

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Privacy Parity

CCN privacy < IP privacy

- What's the "delta"?
 - Interests for same content can be correlated
 - Interest names reveal information about content
 - Content carries explicit names
 - + Location of content not (always) apparent

BTW:

- Anonymity <> privacy
- ... anonymity is out of scope

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Privacy Attacks

 Correlation: learn when two requests correspond to same content

 Identification: learn when specific content was requested

Leakage: learn anything from a request or response

Adversaries

- Eavesdropper: a passive interceptor
- On-path HbC: router that forwards interest and content packets
- Distributed: at least two on-path: one near producer, one near consumer

 Active & Scary: as above, also generates its own probes

Main Questions

 What properties must responses have to prevent privacy attacks?

What about requests?

What about both?

Weak & Strong Privacy

Weak: Adv can not learn anything from a request or response, but can correlate packets

Strong: Adv can not learn, identify, or correlate

Weak Privacy Requirements

Responses must be protected via IND-KPA secure encryption

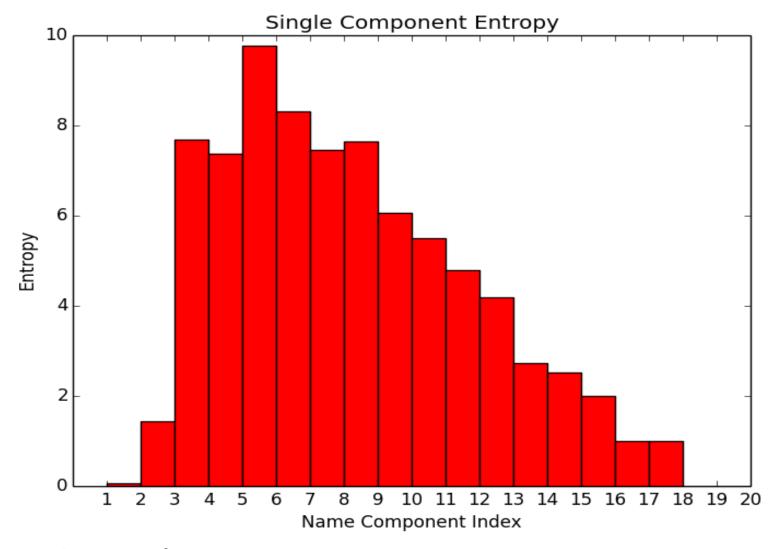
Why? To prevent trivial information leakage

 Requests must be transformed by a deterministic cryptographic PRF that is <u>not</u> length preserving

Why? Length can be used to distinguish requests from one another and the "network" representation must appear random to Adv

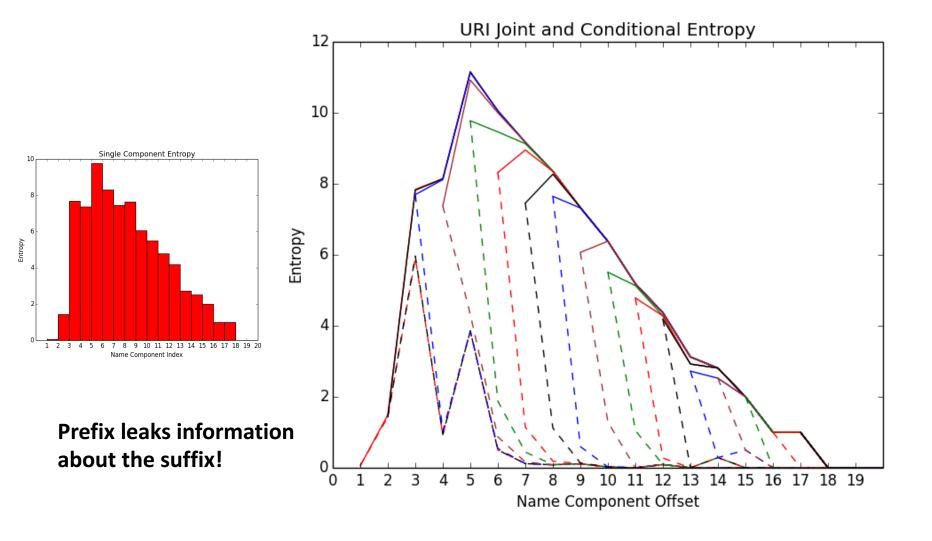
Why deterministic? How to route otherwise? Also: how to preserve the interest collapsing feature?

Hash Functions Are Not Enough



Source: Cisco URI dataset

Hash Functions Are Not Enough



Design Patterns

So... many consumers share the same secret?

Shared Secret?	Strategy
Yes	Consumer and producer: derive ephemeral shared key from secret, use it to encrypt request and response
No	 Consumer: generate random key, encrypt request with producer's public key Producer: decrypt random key, use it to encrypt response
	But what about caches?

Strong Privacy Requirement

Requests and responses must be protected with IND-CCA encryption

Why? To prevent correlation attacks

Design Pattern

Create a secure session (as in TLS) and use it as a pipe to transfer requests and responses

Outcomes?

- Any realistic form of data privacy complicates CCN request-response m.o.
 - It's no longer a simple request-response!
- In most circumstances, privacy inhibits caching
 - How important is caching in CCN?
- To have strong (proper) data privacy, CCN is not very different from IP+TLS
 - So what are we doing here?

Open Question

Will **CCN** privacy remain forever elusive or at least inferior to IPsec/TLS?

Questions?

Thanks!