Practical and Secure Accounting in Content-Centric Networking

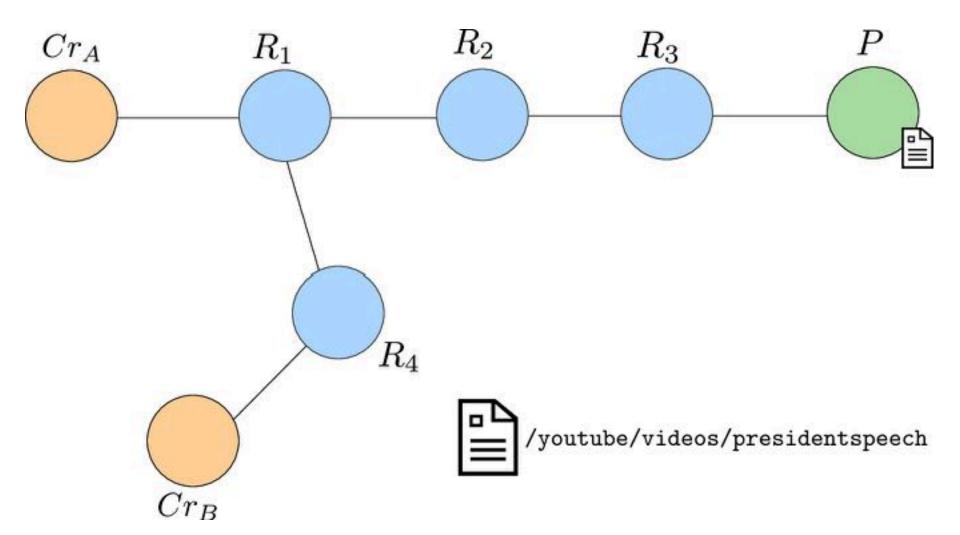
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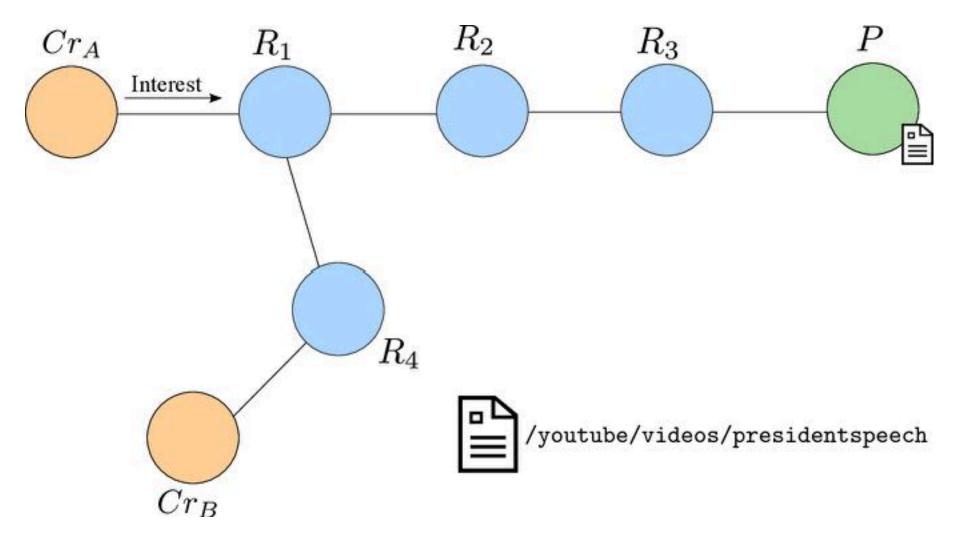
Agenda

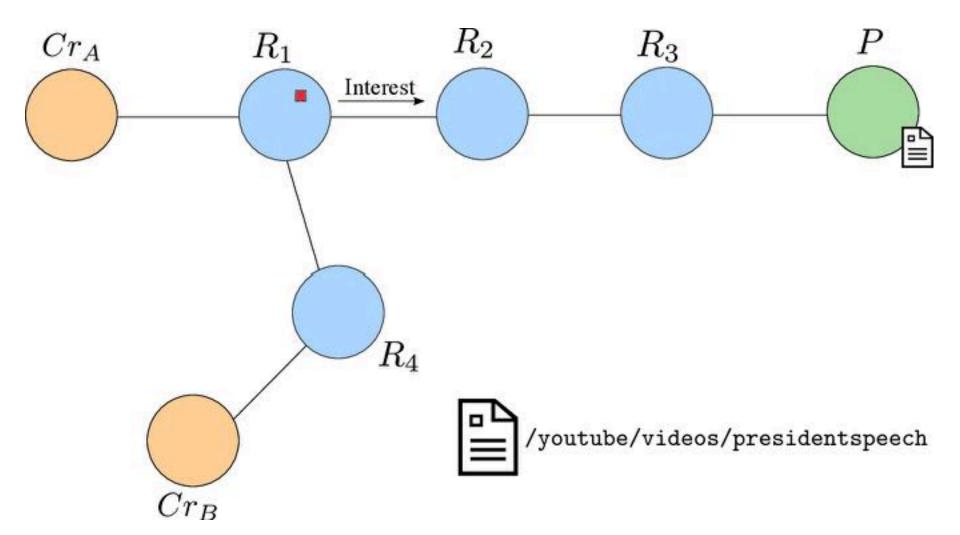
- CCN Overview
- Accounting Information
 - Individual, distinct, and aggregate
 - Cache hits vs content requests
- Accounting techniques
 - Encryption
 - Push Interests
 - Format and processing steps
- Secure accounting
 - Only aggregate and distinct are possible!
- Recommendations
 - Include nonce
- Experimental analysis
 - Show overhead for two scenarios

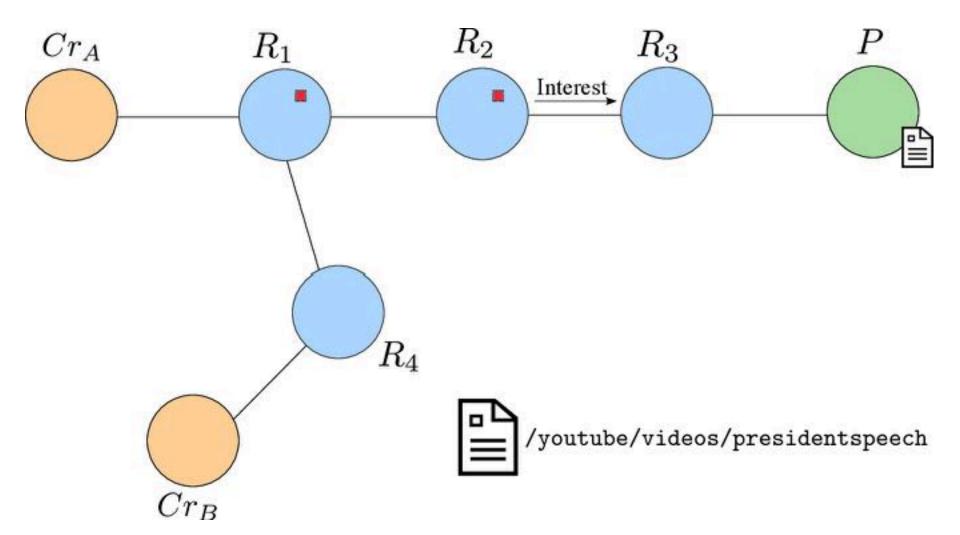
CCN/NDN Overview

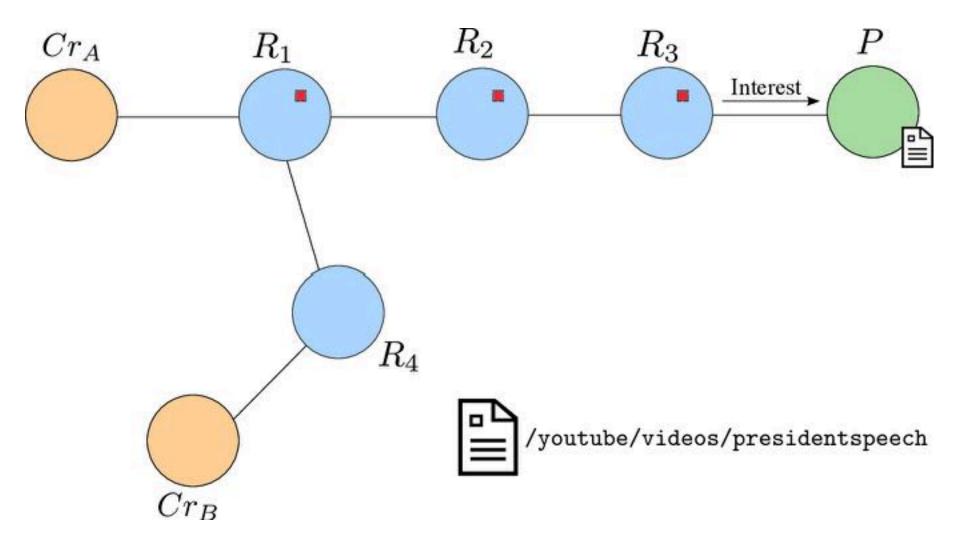
- CCN and NDN are network architectures for transferring named data
 - Data is obtained via an explicit request for the name with an interest
 - Consumers issue interests are routed towards the data producer (using the name)
 - A content object carries the data back to the consumer

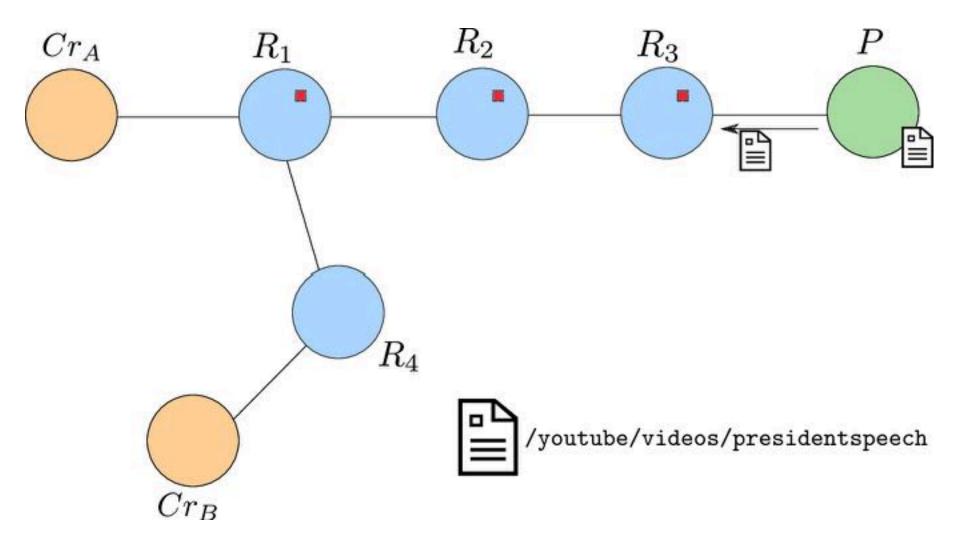


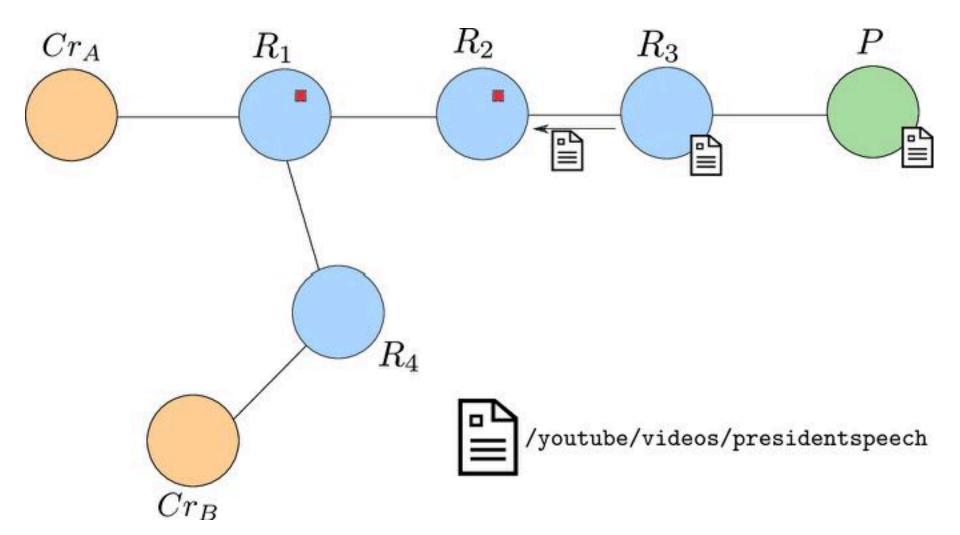


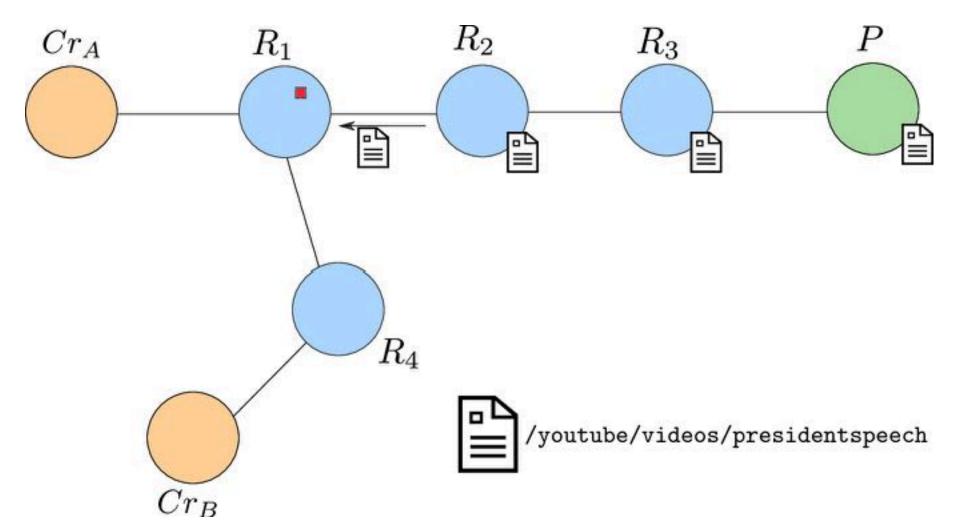


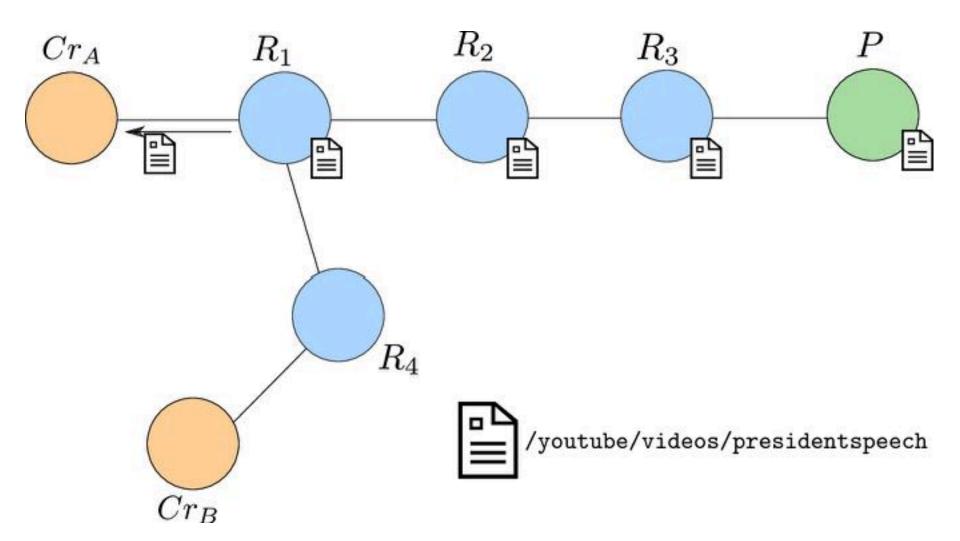


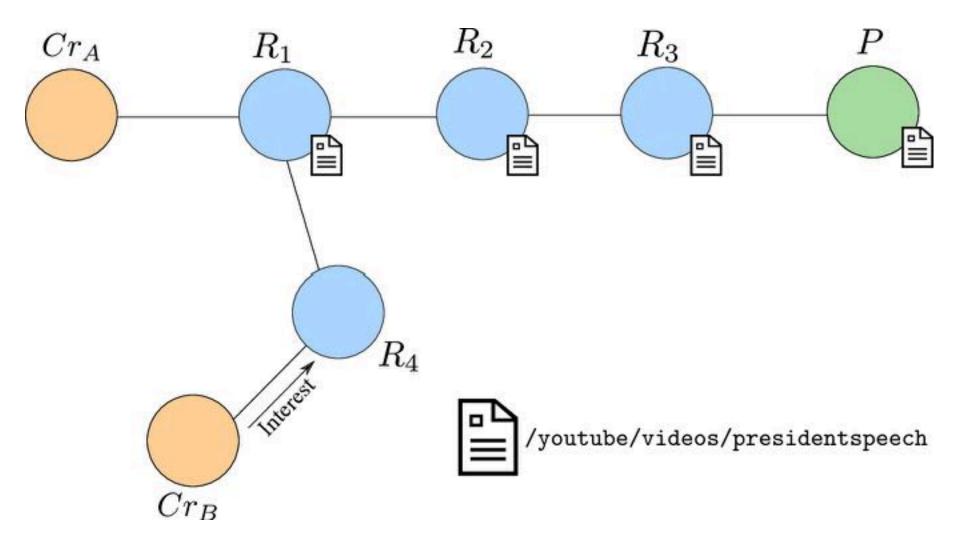


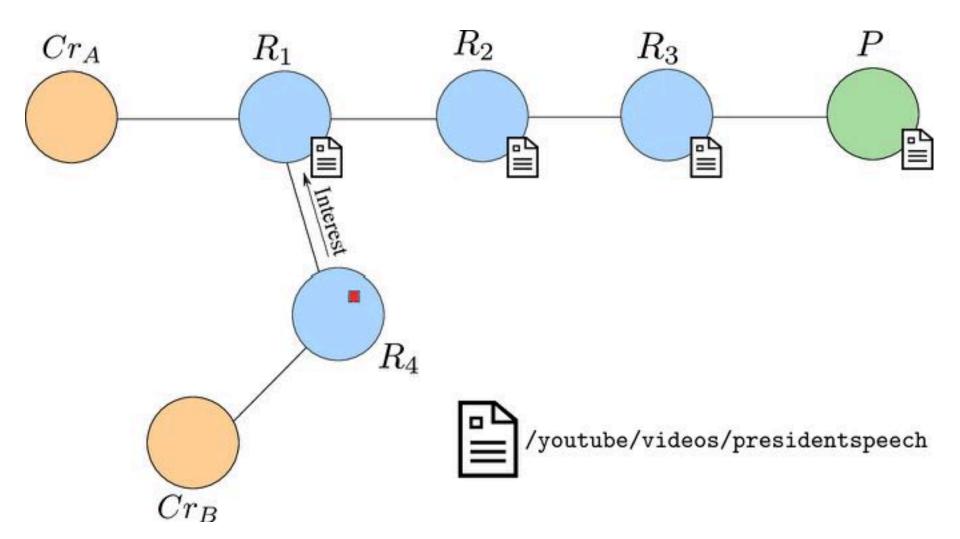












Data Decoupling

- Data is decoupled from its origin
- Routers may cache content to satisfy future duplicate requests
- Benefits:
 - Upstream congestion and bandwidth utilization is reduced via caching
 - Consumer latency is reduced
 - QoS improves (theoretically)
- Drawbacks:
 - Content must be digitally signed (or verified by some other means) to ensure authenticity
 - Producers have no way to know how many times their content was requested

The Accounting Problem(s)

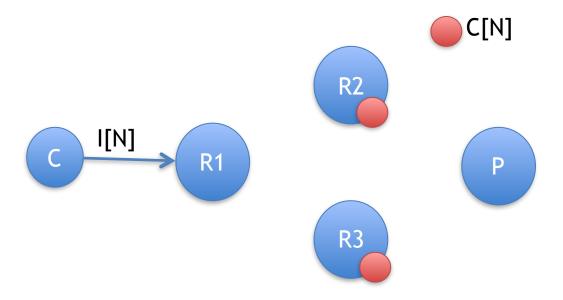
- Problem #1: How can a producer learn the number of times each content object was requested?
- Problem #2: How can a producer learn the above and which consumers requested the content?

Accounting Information

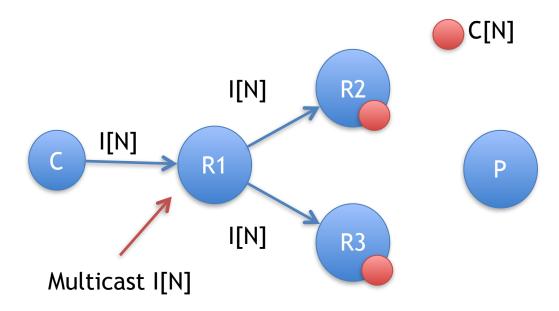
- Varying degrees of accounting information:
 - Individual: information about individual requests (and their sources)
 - Distinct: information about each individual request (but not about their source)
 - Aggregate: a simple count of the number of requests

 Question: Is the number of cache hits equal to the number of content requests?

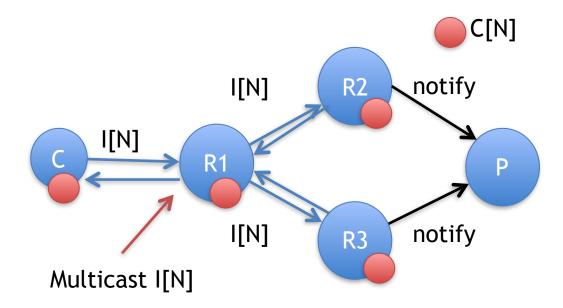
- Question: Is the number of cache hits equal to the number of content requests?
- Answer: No!



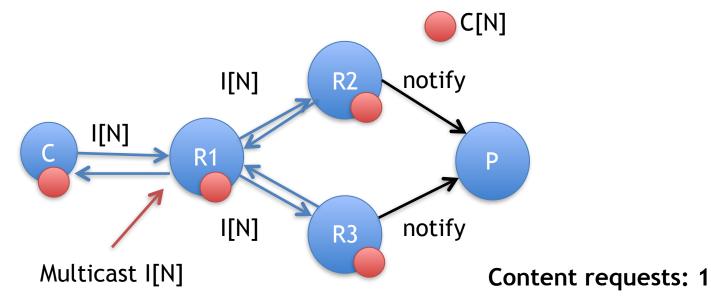
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Cache hits: 2

Goals

- Devise a mechanism to enable the collection of individual, distinct, and aggregate accounting information that is correct (accurate) and secure (when possible)
- Two approaches:
 - Encryption
 - Push Interests (our solution)

Accounting by Encryption

- If content were encrypted, every request would need to be accompanied with a request for the decryption key
 - The producer would learn about every request and could, if desired, enforce consumers to reveal their identities

Problems:

- Enforces two RTTs for a single piece of content
- Mixes two unrelated problems (access control and accounting)

Solution: Use Push Interests

- A push interest (plnt) is an interest that leaves no PIT state in a router because it expects no response
- pInts carry a name and some payload (TBD)
- When a cache hit occurs a router creates and forwards a plnt upstream to the produce [N] plnt[N] plnt[N]

Creating plnts

- The producer indicates the type of accounting information desired with an accounting flag ACCT in content:
 - ACCT ← {INDIVIDUAL, DISTINCT, AGGREGATE}
- plnts echo this flag along with:
 - ORIGIN: an identifier for the originating router (e.g., prefix or public key digest)
 - COUNT: the number of requests served by the catalyst cache hit
 - CDATA: nonce or consumer-specific data

CDATA

Aggregate

- Contents: Random nonce and timestamp
- Rationale: Distinguish between cache hit and content request cases

• Distinct:

- Contents: Random (application-specific) nonce and timestamp
- Rationale: Same as above

Individual:

- Contents: Something to identify the consumer
- Rationale: No other way to identify the consumer!
 (There are no source addresses)

Main Theorems

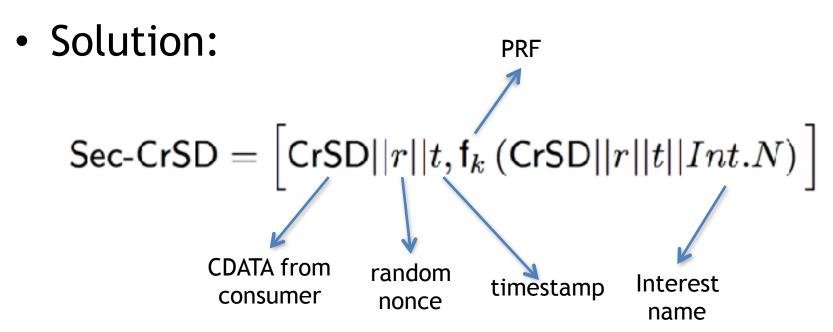
- Correctness: plnts with consumer-supplied nonces and timestamps are sufficient for guaranteeing correctness of the accounting scheme
 - Proof: see the paper
- Secure: coming up next

Secure Accounting

- Adversarial model: malicious router generating plnts for bogus interests when individual accounting is required
- This excludes:
 - Routers that don't generate plnt messages or generate plnt messages without forwarding content
 - A consumer that maliciously tries to inflate accounting information
 - A distributed adversary that can control routers and links

Security Result

 Claim: If a router cannot replay or forge individual accounting information, then the scheme is secure



Security in Practice

- We must assume:
 - Consumers know what CDATA to provide in an interest
 - Consumers have the public key of the producer to generate the secure CDATA
 - Consumers behave honestly

Dishonest Consumers

- Problem: If we assume consumers behave dishonestly (i.e., they generate fake CDATA for individually accountable content), they can bypass the accounting mechanism
- Solution: Routers must verify the secure CDATA contents
- Problem: Routers cannot verify CDATA contents at the network layer
 - Individual accounting via plnts is not possible.

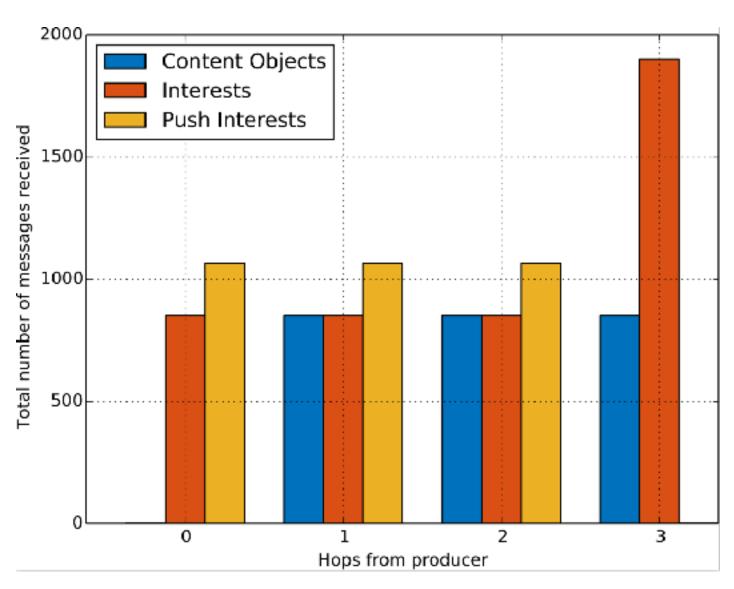
Recommendations

- Consumers always include a random nonce in interests
 - NDN does this, CCN does not!
- Routers copy these nonces into plnts regardless of the ACCT flag for cached content

Performance Assessment

- Interested in the following metrics:
 - Upstream congestion (due to plnts)
 - Forwarder overhead from processing plnts
- Ran simulations in ndnSIM based on the popular DFN and AT&T topologies

Upstream Congestion



Forwarder Overhead

