

Assignment: Codes: Definitions and Theorems, Due Monday, 11:59pm

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1. Non-Perfect Code

Make up an example of a linear code with 16 code words that isn't perfect. List the 16 code words, and give a generating matrix.

```
M <- matrix(c(1,0,0,0,1,1,0, 0,1,0,0,1,1,1, 0,0,1,0,0,1,1, 0,0,0,1,1,1,1), nrow=4, byrow=TRUE)
M
```

	[,1]	[,2]	[,3]	[,4]	[,5]	[,6]	[,7]
[1,]	1	0	0	0	1	1	0
[2,]	0	1	0	0	1	1	1
[3,]	0	0	1	0	0	1	1
[4,]	0	0	0	1	1	1	1

```
C <- generateCode(M)
C
```

	[,1]	[,2]	[,3]	[,4]	[,5]	[,6]	[,7]
[1,]	1	0	0	0	1	1	0
[2,]	0	1	0	0	1	1	1
[3,]	1	1	0	0	0	0	1
[4,]	0	0	1	0	0	1	1
[5,]	1	0	1	0	1	0	1
[6,]	0	1	1	0	1	0	0
[7,]	1	1	1	0	0	1	0
[8,]	0	0	0	1	1	1	1
[9,]	1	0	0	1	0	0	1
[10,]	0	1	0	1	0	0	0
[11,]	1	1	0	1	1	1	0
[12,]	0	0	1	1	1	0	0
[13,]	1	0	1	1	0	1	0
[14,]	0	1	1	1	0	1	1
[15,]	1	1	1	1	1	0	1
[16,]	0	0	0	0	0	0	0

```
isPerfect(C)
```

```
[1] FALSE
```

2. Check bounds

Consider the code generated by the following matrix.

$$B = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix}$$

```
M <- matrix(c(1,0,0,1,1,0, 0,1,0,0,1,1, 0,0,1,1,0,1), nrow=3, byrow=TRUE)
```

a. List all the code words.

```
C <- generateCode(M)
C
```

	[,1]	[,2]	[,3]	[,4]	[,5]	[,6]
[1,]	1	0	0	1	1	0
[2,]	0	1	0	0	1	1
[3,]	1	1	0	1	0	1
[4,]	0	0	1	1	0	1
[5,]	1	0	1	0	1	1
[6,]	0	1	1	1	1	0
[7,]	1	1	1	0	0	0
[8,]	0	0	0	0	0	0

b. Calculate the code rate of this code.

```
log2(nrow(C))/length(C[1, ])
```

```
[1] 0.5
```

c. Is this code a maximum distance separable (MDS) code? Prove or disprove.

```
M <- nrow(C)
n <- length(C[1, ])
d <- codeDistance(C)
isMDS <-function(){
  if(M == 2^(n-d+1))
    return(TRUE)
  else{
    return(FALSE)
  }
}
isMDS()
```

```
[1] FALSE
```

d. Is this code a perfect code? Prove or disprove.

```
isPerfect(C)
```

```
[1] FALSE
```

3. Perfect repetition codes

Prove that a k repetition code is perfect if and only if k is odd.

Repetition codes have two code words, all 0's or all 1's thus $M = 2$.

A code is perfect if $M = \frac{2^k}{\sum_{j=0}^t \binom{n}{j}}$

Since $M = 2$, the numerator must be twice the denominator for a repetition code to be perfect.

They have a minimum distance $d = k$

Thus, if k is odd $t = \frac{k-1}{2}$, and if k is even $t = \frac{k}{2} - 1$

So if k is odd, $\sum_{j=0}^t \binom{k}{j} = 2^{t-1}$ by the binomial theorem.

Thus if, and only if, k is odd, $2 = \frac{2^k}{2^{k-1}}$

4. 8×8 Hadamard code

Here's an 8×8 Hadamard matrix:

$$B = \begin{bmatrix} 1 & 1 & 1 & -1 & 1 & -1 & -1 & -1 \\ 1 & 1 & -1 & -1 & -1 & 1 & -1 & 1 \\ 1 & -1 & 1 & 1 & -1 & -1 & -1 & 1 \\ 1 & -1 & 1 & -1 & -1 & 1 & 1 & -1 \\ 1 & -1 & -1 & -1 & 1 & -1 & 1 & 1 \\ 1 & 1 & -1 & 1 & -1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 & 1 & 1 & -1 & -1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

```
B <- matrix(c(1,1,1,0,1,0,0,0, 1,1,0,0,0,1,0,1, 1,0,1,1,0,0,0,1, 1,0,1,0,0,1,1,0, 1,0,0,0,1,0,1,1, 1,1,
```

- Is the corresponding 8×8 Hadamard code (with -1 's replaced by 0's) a linear code? Explain why or why not.

```
isLinear <- function(B){
  vec3 <- vector()
  nrowB <- nrow(B)
  inv <- vector()
  for(l in 1:length(B[1, ])){
    inv <- c(inv, 1)
  }
  for(j in 1:(nrowB-1)){
    vec1 <- B[j, ]
    for(i in (j+1):nrowB){
      match <- FALSE
      vec2 <- B[i, ]
      vec3 <- as.numeric(xor(vec1, vec2))
      for(k in 1:nrowB){
        if(all(vec3 == B[k, ]) || all((vec3 == xor(inv, B[k, ])) )){
          match <- TRUE
        }
      }
      if(match == FALSE)
    }
  }
}
```

```
        return(FALSE)
    }
    }
    return(TRUE)
}

isLinear(B)
```

```
[1] TRUE
```

This is a linear code, because any two code words added together produce another code word.

b. Is this Hadamard code perfect? Prove or disprove.

```
C <- generateCode(B)
isPerfect(C)
```

```
[1] FALSE
```