

CS 5044

Object-Oriented Programming with Java

Q&A Session

Why (and why not) recursion?

- Why recursion?
 - Certain types of problems are naturally much easier to describe (thus solve) recursively
 - Sorting and searching are the most common examples, but there are many more
 - Domain-specific languages and data formats are often defined and parsed recursively
 - Recursive algorithms can be mathematically proven as correct
 - This makes testing significantly more reliable and convenient
 - The concept is nearly identical to "proof by induction"
 - Converting an *arbitrary* recursive algorithm to its iterative equivalent can be tricky
 - Converting from iterative to recursive is relatively easy (but see "why not" below)
- Why not recursion?
 - Recursion can be slower (and take more memory) than an iterative equivalent, however...
 - ...the differences are often insignificant, and speed is often not critical anyway
 - If the problem is not naturally recursive, a recursive solution is less idiomatic
 - If iteration is more straightforward, or even equivalent, that's what will be expected
- Summary:
 - Feel free to use recursion wherever it makes sense
 - Don't go out of your way to force-fit recursion into naturally iterative processes

Syntax and strategies for recursion

- No special syntax or structures are required for recursion
 - Just like calling any other method to do some work and return some value
 - The method called just happens to be the same method that is making the call
- Typical approach splits input into two groups: **first part** and **all the rest**
 - You get to decide how to split the input, such that:
 - The first part can be processed reasonably easily
 - All the rest is similar to the original input, only a bit simpler and/or smaller
 - Often exactly the same structure but with fewer elements or smaller values
 - Next, you need to actually process the first part
 - You may safely assume that you've *already solved* how to process all the rest
 - This often helps you choose how to split the input
 - This also sometimes helps in processing the first part
 - Finally, you need criteria to detect when there isn't even a first part remaining
 - Usually this is a *degenerate* case (empty String, empty collection, 0, null, etc.)
- Notice that we never actually need to solve how to process all the rest!

Factorials: recursion in practice

- Definition: factorial(n) is the product of all numbers $1 \dots n$ (for all $n \geq 1$)
 - factorial(1) = $1 = 1$
 - factorial(2) = $1 * 2 = 2$
 - factorial(3) = $1 * 2 * 3 = 6$
 - factorial(4) = $1 * 2 * 3 * 4 = 24$
 - factorial(5) = $1 * 2 * 3 * 4 * 5 = 120$
 - factorial(9) = $1 * 2 * 3 * 4 * 5 * 6 * 7 * 8 * 9 = 362,880$
- It's a naturally recursive algorithm, so we'll first solve this iteratively(!)
 - Why? The multiplicands look like they were generated by a loop, and we understand loops
 - We'll use while() loops just for clarity in our examples, rather than for() loops

```
public int factorialI(int n) { // valid for positive n only (we'll ignore any validation here)
    int product = 1;          // need an accumulator for our running total
    int i = 1;                 // start the loop index at 1
    while(i <= n) {            // repeat until we reach n
        product = product * i; // multiply by the next number
        i = i + 1;             // increment the index
    }
    return product;            // return our accumulator
}
```

- That's great, but what about recursion?!


Factorials: recursion in practice, with actual recursion

- Let's look at the examples again, aligned differently, and with parentheses:
 - $\text{factorial}(1) = 1$
 - $\text{factorial}(2) = (1)*2 = 1*2$
 - $\text{factorial}(3) = (1*2)*3 = 1*2*3$
 - $\text{factorial}(4) = (1*2*3)*4 = 1*2*3*4$
 - $\text{factorial}(5) = (1*2*3*4)*5$
- A slightly different pattern emerges here...

Factorials: recursion in practice, with actual recursion


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 - $\text{factorial}(5) = (1 * 2 * 3 * 4) * 5$
- A slightly different pattern emerges here...
 - Each factorial requires only **one new** multiplication, if we can use the prior calculation!

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 - $\text{factorial}(5) = (1 * 2 * 3 * 4) * 5$
- A slightly different pattern emerges here...
 - Each factorial requires only one new multiplication, if we can use the prior calculation!
 - $\text{factorial}(1) = 1$ 
 - $\text{factorial}(2) = \text{factorial}(1) * 2$
 - $\text{factorial}(3) = \text{factorial}(2) * 3$
 - $\text{factorial}(4) = \text{factorial}(3) * 4$
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This seems to be
a special case...

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- A slightly different pattern emerges here...
 - Each factorial requires only one *new* multiplication, if we can use the prior calculation!
 - $\text{factorial}(1) = 1$ 

This seems to be a special case...
 - $\text{factorial}(2) = \text{factorial}(1) * 2$
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 - $\text{factorial}(4) = \text{factorial}(3) * 4$
 - $\text{factorial}(5) = \text{factorial}(4) * 5$
- Let's redefine the factorial function more generally, with a recursive definition:
 - $\text{factorial}(n) = 1$ (*special case, for $n == 1$*)
 - $\text{factorial}(n) = \text{factorial}(n-1) * n$ (*for all $n > 1$*)

Recursion walkthrough

$\text{factorial}(n) = 1$ *(for $n == 1$)*

$\text{factorial}(n) = \text{factorial}(n-1) * n$ *(for all $n > 1$)*

Notice there are no loops suggested by this definition

```
public int factorialR(int n) {  
    if (n == 1) {           // detect terminating condition  
        return 1;          // return special case  
    }  
    return factorialR(n - 1) * n; // recursion!  
}
```

Currently
Running



CALL STACK

someMethod

int value: 4
int result: pending

Let's trace through a call with argument 4

```
public void someMethod() {  
    int value = 4;  
    int result = factorialR(value);  
    System.out.println(result);  
}
```

Recursion walkthrough

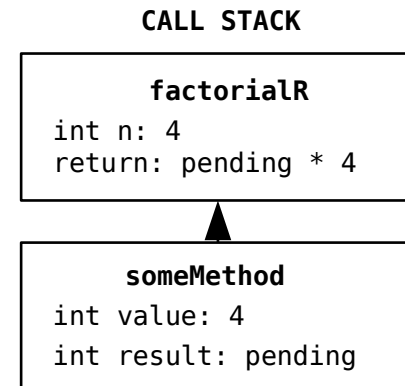
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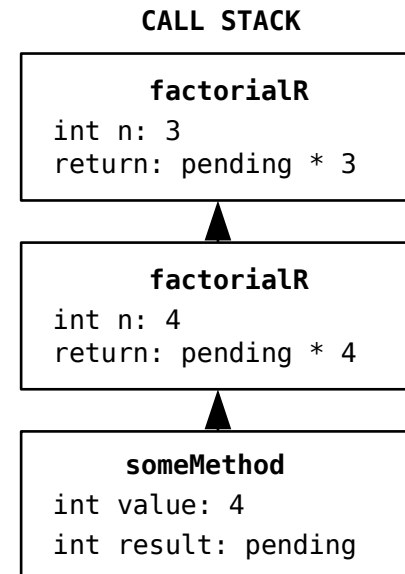
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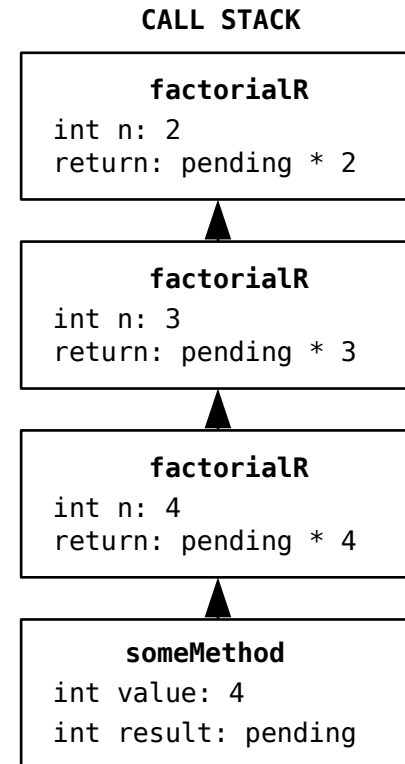
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Currently
Running



Recursion walkthrough

$\text{factorial}(n) = 1$ (for $n == 1$)

$\text{factorial}(n) = \text{factorial}(n-1) * n$ (for all $n > 1$)

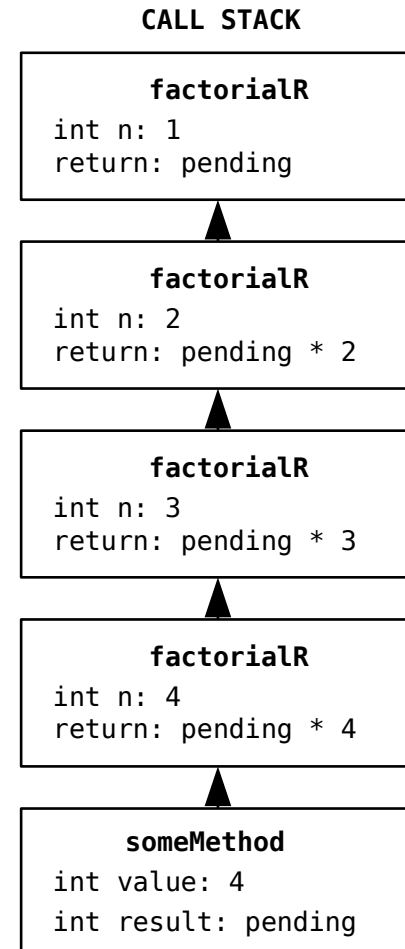
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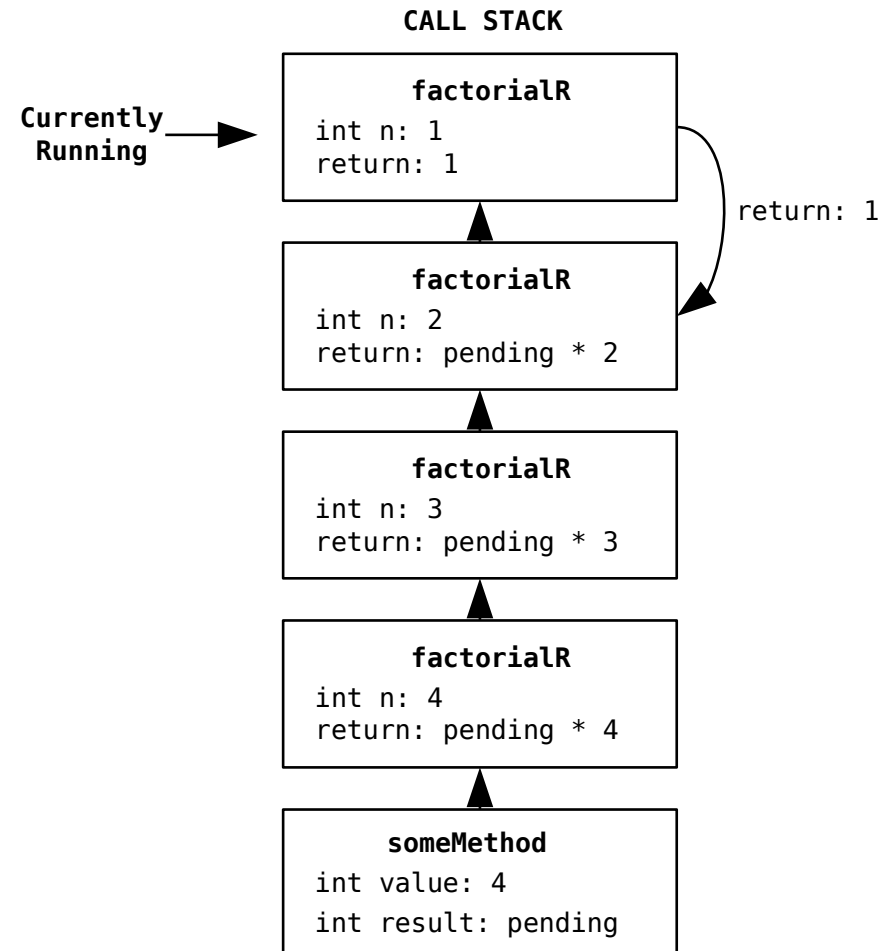
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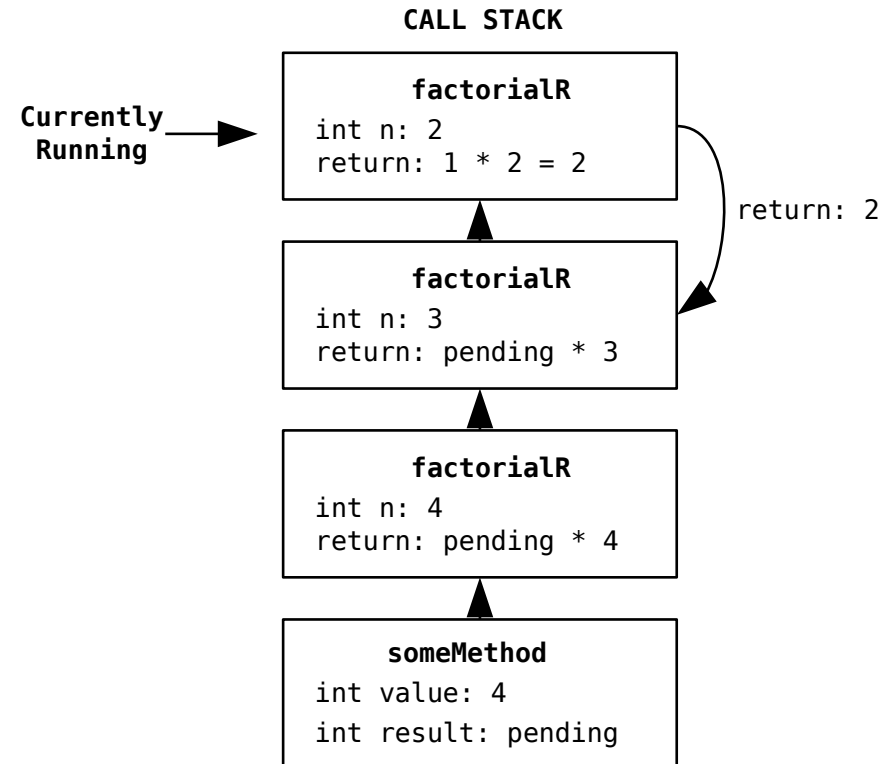
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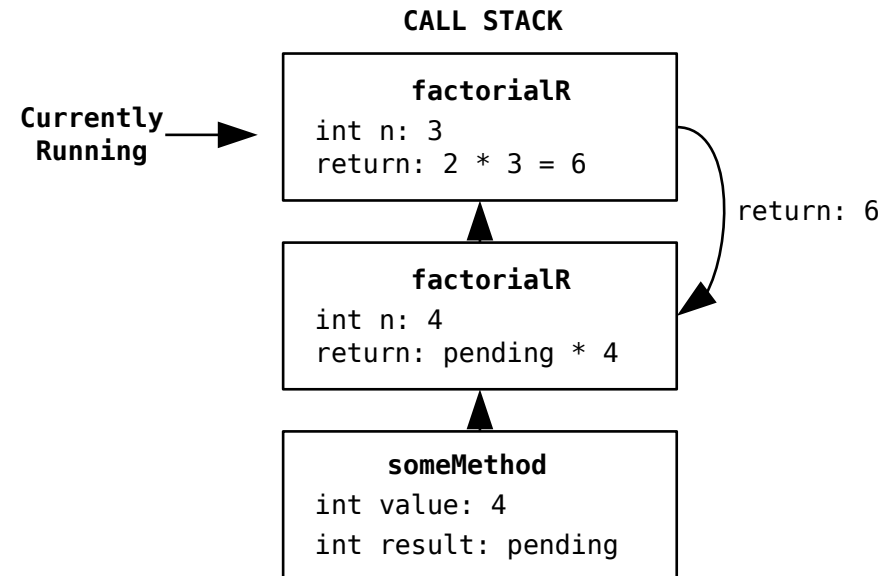
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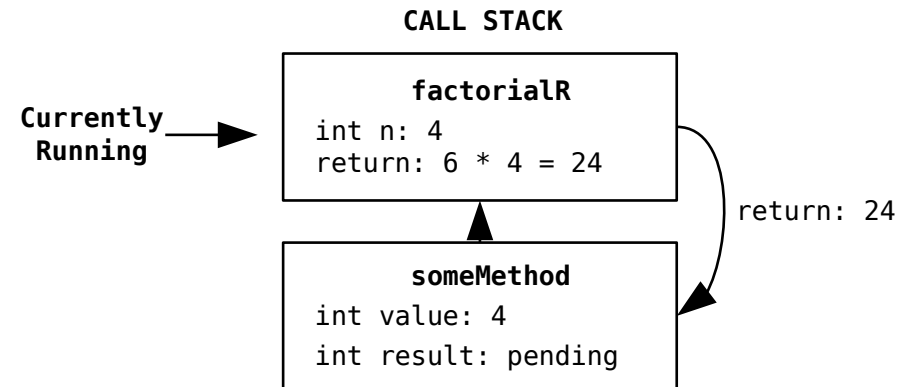
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```

Currently
Running →

CALL STACK

someMethod

int value: 4
int result: 24

Let's trace through a call with argument 4

```
public void someMethod() {  
    int value = 4;  
    int result = factorialR(value);  
    System.out.println(result);  
}
```

Recursion vs iteration by example: find minimum

- Consider a method to find the minimum of a collection of numbers
 - The iterative approach, using a while() loop for clarity:

```
private int findMinIterative(List<Integer> list) {  
    int currMin = Integer.MAX_VALUE;  
    int index = 0;  
    while (index < list.size()) {  
        currMin = Math.min(currMin, list.get(index));  
        index = index + 1;  
    }  
    return currMin;  
}
```

- Convert to the recursive equivalent, using a recursive helper:

```
private int findMinRecursive(List<Integer> list) {  
    return findMinHelper(list, Integer.MAX_VALUE, 0);  
}  
  
private int findMinHelper(List<Integer> list, int currMin, int index) {  
    if (index >= list.size()) {  
        return currMin;  
    }  
    return findMinHelper(list, Math.min(currMin, list.get(index)), index + 1);  
}
```

Recursion vs iteration by example: find minimum

- Consider a method to find the minimum of a collection of numbers
 - The iterative approach, using a while() loop for clarity:

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private int findMinIterative(List<Integer> list) {  
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        ⑤ index = index + 1;  
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    ③ if (index >= list.size()) {  
        ⑥ return currMin;  
    }  
    return findMinHelper(list, Math.min(currMin, list.get(index)), index + 1);  
}
```

```
graph TD
    subgraph findMinRecursive
        L1[Line 1: return findMinHelper(list, Integer.MAX_VALUE, 0);]
    end
    subgraph findMinHelper
        L3[Line 3: if (index >= list.size()) {]
        L6[Line 6: return currMin;]
        L4[Line 4: return findMinHelper(list, Math.min(currMin, list.get(index)), index + 1);]
    end
    L1 -- ① --> L3
    L3 -- ② --> L4
    L4 -- ④ --> L6
    L6 -- ⑤ --> L4
    L4 -- ⑥ --> L1
```

Recursion vs iteration by example: find minimum

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- The iterative approach, using a while() loop for clarity:

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}
```

```
graph TD  
  A[findMinRecursive] -- 1 --> B[findMinHelper list, MAX_VALUE, 0]  
  B -- 2 --> C[findMinHelper list, min(MAX_VALUE, list.get(0)), 1]  
  C -- 3 --> D[findMinHelper list, min(min(MAX_VALUE, list.get(0)), list.get(1)), 2]  
  D -- 4 --> E[findMinHelper list, min(min(min(MAX_VALUE, list.get(0)), list.get(1)), list.get(2)), 3]  
  E -- 5 --> F[findMinHelper list, min(min(min(min(MAX_VALUE, list.get(0)), list.get(1)), list.get(2)), list.get(3)), 4]  
  F -- 6 --> G[return currMin to previous call]
```

Note that the recursive approach doesn't require any local variables, nor does it mutate anything!

These can be extremely important properties in certain situations, such as in multi-threaded and/or distributed systems.

However, the price we pay is more memory usage, plus some added processing overhead, so there are always trade-offs to consider!

It's also arguably a somewhat less elegant solution than iteration; can we do any better?

Recursion vs iteration by example: find minimum

- Consider a method to find the minimum of a collection of numbers
 - The iterative approach, using a while() loop for clarity:

```
private int findMinIterative(List<Integer> list) {  
    int currMin = Integer.MAX_VALUE;  
    int index = 0;  
    while (index < list.size()) {  
        currMin = Math.min(currMin, list.get(index));  
        index = index + 1;  
    }  
    return currMin;  
}
```

- Normally we wouldn't retain the index, and would instead just iterate with a smaller list:

```
private int findMinRecursive(List<Integer> list) {  
    return findMinHelper(list, Integer.MAX_VALUE);  
}  
  
private int findMinHelper(List<Integer> list, int currMin) {  
    if (list.isEmpty()) {  
        return currMin;  
    }  
    return findMinHelper(list.subList(1, list.size()), Math.min(currMin, list.get(0)));  
}
```

- This is more fitting of the concept of separating the "first" from the "rest"

Recursion vs iteration by example: find minimum

- Consider a method to find the minimum of a collection of numbers
 - The iterative approach, using a while() loop for clarity:

```
private int findMinIterative(List<Integer> list) {  
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    int index = 0;  
    while (index < list.size()) {  
        currMin = Math.min(currMin, list.get(index));  
        index = index + 1;  
    }  
    return currMin;  
}
```

- Further, we typically also remove the accumulator, so we no longer need the helper:

```
private int findMinRecursive(List<Integer> list) {  
    if (list.isEmpty()) {  
        return Integer.MAX_VALUE;  
    }  
    return Math.min(list.get(0), findMinRecursive(list.subList(1, list.size())));  
}
```

- The code is now slightly denser, yet somewhat *more* elegant than iteration...
 - However, unless the code is in a multi-threaded or distributed environment...
 - » ...the iterative approach would be a much less surprising way to find a minimum

Thinking recursively (with an iterative start)

- Let's try to multiply two non-negative integers, by only using addition
 - For example: $6*8 = 8 + 8 + 8 + 8 + 8 + 8$
- Iterative approach ($A*B$):
 - Start with an accumulator set to zero, repeat A times of adding B each time

```
public int multiplyI(int a, int b) {  
    int sum = 0;  
    int index = a;  
    while (index > 0) {  
        sum = sum + b;  
        index = index - 1;  
    }  
    return sum;  
}
```

- Recursive approach ($A*B$):
 - Notice that $6*8 = 8 + (5*8)$ so we have an answer in terms of a simplified problem
 - Eventually we reach $1*8 = 8 + (0*8)$ then finally $0*8 = 0$ (the terminating case)
 - More generally $A*B = B + (A - 1)*B$ (for all $A > 0$); $A*B = 0$ (for $A == 0$)

```
public int multiplyR(int a, int b) {  
    if (a == 0) {  
        return 0;  
    }  
    return b + multiplyR(a - 1, b);  
}
```


Count on recursions (yet another example)

- Suppose we need to count the number of spaces in a string
 - Example: `countSpaces("Suppose we need to count the number of spaces in a string") = 11`
- Iterative approach:
 - Start with `index=0` and `count=0`, repeat for every char, increment count if char is space

```
public int countSpacesI(String s) {  
    int count = 0;  
    int index = 0;  
    while (index < s.length()) {  
        if (s.charAt(index) == ' ') {  
            count = count + 1;  
        }  
        index = index + 1;  
    }  
    return count;  
}
```

- Recursive approach:
 - Total spaces = number of spaces in first char + number of spaces in all remaining chars
 - `countSpaces(...)` = ((firstChar is space?) 1 else 0) + `countSpaces(substring after firstChar)`
 - Eventually we reach an empty string, where `count("") = 0` as the terminating case

```
public int countSpacesR(String s) {  
    if (s.isEmpty()) {  
        return 0;  
    }  
    return (s.charAt(0) == ' ' ? 1 : 0) + countSpacesR(s.substring(1));  
}
```

Sorting and searching

- Sorting problems tend to be naturally recursive algorithms
 - No single sorting algorithm is "best" for all cases
 - Best may involve several factors: fastest, least memory, most interruptible, etc.
 - Even on a single factor, algorithms vary by best case, worst case, average, etc.
- Searching also tends to be recursive, especially for common data structures
 - There are dozens of distinct search tree algorithms, with varying properties
 - Java's `TreeMap` and `TreeSet` happen use use a "Red-Black" binary tree structure
- Order-of-the-function ("Big O Notation") performance is usually all we need to know
 - Still, depending on usage patterns, specific implementations may vary in several ways:
 - Memory usage, index size, speed of insertion, speed of retrieval, etc.
 - However, for the vast majority of applications, we simply don't care!
 - Selecting any appropriate collection class is normally the best approach
 - For example, choose `Set` or `Map` rather than `List` when `contains()` is important
 - Unless you have exceptional performance needs, just use what the language provides
 - Collections typically mention performance orders in their API documentation