

COMP ENG 4DS4

Lab 0 - Introduction Lab

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As a future member of the engineering profession, the student is responsible for performing the required work in an honest manner, without plagiarism and cheating. Submitting this work with my name and student number is a statement and understanding that this work is my own and adheres to the Academic Integrity Policy of McMaster University and the Code of Conduct of the Professional Engineers of Ontario. **Submitted by Aidan Mathew, Aaron Rajan, Sameer Shakeel, and Chris Jiang.**

Declaration of Contributions:

Student	Contributions
Aidan Mathew - 400306142	Worked on Report, participated in all lab experiments, and worked on Problems 1, 2, 3, and 5.
Aaron Rajan -	Participated in all lab experiments and worked on Problems 1, 3, 4, 5.
Chris Jiang -	Helped with lab experiments and edited lab report.
Sameer Shakeel -	Participated in all lab experiments and worked on Problems 1, 3, 4, 5.

Lab Goals

- Get familiar with the tools (IDE, Debugger, board connection)
- Revise frequently used concepts in C
- Deal with memory-mapped registers
- Deal with the MCU's datasheet and the board's schematics
- Learn how to use GPIOs, UART and PWM

Experiments

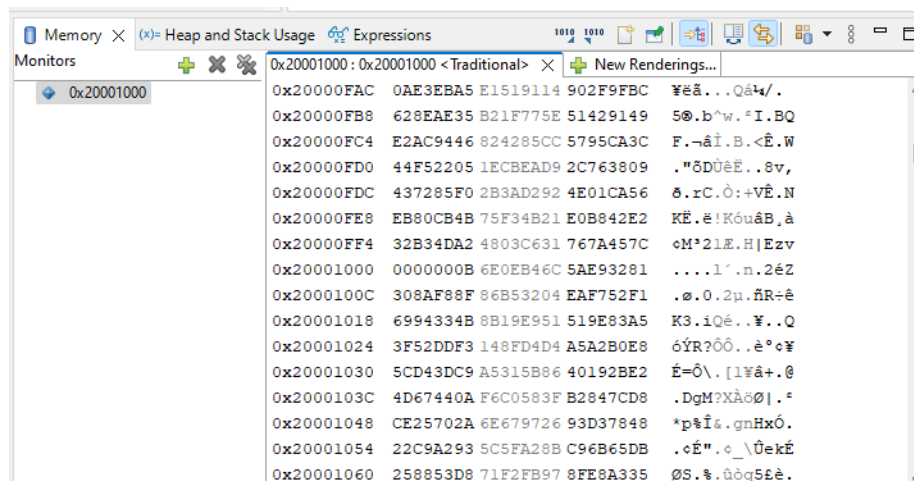
Experiment 1 - Hello World

In this first experiment we learned how to setup our environment and ran the hello world example.

Experiment 2 - C Pointers and Structures

During Part A of this experiment, we explored the use of pointers to access and manipulate memory on the microcontroller. A test function was implemented to demonstrate pointer behavior. An integer variable `x` was declared and initialized to `0`, and a pointer `ptr` was assigned the address of `x`. Additionally, a second pointer, `ptr_location`, was assigned a direct memory address (`0x20001000`). The values at these locations were modified by dereferencing the pointers: `*ptr` was set to `10`, and `*ptr_location` was set to `11`. Debugging tools were used to verify the values. The expressions panel showed `x` with a value of `10`, confirming that the pointer `ptr` correctly referenced `x`. Similarly, the value at memory address `0x20001000` was updated to `11` (the value `B` in Hex = `11`), as shown in the memory panel screenshot. This experiment demonstrated the ability to access and modify specific memory locations on the MCU and the importance of debugging tools like the watch and memory panels for visualizing memory state and verifying pointer operations. See the below screenshots of the watch panel and the memory.

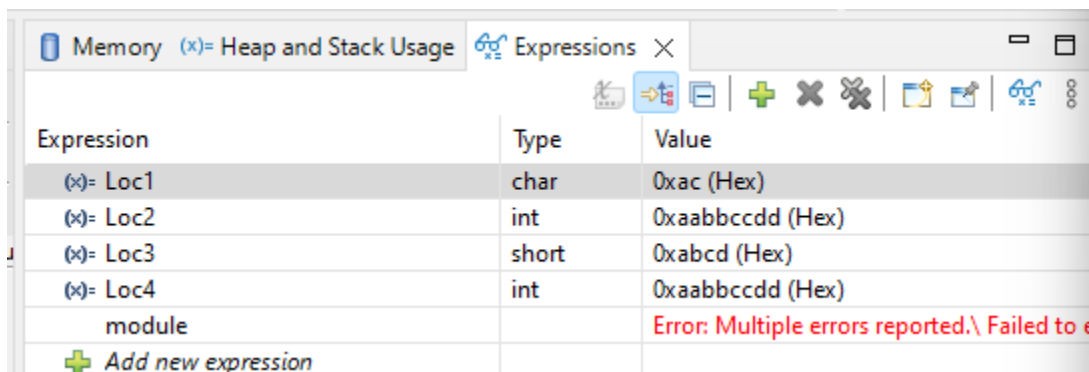
Expression	Type	Value
(x)= x	int	10
▼ ➔ ptr	int *	0x2002ffdc
(x)= *ptr	int	10
▼ ➔ ptr_location	int *	0x20001000
(x)= *ptr_locat	int	11



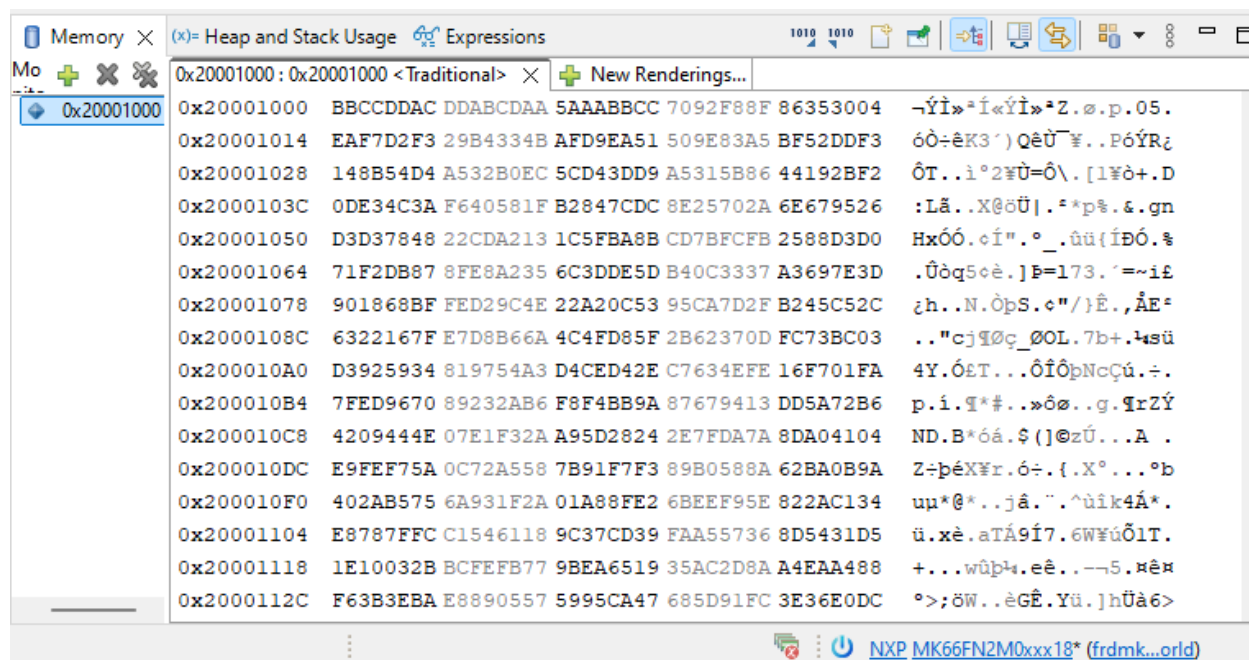
During part B of this experiment, we demonstrated how to efficiently access and manipulate specific memory locations using direct typecasting and dereferencing. Initially, a pointer (`ptr_location`) was declared and set to the memory address `0x20001000`, where the value

11 was written. Using direct typecasting, the value 12 was written to another memory location, 0x20001004, without creating a separate pointer. The value at this location was then read and stored in the variable x. The screenshots validate the correctness of our implementation: the Expressions Panel shows x = 12, confirming that the value at 0x20001004 was successfully read, and ptr_location = 0x20001000, confirming the pointer's correct assignment. Additionally, the Memory Panel displays 11 at address 0x20001000 and 12 at 0x20001004, verifying that both memory locations were modified as intended. These results align with the experiment's objectives, demonstrating the efficiency of typecasting and the ability to manipulate memory on the MCU directly.

Problem 1:



Expression	Type	Value
(x)= Loc1	char	0xac (Hex)
(x)= Loc2	int	0xaabbccdd (Hex)
(x)= Loc3	short	0xabcd (Hex)
(x)= Loc4	int	0xaabbccdd (Hex)
module		Error: Multiple errors reported.\ Failed to e
+ Add new expression		



Address	Hex Data	ASCII Data
0x20001000	BBCCDDAC DDABCDAA 5AAABBC 7092F88F 86353004	-Yi»=i«Yi»=Z.ø.p.05.
0x20001014	EAF7D2F3 29B4334B AFD9EA51 509E83A5 BF52DDF3	ôð÷êK3') QêÛ=...PóYRç
0x20001028	148B54D4 A532B0EC 5CD43DD9 A5315B86 44192BF2	ÔT...i°2¥Û=Ô\.[1¥ð+.D
0x2000103C	0DE34C3A F640581F B2847CDC 8E25702A 6E679526	:Lă...X@öü .°*p%.&.gn
0x20001050	D3D37848 22CDA213 1C5FBA8B CD7BFCFB 2588D3D0	HxÓÓ.°í".°_üü{ÍĐÓ.°
0x20001064	71F2DB87 8FE8A235 6C3DDE5D B40C3337 A3697E3D	.Ûðq5çè.}B=173.'=~if
0x20001078	901868BF FED29C4E 22A20C53 95CA7D2F B245C52C	çh...N.òpS.°"/}Ê.,ÅE=
0x2000108C	6322167F E7D8B66A 4C4FD85F 2B62370D FC73BC03	.."c]qç_00L.7b+.4sü
0x200010A0	D3925934 819754A3 D4CED42E C7634EFE 16F701FA	4Y.ÓêT...ÔîôpNcÇú.÷.
0x200010B4	7FED9670 89232AB6 F8F4BB9A 87679413 DD5A72B6	p.í.¶*#...»ðø..g.¶rZÝ
0x200010C8	4209444E 07E1F32A A95D2824 2E7FDA7A 8DA04104	ND.B*óá.\$(@zÛ...A .
0x200010DC	E9FEF75A 0C72A558 7B91F7F3 89B0588A 62BA0B9A	Z÷péX¥r.ó÷.{.X°...°b
0x200010F0	402AB575 6A931F2A 01A88FE2 6BEEF95E 822AC134	up*@*...já...^ûik4Ä*
0x20001104	E8787FFC C1546118 9C37CD39 FAA55736 8D5431D5	ü.xè.aTÁ9Í7.6W¥úÔ1T.
0x20001118	1E10032B BCFEFB77 9BEA6519 35AC2D8A A4EAA488	+...wûp4.eê...→5.xêx
0x2000112C	F63B3EBA E8890557 5995CA47 685D91FC 3E36E0DC	°>;öW...èGÊ.Yü.jhÜä6>

In Problem 1, we addressed writing specific values to designated memory locations in the microcontroller's memory. To achieve this, we used macros to define memory locations (MEM_LOC, MEM_LOC_short, and MEM_LOC_char) with data types corresponding to the size of the required values. Each memory location (Loc1, Loc2, Loc3, and Loc4) was assigned its respective value as specified in the table. By leveraging macros, the code is concise and

improves readability while ensuring proper handling of data types for memory alignment. The results, validated through memory inspection tools in the IDE, confirmed that the values were written correctly to the specified addresses with the appropriate configuration of Endianness, Cell Size, and Radix. This solution demonstrates efficient memory manipulation in embedded systems programming.

Problem 2:

The memory configuration used for analyzing the structures is Little Endian with a 4-byte Cell Size and values displayed in Hexadecimal Radix.

For **struct1**, the **char x2** occupies 1 byte at offset 0. Since the following member, **int x1**, requires 4-byte alignment, the compiler adds 3 bytes of padding between **x2** and **x1**, ensuring that **x1** starts at offset 4. The total size of **struct1** is 8 bytes, with 3 bytes of padding.

For **struct2**, the **short x2** occupies 2 bytes at offset 0. The **int x1** requires alignment to a 4-byte boundary, so the compiler adds 2 bytes of padding after **x2**, allowing **x1** to start at offset 4. The total size of **struct2** is also 8 bytes, with 2 bytes of padding.

For **struct3**, the **int x1** occupies 4 bytes at offset 0, while the **short x2** takes 2 bytes at offset 4. Since structures must be aligned to the size of their largest member (4 bytes in this case), the compiler adds 2 bytes of padding at the end of the structure, resulting in a total size of 8 bytes with 2 bytes of padding.

Finally, for **struct4**, the nested **inner_struct** comprises three members: **char x1** (1 byte), **short x2** (2 bytes with 1 byte of padding for alignment), and **int x3** (4 bytes), making its size 8 bytes. The outer **struct4** includes the **inner_struct** (8 bytes) and an **int x1** (4 bytes) at offset 8, with no additional padding required. The total size of **struct4** is 12 bytes, with no extra padding. See the table below for a clear definition of the structure, padding, and padding locations.

Structure	Size (bytes)	Padding (bytes)	Padding Locations
struct1	8	3	Between x2 (char) and x1 (int)
struct2	8	2	Between x2 (short) and x1 (int)
struct3	8	2	At the end of the structure
struct4	12	0	No padding needed

Problem 3:

In Problem 3, we modified the code to toggle three LEDs—`LEDRGB_BLUE`, `LEDRGB_GREEN`, and `LEDRGB_RED`—sequentially in the order BLUE → GREEN → RED, repeating in a loop. The changes involved configuring the necessary GPIO pins and toggling them using a delay loop to achieve a visible blinking effect.

In the `BOARD_InitPins` function, we enabled the clock for both Port C and Port D using the `CLOCK_EnableClock` function. We then configured the pins `PTC8`, `PTC9`, and `PTD1` to operate as GPIO using the `PORT_SetPinMux` function. These pins correspond to the blue, green, and red LEDs, respectively.

In the `main` function, we initialized the LEDs as digital output pins using the `GPIO_PinInit` function. A simple delay function (`delay()`) was implemented using a for-loop with `NOP` instructions to introduce a delay between each LED toggle. The LEDs were toggled sequentially using the `GPIO_PortToggle` function, ensuring that each LED blinked on and off before moving to the next LED in the sequence. This process was repeated indefinitely in the main loop.

Problem 4:

In Problem 4, a custom GPIO driver was developed to replace the default SDK driver (`fsl_gpio.h`) for managing GPIO functionality. This driver consists of two key components: a header file (`gpio_led_output.h`) and a source file. The header file defines the `GPIO_Types` structure, which represents six key GPIO registers: `PDOR`, `PSOR`, `PCOR`, `PTOR`, `PDIR`, and `PDDR`. These registers are used to control GPIO operations such as setting, clearing, toggling, and reading pin values, as well as configuring pin directions. The header file also includes the prototypes for three helper functions: `GPIONTogglePin`, `GPIOClearPin`, and `GPIONPinInit`.

The `GPIONTogglePin` function toggles the logic state of a specific GPIO pin by writing to the `PTOR` register, while `GPIOClearPin` clears all pins in a given port by writing `0xFFFF` to the `PCOR` register. The `GPIONPinInit` function configures a pin as input or output by modifying the `PDDR` register, using a configuration structure to specify the desired direction. These functions provide direct and efficient control over GPIO operations, bypassing the overhead of the higher-level SDK driver.

In the main function, the custom driver was used to toggle three LEDs—blue, green, and red—sequentially in the order BLUE (`PTC8`) → GREEN (`PTC9`) → RED (`PTD1`). The pins were first initialized as digital output using the `GPIONPinInit` function. A delay function was implemented to introduce a visible blinking effect, and the LEDs were toggled sequentially using `GPIONTogglePin`, with `GPIOClearPin` resetting the pins after each toggle. The process repeated in an infinite loop, demonstrating the effectiveness of the custom driver in managing GPIO operations.

This implementation highlights the importance of low-level programming in embedded systems by directly manipulating memory-mapped registers. The custom GPIO driver provides greater control over hardware, reduces abstraction layers, and enhances understanding of how embedded systems interact with peripherals. This approach not only optimizes performance but also serves as a practical learning experience for designing hardware-specific drivers.

Problem 5:

In Problem 5, we extended the use of Pulse-Width Modulation (PWM) to control the brightness of three LEDs—red, green, and blue—simultaneously, based on an RGB hex-code input from the user. The FlexTimer Module (FTM) was configured to generate PWM signals for each LED, allowing precise control of their brightness levels to produce a wide range of colors. The pins corresponding to the LEDs (PTD1, PTC8, and PTC9) were configured in the `BOARD_InitPins` function using the `PORT_SetPinMux` function to assign their alternative functionalities as PWM outputs. Each pin was mapped to a specific FTM channel: `FTM3_CH1` for the red LED, `FTM3_CH5` for the green LED, and `FTM3_CH4` for the blue LED.

The `pwm_setup` function initialized the FTM for each LED by configuring the PWM parameters using the `ftm_chnl_pwm_signal_param_t` structure. The parameters included the channel number, a high-true PWM signal level, an initial duty cycle of 0%, and a frequency of 5 kHz, which ensured the LEDs did not flicker. The PWM signals were set up and started using the `FTM_SetupPwm` and `FTM_StartTimer` functions.

In the `main` function, the RGB hex-code input was read from the user, parsed into three separate integer values, and scaled from a range of 0 to 255 into percentages (0 to 100) to match the PWM duty cycle. The `FTM_UpdatePwmDutycycle` function was used to update the duty cycle of each LED, and a software trigger ensured the changes were applied immediately. For example, an input of `FFFF00` resulted in a 100% duty cycle for the red and green LEDs and 0% for the blue LED, producing a yellow color. The program executed continuously in an infinite loop, maintaining the LED brightness levels based on the user-provided input.

This experiment demonstrated the practical application of PWM for controlling LED brightness and dynamically adjusting colors. It highlighted the use of the FTM module for generating independent PWM signals for multiple channels and showcased how hardware configurations, such as pin multiplexing and alternative functionalities, enable efficient control of peripheral devices in embedded systems. By incorporating an intuitive RGB hex-code interface, the implementation provided a user-friendly way to manage LED colours dynamically.