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## Lab 3

### Introduction:

- For this lab, we wrote a program that reads a grammar and finds its canonical LR(0) sets. This program outputs the list of all productions in G and the canonical LR(0) collection of sets for G. In addition to numbering the sets, it identifies the symbol X that generates the set through the function GOTO(I,X). The test results used for this program are grammar G419 and at least one other grammar.

### Source Code:

- **lab3.cpp**

```
/*
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 *
 **/

#include "Parser.h"
#include <iostream>

int main(int argc, char *argv[])
{
    if (argc < 2)
    {
        std::cout << "Please run the program in this format: ./lab3
<filename>\n";
        return 0;
    }

    std::string grammer = argv[1];

    Parser *parser = new Parser(grammer);

    parser->parse();

    //parser->printFirst();

    //parser->printFollow();

    parser->printCanonicalSet();
```

```

    delete parser;

    return 0;
}

```

#### - Parser.cpp

```

#include "Parser.h"
#include "Utility.h"
#include "Item.h"

#include <iostream>
#include <string>
#include <map>
#include <set>
#include <fstream>
#include <algorithm>

#define EPLISON 'e'

Parser::Parser(std::string &grammerFile) : grammer(grammerFile){};

void Parser::parse()
{
    std::ifstream inputGrammerFile;

    inputGrammerFile.open(grammer.c_str());

    if (!inputGrammerFile.is_open())
    {
        std::cerr << "Error: " << grammer << " not found." << '\n';
        exit(-1);
    }

    std::string input = "";

    int productionsPositionInTheGrammer = 0;

    while (inputGrammerFile >> input)
    {
        if (input == "$") // End of non-terminals
        {
            productionsPositionInTheGrammer =
inputGrammerFile.tellg();
            break;
        }
        else
        {
            // Add FIRST(X) = {X}

```

```

        first[input[0]].insert(input[0]);
        terminals.insert(input[0]);
    }
}

inputGrammarFile.seekg(productionsPositionInTheGrammar);
input.clear();

// Now add the productions
while (inputGrammarFile >> input)
{
    productions.insert(input);
}

// Do not need the $ in our productions
// since we just want the productions
productions.erase(productions.find("$"));

for (auto p : productions)
{
    std::string currProd = p;
    char nonTerminal = currProd[0];
    nonTerminals.insert(nonTerminal);

    std::string rhs = currProd.substr(3);
    size_t index = 0;

    // If  $X \rightarrow \hat{\mu}$  is production, add  $\hat{\mu}$  to FIRST(X)
    while (index < rhs.length())
    {
        if (rhs[index] == EPLISON)
        {
            first[nonTerminal].insert(EPLISON);
        }
        index++;
    }
}

//computeFirst();
//computeFollow();
findCanonicalSet();
}

// Computes the FIRST set following the algorithm
// given in the CSE 570/670 website
void Parser::computeFirst()
{
    do
    {
        changed = false;
    }
}

```

```

// Go through the productions
for (auto p : productions)
{
    std::string currProduction = p;
    char nonTerminal = currProduction[0];
    std::string rhs = currProduction.substr(3);

    for (size_t i = 0; i < rhs.length(); i++)
    {
        // Retrieve a set from FIRST(X)
        std::set<char> temp = first[rhs[i]];

        // Add everything in FIRST(Y_i) except for Îµ and
        // increment i
        if (util.hasEpsilon(temp))
        {
            temp.erase(temp.find(EPLISON));
            std::set<char> unionSet =
util.setUnion(first[nonTerminal], temp);

            for (auto s : unionSet)
            {
                auto checkInsert =
first[nonTerminal].insert(s);

                // Check if a set in FIRST has been changed
after inserting
                // a new nonTerminal key
                if (checkInsert.second)
                {
                    changed = true;
                }
            }

            i++;
        }
        // Else FIRST(X) = {X}
        else
        {
            std::set<char> unionSet =
util.setUnion(first[nonTerminal], temp);

            for (auto s : unionSet)
            {
                auto checkInsert =
first[nonTerminal].insert(s);

                if (checkInsert.second)
                {

```

```

        changed = true;
    }
    }
    break;
}
// If i > k then add  $\hat{\mu}$  to FIRST(X)
if (i >= rhs.length())
{
    auto checkInsert =
first[nonTerminal].insert(EPLISON);

    if (checkInsert.second)
    {
        changed = true;
    }
}
i++;
}
}
} while (changed);
}

// Computes the FOLLOW set following the algorithm
// given in the CSE 570/670 website
void Parser::computeFollow()
{
    do
    {
        changed = false;

        // Go through the productions of G
        for (auto p : productions)
        {
            std::string currProduction = p;
            char nonTerminal = currProduction[0];
            std::string rhs = currProduction.substr(3);

            // Go through the characters in the right hand side of
the production
            // e.g. E -> T + F where T + F is the RHS
            for (size_t i = 0; i < rhs.length(); i++)
            {
                // Check if the current char in the RHS is a
non-terminal
                // and making sure we still have input to process on
the RHS
                if (util.isNonTerminal(rhs[i]) && i < rhs.length() -
1)
                {
                    // Retrieve a set in FIRST based the non-terminal

```

```

// we get from the RHS of the production
std::set<char> temp = first[rhs[i + 1]];

// If  $A \rightarrow \hat{A}B\hat{A}^2$  and  $\hat{A}$  is in  $FIRST(\hat{A}^2)$  then
// add everything in  $FOLLOW(A)$  to  $FOLLOW(B)$ 
if (util.hasEpsilon(temp))
{
    std::set<char> nonTerminalFollowSet =
follow[nonTerminal];
    for (auto n : nonTerminalFollowSet)
    {
        auto checkInsert =
follow[rhs[i]].insert(n);
        if (checkInsert.second)
        {
            changed = true;
        }
    }
    // Else, just add everything in  $FIRST(\hat{A}^2)$  except
 $\hat{A}$  to  $FOLLOW(B)$ 
    for (auto t : temp)
    {
        auto checkInsert = follow[rhs[i]].insert(t);
        if (checkInsert.second)
        {
            changed = true;
        }
    }
    // If we reached the last character in the RHS
    // then just add everything in  $FOLLOW(A)$  to  $FOLLOW(B)$ 
else if (util.isNonTerminal(rhs[i]) && i ==
rhs.length() - 1)
{
    std::set<char> temp = follow[nonTerminal];
    for (auto t : temp)
    {
        auto checkInsert = follow[rhs[i]].insert(t);
        if (checkInsert.second)
        {
            changed = true;
        }
    }
}
i++;
}
} while (changed);

```

```

    // Add the $ symbol to FOLLOW(S)
    follow['S'].insert('$');
}

std::set<Item> Parser::closure(std::set<Item> items)
{
    std::set<Item> closure;

    // Add every item to Closure(i)
    for (auto i : items)
    {
        closure.insert(i);
    }

    do
    {
        changed = false;
        for (auto c : closure)
        {
            std::string currProduction = c.getProduction();
            int dotPos = c.getDotPos();

            char lookAhead = currProduction[dotPos];

            // If A -> Î±BÎ² is in Closure(j)
            if (util.isNonTerminal(lookAhead))
            {
                for (auto p : productions)
                {
                    // If B is a production by using a look ahead
                    if (p[0] == lookAhead)
                    {
                        Item temp(p, 3);

                        // Add B to Closure(j)
                        auto checkInsert = closure.insert(temp);
                        if (checkInsert.second)
                        {
                            changed = true;
                        }
                    }
                }
            }
        }
    } while (changed);

    return closure;
}

```

```

std::set<Item> Parser::getGoto(std::set<Item> items, char symbol)
{
    // The empty set j
    std::set<Item> j;

    for (auto item : items)
    {
        std::string currProduction = item.getProduction();
        int dotPos = item.getDotPos();

        char lookAhead = currProduction[dotPos];

        if (symbol == lookAhead)
        {
            Item temp(currProduction, dotPos + 1);
            j.insert(temp);
        }
    }
    return closure(j);
}

void Parser::findCanonicalSet()
{
    std::set<char> symbols;

    std::string startProduction = "S->E";

    // Add the first production S->E to the canonical set to
    // generate the correct item sets
    auto firstItem = Item(startProduction, 3);
    std::set<Item> tempItem;
    tempItem.insert(firstItem);
    auto firstClosure = closure(tempItem);
    canonicalSet.insert(LRSet(firstClosure, 0, '\0'));

    for (auto p : productions)
    {
        auto currentItem = Item(p, 3);
        std::set<Item> tempItem;
        tempItem.insert(currentItem);
        auto currentClosure = closure(tempItem);

        symbols = util.setUnion(terminals, nonTerminals);

        int id = 1;
        do
        {
            changed = false;
            for (auto item : canonicalSet)
            {

```



```

        for (auto symbol : symbols)
        {
            auto temp = getGoto(item.getClosure(), symbol);

            // If GOTO(I, X) is not empty && not in the
canonicalSet
            // then add the GOTO(I, X) to the canonicalSet
            if (!temp.empty() && !isIn(canonicalSet, temp))
            {
                LRSet lrset(temp, id, symbol);
                canonicalSet.insert(lrset);
                changed = true;
                ++id;
            }
        }
    } while (changed);
}

// Checks if a set of items is in the LRSet
bool Parser::isIn(std::set<LRSet> lrset, std::set<Item> items)
{
    for (auto curr : lrset)
    {
        if (curr.getClosure() == items)
            return true;
    }

    return false;
}

void Parser::printFirst() const
{
    std::cout << "FIRST = " << '\n';
    for (auto f : first)
    {
        std::cout << f.first << " -> ";
        std::cout << "{ ";
        for (auto s : f.second)
        {
            std::cout << s << ' ';
        }
        std::cout << "}";
        std::cout << '\n';
    }
}

void Parser::printFollow() const
{

```

```

std::cout << "FOLLOW = " << '\n';
for (auto f : follow)
{
    std::cout << f.first << " -> ";
    std::cout << "{ ";
    for (auto s : f.second)
    {
        std::cout << s << ' ';
    }
    std::cout << "}";
    std::cout << '\n';
}
}

void Parser::printProductions() const
{
    std::cout << "The productions in " << grammer << " are:" << '\n';
    for (auto p : productions)
    {
        std::cout << p << '\n';
    }
    std::cout << '\n';
}

void Parser::printCanonicalSet() const
{
    for (auto c : canonicalSet)
    {
        auto temp = c.getClosure();
        std::cout << "Item: " << c.getId() << '\n';
        std::cout << "Symbol: " << c.getSymbol() << '\n';
        for (auto t : temp)
        {
            std::string currProd = t.getProduction();
            currProd.insert(currProd.begin() + t.getDotPos(), '.');
            std::cout << "Production: [" << t.getDotPos() << "] = "
<< currProd << '\n';
        }
        std::cout << '\n';
    }
}

```

#### - Item.cpp

```

#include "Item.h"
#include <string>

std::string Item::getProduction()
{
    return production;
}

```

```

}

int Item::getDotPos()
{
    return dotPos;
}

bool Item::operator==(const Item &other) const
{
    return other.production == production;
}

bool Item::operator<(const Item &other) const
{
    return other.production > production || dotPos < other.dotPos;
}

Item &Item::operator=(const Item &other)
{
    production = other.production;
    dotPos = other.dotPos;
    return *this;
}

```

#### - Utility.cpp

```

#include "Utility.h"

#include <set>
#include <cctype>

#define EPLISON 'e'

// Checks if a set has an eplison
bool Utility::hasEplison(std::set<char> &s)
{
    return s.find(EPLISON) != s.end();
}

// Combines two sets which one set will contain
// all of the elements from the two sets
std::set<char> Utility::setUnion(std::set<char> &s1, std::set<char> &s2)
{
    std::set<char> result;
    result.insert(s1.begin(), s1.end());
    result.insert(s2.begin(), s2.end());
    return result;
}

```

```

// Checks if a character is an non-terminal
// where a non-terminal is uppercase and alphanumeric
bool Utility::isNonTerminal(char input)
{
    return std::isalpha(input) && std::isupper(input);
}

```

#### - LRSet.cpp

```

#include "LRSet.h"
#include "Item.h"

#include <set>

std::set<Item> LRSet::getClosure() const
{
    return closure;
}

int LRSet::getId() const
{
    return id;
}

char LRSet::getSymbol() const
{
    return symbol;
}

bool LRSet::operator<(const LRSet &other) const
{
    return id < other.id;
}

```

#### - Parser.h

```

#pragma once
#ifndef PARSER_H
#define PARSER_H

#include "Utility.h"
#include "Item.h"
#include "LRSet.h"

#include <string>
#include <map>
#include <set>
#include <fstream>

class Parser

```

```

{
public:
    Parser(std::string &grammarFile);

    void parse();
    void computeFirst();
    void computeFollow();
    void findCanonicalSet();

    bool isIn(std::set<LRSet> lrset, std::set<Item> items);

    std::set<Item> closure(std::set<Item> items);
    std::set<Item> getGoto(std::set<Item> items, char symbol);

    void printFirst() const;
    void printFollow() const;
    void printProductions() const;
    void printCanonicalSet() const;

private:
    std::string grammar;

    std::set<std::string> productions;
    std::set<char> nonTerminals;
    std::set<char> terminals;
    std::map<char, std::set<char> > first;
    std::map<char, std::set<char> > follow;
    std::set<LRSet> canonicalSet;

    Utility util;

    bool changed;
};

#endif

```

#### - Item.h

```

#pragma once
#ifndef ITEM_H
#define ITEM_H

#include <string>

class Item
{
public:
    Item(std::string &production, int dotPos) : production(production),
    dotPos(dotPos) {};

```

```

std::string getProduction();
int getDotPos();

bool operator==(const Item &) const;
bool operator<(const Item &) const;
Item &operator=(const Item &);

private:
    std::string production;
    int dotPos;
};

#endif

```

#### - Utility.h

```

#pragma once
#ifndef UTILITY_H
#define UTILITY_H

#include <set>

class Utility
{
public:
    Utility(){};
    bool hasEpsilon(std::set<char> &s);
    std::set<char> setUnion(std::set<char> &s1, std::set<char> &s2);
    bool isNonTerminal(char input);

};

#endif

```

#### - LRSet.h

```

#pragma once
#ifndef LRSET_H
#define LRSET_H

#include "Item.h"
#include <set>

class LRSet
{
public:
    LRSet(std::set<Item> closure, int id, char symbol)
        : closure(closure), id(id), symbol(symbol){};

    std::set<Item> getClosure() const;

```

```
int getId() const;
char getSymbol() const;

bool operator<(const LRSet &) const;

private:
    std::set<Item> closure;
    int id;
    char symbol;
};

#endif
```

- G419.txt

**i**  
**+**  
**\***  
**(**  
**)**  
**\$**  
**S->E**  
**E->E+T**  
**E->T**  
**T->T\*F**  
**T->F**  
**F->(E)**  
**F->i**  
**\$**

- samplegrammar.txt

**(**  
**)**  
**+**  
**\***  
**e**  
**i**  
**\$**  
**E->TX**  
**X->+E**  
**T->iY**  
**Y->\*T**  
**X->e**  
**T->(E)**  
**Y->e**  
**\$**



## Outputs:

- Test results for grammar G419:

```
[005319687@csusb.edu@jlb359-2 Lab3]$ ./lab3 G419.txt
Item: 0
Symbol:
Production: [3] = E->.E+T
Production: [3] = E->.T
Production: [3] = F->.(E)
Production: [3] = F->.i
Production: [3] = S->.E
Production: [3] = T->.F
Production: [3] = T->.T*F

Item: 1
Symbol: (
Production: [3] = E->.E+T
Production: [3] = E->.T
Production: [3] = F->.(E)
Production: [3] = F->.i
Production: [3] = T->.F
Production: [3] = T->.T*F
Production: [4] = F->(.E)

Item: 2
Symbol: E
Production: [4] = E->E.+T
Production: [4] = S->E.

Item: 3
Symbol: F
Production: [4] = T->F.

Item: 4
Symbol: T
Production: [4] = E->T.
Production: [4] = T->T.*F

Item: 5
Symbol: i
Production: [4] = F->i.

Item: 6
Symbol: E
Production: [4] = E->E.+T
Production: [5] = F->(E.)
```

Item: 7  
Symbol: +  
Production: [3] =  $F \rightarrow \cdot (E)$   
Production: [3] =  $F \rightarrow \cdot i$   
Production: [3] =  $T \rightarrow \cdot F$   
Production: [3] =  $T \rightarrow \cdot T * F$   
Production: [5] =  $E \rightarrow E + \cdot T$

Item: 8  
Symbol: \*  
Production: [3] =  $F \rightarrow \cdot (E)$   
Production: [3] =  $F \rightarrow \cdot i$   
Production: [5] =  $T \rightarrow T * \cdot F$

Item: 9  
Symbol: )  
Production: [6] =  $F \rightarrow (E) \cdot$

Item: 10  
Symbol: T  
Production: [4] =  $T \rightarrow T \cdot * F$   
Production: [6] =  $E \rightarrow E + T \cdot$

Item: 11  
Symbol: F  
Production: [6] =  $T \rightarrow T * F \cdot$

- Test results for one other sample grammar:

```
[005319687@csusb.edu~]$ g++ -o lab3 lab3.cpp Item.cpp Parser.cpp Utility.cpp LRSet.cpp
[005319687@csusb.edu~]$ ./lab3 samplegrammar.txt
Item: 0
Symbol:
Production: [3] = E->.TX
Production: [3] = S->.E
Production: [3] = T->.(E)
Production: [3] = T->.iY

Item: 1
Symbol: (
Production: [3] = E->.TX
Production: [3] = T->.(E)
Production: [3] = T->.iY
Production: [4] = T->(.E)

Item: 2
Symbol: E
Production: [4] = S->E.

Item: 3
Symbol: T
Production: [3] = X->.+E
Production: [3] = X->.e
Production: [4] = E->T.X

Item: 4
Symbol: i
Production: [3] = Y->.*T
Production: [3] = Y->.e
Production: [4] = T->i.Y

Item: 5
Symbol: E
Production: [5] = T->(E.)

Item: 6
Symbol: +
Production: [3] = E->.TX
Production: [3] = T->.(E)
Production: [3] = T->.iY
Production: [4] = X->+.E

Item: 7
Symbol: X
Production: [5] = E->TX.
```

Item: 8  
Symbol: e  
Production: [4] =  $X \rightarrow e$ .

Item: 9  
Symbol: \*  
Production: [3] =  $T \rightarrow (E)$   
Production: [3] =  $T \rightarrow iY$   
Production: [4] =  $Y \rightarrow *T$

Item: 10  
Symbol: Y  
Production: [5] =  $T \rightarrow iY$ .

Item: 11  
Symbol: e  
Production: [4] =  $Y \rightarrow e$ .

Item: 12  
Symbol: E  
Production: [5] =  $X \rightarrow +E$ .

Item: 13  
Symbol: T  
Production: [5] =  $Y \rightarrow *T$ .