## Math 320 Homework 1.9 part 2

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## Algorithm description

Consider the following problem. An individual, Charlie, is on a treasure-hunting expedition. He stumbles upon a vast treasure trove, where there is a large, but finite, amount of jewels and valuables. Charlie is faced with a problem, though. He can only carry a certain weight W before his knapsack breaks. Every item has its weight  $w_i$  and value  $v_i$  etched on the side. How does Charlie figure out what to take? He employs the following algorithm, implemented in python on his MacBook Pro.

He fills his knapsack with n random objects  $x_i$  such that  $\sum_{i=1}^n x_i < W$ . He then checks the ratio  $\frac{v_i}{w_i}$  to see if this ratio is higher than the objects already in the backpack. This algorithm will have temporal complexity of O(2n), since he will have to store every items ratio and replace the items in his pack a potentially huge number of times.

He then realizes that there is a better way! He decides not to store every ratio in memory. Rather, he only stores a ratio if it is a new, unseen ratio. This way, he stores a fraction of what he once stored, and only places the most valuable items in his knapsack. This algorithm has a temporal complexity of O(n), since he only has to store a few values and only places the most valuable items in his knapsack, and Charlie goes home rich!

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