Question 1)

**(a)** The array A and the auxiliary array C after line 5 of the COUNTING-SORT algorithm.

A: (indexes range from 0 to 11, left to right)

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| 6 | 0 | 2 | 0 | 1 | 3 | 4 | 6 | 1 | 3 | 2 |

C: (indexes range from 0 to 6, left to right)

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| 2 | 2 | 2 | 2 | 1 | 0 | 2 |

-----------------------------------------------------------------------------------------------------------------

**(b)** The array C after line 8 of the COUNTING-SORT algorithm.

C: (indexes range from 0 to 6, left to right)

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| 2 | 4 | 6 | 8 | 9 | 9 | 11 |

-----------------------------------------------------------------------------------------------------------------

**(c)** The array B and the auxiliary array C after *one* iteration of the loop in lines 10-12 of the COUNTING-SORT algorithm.

B: (indexes range from 0 to 11, left to right)

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
|  |  |  |  |  |  | 2 |  |  |  |  |

C: (indexes range from 0 to 6, left to right)

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| 2 | 4 | 5 | 8 | 9 | 9 | 11 |

-----------------------------------------------------------------------------------------------------------------

**(d)** The array B and the auxiliary array C after *two* iterations of the loop in lines 10-12 of the COUNTING-SORT algorithm.

B: (indexes range from 0 to 11, left to right)

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
|  |  |  |  |  | 2 |  | 3 |  |  |  |

C: (indexes range from 0 to 6, left to right)

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| 2 | 4 | 5 | 7 | 9 | 9 | 11 |

-----------------------------------------------------------------------------------------------------------------

**(e)** The array B and the auxiliary array C after *three* iterations of the loop in lines 10-12 of the COUNTING-SORT algorithm.

B: (indexes range from 0 to 11, left to right)

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
|  |  |  | 1 |  | 2 |  | 3 |  |  |  |

C: (indexes range from 0 to 6, left to right)

|  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- |
| 2 | 3 | 5 | 7 | 9 | 9 | 11 |

-----------------------------------------------------------------------------------------------------------------

**(f)** The final sorted output array B

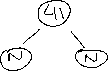
A: (indexes range from 0 to 11, left to right)

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| 0 | 0 | 1 | 1 | 2 | 2 | 3 | 3 | 4 | 6 | 6 |

-----------------------------------------------------------------------------------------------------------------

Question 2)

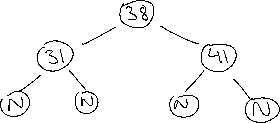
(Insert 41)



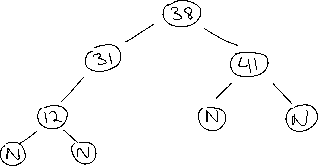
(Insert 38)



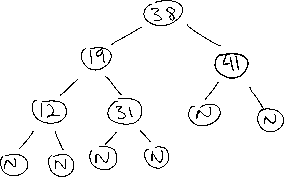
(Insert 31)



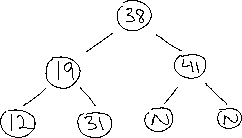
(Insert 12)



(Insert 19)



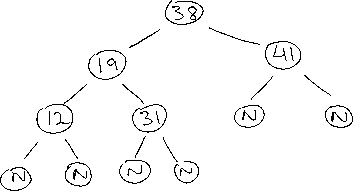
(Insert 8)



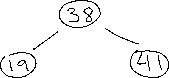
-----------------------------------------------------------------------------------------------------------------

Question 3)

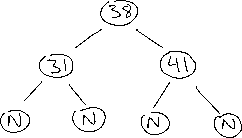
(Delete 8)



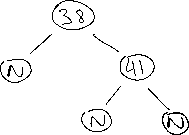
(Delete 12)



(Delete 19)



(Delete 31)



(Delete 38)



(Delete 41)



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Question 4)

Question 5)

We know that a full binary tree has exactly 2*n*-1 nodes. The following pseudo-code describes how to represent any optical prefix code using only 2*n*-1+*n*[lg *n*] bits.

1. Encode the full binary tree structure by performing a pre-order traversal of T
2. FOR EACH node recorded in the traversal {

IF node is an internal node THEN write 0

ELSE (node is a leaf) write 1

}

Since we know the binary tree is full its structure is now uniquely determined. As a property of this tree, we can encode any character of *C* in [lg *n*] bits. Since there are *n* characters, we may encode them in pre-order traversal order using *n* [lg *n*] bits.

Question 6)

We will use a proof by strong induction to prove that every node has rank at most ⌊lg *n*⌋. In the base case where *n* = 1, the nodes rank is 0 = ⌊lg 1⌋. Let’s suppose that the claim holds for 1 to *n* nodes. Given *n*+1 nodes, we may perform a UNION operation on two different sets with *x* and *y* nodes respectively where *x, y* ≤ *n*. With this we know that the root of the set containing *x* nodes has a rank of at most ⌊lg *x*⌋, and that the root of the set containing *y* nodes has a rank of at most ⌊lg *y*⌋. If the ranks of sets *x* and *y* are unequal, then the UNION operation will result in a set that has the same rank as the highest of the two individual sets. If the ranks are indeed equal then the rank of the UNION operation increases the rank by 1, and the resulting set has a rank of ⌊lg *x*⌋ + 1 ≤ ⌊lg (*n*+1)/2⌋ + 1 = ⌊lg (*n*+1)⌋.

Question 7)

We may determine the number of bits necessary to store *x.rank* for each node *x* given that their value will be at most ⌊lg *n*⌋. We may represent *x.rank* using Θ(lg(lg(n))) bits and may need to use Θ(lg(lg(n))) many bits to represent a number that can take that many values.