

## 4th Slide Set Operating Systems

Prof. Dr. Christian Baun

Frankfurt University of Applied Sciences  
(1971–2014: Fachhochschule Frankfurt am Main)  
Faculty of Computer Science and Engineering  
[christianbaun@fb2.fra-uas.de](mailto:christianbaun@fb2.fra-uas.de)

# Learning Objectives of this Slide Set

- At the end of this slide set You know/understand. . .
  - the structure, functioning and characteristics of **Hard Disk Drives**
  - the structure, functioning and characteristics of **Solid State Drives**
  - the functioning and the most commonly implemented variants of Redundant Array of Independent Disks (**RAID**)

Exercise sheet 4 repeats the contents of this slide set which are relevant for these learning objectives

# Hard Disk Drives

- HDDs are approx. 100 times less expensive per bit versus main memory and they offer approx. 100 times more capacity
  - Drawback: Accessing data on HDDs is approx. 1000 times slower
- Reason for the poorer **access time**:
  - HDDs are mechanical devices
    - They contain one or more discs, rotating with 4200, 5400, 7200, 10800, or 15000 revolutions per minute (RPM)
- For each side of each disk, an arm exists with a **read-and-write head**
  - The read-and-write head is used to detect and modify the magnetization of the material
  - The distance between disk and head is approx. 20 nanometers
- Also, HDDs have a cache (usually  $\leq 32$  MB)
  - This cache buffers read and write operations

# Logical Structure of Hard Disk Drives (1/2)

- The disc surfaces are magnetized in circular **tracks** by the heads
- All tracks on all disks at a specific arm position are part of a **cylinder**
- The tracks are divided into logical units (segments of a circle), which are called **blocks** or **sectors**
  - Typically, a sector contains 512 bytes payload
  - Sectors are the smallest addressable units of HDDs
  - If data need be modified, the entire sector must be read and rewritten
- Today, **clusters** are addressed by the software
  - Clusters are groups of sectors with a fixed size, e.g. 4 or 8 kB
  - In modern operating systems, clusters are the smallest addressable unit of HDDs

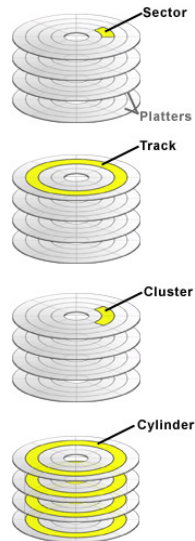
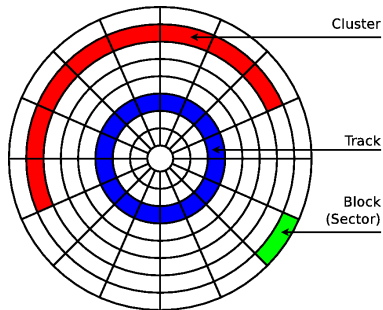
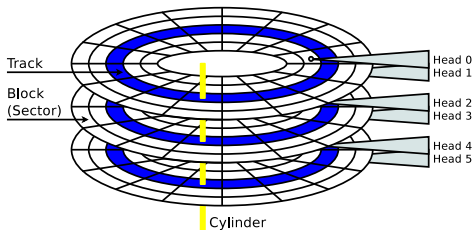


Image source: SweetScape

# Logical Structure of Hard Disk Drives (2/2)



Image Source: <http://www.hitechreview.com>



# Addressing Data on Hard Disk Drives (1/4)

- HDDs with a capacity  $\leq 8$  GB use the **Cylinder-Head-Sector** addressing
  - CHS faces several limitations:
    - The Parallel ATA interface uses 28 bits for CHS addressing and thereof...
      - 16 bits for the cylinders (65,536)
      - 4 bits for the heads (up to 16)
      - 8 bits for the sectors/track (up to 255, because counting starts at 1)
    - The BIOS uses 24 bits for CHS addressing and thereof...
      - 10 bits for the cylinders (1,024)
      - 8 bits for the heads (up to 256)
      - 6 bits for the sectors/track (up to 63, because counting starts at 1)
  - For all these limits, the lowest value is important
    - For this reason, old BIOS revisions can address a maximum of 504 MB
- 1,024 cylinders \* 16 heads \* 63 sectors/track \* 512 bytes/sector = 528,482,304 bytes
  - 528,482,304 bytes / 1024 / 1024 = 504 MB

# Addressing Data on Hard Disk Drives (2/4)



Image Source: <http://www.eak-computers.com>

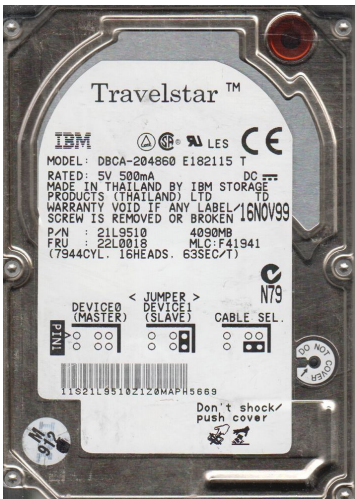
- $1,024 \text{ cylinders} * 16 \text{ heads} * 63 \text{ sectors/track} * 512 \text{ bytes/sector} = 528,482,304 \text{ bytes}$
- $528,482,304 \text{ bytes} / 1024 / 1024 = 504 \text{ MB}$

- Issue: No 2.5" or 3.5" HDD contains  $> 16$  heads
- Solution: Logical heads
  - HDDs usually implement 16 logical heads

⇒ **Extended CHS (ECHS)**

# Addressing Data on Hard Disk Drives (3/4)

Image Source: eBay



- Later BIOS revisions implemented **Extended CHS (ECHS)**
  - Increases (multiplies) the number of heads up to 255 and decreases the number of cylinders by the same factor
  - Result:  $\leq 7.844$  GB can be addressed
- 1,024 cylinders \* 255 heads \* 63 sectors/track \* 512 bytes/sector = 8,422,686,720 bytes
- 8,422,686,720 bytes / 1,024 / 1,024 / 1,024 = 7.844 GB

Better explanation  $\Rightarrow$  see next slide



# Extended CHS (ECHS) – Better Explanation (1/2)

Standard	Max. Cylinders	Max. Heads	Max. Sectors	Max. Capacity
IDE/ATA	65,536	16	256	128 GB
BIOS	1,024	256	63	7.88 GB
Combination (Smaller of Each)	1,024	16	63	504 MB

- The IDE/ATA standard allows more cylinders than the BIOS does, and the BIOS allows more heads than IDE/ATA does
- Remember: These are logical disk parameters, not physical ones
- The BIOS takes the logical geometry that the hard disk specifies according to the IDE/ATA standard, and translates it into an equivalent geometry that will „fit“ into the maximums allowed by the BIOS
- This is done by **dividing the number of logical cylinders by an integer**, and then **multiplying the number of logical heads by the same number**

Source: <http://www.pcguides.com/ref/hdd/bios/modesECHS-c.html>

## Extended CHS (ECHS) – Better Explanation (2/2)

Standard	Max. Cylinders	Max. Heads	Max. Sectors	Max. Capacity
IDE/ATA	65,536	16	256	128 GB
BIOS	1,024	256	63	7.88 GB
Combination (Smaller of Each)	1,024	16	63	504 MB

- Let's take the case of a 3.1 GB Western Digital Caviar hard drive, AC33100
- This drive actually has a capacity of 2.95 binary GB, and logical geometry of 6,136 cylinders, 16 heads and 63 sectors. This is well within the bounds of the IDE/ATA limitations, but exceeds the BIOS limit of 1,024 cylinders
- The BIOS picks a translation factor such that dividing the logical number of cylinders by this number will produce a number of cylinders below 1,024
- Usually one of 2, 4, 8, or 16 are selected; in this case the optimal number is 8
- The BIOS divides the number of cylinders by 8 and multiplies the number of heads by 8
- This results in a translated geometry of 767 cylinders, 128 heads and 63 sectors. The capacity is of course unchanged, and the new geometry fits quite nicely into the BIOS limits

Standard	Max. Cylinders	Max. Heads	Max. Sectors	Max. Capacity
IDE/ATA	65,536	16	256	128 GB
Hard Disk Logical Geometry	6,136	16	63	2.95 GB
BIOS Translation Factor	divide by 8	multiply by 8	—	—
BIOS Translated Geometry	767	128	63	2.95 GB
BIOS	1,024	256	63	7.88 GB

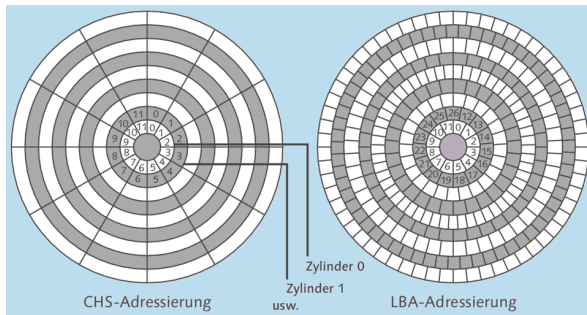
# Addressing Data on Hard Disk Drives (4/4)



- HDDs > 7.844 GB use **Logical Block Addressing (LBA)**
  - All sectors are numbered consecutively beginning with 0
- To ensure compatibility, the first 7.844 GB of all HDDs > 7.844 GB can be addressed via CHS

# Logical Block Addressing (LBA)

Image Source: Sascha Kersken (Rheinwerk Verlag)



- When CHS addressing is used, all tracks contain **the same number of sectors**
  - Each sector stores stores 512 bytes of payload
- Drawback: **Storage capacity is wasted**, because the data density decreases from the inner tracks to the outer tracks
- When LBA is implemented, this drawback does not exist

# Required Time to access Data on HDDs

- The access time is an important performance factor
- 2 factors influence the access time of HDDs
  - ① **Average Seek Time**
    - The time that it takes for the arm to reach a desired track
    - Is for modern HDDs between 5 and 15 ms
  - ② **Average Rotational Latency Time**
    - Delay of the rotational speed, until the required disk sector is located under the head
    - Depends entirely on the rotational speed of the disks
    - Is for modern HDDs between 2 and 7.1 ms

$$\text{Average Rotational Latency Time [ms]} = \frac{30.000}{\text{rotational speed [r/min]}}$$

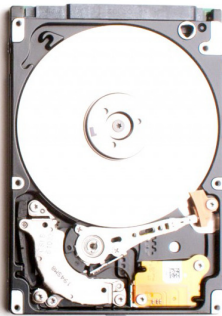
# Solid State Drives (SSD)

Image Source: <http://hardwrk.com>

- Are sometimes falsely called Solid State Disks
- Do not contain moving parts
- Benefits:
  - Fast access time
  - Low power consumption
  - No noise generation
  - Mechanical robustness
  - Low weight
  - The location of data does not matter  $\implies$  defragmenting makes no sense
- Drawbacks:
  - Higher price compared with HDDs of the same capacity
  - Secure delete or overwrite is hard to implement
  - Limited number of program/erase cycles



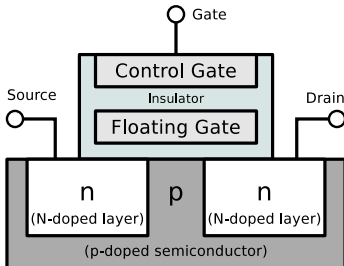
Left image: SSD



Right image: HDD

# Functioning of Flash Memory

- Data is stored as electrical charges
- In contrast to main memory, no electricity is required to keep the data



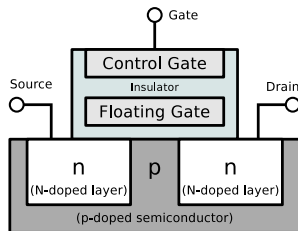
- Each flash memory cell is a transistor and has 3 connectors
  - **Gate** = control electrode
  - **Drain** = electrode
  - **Source** = electrode
- The floating gate stores electrons (data)
  - Completely surrounded by an insulator
  - Electrical charge remains stable for years

Well written explanation about the functioning of flash memory

Benjamin Benz. *Die Technik der Flash-Speicherkarten*. c't 23/2006

## Reading Data from Flash Memory Cells

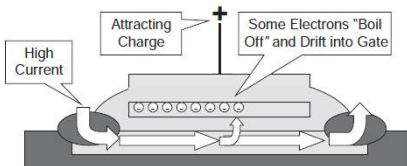
- A positive doped (p) semiconductor separates the 2 negative doped (n) electrodes drain and source
  - Equal to a npn transistor, the npn passage does not conduct without a base current
- Above a certain positive voltage (5V) at the gate (**threshold**) a n-type channel is created in the p-area
  - Current can flow between source and drain through this channel
- If the floating gate contains electrons, the threshold gets changed
  - A higher positive voltage at the gate is required, so that current can flow between source and drain
  - **This way the stored value of the flash memory cell is read out**





# Writing Data from Flash Memory Cells

Image Source: Thomas Coughlin



- Data is stored inside flash memory cells by using **Fowler-Nordheim tunneling**

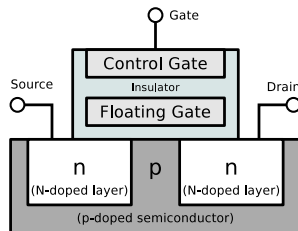
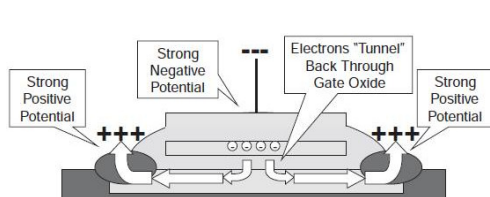
- A positive voltage (5V) is applied to the control gate
  - As a result, electrons can flow between source and drain
- If the high positive voltage is sufficient high (6–20V), some electrons are tunneled ( $\implies$  Fowler-Nordheim tunneling) through the insulator into the floating gate
- This method is also called **Channel Hot Electron Injection**

## Recommended Source

Flash memory. Alex Paikin. 2004. [http://www.hitequest.com/Kiss/Flash\\_terms.htm](http://www.hitequest.com/Kiss/Flash_terms.htm)

# Erasing Data in Flash Memory Cells

Image Source: Thomas Coughlin



- For erasing a flash memory cell, a negative voltage (-6–20V) is applied at the control gate
  - As a result, electrons are tunneled in the reverse direction from the floating gate
  - The insulating layer, which surrounds the floating gate, suffers from each erase cycle
    - At some point the insulating layer is no longer sufficient to hold the charge in the floating gate
    - For this reason, flash memory survives only a limited number of program/erase cycles

# Functioning of Flash Memory

- Memory cells are connected to **blocks** and (depending on the structure also in) **pages**
  - A block always contains a fixed number of pages
- Write and erase operations can only be carried out for entire pages or blocks
  - For this reason, write and erase operations are more complex compared with read operations
- If data inside a page need to be modified, the entire block must be erased
  - To do this, the block is copied to a buffer memory
  - In the buffer memory, the data is modified
  - Next, the block is erased from the flash memory
  - Finally, the modified block is written into the flash memory

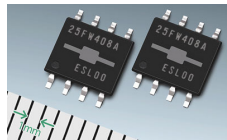
## Different Types of Flash Memory

- 2 types of flash memory exist:
  - **NOR memory**
  - **NAND memory**
- The circuit symbol indicates the way, the memory cells are connected
  - This influences the capacity and latency

# NOR Memory

Image Source: Sanyo

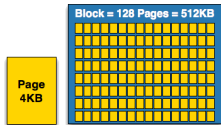
- Each memory cell has its own data line
  - Benefit:
    - Random access for read and write operations  
⇒ Better latency compared with NAND memory
  - Drawback:
    - More complex (⇒ expensive) construction
    - Higher power consumption compared with NAND memory
    - Typically small capacities ( $\leq 32$  MB)
- Does not contain pages
  - The memory cells are grouped together to blocks
    - Typical block sizes: 64, 128 or 256 kB
- No random access for erase operations
  - Erase operations can only be carried out for entire blocks
- Fields of application:
  - Industrial environment
  - Storing the firmware of a computer system



# NAND Memory

Image Source: engadget.com and Samsung

- The memory cells are grouped to pages
  - Typical page size: 512-8192 bytes
    - Each page has it's own data line
  - Each block consists of a number of pages
    - Typical block sizes: 32, 64, 128 or 256 pages
- Benefit:
  - Lesser data lines  $\implies$  requires  $< 50\%$  of the surface area of NOR memory
  - Lower manufacturing costs compared with NOR flash memory
- Drawback:
  - No random access
    - $\implies$  Poorer latency compared with NOR memory
  - Read and write operations can only be carried out for entire pages
  - Erase operations can only be carried out for entire blocks
- Fields of application: USB flash memory drives, SSDs, memory cards



# Single-Level Cell, Multi-Level Cell, Triple-Level Cell

- 3 types of NAND flash memory exist
  - TLC memory cells store 3 bits
  - MLC memory cells store 2 bits
  - SLC memory cells store 1 bit
- SLC storage. . .
  - is most expensive
  - provides the best write speed
  - survives most program/erase cycles

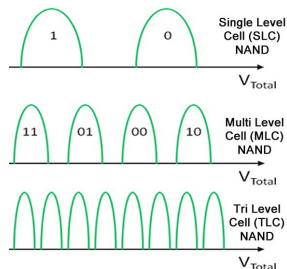


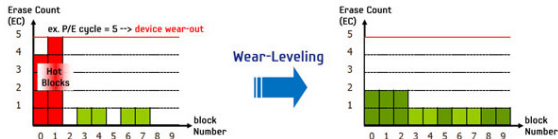
Image Source:  
<http://embedded-computing.com>

- SLC memory survives approx. 100,000 - 300,000 program/erase cycles
- MLC memory survives approx. 10,000 program/erase cycles
- TLC memory survives approx. 1,000 program/erase cycles
- Also memory cells exist, which survive millions of program/erase cycles

# Wear Leveling

Image Source: <http://notebookitalia.it>

- Wear leveling algorithms evenly distribute write operations



- File systems, which are designed for flash memory, and therefore minimize write operations, are e.g. JFFS, JFFS2, YAFFS and LogFS
  - JFFS contains its own wear leveling algorithm
    - This is often required in embedded systems, where flash memory is directly connected



# Latency of Hard Disk Drives

- The performance of CPUs, cache and main memory is growing faster than the data access time (*latency*) of HDDs:
  - **HDDs**
    - 1973: IBM 3340, 30 MB capacity, 30 ms data access time
    - 1989: Maxtor LXTI00S, 96 MB capacity, 29 ms data access time
    - 1998: IBM DHEA-36481, 6 GB capacity, 16 ms data access time
    - 2006: Maxtor STM320820A, 320 GB capacity, 14 ms data access time
    - 2011: Western Digital WD30EZRSDDL, 3 TB capacity, 8 ms data access time
  - **CPUs**
    - 1971: Intel 4004, 740 kHz clock speed
    - 2007: AMD Opteron Santa Rosa F3, 2.8 GHz clock speed
    - 2010: Core i7 980X Extreme (6 Cores), 3.33 GHz clock speed
- The latency of SSDs is  $\leq 1 \mu s \implies \approx$  factor 100 better than HDDs
  - But the gap grows because of interface limitations and multiple CPU cores
- Further challenge
  - Storage drives can fail  $\implies$  risk of data loss
- Enhance **latency** and **reliability** of HDDs and SSDs  $\implies$  **RAID**

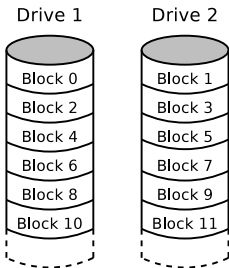
# Redundant Array of independent Disks (RAID)

- The performance of the HDDs can not be improved infinitely
  - HDDs contain moving parts
    - Physical boundaries must be accepted
- One way to avoid the given limitations in terms of speed, capacity and reliability, is the parallel use multiple components
- A RAID consists of multiple drives (HDDs or SSDs)
  - For users and processes, a RAID behaves like a single large drive
- Data is distributed across the drives of a RAID system
  - The RAID level specifies how the data is distributed
    - The most commonly used RAID levels are RAID 0, RAID 1 and RAID 5

Patterson, David A., Garth Gibson, and Randy H. Katz, **A Case for Redundant Arrays of Inexpensive Disks (RAID)**, Vol. 17. No. 3, ACM (1988)

## RAID 0 – Striping – Acceleration without Redundancy

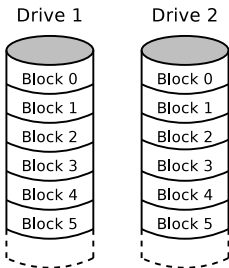
- No redundancy
  - Increases only the data rate
- Drives are split into blocks of equal size
- If read/write operations are big enough ( $> 4$  or  $8$  kB), the operations can be carried out in parallel on multiple drives or on all drives



- In the event of a drive failure, not the entire data can be reconstructed
  - Only small files, which are stored entirely on the remaining drives, can be reconstructed (in theory)
- RAID 0 should only be used when security is irrelevant or backups are created at regular intervals

## RAID 1 – Mirroring

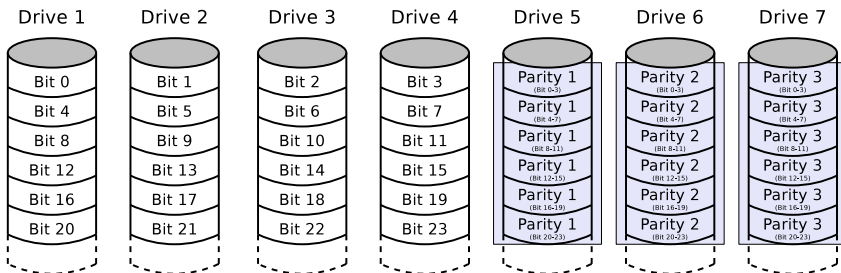
- At least 2 drives of the same capacity store identical data
  - If the drives are of different sizes, RAID 1 provides only the capacity of the smallest drive
- Failure of a single drive does not cause any data loss
  - Reason: The remaining drives store the identical data
- A total loss occurs only in case of the failure of all drives



- Any change of data is written on all drives
- Not a backup replacement
  - Wrong file operations or virus attacks take place on all drives
- The read performance can be increased by intelligent distribution of accesses to the attached drives

## RAID 2 – Bit-Level Striping with Hamming Code for Error Correction

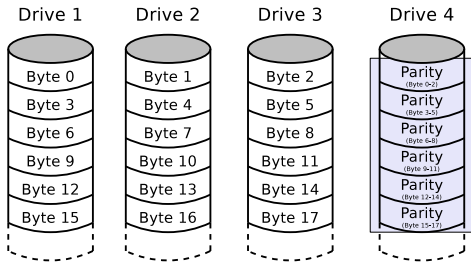
- Each sequential bit is stored on a different drive
  - Bits, which are powers of 2 (1, 2, 4, 8, 16, etc.) are parity bits



- The individual parity bits are distributed across multiple drives
  - ⇒ increases the throughput
- Was used only in mainframes
  - Is no longer relevant

# RAID 3 – Byte-level Striping with Parity Information

- Parity information are stored on a dedicated parity drive

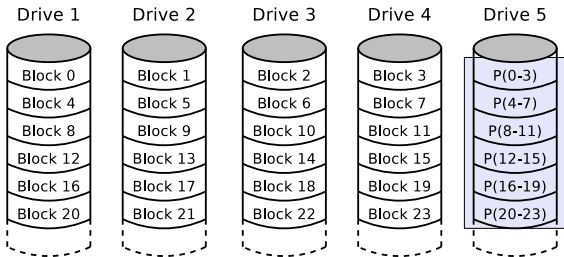


- Each write operation on the RAID causes write operations on the dedicated parity drive  
⇒ bottleneck
- Was replaced by RAID 5

Payload drives		Sum		even/odd		Parity drives
Bits are 0 + 0 + 0	⇒	0	⇒	Sum is even	⇒	Sum bit 0
Bits are 1 + 0 + 0	⇒	1	⇒	Sum is odd	⇒	Sum bit 1
Bits are 1 + 1 + 0	⇒	2	⇒	Sum is even	⇒	Sum bit 0
Bits are 1 + 1 + 1	⇒	3	⇒	Sum is odd	⇒	Sum bit 1
Bits are 1 + 0 + 1	⇒	2	⇒	Sum is even	⇒	Sum bit 0
Bits are 0 + 1 + 1	⇒	2	⇒	Sum is even	⇒	Sum bit 0
Bits are 0 + 1 + 0	⇒	1	⇒	Sum is odd	⇒	Sum bit 1
Bits are 0 + 0 + 1	⇒	1	⇒	Sum is odd	⇒	Sum bit 1

# RAID 4 – Block-level Striping with Parity Information

- Parity information are stored at a dedicated parity drive
- Difference to RAID 3:
  - Not individual bits or bytes, but blocks (**chunks**) are written



- Each write operation on the RAID causes write operations on the dedicated parity drive

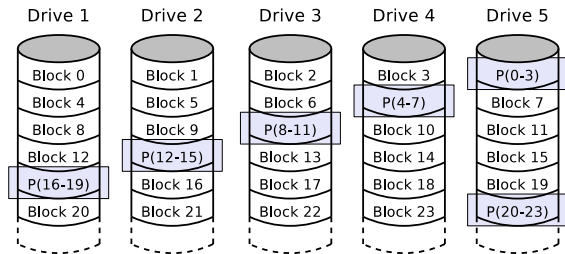
- Drawbacks:
  - Bottleneck
  - Dedicated parity drive fails more frequently

$P(16-19) = \text{Block 16 XOR Block 17 XOR Block 18 XOR Block 19}$

- Seldom implemented, because RAID 5 does not face these drawbacks
- The company NetApp implements RAID 4 in their NAS servers
  - e.g. NetApp FAS2020, FAS2050, FAS3040, FAS3140, FAS6080

## RAID 5 – Block-level Striping with distributed Parity Information

- Payload and parity information are distributed across all drives
- Benefits:
  - High throughput
  - High security level against data loss
  - No bottleneck

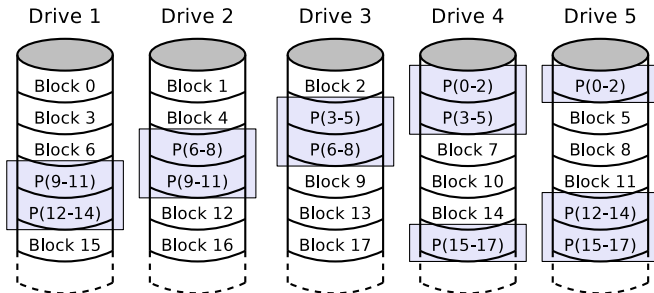


$P(16-19) = \text{block 16 XOR block 17 XOR block 18 XOR block 19}$



## RAID 6 – Block-level Striping with double distributed Parity Information

- Functioning is similar to RAID 5
  - But it can handle the simultaneous failure of up to 2 drives
- In contrast to RAID 5...
  - is the availability better, but the write performance is lower
  - is the effort to write the parity information higher



# Summary of the RAID Levels

RAID	$n$ (number of drives)	$k$ (net capacity)	Allowed to fail	Performance (read)	Performance (write)
0	$\geq 2$	$n$	0 (none)	$n * X$	$n * X$
1	$\geq 2$	capacity of the smallest drive	$n - 1$ drives	$n * X$	$X$
2	$\geq 3$	$n - \lceil \log_2 n \rceil$	1 drive	variable	variable
3	$\geq 3$	$n - 1$	1 drive	$(n - 1) * X$	$(n - 1) * X$
4	$\geq 3$	$n - 1$	1 drive	$(n - 1) * X$	$(n - 1) * X$
5	$\geq 3$	$n - 1$	1 drive	$(n - 1) * X$	$(n - 1) * X$
6	$\geq 4$	$n - 2$	2 drive	$(n - 2) * X$	$(n - 2) * X$

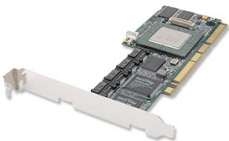
- $X$  is the performance of a single drive during read or write
- The maximum possible performance in theory is often limited by the controller and the computing power of the CPU

# RAID Combinations

- Usually RAID 0, 1 or 5 is used
- In addition to the popular RAID levels, several RAID combinations exist
  - At least 2 RAIDs are combined to a bigger RAID
- Examples:
  - RAID 00: Multiple RAID 0 are connected to a RAID 0
  - RAID 01: Multiple RAID 0 are connected to a RAID 1
  - RAID 05: Multiple RAID 0 are connected to a RAID 5
  - RAID 10: Multiple RAID 1 are connected to a RAID 0
  - RAID 15: Multiple RAID 1 are connected to a RAID 5
  - RAID 50: Multiple RAID 5 are connected to a RAID 0
  - RAID 51: Multiple RAID 5 are connected to a RAID 1

# Hardware / Host / Software RAID (1/2)

Image Source: Adaptec



Adaptec SATA RAID 2410SA



Adaptec SATA II RAID 1220SA

## • Hardware RAID

- A RAID controller with a processor does the calculation of the parity information and monitors the state of the RAID

Benefit: Operating system independent  
No additional CPU load

Drawback: High price (approx. € 200)

## • Host RAID

- Either an inexpensive RAID controller or the chipset provide the RAID functionality
- Usually only supports RAID 0 and RAID 1

Benefit: Operating system independent  
Low price (approx. € 50)

Drawback: Additional CPU load  
Possible dependence of rare hardware

## Hardware / Host / Software RAID (2/2)

- **Software RAID**

- Linux, Windows and MacOS allow to connect drives to a RAID without a RAID controller

Benefit: No cost for additional hardware

Drawback: Operating system dependent  
Additional CPU load

- Example: Create a RAID 1 (md0) with the partitions sda1 and sdb1:

```
mdadm --create /dev/md0 --auto md --level=1  
--raid-devices=2 /dev/sda1 /dev/sdb1
```

- Obtain information about any software RAID in the system:

```
cat /proc/mdstat
```

- Obtain information about a specific software RAID (md0):

```
mdadm --detail /dev/md0
```

- Remove partition sdb1 and add partition sdc1 to the RAID:

```
mdadm /dev/md0 --remove /dev/sdb1  
mdadm /dev/md0 --add /dev/sdc1
```