

Hard Disk Drives (HDD)  
oooooo

Solid State Drives (SSD)  
oooooooooooo

RAID  
ooooooooooooooo

## 4th Slide Set Operating Systems

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# Learning Objectives of this Slide Set

- At the end of this slide set You know/understand...
  - the structure, functioning and characteristics of **Hard Disk Drives**
  - the structure, functioning and characteristics of **Solid State Drives**
  - the functioning and the most commonly implemented variants of Redundant Array of Independent Disks (**RAID**)

Exercise sheet 4 repeats the contents of this slide set which are relevant for these learning objectives

# Hard Disk Drives

- HDDs are approx. 100 times less expensive per bit versus main memory and they offer approx. 100 times more capacity
  - Drawback: Accessing data on HDDs is approx. 1000 times slower
- Reason for the poorer **access time**:
  - HDDs are mechanical devices
    - They contain one or more discs, rotating with 4200, 5400, 7200, 10800, or 15000 revolutions per minute (RPM)
- For each side of each disk, an arm exists with a **read-and-write head**
  - The read-and-write head is used to detect and modify the magnetization of the material
  - The distance between disk and head is approx. 20 nanometers
- Also, HDDs have a cache (usually  $\leq$  32 MB)
  - This cache buffers read and write operations

# Logical Structure of Hard Disk Drives (1/2)

- The surfaces of the **platters** (discs) are magnetized in circular **tracks** by the heads
- All tracks on all disks at a specific arm position are part of a **cylinder**
- The tracks are divided into logical units (segments of a circle), which are called **blocks** or **sectors**
  - Typically, a sector contains 512 bytes payload
  - Sectors are the smallest addressable units of HDDs

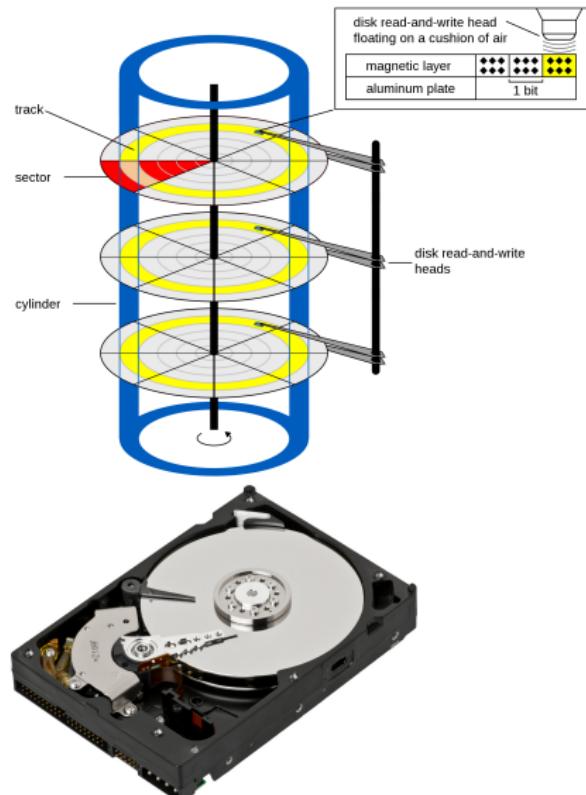


Image source (structure): Henry Mühlfordt. Wikimedia  
(CC-BY-SA-1.0)

Image source (HDD): purepng.com (CC0)

## Logical Structure of Hard Disk Drives (2/2)

- If data need be modified, the entire sector must be read and rewritten
- Today, **clusters** are addressed by the software (see slide set 6)
  - Clusters are groups of sectors with a fixed size, e.g. 4 or 8 kB
  - In modern operating systems, clusters are the smallest addressable unit of HDDs

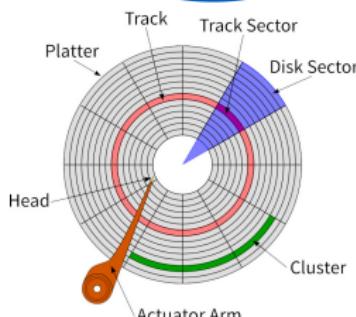
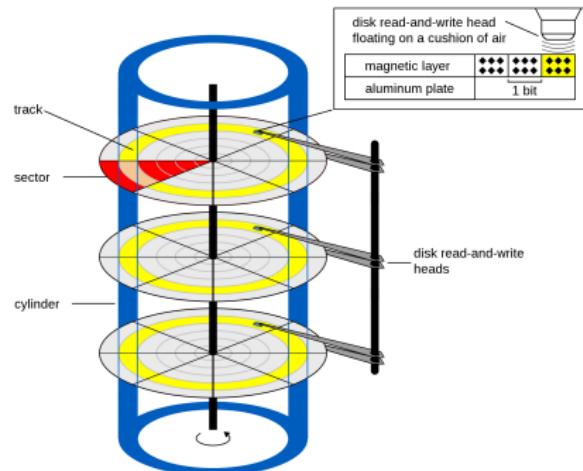


Image source (structure): Henry Mühlfordt. Wikimedia  
(CC-BY-SA-1.0)

Image source (platter): Tim Bielawa. The Linux Sysadmins Guide to  
Virtual Disks (CC-BY-SA-4.0)

# Addressing Data on Hard Disk Drives

- HDDs with a capacity  $\leq$  8 GB use the **Cylinder-Head-Sector addressing**
- CHS faces several limitations:
  - The Parallel ATA interface and the BIOS offer just...
    - 16 bits for the cylinders (up to 65,536)
    - 8 bits for the heads (up to 255. Head number 0 is not used)
    - 8 bits for the sectors/track (up to 255. Sector number 0 is not used)
- $\leq 7.844$  GB can be addressed this way
  - $1024 \text{ cylinders} * 255 \text{ heads} * 63 \text{ sectors/track} * 512 \text{ bytes/sector} = 8,422,686,720 \text{ bytes}$
  - $8,422,686,720 \text{ bytes} / 1024 / 1024 / 1024 = 7844 \text{ GB}$
- No 2.5" or 3.5" HDD contains  $> 16$  heads!!!
  - Logical heads were used
- HDDs with a capacity  $> 7.844$  GB use **Logical Block Addressing (LBA)**
  - All sectors are numbered consecutively beginning with 0

# Logical Block Addressing (LBA)

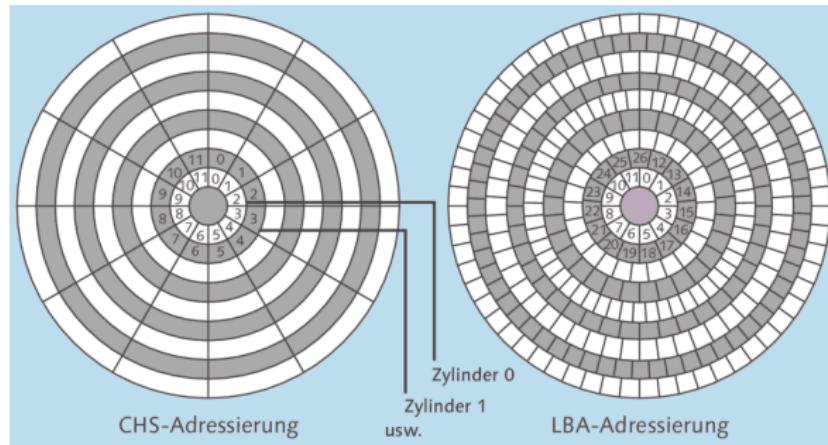


Image source

IT-Handbuch für  
Fachinformatiker.  
Sascha Kersken.  
6th edition.  
Rheinwerk Verlag

- When CHS addressing is used, all tracks contain **the same number of sectors**
  - Each sector stores 512 bytes of payload
- Drawback: **Storage capacity is wasted**, because the data density decreases from the inner tracks to the outer tracks
- When LBA is implemented, this drawback does not exist

# Required Time to access Data on HDDs

- The access time is an important performance factor
- 2 factors influence the access time of HDDs

## ① Average Seek Time

- The time that it takes for the arm to reach a desired track
- Is for modern HDDs between 5 and 15 ms

## ② Average Rotational Latency Time

- Delay of the rotational speed, until the required disk sector is located under the head
- Depends entirely on the rotational speed of the disks
- Is for modern HDDs between 2 and 7.1 ms

$$\text{Average Rot. Lat. Time [ms]} = \frac{1000 \frac{\text{[ms]}}{\text{[sec]}} \times 60 \frac{\text{[sec]}}{\text{[min]}} \times 0.5}{\frac{\text{revolutions}}{\text{[min]}}} = \frac{30,000 \frac{\text{[ms]}}{\text{[min]}}}{\frac{\text{revolutions}}{\text{[min]}}}$$

Why does the equation contain 0.5 ?

Once the head has reached the right track, on average a half rotation of the disk must be waited for the correct sector to be under the head  $\implies$  Average Rotational Latency Time = half Rotational Latency Time

# Solid State Drives (SSD)

- Are sometimes falsely called Solid State Disks

- Do not contain moving parts

- Benefits:

- Fast access time
- Low power consumption
- No noise generation
- Mechanical robustness
- Low weight
- The location of data does not matter  $\Rightarrow$  defragmenting makes no sense



Image (SSD): Thomas Springer. Wikimedia (CC0)

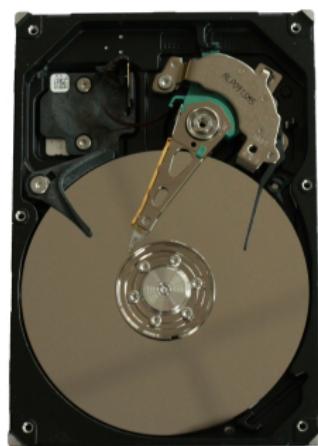


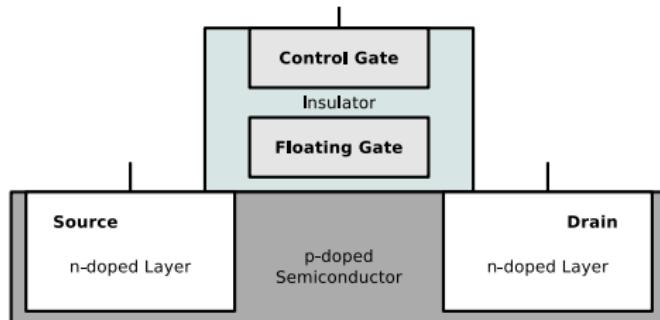
Image (HDD): Erwan Velu. Wikimedia (CC-BY-SA-1.0)

- Drawbacks:

- Higher price compared with HDDs of the same capacity
- Secure delete or overwrite is hard to implement
- Limited number of program/erase cycles

# Functioning of Flash Memory

- Data is stored as electrical charges
- In contrast to main memory, no electricity is required to keep the data
- Each flash memory cell is a transistor and has 3 connectors
  - **Gate** = control electrode
  - **Drain** = electrode
  - **Source** = electrode
- The floating gate stores electrons (data)
  - Completely surrounded by an insulator
  - Electrical charge remains stable for years

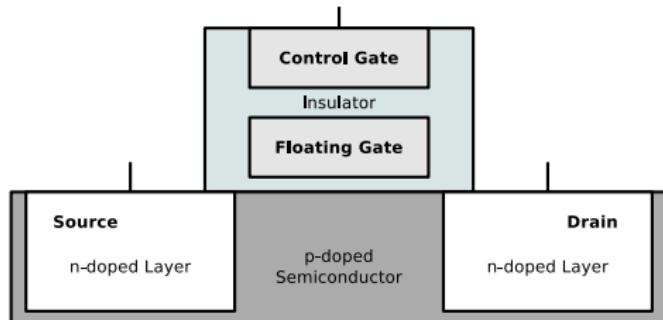


Well written explanation about the functioning of flash memory

Benjamin Benz. *Die Technik der Flash-Speicherkarten.* c't 23/2006

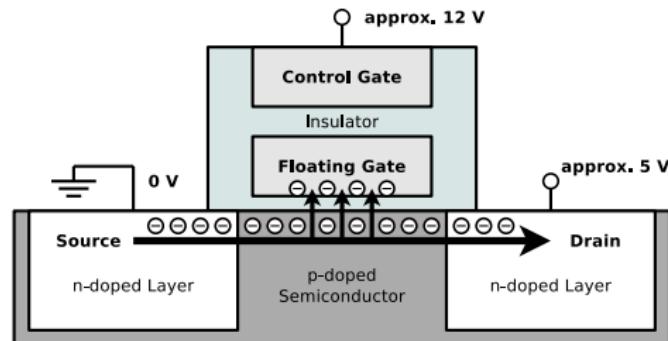
# Reading Data from Flash Memory Cells

- A positively doped (p) semiconductor separates the 2 negatively doped (n) electrodes drain and source
  - Equal to a npn transistor, the npn passage is not conductive without a base current
- Above a certain positive voltage (5V) at the gate (**threshold**) a n-type channel is created in the p-area
  - Current can flow between source and drain through this channel
- If the floating gate contains electrons, the threshold is different
  - A higher positive voltage at the gate is required, so that current can flow between source and drain
    - **This way the stored value of the flash memory cell is read out**



# Writing Data into Flash Memory Cells

- Data is stored inside flash memory cells by using **Fowler-Nordheim tunneling**



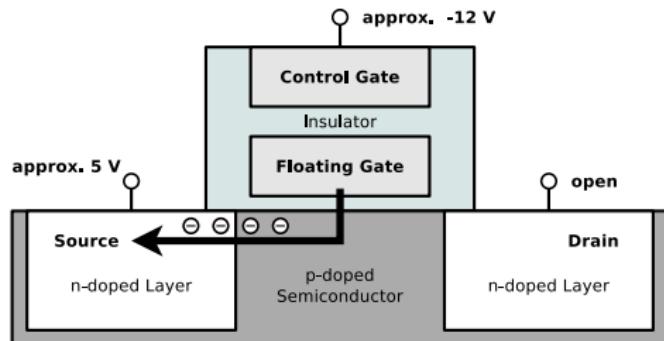
- A positive voltage (5V) is applied to the control gate
  - As a result, electrons can flow between source and drain
- If the high positive voltage is sufficient high (6 to 20V), some electrons are tunneled ( $\Rightarrow$  Fowler-Nordheim tunneling) through the insulator into the floating gate
- This method is also called **Channel Hot Electron Injection**

## Recommended Source

Flash memory. Alex Paikin. 2004. [http://www.hitequest.com/Kiss/Flash\\_terms.htm](http://www.hitequest.com/Kiss/Flash_terms.htm)

# Erasing Data in Flash Memory Cells

- For erasing a flash memory cell, a negative voltage (-6 to -20V) is applied at the control gate
  - As a result, electrons are tunneled in the reverse direction from the floating gate
- The insulating layer, which surrounds the floating gate, suffers from each erase cycle
  - At some point the insulating layer is no longer sufficient to hold the charge in the floating gate
  - For this reason, flash memory survives only a limited number of program/erase cycles



# Functioning of Flash Memory

- Memory cells are connected to **blocks** and (depending on the structure also in) **pages**
  - A block always contains a fixed number of pages
- Write and erase operations can only be carried out for entire pages or blocks
  - For this reason, write and erase operations are more complex than read operations
- If data in a page need to be modified, the entire block must be erased
  - To do this, the block is copied into a buffer memory
  - In the buffer memory, the data is modified
  - Next, the block is erased from the flash memory
  - Finally, the modified block is written into the flash memory

# Different Types of Flash Memory

- 2 types of flash memory exist:
  - **NOR memory**
  - **NAND memory**
- The circuit symbol indicates the way, the memory cells are connected
  - This influences the capacity and access time (latency)

# NOR Memory

- Each memory cell has its data line
  - Benefit:
    - Random access for read and write operations  
⇒ Better latency compared with NAND memory
  - Drawback:
    - More complex (⇒ expensive) construction
    - Higher power consumption than NAND memory
    - Typically small capacities ( $\leq$  32 MB)
- Does not contain pages
  - The memory cells are grouped together to blocks
    - Typical block sizes: 64, 128 or 256 kB
- No random access for erase operations
  - Erase operations can only be done for entire blocks

## Fields of application

Industrial environment (e.g. automotive), storing the firmware of a computer system



NOR flash memory (top image) on the iPhone 3G mainboard (bottom image)

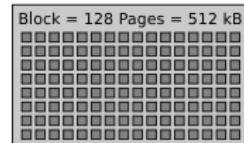


Images: Raimond Spekking.  
Wikimedia (CC-BY-SA-4.0)

# NAND Memory

- The memory cells are grouped to pages
  - Typical page size: 512-8192 bytes
    - Each page has its data line
  - Each block consists of a number of pages
    - Typical block sizes: 32, 64, 128 or 256 pages
- Benefit:
  - Lesser data lines  $\Rightarrow$  requires < 50% of the surface area of NOR memory
  - Lower manufacturing costs compared with NOR flash memory
- Drawback:
  - No random access  $\Rightarrow$  Poorer latency compared with NOR memory
  - Read and write operations can only be carried out for entire pages
  - Erase operations can only be carried out for entire blocks

Page  
4 kB



## Fields of application

USB flash memory drives,  
SSDs, memory cards



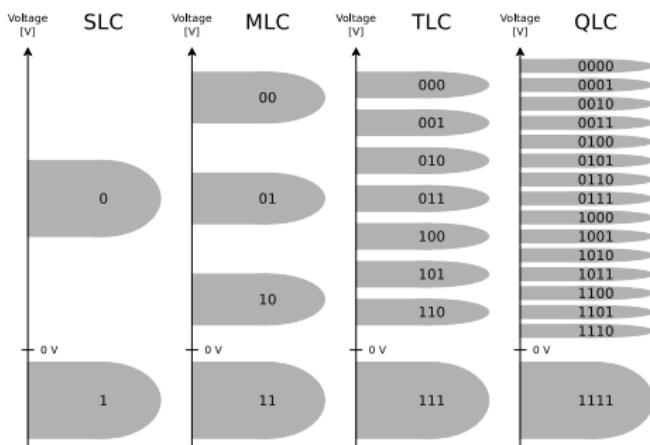
# Single/Multi/Triple/Quad-Level Cell

- 4 types of NAND flash memory exist

- QLC memory cells store 4 bits
- TLC memory cells store 3 bits
- MLC memory cells store 2 bits
- SLC memory cells store 1 bit

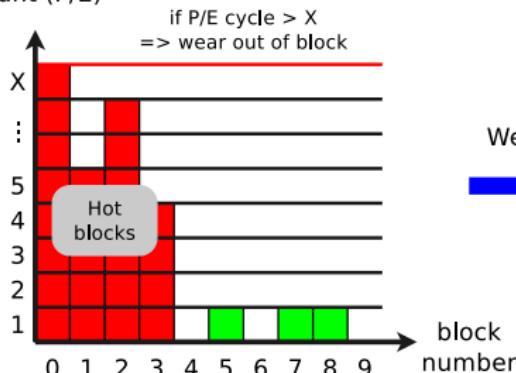
- SLC storage...

- is most expensive
- provides the best write speed
- survives most program/erase cycles

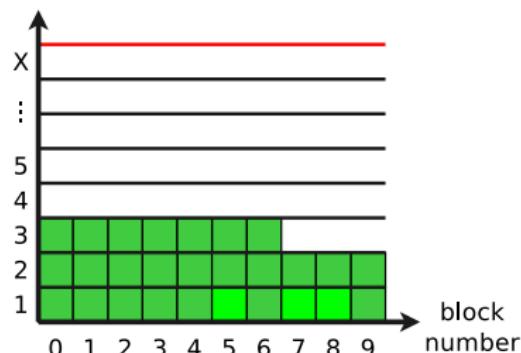


- SLC memory survives approx. 100,000 - 300,000 program/erase cycles
- MLC memory survives approx. 10,000 program/erase cycles
- TLC and QLC memory survives approx. 1,000 program/erase cycles
- Also memory cells exist, which survive millions of program/erase cycles

# Wear Leveling

Program/Erase  
count (P/E)Program/Erase  
count (P/E)

Wear Leveling



- Wear leveling algorithms evenly distribute write operations
- File systems, which are designed for flash memory, and therefore minimize write operations, are e.g. JFFS, JFFS2, YAFFS and LogFS
  - JFFS contains its own wear leveling algorithm
    - This is often required in embedded systems, where flash memory is directly connected

# Latency of Hard Disk Drives

- The performance of CPUs, cache and main memory is growing faster than the data access time (*latency*) of HDDs:

- HDDs**

1973: IBM 3340, 30 MB capacity, 30 ms data access time  
1989: Maxtor LXT100S, 96 MB capacity, 29 ms data access time  
1998: IBM DHEA-36481, 6 GB capacity, 16 ms data access time  
2006: Maxtor STM320820A, 320 GB capacity, 14 ms data access time  
2011: Western Digital WD30EZRSCTL, 3 TB capacity, 8 ms data access time  
2018: Seagate BarraCuda Pro ST14000DM001, 14 TB capacity, 4-5 ms data access time

- CPUs**

1971: Intel 4004, 740 kHz clock speed  
1989: Intel 486DX, 25 Mhz clock speed  
1997: AMD K6-2, 550 Mhz clock speed  
2007: AMD Opteron Santa Rosa F3, 2.8 GHz clock speed  
2010: Core i7 980X Extreme (6 Cores), 3.33 Ghz clock speed  
2018: Ryzen Threadripper 2990WX (32 Cores), 3 Ghz clock speed

- The latency of SSDs is  $\leq 1 \mu\text{s}$   $\Rightarrow \approx$  factor 100 better than HDDs
  - But the gap grows because of interface limitations and multiple CPU cores
- Further challenge
  - Storage drives can fail  $\Rightarrow$  risk of data loss
- Enhance **latency** and **reliability** of HDDs and SSDs  $\Rightarrow$  **RAID**

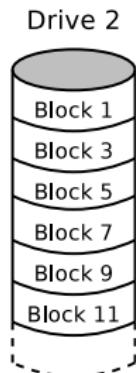
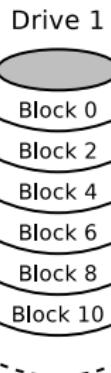
# Redundant Array of independent Disks (RAID)

- The performance of the HDDs can not be improved infinitely
  - HDDs contain moving parts
    - Physical boundaries must be accepted
- One way to avoid the given limitations in terms of speed, capacity and reliability, is the parallel use multiple components
- A RAID consists of multiple drives (HDDs or SSDs)
  - For users and their processes, a RAID behaves like a single large drive
- Data is distributed across the drives of a RAID system
  - The RAID level specifies how the data is distributed
    - The most commonly used RAID levels are RAID 0, RAID 1 and RAID 5

Patterson, David A., Garth Gibson, and Randy H. Katz, **A Case for Redundant Arrays of Inexpensive Disks (RAID)**, Vol. 17. No. 3, ACM (1988)

# RAID 0 – Striping – Acceleration without Redundancy

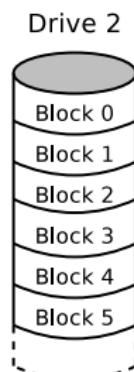
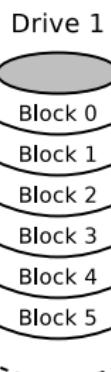
- No redundancy
  - Increases only the data rate
- Drives are partitioned into blocks of equal size
- If read/write operations are big enough ( $> 4$  or  $8$  kB), the operations can be carried out in parallel on multiple drives or on all drives



- In the event of a drive failure, not the entire data can be reconstructed
  - Only small files, which are stored entirely on the remaining drives, can be reconstructed (in theory)
- RAID 0 should only be used when security is irrelevant or backups are created at regular intervals

# RAID 1 – Mirroring

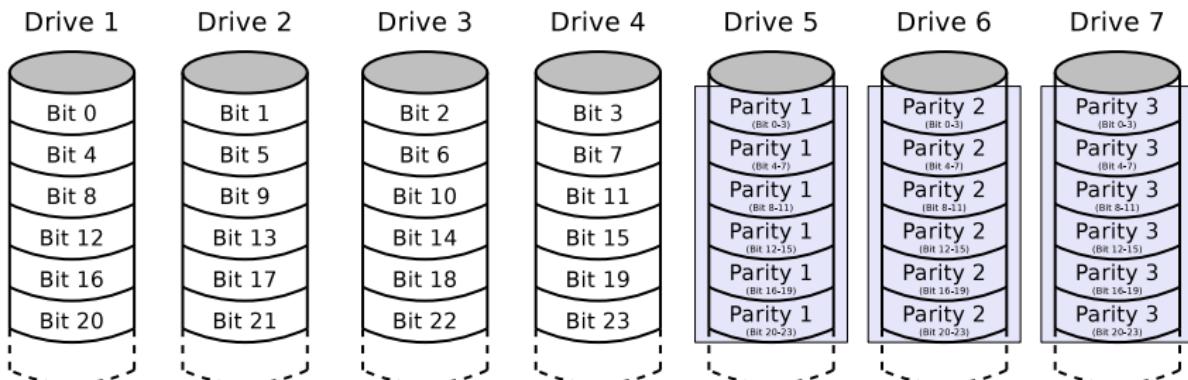
- At least 2 drives of the same capacity store identical data
  - If the drives are of different sizes, RAID 1 provides only the capacity of the smallest drive
- Failure of a single drive does not cause any data loss
  - Reason: The remaining drives store the identical data
- A total loss occurs only in case of the failure of all drives



- Any change of data is written on all drives
- Not a backup replacement
  - Corrupted file operations or virus attacks take place on all drives
- The read performance can be increased by intelligent distribution of requests to the attached drives

## RAID 2 – Bit-Level Striping with Hamming Code Error Correction

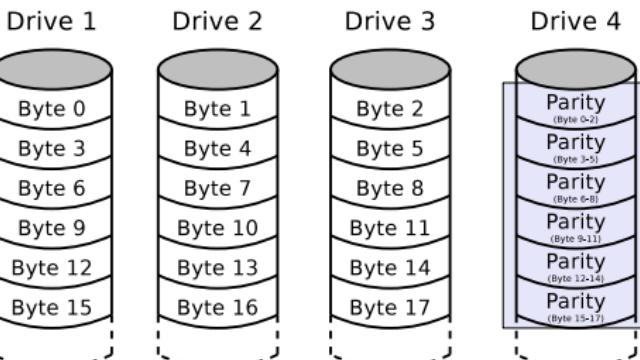
- Each sequential bit is stored on a different drive
  - Bits, which are powers of 2 (1, 2, 4, 8, 16, etc.) are parity bits



- The individual parity bits are distributed over multiple drives  
⇒ increases the throughput
- Was used only in mainframes
  - Is no longer relevant

# RAID 3 – Byte-level Striping with Parity Information

- Parity information is stored on a dedicated parity drive

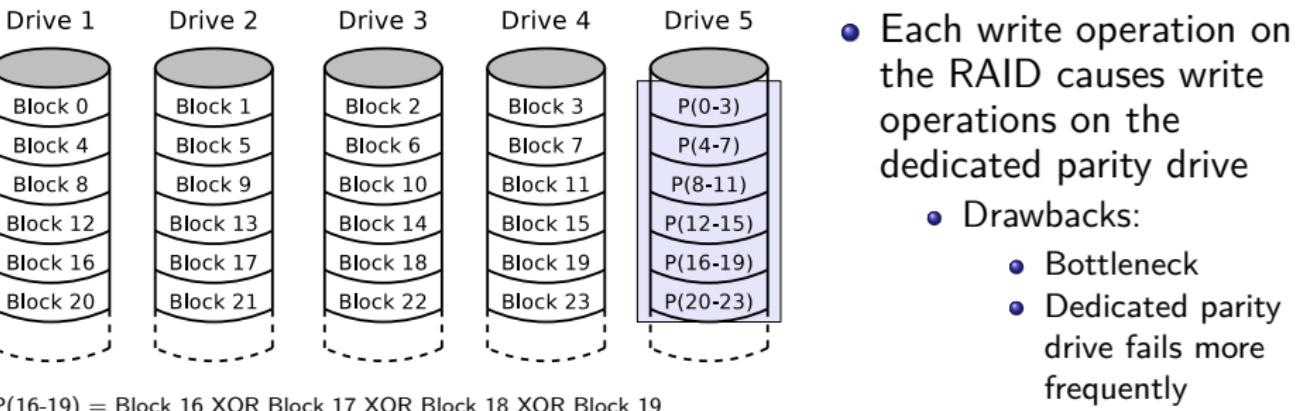


- Each write operation on the RAID causes write operations on the dedicated parity drive  
⇒ bottleneck
- Was replaced by RAID 5

Payload drives	Sum	even/odd	Parity drive
Bits are 0 + 0 + 0	⇒ 0	⇒ Sum is even	⇒ Sum bit 0
Bits are 1 + 0 + 0	⇒ 1	⇒ Sum is odd	⇒ Sum bit 1
Bits are 1 + 1 + 0	⇒ 2	⇒ Sum is even	⇒ Sum bit 0
Bits are 1 + 1 + 1	⇒ 3	⇒ Sum is odd	⇒ Sum bit 1
Bits are 1 + 0 + 1	⇒ 2	⇒ Sum is even	⇒ Sum bit 0
Bits are 0 + 1 + 1	⇒ 2	⇒ Sum is even	⇒ Sum bit 0
Bits are 0 + 1 + 0	⇒ 1	⇒ Sum is odd	⇒ Sum bit 1
Bits are 0 + 0 + 1	⇒ 1	⇒ Sum is odd	⇒ Sum bit 1

# RAID 4 – Block-level Striping with Parity Information

- Parity information is stored at a dedicated parity drive
- Difference to RAID 3:
  - Not individual bits or bytes, but blocks (**chunks**) are stored

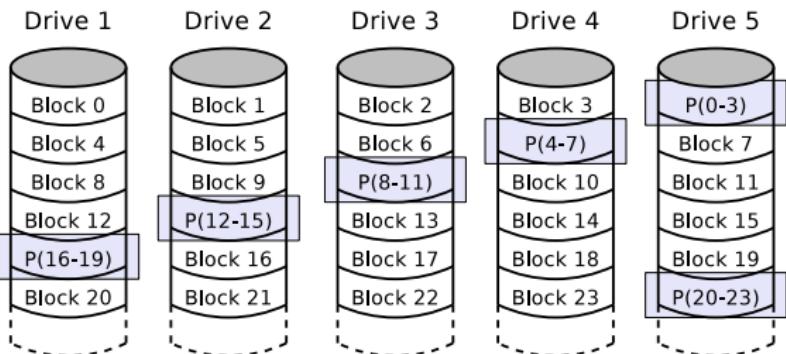


- Each write operation on the RAID causes write operations on the dedicated parity drive
- Drawbacks:
  - Bottleneck
  - Dedicated parity drive fails more frequently

- Seldom implemented, because RAID 5 does not face these drawbacks
- The company NetApp implements RAID 4 in their NAS servers
  - e.g. NetApp FAS2020, FAS2050, FAS3040, FAS3140, FAS6080

## RAID 5 – Block-level Striping with distributed Parity Information

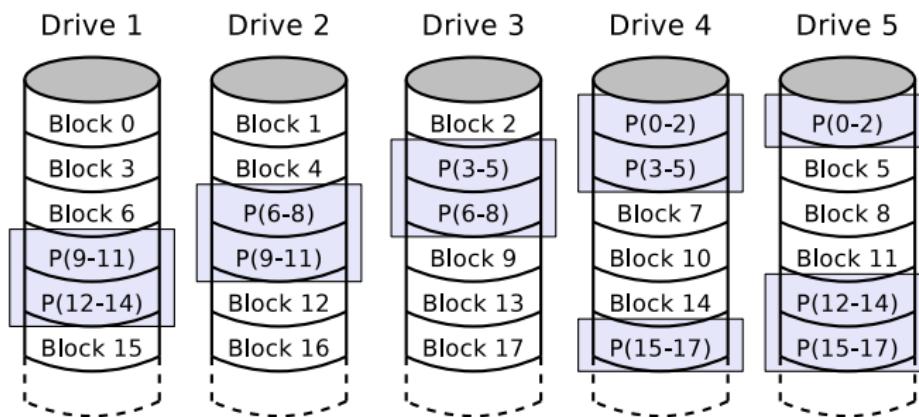
- Payload and parity information are distributed to all drives
- Benefits:
  - High throughput
  - High security level against data loss
  - No bottleneck



$P(16-19) = \text{block 16 XOR block 17 XOR block 18 XOR block 19}$

## RAID 6 – Block-level Striping with double distributed Parity Information

- Functioning is similar to RAID 5
  - But it can handle the simultaneous failure of up to 2 drives
- In contrast to RAID 5...
  - is the availability better, but the write performance is lower
  - is the effort to write the parity information higher



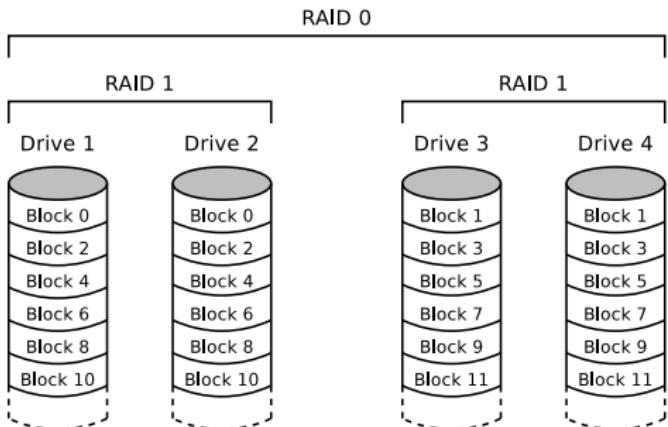
# Summary of the RAID Levels

RAID	$n$ (number of drives)	$k$ (net capacity)	Allowed to fail	Performance (read)	Performance (write)
0	$\geq 2$	$n$	0 (none)	$n * X$	$n * X$
1	$\geq 2$	1	$n - 1$ drives	$n * X$	$X$
2	$\geq 3$	$n - [\log_2 n]$	1 drive	variable	variable
3	$\geq 3$	$n - 1$	1 drive	$(n - 1) * X$	$(n - 1) * X$
4	$\geq 3$	$n - 1$	1 drive	$(n - 1) * X$	$(n - 1) * X$
5	$\geq 3$	$n - 1$	1 drive	$(n - 1) * X$	$(n - 1) * X$
6	$\geq 4$	$n - 2$	2 drives	$(n - 2) * X$	$(n - 2) * X$

- $X$  is the performance of a single drive during read or write
- The maximum possible performance in theory is often limited by the controller and the computing power of the CPU

If the drives of a RAID 1 have different capacities, the net capacity of a RAID 1 is equal to the capacity of its smallest drive

# RAID Combinations



- Usually RAID 0, 1 or 5 is used
- In addition to the popular RAID levels, several RAID combinations exist
  - At least 2 RAIDs are combined to a bigger RAID

## Examples

- RAID 00: Multiple RAID 0 are connected to a RAID 0
- RAID 01: Multiple RAID 0 are connected to a RAID 1
- RAID 05: Multiple RAID 0 are connected to a RAID 5
- RAID 10: Multiple RAID 1 are connected to a RAID 0 (see figure)
- RAID 15: Multiple RAID 1 are connected to a RAID 5
- RAID 50: Multiple RAID 5 are connected to a RAID 0
- RAID 51: Multiple RAID 5 are connected to a RAID 1

# Hardware / Host / Software RAID (1/2)

Image Source: Adaptec



Adaptec SATA RAID 2410SA

## • Hardware RAID

- A RAID controller with a processor does the calculation of the parity information and monitors the state of the RAID

Benefit: Operating system independent  
No additional CPU load

Drawback: High price (approx. € 200)



Adaptec SATA II RAID 1220SA

## • Host RAID

- Either an inexpensive RAID controller or the chipset provide the RAID functionality
- Usually only supports RAID 0 and RAID 1

Benefit: Operating system independent  
Low price (approx. € 50)

Drawback: Additional CPU load  
Possible dependence of rare hardware

## Hardware / Host / Software RAID (2/2)

- **Software RAID**

- Linux, Windows and MacOS allow to connect drives to a RAID without a RAID controller

Benefit: No cost for additional hardware

Drawback: Operating system dependent  
Additional CPU load

- Example: Create a RAID 1 (md0) with the partitions sda1 and sdb1:

```
mdadm --create /dev/md0 --auto md --level=1  
--raid-devices=2 /dev/sda1 /dev/sdb1
```

- Obtain information about any software RAID in the system:

```
cat /proc/mdstat
```

- Obtain information about a specific software RAID (md0):

```
mdadm --detail /dev/md0
```

- Remove partition sdb1 and add partition sdc1 to the RAID:

```
mdadm /dev/md0 --remove /dev/sdb1  
mdadm /dev/md0 --add /dev/sdc1
```