MS-E2122 - Nonlinear Optimization Lecture 5

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Outline of this lecture

Line search methods - univariate optimisation

Line searches without derivatives

Line searches with derivatives

Methods for unconstrained optimisation

Coordinate descent

Gradient method

Newton's method

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Most optimisation methods can be represented by this pseudocode:

Algorithm Conceptual optimisation algorithm

- 1: **initialise.** iteration count k = 0, starting point x_0
- 2: while stopping criteria are not met do
- 3: compute direction d_k
- 4: compute step size λ_k
- 5: $x_{k+1} = x_k + \lambda_k d_k$
- 6: k = k + 1
- 7: end while
- 8: return x_k

where

- ▶ *k* is an iteration counter;
- $\triangleright \lambda_k$ is a suitable step size;
- $ightharpoonup d_k$ is a direction vector;

Finding an optimal step size λ_k is in itself an optimisation problem called line search due to its unidimensional nature.

Line searches are the backbone of most optimisation methods.

Let $\theta(\lambda) = f(x + \lambda d)$. If f is differentiable, a straightforward approach is to find an optimal setup size λ is

$$\theta'(\lambda) = d^{\top} \nabla f(x + \lambda d) = 0$$

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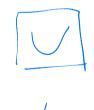
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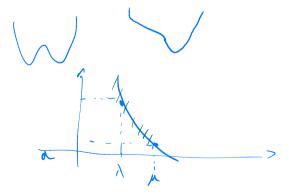
However, one must bear in mind that:

- $ightharpoonup d^{\top}f(x+\lambda d)$ is often nonlinear in λ ;
- $\overline{\lambda} = \operatorname{argmin}_{\lambda} d^{\top} \nabla f(x + \lambda d) = 0$ is not necessarily optimal.



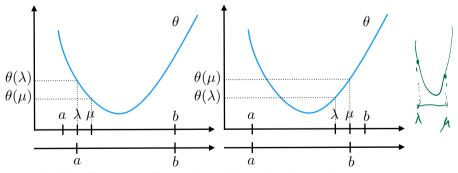
Theorem 1 (Line search reduction) _ vnimodal

Let $\theta: \mathbb{R} \to \mathbb{R}$ be strictly quasiconvex over the interval [a,b], and let $\lambda, \mu \in [a, b]$ such that $\lambda < \mu$. If $\theta(\lambda) > \theta(\mu)$, then $\theta(z) > \theta(\mu)$ for all $z \in [a, \lambda]$. If $\theta(\lambda) \leq \theta(\mu)$, then $\theta(z) \geq \theta(\lambda)$ for all $z \in [\mu, b]$.



Theorem 1 (Line search reduction)

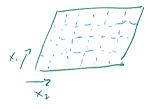
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Applying Theorem 1 allows to iteratively reduce the search space.

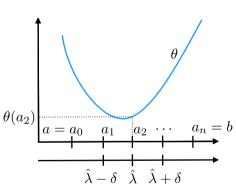
Line search methods - uniform search

Break [a,b] into n uniform intervals of size δ , which leads to n+1 grid points $a_k=a_0+k\delta$, with $a=a_0,b=a_n$, and $k=0,\ldots,n$.



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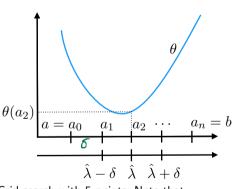


Let $\hat{\lambda} = \operatorname{argmin}_{i=0,\dots,n} \theta(a_i)$. We know that the optimal $\overline{\lambda} \in [\hat{\lambda} - \delta, \hat{\lambda} + \delta]$.

Grid search with 5 points; Note that $\theta(a_2) = \min_{i=0,...,n} \theta(a_i)$.

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Remarks:

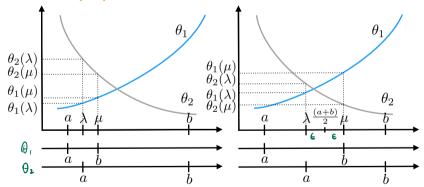
- The search can be repeated making $a = \hat{\lambda} \delta$ and $b = \hat{\lambda} + \delta$.
- The number of grid points can increase dynamically, saving function evaluations.

More efficient methods can be devised by using information of the previous evaluation of θ . These are known as sequential searches.

1. Dichotomous search: we place two points, λ and μ , around the midpoint of [a,b] at a small distance ϵ .

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1. Dichotomous search: we place two points, λ and μ , around the midpoint of [a,b] at a small distance ϵ .



Using the midpoint (a+b)/2 and Theorem 1 to reduce the search space.

Algorithm Dichotomous search

```
1: initialise. distance \epsilon > 0, tolerance l > 0, [a_0, b_0] = [a, b], k = 0
 2: while b_k - a_k > l do
        \lambda_k = \frac{a_k + b_k}{2} - \epsilon, \ \mu_k = \frac{a_k + b_k}{2} + \epsilon
 3:
                                                                  AMZ OUM
        if \theta(\lambda_k) < \theta(\mu_k) then
 5:
            a_{k+1} = a_k, b_{k+1} = \mu_k
 6:
        else
          a_{k+1} = \lambda_k, b_{k+1} = b_k
        end if
 8:
          k = k + 1
10: end while
11: return \overline{\lambda} = \frac{a_k + b_k}{2}
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4: if \theta(\lambda_k) < \theta(\mu_k) then
5: a_{k+1} = a_k, b_{k+1} = \mu_k
6: else
7: a_{k+1} = \lambda_k, b_{k+1} = b_k
8: end if
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Remark: The number of steps (and evaluations of θ) can be predicted beforehand:

$$\int b_{k+1} - a_{k+1} = \frac{1}{2^k} (b_0 - a_0) + 2\epsilon \left(1 - \frac{1}{2^k} \right).$$

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 - 3. if $\theta'(\lambda_k) < 0$, we have $\theta'(\lambda_k)(\lambda \lambda_k) \ge 0$ for $\lambda < \lambda_k$. Thus, the new search interval becomes $[a_{k+1}, b_{k+1}] = [\lambda_k, b_k]$.

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 - 3. if $\theta'(\lambda_k) < 0$, we have $\theta'(\lambda_k)(\lambda \lambda_k) \ge 0$ for $\lambda < \lambda_k$. Thus, the new search interval becomes $[a_{k+1}, b_{k+1}] = [\lambda_k, b_k]$.
 - 4. As in the dichotomous search, to maximise overall interval reduction, we set $\lambda_k = \frac{1}{2}(b_k + a_k)$.

Algorithm Bisection method

```
1: initialise. tolerance l > 0, [a_0, b_0] = [a, b], k = 0
 2: while b_k - a_k > l do
        \lambda_k = \frac{(b_k + a_k)}{2} and evaluate \theta'(\lambda_k)
        if \theta'(\lambda_k) = 0 then return \lambda_k
 5:
        else if \theta'(\lambda_k) > 0 then
            a_{k+1} = a_k, b_{k+1} = \lambda_k
       else
            a_{k+1} = \lambda_k, b_{k+1} = a_k
        end if
      k = k + 1.
10:
11: end while
12: return \overline{\lambda} = \frac{a_k + b_k}{2}
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Inexact line searches - Armijo rule

Often the use of nonoptimal (i.e., inexact) step sizes λ_k is enough to guarantee a good performance.

Armijo's rule: find acceptable step sizes by balancing the trade-off between convergence and numerical performance.

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$$f(\overline{x} + d\overline{\lambda}) - f(\overline{x}) \leq \alpha \overline{\lambda} \nabla f(\overline{x})^{\top} d$$
 which, at $\lambda = 0$, is the same as
$$\theta(\overline{\lambda}) - \theta(0) \leq \alpha \overline{\lambda} \theta'(0)$$

$$\theta(\overline{\lambda}) \leq \theta(0) + \alpha \overline{\lambda} \theta'(0) : \text{Armijo's rule (AR)}$$

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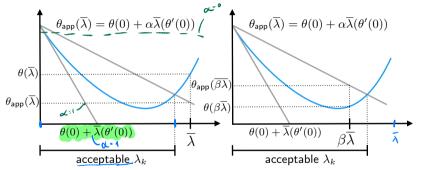
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 $\theta(\overline{\lambda}) \le \theta(0) + \alpha \overline{\lambda} \theta'(0)$: Armijo's rule (AR)

If $\overline{\lambda}$ does not satisfy AR, $\overline{\lambda}$ is reduced by a factor $\beta \in (0,1)$ and the test is repeated until (AR) is satisfied.

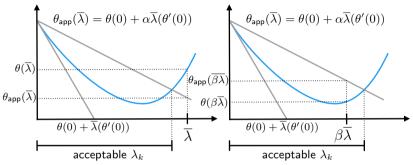
Inexact line searches - Armijo's rule

Armijo's rule has a nice graphical interpretation: $\overline{\lambda}$ is accepted if it is in an interval where the function $\theta(\lambda)$ is below a deflected linear extrapolation (from 0).



At first $\lambda_0 = \overline{\lambda}$ is not acceptable; after reducing the step size to $\lambda_1 = \beta \overline{\lambda}$, it enters the acceptable range where $\theta(\lambda_k) \leq \theta_{\rm app}(\lambda_k) = \theta(0) + \alpha \lambda_k(\theta'(0))$.

Inexact line searches - Armijo's rule



Remarks:

- 1. Some variants also consider rules to guarantee that $\overline{\lambda}$ is not too small, such as $\theta(\delta\overline{\lambda}) \leq \theta(0) + \alpha \delta \overline{\lambda} \theta'(0)$, with $\delta > 1$.
- Wolfe Note.

- 2. Also known in the literature as backtracking.
- 3. Typical values: $\alpha \in [0.1, 0.5]$ and $\beta \in [0.6, 0.99]$. Very small α , e.g. 10^{-4} is often used as well.

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Next, we focus on optimising functions $f: \mathbb{R}^n \to \mathbb{R}$ with more than one dimension.

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1. Coordinate descent: the search direction is one coordinate axis per iteration. That is, $d_i = 1$ for coordinate i and $d_{j \neq i} = 0$, for $i, j \in \{1, \ldots, n\}$.

$$\mathbb{R}^2$$
: $d_1 = [1,0]$ \longrightarrow $d_2 = [0,1]$

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Several variants:

- 1. **Cyclic:** coordinates are considered in order $1, \ldots, n$;
- Double-sweep: swap the coordinate order at each iteration;
- 3. **Gauss-Southwell:** choose components with largest $\frac{\partial f(x)}{\partial x_i}$;
- 4. Stochastic: coordinates are selected at random

Algorithm Coordinate descent method (cyclic)

```
1: initialise. tolerance \epsilon>0, initial point x^0, iteration count k=0

2: while ||x^{k+1}-x^k||>\epsilon do

3: for j=1,\dots n do

4: d=\{d_i=1, \text{ if } i=j; d_i=0, \text{ if } i\neq j\}

5: \overline{\lambda_j}= \underset{\lambda\in\mathbb{R}}{\operatorname{argmin}}_{\lambda\in\mathbb{R}}\{f(x_j^k+\lambda d_j)\}

6: x_j^{k+1}=x_j^k+\overline{\lambda_j}d_j

7: end for

8: k=k+1

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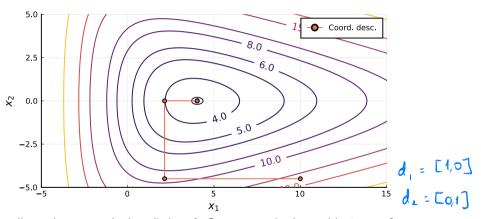
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3: for j = 1, \dots n do
4: d = \{d_i = 1, \text{ if } i = j; d_i = 0, \text{ if } i \neq j\}
5: \overline{\lambda}_j = \operatorname{argmin}_{\lambda \in \mathbb{R}} \{f(x_j^k + \lambda d_j)\}
6: x_j^{k+1} = x_j^k + \overline{\lambda}_j d_j
7: end for
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```

Remarks:

- 1. The one-dimensional minimisation is called Gauss-Seidel step;
- Block-coordinate methods use subgroups (blocks) of coordinates to define directions.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



Coordinate descent method applied to f. Convergence is observed in 4 steps for a tolerance $\epsilon=10^{-4}$

Recall that if d is a descent direction, there exists $\delta>0$ such that $f(x+\lambda d)< f(x)$ for all $\lambda\in(0,\delta)$. The following result provides directions of steepest descent.

Lemma 2 (Steepest descent direction)

Suppose that $f: \mathbb{R}^n \to \mathbb{R}$ is differentiable at $x \in \mathbb{R}^n$ and $\nabla f(x) \neq 0$.

Then $\overline{d} = -\frac{\nabla f(x)}{||\nabla f(x)||}$ is the direction of steepest descent of f at x.



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Proof.

From differentiability of f, we have

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$$f$$
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$$f'(x;d) = \lim_{\lambda \to 0^+} \frac{f(x+\lambda d) - f(x)}{\lambda} = \nabla f(x)^\top d.$$

$$f = \operatorname{argmin}_{||d|| \le 1} \left\{ \nabla f(x)^\top d \right\} = -\frac{\nabla f(x)}{||\nabla f(x)||}.$$

Thus,
$$\overline{d} = \operatorname{argmin}_{||d|| \le 1} \left\{ \nabla f(x)^{\top} d \right\} = -\frac{\nabla f(x)}{||\nabla f(x)||}.$$

Algorithm Gradient method

```
1: initialise. tolerance \epsilon > 0, initial point x_0, iteration count k = 0.
```

```
2: while ||\nabla f(x_k)|| > \epsilon do
3: d = -\frac{|\nabla f(x_k)|}{||\nabla f(\overline{x})||}
4: \overline{\lambda} = \operatorname{argmin}_{\lambda \in \mathbb{R}} \{ f(x_k + \lambda d) \}
5: x_{k+1} = x_k + \overline{\lambda} d_j
6: k = k+1
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- 7: end while
- 8: return x_k

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5: x_{k+1} = x_k + \overline{\lambda} d_i
       k = k + 1
```

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7: end while

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Remarks:

1. Steepest descent and gradient methods are different. When ||d|| < 1 uses 2-norm (in Lemma 2), they are equivalent;

Gradient method

Algorithm Gradient method

- 1: **initialise.** tolerance $\epsilon > 0$, initial point x_0 , iteration count k = 0.
- 2: while $||\nabla f(x_k)|| > \epsilon$ do

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$$d = -\frac{\nabla f(x_k)}{||\nabla f(\overline{x})||}$$

4:
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- 5: $x_{k+1} = x_k + \overline{\lambda} d_j$
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- 1. Steepest descent and gradient methods are different. When $||d|| \le 1$ uses 2-norm (in Lemma 2), they are equivalent;
- 2. Poor convergence and zigzagging can be observed due to imprecise linear approximations (more on this later);

Gradient method

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$

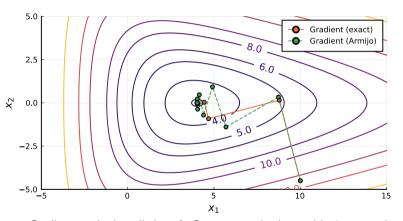
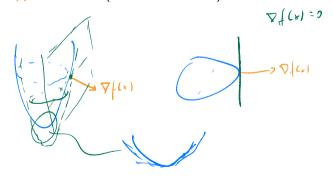


Figure: Gradient method applied to f. Convergence is observed in 9 steps using exact line search and 15 using Armijo's rule ($\epsilon=10^{-4}$)

Same idea as in the univariate case. Can also be seen as deflected steepest descent.

Deflection is achieved using the Hessian, which is equivalent to relying on quadratic approximations (rather than linear).



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Consider the 2nd-order approximation of
$$f$$
 at x_k :
$$q(x) = f(x_k) + \nabla f(x_k)^\top (x - x_k) + \frac{1}{2} (x - x_k)^\top \underline{H(x_k)}(x - x_k),$$

where $H(x_k)$ is the Hessian at x_k .

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Consider the 2nd-order approximation of f at x_k :

$$q(x) = f(x_k) + \nabla f(x_k)^{\top} (x - x_k) + \frac{1}{2} (x - x_k)^{\top} H(x_k) (x - x_k),$$

where $H(x_k)$ is the Hessian at x_k . We require that $\nabla q(x_{k+1}) = 0$, which leads to

$$\nabla f(x_k) + H(x_k)(x - x_k) = 0.$$

Assuming that $H^{-1}(x_k)$ exists, we obtain the update rule

$$x_{k+1} = x_k - H^{-1}(x_k)\nabla f(x_k).$$

Algorithm Newton's method

```
1: initialise. tolerance \epsilon>0, initial point x_0, iteration count k=0
2: while ||\nabla f(x_k)|| > \epsilon do
3: d=-H^{-1}(x_k)\nabla f(x_k)
4: \overline{\lambda}= \underset{k=k+1}{\operatorname{argmin}}_{\lambda\in\mathbb{R}}\{f(x_k+\lambda d)\}
5: x_{k+1}=x_k+\overline{\lambda}d
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5: x_{k+1} = x_k + \lambda d
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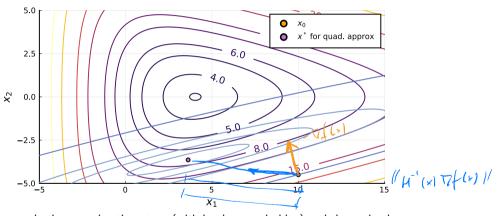
- 1. Setting $\overline{\lambda} = 1$ recovers the "pure" Newton's method;
- 2. As $\nabla f(x_k)$ gets close to 0, $H^{-1}(x_k)$ becomes singular;
- 3. It might not converge if x_0 is too far from optimal and fixed step size is used;

Algorithm Newton's method

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1: initialise. tolerance \epsilon>0, initial point x_0, iteration count k=0
2: while ||\nabla f(x_k)|| > \epsilon do
3: d=-H^{-1}(x_k)\nabla f(x_k)
4: \overline{\lambda}= \operatorname*{argmin}_{\lambda\in\mathbb{R}}\{f(x_k+\lambda d)\}
5: x_{k+1}=x_k+\overline{\lambda}d
6: k=k+1
7: end while
8: return x_k
```

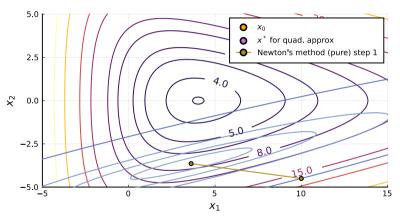
- 1. Setting $\overline{\lambda} = 1$ recovers the "pure" Newton's method;
- 2. As $\nabla f(x_k)$ gets close to 0, $H^{-1}(x_k)$ becomes singular;
- 3. It might not converge if x_0 is too far from optimal and fixed step size is used;
- 4. Levenberg-Marquardt method and other trust-region method variants also address convergence issues of Newton's method.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



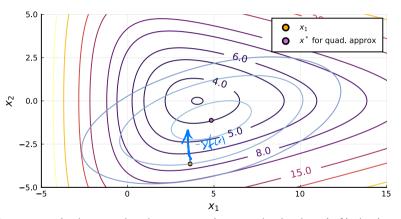
The quadratic approximation at x_0 (with level curves in blue) and the optimal point x^{\ast} .

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



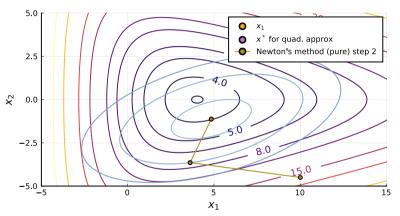
The new point x_1 becomes x^* .

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



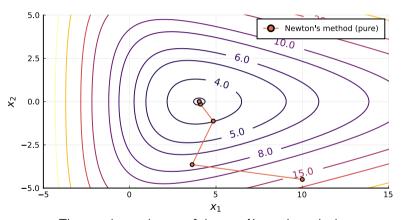
The new quadratic approximation at x_1 and new optimal point x^* . Notice how the approximation improved.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



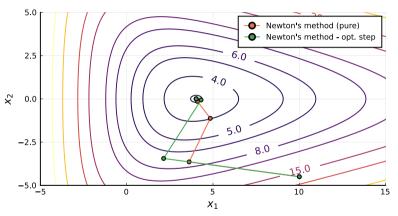
The new point x_2 becomes x^* .

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



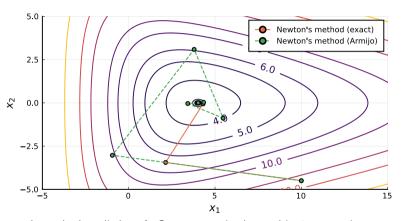
The complete trajectory of the pure Newton's method.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



When employing line searches, the direction $x_k - x_{k-1}$ from the pure method is used, but the actual step is optimised.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



Newton's method applied to f. Convergence is observed in 4 steps using exact line search and 27 using Armijo's rule ($\epsilon=10^{-4}$)