#### Searching:

- Searching is the most common operation performed by a database system
- In SQL, the SELECT statement is arguably the most versatile / complex (they can be recursive and there can even be select statements in select statements)

#### General Vocab:

- Record: a collection of values for attributes of a single entity instance; a row of a table
- <u>Collection:</u> a set of records of the same entity type; a table (trivially stored in some sequential order like a list)
- Search Key: a value for an attribute from the entity type (could be >= 1 attribute)

id	specVal
1	55
2	87
3	50
4	108

- Assume data is stored on disk by column id's value, searching for a specific id is fast
- But searching for specific specialVal since data is unsorted, the only option is linear scan the column
- Can't store data on disk sorted by both id and specialVal at the same time so the data would have to be duplicated and there'd be inefficient space
  - Therefore we need an external data structure to support faster searching by specialVal than a linear search

## What to do?

- An array of tuples (specialVal, rowNumber) sorted by specialVal
  - We could use Binary Search to quickly locate a particular specialVal and find its corresponding row in the table
  - But, every insert into the table would be like inserting into a sorted array slow...
- OR A linked list of tuples (specialVal, rowNumber) sorted by specialVal
  - searching for a specialVal would be slow linear scan required
  - But inserting into the table would theoretically be quick to also add to the list... INSTEAD USE BINARY SEARCH TREE

# **Linear Search:**

- Baseline for efficiency where you start at the beginning of a list and proceed element element by element until you either find what you're looking for or get to the last element and haven't found it aka O(n)
- If each record takes up x bytes of memory, then for n records, we need n\*x bytes of memory
- There are 2 different Data Structures for Linear Search:
  - Contiguously Allocated List (aka Array): all n\*x are allocated as a single "chunk" of memory

- Pro: Faster for random access
- Con: Slow for inserting anywhere but the end
- <u>Linked List:</u> each record needs x bytes and additional space for 1 or 2 memory addresses - individual records are linked together in a type of chain using memory addresses
  - Pro: Faster for inserting anywhere in the list
  - Con: Slower for random access
- Best Case: target is found at the first element where only 1 comparison is needed
- Worst Case: target is not in the array; n comparisons O(n) time complexity is for worst case

## **Binary Search:**

- <u>Input:</u> array of values in sorted order, target value
- <u>Output:</u> location (index) of where target is located or some value indicating target was not found
- Best Case: target is found at mid; 1 comparison (inside the loop)
- <u>Worst Case:</u> target is not in the array; log2n comparisons O(log2n) time complexity is for worse case

# **Binary Search Tree:**

- A binary tree where every node in the left subtree is less than its parent and every node in the right subtree is greater than its parent