

1 Introduction

This document contains detailed notes related to the slides.

2 Polynomials

This is the general form of a polynomial of degree k .

$$f(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + \dots + a_kx^k$$

Or, more generally...

$$f(x) = \sum_{k=0}^n a_kx^k$$

3 Curve Fitting

A polynomial of degree n may be uniquely defined given $n + 1$ points on the line. An infinite number of lines will intersect only n points. These lines may be found by solving the simultaneous equations for the points provided. These are the points used in the slides.

3.1 Linear Polynomial

A first degree polynomial is represented by the equation

$$f(x) = ax + b$$

Given our secret 42 and a random $a = 4$ we can represent a straight line

$$f(x) = 4x + 42 \tag{1}$$

Suppose we have a single point on this line, say $P = (-17, -26)$, we want to want to find an equation that describes all possible lines that include this point. Given point P , we solve for b :

$$\begin{aligned} -26 &= -17a + b \\ b &= -26 + 17a \end{aligned} \tag{2}$$

Now we may substitute b from equation 2 into equation 1.

3.2 Quadratic Polynomial

$$\begin{aligned}f_a(x) &= ax - 26 + 17a \\f_a(x) &= a(x + 17) - 26\end{aligned}\tag{3}$$

Equation 3 describes all lines passing through point $P = (-17, -26)$.

3.2 Quadratic Polynomial

A second degree, or quadratic, polynomial is represented as

$$f(x) = ax^2 + bx + c\tag{4}$$

Given a secret $S = 42$, we can choose random values for a and b . Suppose we choose $a = 7$ and $b = 3$. We may then write

$$f(x) = 7x^2 + 3x + 42\tag{5}$$

Three points must be given to describe a quadratic equation. Given only two points, we want to find all quadratic curves that pass through the two given points.

Suppose we have points $P_1 = (2, 76)$ and $P_2 = (5, 232)$ discovered by choosing two x values randomly and solving equation 5 for $f(x)$. We may find the equation describing all curves passing through P_1 and P_2 by solving the system of equations as we did in the case of the linear polynomial.

Substituting P_1 and P_2 into equation 4 we can write

$$\begin{aligned}f(2) &= 4a + 2b + c \\76 &= 4a + 2b + c \\c &= -4a - 2b + 76\end{aligned}\tag{6}$$

and

$$\begin{aligned}f(5) &= 25a + 5b + c \\232 &= 25a + 5b + c \\c &= -25a - 5b + 232\end{aligned}\tag{7}$$

Since we know equations 6 and 7 are equal, we can solve for b in terms of a as follows:

3.3 Cubic Polynomial

$$\begin{aligned} -4a - 2b + 76 &= -25a - 5b + 232 \\ 5b - 2b &= 232 - 76 + 4a - 25a \\ 3b &= -21a + 156 \\ b &= -7a + 52 \end{aligned} \tag{8}$$

We can now substitute equation 8 into equation 6 to write:

$$\begin{aligned} c &= -4a - 2(-7a + 52) + 76 \\ &= -4a + 14a - 104 + 76 \\ c &= 10a - 28 \end{aligned} \tag{9}$$

Now that we have solved for both b and c in terms of a , we may rewrite equation 4 as follows:

$$\begin{aligned} f_a(x) &= ax^2 + (-7a + 52)x + 10a - 28 \\ &= ax^2 - 7ax + 52x + 10a - 28 \\ f_a(x) &= a(x^2 - 7x + 10) + 52x - 28 \end{aligned} \tag{10}$$

Equation 10 is the set of quadratic curves that pass through points $P_1 = (2, 76)$ and $P_2 = (5, 232)$.

3.3 Cubic Polynomial

A third degree, or cubic, polynomial is represented as:

$$f(x) = ax^3 + bx^2 + cx + d \tag{11}$$

As before, our secret $S = 42$. We then select random values for a , b , and c . Suppose we have $a = 1$, $b = 3$, and $c = 13$. We may write:

$$f(x) = x^3 + 3x^2 + 13x + 42 \tag{12}$$

Four points are required to identify a specific cubic equation. Suppose we have three points $P_1 = (-6, -144)$, $P_2 = (1, 59)$, $P_3 = (4, 206)$. We can find the set of cubic polynomials that contains each of P_1 , P_2 , and P_3 by solving the system of equations as before. We can write:

$$\begin{aligned} f(-6) &= a(-6)^3 + b(-6)^2 + c(-6) + d \\ -144 &= -216a + 36b - 6c + d \\ d &= 216a - 36b + 6c - 144 \end{aligned} \tag{13}$$

3.3 Cubic Polynomial

$$\begin{aligned}f(1) &= a + b + c + d \\59 &= a + b + c + d \\d &= -a - b - c + 59\end{aligned}\tag{14}$$

$$\begin{aligned}f(4) &= 4^3a + 4^2b + 4c + d \\206 &= 64a + 16b + 4c + d \\d &= -64a - 16b - 4c + 206\end{aligned}\tag{15}$$

We may now represent c in terms of a and b by simplifying the equality of equations 13 and 14:

$$\begin{aligned}216a - 36b + 6c - 144 &= -a - b - c + 59 \\7c &= -217a + 35b + 203 \\c &= -31a + 5b + 29\end{aligned}\tag{16}$$

And now repeat for equations 14 and 15:

$$\begin{aligned}-a - b - c + 59 &= -64a - 16b - 4c + 206 \\3c &= -63a - 15b + 147 \\c &= -21a - 5b + 49\end{aligned}\tag{17}$$

Repeating once again to solve for b by simplifying the equality represented by the equations 16 and 17:

$$\begin{aligned}-31a + 5b + 29 &= -21a - 5b + 49 \\10b &= 10a + 20 \\b &= a + 2\end{aligned}\tag{18}$$

At this point, we may substitute equations 18, 16, and 14 into the general cubic polynomial formula 11. Terms with b and c must be replaced as needed in equations 16 and 14 so that we have an equation in terms of a and x only. The algebra is simple but long so here is the result without further ceremony:

$$f_a(x) = a(x^3 + x^2 - 26x + 24) + 2x^2 + 39x + 18\tag{19}$$

4 Shamir Secret Sharing

4.1 Mathematical Definition

Divide secret S into n parts S_1, \dots, S_n such that

- Knowledge of any k or more S_i makes secret S computable.
- Knowledge of any $k - 1$ or less S_i leaves secret S completely undetermined.

This is called a (k, n) threshold scheme.

The essential idea is that k points are required to define a polynomial of degree $k - 1$.

For our purposes we can define a polynomial as follows:

$$f(x) = a_0 + a_1x + a_2x^2 + \dots + a_{k-1}x^{k-1}$$

$$f(x) = \sum_{i=0}^{k-1} a_i x^i$$

Suppose we want to use a (k, n) threshold scheme to share our secret S , an element in a finite field \mathbb{F} of size P where $0 < k \leq n < P$; $S < P$ and P is a prime number.

- Choose at random $k - 1$ positive integers a_1, \dots, a_{k-1} with $0 < a_i < P$; $a_i \in \mathbb{N}$ and let $a_0 = S$; $S < P$.
- Build the polynomial $f(x) = a_0 + a_1x + a_2x^2 + \dots + a_{k-1}x^{k-1} \pmod{P}$
- Construct any n points, *for instance*, set $i = 1, \dots, n$ to retrieve $(i, f(i))$.

Every participant is given a point, a point, e.g. in integer input to the polynomial and the corresponding integer output.

Given any subset k of these pairs, we can find the coefficients of the polynomial. The secret is the constant term a_0 .

4.2 Examples

4.2.1 The Problem

In this example we will omit the requirement that \pmod{P} is applied to the polynomial.

Suppose Eve has managed to obtain a share, say $P = (16, 106)$ which is a point on the curve defined by 1. Eve knows $f(x) = ax + b$.

4.2 Examples

She can solve for b in terms of a .

$$\begin{aligned} f(x) &= ax + b \\ P \rightarrow f(16) &= 106 = a(16) + b \\ b &= -16a + 106 \end{aligned} \tag{20}$$

Recall the requirement that $a, b \in \mathbb{N}$.

Eve may then substitute values for unknown a in equation 20.

$$\begin{aligned} a = 0 &\rightarrow b = -16(0) + 106 = 106 \\ a = 1 &\rightarrow b = -16(1) + 106 = 90 \\ a = 2 &\rightarrow b = -16(2) + 106 = 74 \\ a = 3 &\rightarrow b = -16(3) + 106 = 58 \\ a = 4 &\rightarrow b = -16(4) + 106 = 42 \\ a = 5 &\rightarrow b = -16(5) + 106 = 26 \\ a = 6 &\rightarrow b = -16(6) + 106 = 10 \\ a = 7 &\rightarrow b = -16(7) + 106 = -6 \end{aligned}$$

Since the requirement is that $a, b \in \mathbb{N}$, a cannot be negative. Therefore, Eve can conclude

$$\begin{aligned} a &\in [0, 1, 2, 3, 4, 5, 6] \\ b &\in [106, 90, 74, 58, 42, 26, 10] \end{aligned}$$

4.2.2 The Solution

The solution is to require that S is an element in a finite field \mathbb{F} of size P where $S < P$ and P is prime.

Since $S = 42$, choosing $P = 43$ satisfies the requirement that $S < P$ and P is prime.

$$\begin{aligned} f(x) &= 4x + 42 \pmod{43} \\ f(16) &= 4(16) + 42 \pmod{43} \\ &= 106 \pmod{43} \\ f(16) &= 20 \end{aligned}$$

Recall from the definition of a modulus that

$$a \pmod{P} = a - Pm \mid 0 \leq a - pm \leq P$$

4.2 Examples

In other words, m is a multiplier.

Eve knows $P = (16, 20)$, $f(x) = ax + b \pmod{P}$, and modulus $P = 43$.

So she can substitute and write

$$\begin{aligned} f(x) &= ax + b \pmod{P} \\ &= ax + b - pm \\ 20 &= 16a + b - 43m \\ b &= -16a + 20 + 43m \end{aligned} \tag{21}$$

As before, Eve can then substitute values for a into equation 21.

$$\begin{aligned} a = 0 &\rightarrow b = -16(0) + 20 + 43m = 20 + 43m \\ a = 1 &\rightarrow b = -16(1) + 20 + 43m = 4 + 43m \\ a = 2 &\rightarrow b = -16(2) + 20 + 43m = -12 + 43m \\ a = 3 &\rightarrow b = -16(3) + 20 + 43m = -28 + 43m \\ a = 4 &\rightarrow b = -16(4) + 20 + 43m = -44 + 43m \\ a = 5 &\rightarrow b = -16(5) + 20 + 43m = -60 + 43m \\ a = 6 &\rightarrow b = -16(6) + 20 + 43m = -76 + 43m \\ a = 7 &\rightarrow b = -16(7) + 20 + 43m = -92 + 43m \end{aligned} \tag{22}$$

This time since m is unknown, Eve can gain no information about the secret. But since we know $a = 4$ is the correct value, we can see that equation 22 is the correct result. In other words, $m = 2$.

The secret S is equally likely to be any element of the finite field \mathbb{F} . No information may be learned about S unless $S_1 \dots S_k$ are known.