## Utils.v

```
Require Import Coq.Classes.EquivDec.
   Require Import Coq.Relations.Relation Operators.
   Require Import Coq.Lists.List.
   Require Import Coq.omega.Omega.
   Require Import Datatypes.
   Require Import Omega.
   Import ListNotations.
   Import EqNotations.
   Lemma in_tuple_snd {A} {eq_dec : EqDec A eq} : forall (l : list (A*A)) a, {b &
11
       In (a, b) l} + {forall b, ~In (a,b) l}.
   Proof.
12
     intros.
13
     pose proof prod_eqdec eq_dec eq_dec as tuple_eq_dec.
14
     induction 1.
15
     - right. auto.
16

    destruct IHl.

17
       + destruct s. left. intros. exists x. constructor 2. assumption.
18
       + destruct a0. destruct (a = a0).
19
          * rewrite e. left. exists a1. constructor. auto.
          * right. intros. intros F. inversion F.
21
            ** apply c. injection H. intros. subst. reflexivity.
22
            ** pose proof n b. contradiction.
23
   Qed.
24
25
   Lemma ge_0eq: forall m, \emptyset \geqslant m \rightarrow \emptyset = m. Proof. intros. omega. Qed.
27
28
   Lemma Odd plus Even is Odd : forall n m, Nat.Odd n \rightarrow Nat.Even m \rightarrow Nat.Odd (n +
29
    \rightarrow m).
   Proof.
     intros n m [n' Hodd] [m' Heven].
31
     unfold Nat.Odd.
32
     exists (n' + m').
33
     omega.
34
   Qed.
35
36
   Lemma count_occ_split {A} : forall eq_dec (x : A) l1 l2, count_occ eq_dec (l1 ++
38
    → l2) x = count_occ eq_dec l1 x + count_occ eq_dec l2 x.
   Proof.
39
     induction 11.
40
     reflexivity.
41

    intros. rewrite ← app_comm_cons. simpl. destruct (eq_dec a x) eqn:Heq.

       + rewrite IHl1. firstorder.
43
       + apply IHl1.
   Qed.
45
46
   Lemma count_occ_head_last {A} : forall eq_dec (x: A) y l, count_occ eq_dec (y ::
    \rightarrow l) x = count_occ eq_dec (l ++ [y]) x.
   Proof.
48
     intros.
49
     rewrite count_occ_split.
50
```

```
simpl.
51
      destruct (eq_dec y x) eqn:Heq.
52

    firstorder.

53
      - firstorder.
54
    Qed.
55
56
    Lemma count_occ_last_neq \{A\} : forall eq_dec (x: A) y l, y \Leftrightarrow x \rightarrow count_occ
57
     \rightarrow eq_dec (l ++ [y]) x = count_occ eq_dec l x.
    Proof.
58
      intros.
59
      rewrite ← count_occ_head_last.
      apply count_occ_cons_neq.
      assumption.
62
    Qed.
63
64
    Lemma In_head_set {A} {eqdec : EqDec A eq} : forall l (a x : A), In x (a :: l)
     \rightarrow {a = x} + {In x l}.
    Proof.
      intros.
67
      destruct (a = x).
68
      - left. assumption.
69
      - right. destruct H.
70
         + contradiction.
71
         + assumption.
72
    Qed.
73
74
    Lemma In_app_sumbool {A} {eqdec : EqDec A eq}: forall (a : A) l1 l2, In a (l1
75
     \rightarrow ++ l2) \rightarrow {In a l1} + {In a l2}.
    Proof.
      intros. apply in_app_or in H. destruct (in_dec eqdec a l1).
77
      left. assumption. right. destruct H. contradiction. assumption.
79
80
    Lemma in_map_cons \{A\}: forall (x : A) xs ys, In (x :: xs) (map\ (cons\ x)\ ys) \rightarrow In
     \hookrightarrow XS VS.
    Proof.
82
      induction ys.
83
      - auto.
84
      - simpl in *. intros. destruct H.
85
         + left. inversion H. reflexivity.
         + right. apply IHys. assumption.
    Qed.
88
    Lemma in_map_cons_not \{A\} : forall (x \ y : A) \ xs \ ys, x \Leftrightarrow y \to \sim (In \ (x :: xs) \ (map_x := xs))
     \hookrightarrow (cons y) ys)).
    Proof.
      induction ys.
92
      - intros Heq F. exact F.
93
      - intros Heq F. destruct F.
94
         + inversion H. symmetry in H1. contradiction.
95
         + apply IHys in Heq. contradiction.
96
    Qed.
    Lemma in_map_cons_iff {A} : forall (x : A) xs ys, In (x::xs) (map (cons x) ys)
     \rightarrow \longrightarrow In xs ys.
    Proof.
100
```

```
intros. split.
101
             - apply in_map_cons.
102
             intros. induction ys.
103
                 + inversion H.
104
                 + simpl. destruct H.
                      * left. subst. reflexivity.
                      * right. apply IHys. apply H.
107
        Qed.
108
109
        Lemma some eq : forall (T : Type) (a b : T), a = b \rightarrow Some \ a = Some \ b.
110
        Proof. intros. split.

    intros Heq. subst. reflexivity.

112

    intros Heq. inversion Heq. reflexivity.

113
        Qed.
114
115
        Inductive Forall_T \{A : Type\} (P : A \rightarrow Type) : list A \rightarrow Type :=
116
        | Forall_T_nil : Forall_T P nil
        | Forall_T_cons : forall (x : A) (l : list A), P x \rightarrow Forall_T P l \rightarrow Forall_T P
          \rightarrow (x :: 1)
119
120
        Lemma Forall T forall \{A P\} {eqdec : EqDec A eq} \{l: list A\} : Forall T P l \rightarrow
121
          \rightarrow (forall x : A, In x l \rightarrow P x).
        Proof.
122
             induction 1; intros.
123
             inversion H.
124
             - destruct (x = x0).
125
                 + rewrite e in p. assumption.
126
                 + apply IHX.
                      inversion H.
                      * pose proof (Equivalence.equiv_reflexive_obligation_1 _ x0).
129

→ contradiction.

                      * assumption.
130
        Qed.
        Inductive Forall2_T {A B : Type} (R : A \rightarrow B \rightarrow Type) : list A \rightarrow list B \rightarrow Type
133
             | Forall2 T nil : Forall2 T R [] []
134
             | Forall2 T cons : forall (x : A) (y : B) (l : list A) (l' : list B),
135
                                                     R \times y \rightarrow Forall2_T R l l' \rightarrow Forall2_T R (x :: l) (y :: l').
137
        Inductive Forall2_T_idx {A B : Type} (R : nat \rightarrow A \rightarrow B \rightarrow Type) : nat \rightarrow list A
138
          \rightarrow list B \rightarrow Type :=
             | Forall2_T_idx_nil : Forall2_T_idx R 0 [] []
139
             | Forall2_T_idx_cons : forall (x : A) (y : B) (l : list A) (l' : list B) (n : A) (l' : list B) (n : A) (l' : list B) (n : B) (l' : B) (
140
               \rightarrow nat),
                                                     R n x y \rightarrow Forall2\_T\_idx R n l l' \rightarrow Forall2\_T\_idx R (S n) (x)
141
                                                       → :: l) (y :: l').
142
        Ltac dec eq one :=
143
                 match goal with
144
                  | H : ?x \neq ?x \vdash \_ \Rightarrow elimtype False; apply H; reflexivity
                  | H : ?x \implies ?y \vdash \_ \Rightarrow red in H; subst
146
                  | \vdash \{ ?x \equiv ?x \} + \{ \_ \} \Rightarrow left; reflexivity
147
                  | \vdash \{ \_ \} + \{ ?x \neq ?x \} \Rightarrow left; f_equal
148
                  | \vdash \{ ?x \equiv ?y \} + \{ \_ \} \Rightarrow right; let H := fresh in intro H; red in H;
149
```

```
(injection H || discriminate); intros; subst
150
         end.
151
152
    Ltac dec_eq := try solve [ repeat dec_eq_one ];
153
         repeat match goal with H : \_ \equiv \_ \vdash \_ \Rightarrow \text{red in } H; subst end.
    Ltac isfalse := intros F; inversion F; try contradiction.
156
157
158
    Definition kleisli option {A B C : Type} (f : A \rightarrow (option B)) (g : B \rightarrow option
159
     \rightarrow C) x :=
         match f \times with
160
         | None ⇒ None
161
         | Some y \Rightarrow g(y)
162
         end.
163
164
    Definition fmap_option {A B : Type} (f : A \rightarrow B) (a : option A) : option B :=
165
         match a with
166
         | None ⇒ None
167
         | Some x \Rightarrow Some (f x)
168
         end.
169
170
    Definition bind_option {A B : Type} (m : option A) (f : A \rightarrow option B) : option
         match m with
172
         | None ⇒ None
173
         | Some x \Rightarrow f x
174
         end.
175
    Notation "A \Longrightarrow B" := (kleisli_option A B) (at level 50, left associativity).
177
    Notation "m > f" := (bind_option m f) (at level 50, left associativity).
178
179
    Lemma kleisli_to_bind_option {A B C : Type} :
180
         forall (m : A \rightarrow option B) (n : B \rightarrow option C) x,
              (m \implies n) x = m x \gg (fun y \implies n y).
    Proof.
183
         intros. unfold kleisli_option. destruct (m x); reflexivity.
184
    Qed.
185
186
    Lemma monad_law_option_1 {A B: Type} : forall (f: A \rightarrow option B) a , Some a \gg
     \rightarrow f = f a.
    Proof. reflexivity. Qed.
188
189
    Lemma monad_law_option_2 {A : Type} : forall (m : option A) , m >= Some = m.
190
    Proof.
191
         destruct m; reflexivity.
    Qed.
193
194
    Lemma monad_law_option_3 {A B C: Type} :
195
         forall m (f : A \rightarrow option B) (g : B \rightarrow option C),
196
              (m \gg f) \gg g = m \gg (fun x \Rightarrow f x \gg g).
197
    Proof.
198
         destruct m; reflexivity.
199
    Qed.
200
```

201

```
Lemma split_rev : forall A (ts1:list A) ts2, rev (ts1 ++ ts2) = rev ts2 ++ rev

→ ts1.

    Proof.
203
      induction ts1.
204

    intros. simpl. rewrite app nil r. reflexivity.

      - intros. simpl. rewrite IHts1. rewrite app_assoc. reflexivity.
    Qed.
207
208
    Lemma rev head last : forall A (ts : list A) a, rev (a :: ts) = rev ts ++ [a].
209
    Proof.
210
      intros.
      assert (a :: ts = [a] ++ ts). reflexivity.
212
      rewrite H. apply split_rev.
213
214
215
216
    Inductive Forall2_rev {A B} R : list A \rightarrow list B \rightarrow Prop :=
    | Forall2_rev_nil : Forall2_rev R [] []
218
    | Forall2_rev_cons : forall x y l l',
219
                    R \times y \rightarrow Forall2\_rev R l l' \rightarrow Forall2\_rev R (rev (x :: l)) (rev
220
                     \rightarrow (y :: l')).
221
    Hint Constructors Forall2_rev.
223
    Lemma Forall2_head : forall A B R (l1 : list A) (l2 : list B) a b ,
224
      Forall2 R (a :: l1) (b :: l2) \rightarrow R a b.
225
    Proof.
226
      intros. inversion H. assumption.
227
    Qed.
229
    Lemma Forall2_head_T : forall A B R (l1 : list A) (l2 : list B) a b ,
230
      Forall2_T R (a :: l1) (b :: l2) \rightarrow R a b.
231
    Proof.
232
      intros. inversion X. assumption.
    Qed.
235
    Lemma app_eq_length_eq : forall A (l1 l2 : list A),
236
      l1 = l2 \rightarrow length l1 = length l2.
237
    Proof.
238
      intros. subst. reflexivity.
    Qed.
240
241
    Lemma app_singl_eq_singl_nil : forall A (l: list A) a b,
242
      [a] = l + + [b] \rightarrow l = [] / a = b.
243
    Proof.
244
      intros.
      induction l.
246

    simpl in H. inversion H. auto.

247
      apply app_eq_length_eq in H. simpl in H. inversion H.
248
        rewrite app_length in H1. simpl in H1. rewrite ← plus_n_Sm in H1. inversion
249
          → H1.
    Qed.
251
    Lemma l1_le_length_split : forall A (l: list A),
252
      1 \leq \text{length } l \rightarrow \text{exists } a l', l = a::l'.
253
    Proof.
254
```

```
destruct l.
255
      - isfalse.
256

    intros. exists a. exists l. reflexivity.

257
258
    Lemma Sn_impl_1_lt_m : forall m n, S n = m \rightarrow 1 \leq m.
260
261
      intros. omega.
262
263
264
    Lemma Forall2_idx {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2 R ns ms
     \rightarrow forall i, nth_error ns i = Some n \rightarrow
      nth\_error ms i = Some m \rightarrow R n m.
266
267
      induction 1.
268

    intros. destruct i; inversion H.

269
      - intros. destruct i.
        + simpl in *. inversion H1. inversion H2. subst. assumption.
271
        + simpl in *. eapply IHForall2.
272
           { apply H1. }
273
           { apply H2. }
274
    Qed.
275
    Lemma Forall2_idx_T {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2_T R
277
     \rightarrow ns ms \rightarrow forall i, nth_error ns i = Some n \rightarrow
      nth error ms i = Some m \rightarrow R n m.
278
    Proof.
279
      induction 1.
280

    intros. destruct i; inversion H.

      - intros. destruct i.
        + simpl in *. inversion H. inversion H0. subst. assumption.
283
        + simpl in *. eapply IHX.
284
           { apply H. }
285
           { apply H0. }
    Qed.
287
288
    Lemma Forall2_length : forall A B R (l : list A) (l' : list B), Forall2 R l l'
289
     \rightarrow length l = length l'.
    Proof.
290
      induction 1.
291
      reflexivity.
292

    simpl. rewrite IHForall2. reflexivity.

293
    Qed.
294
295
    Lemma Forall2_length_T : forall A B R (l : list A) (l' : list B), Forall2_T R l
296
     \rightarrow l' \rightarrow length l = length l'.
    Proof.
297
      induction 1.
298

    reflexivity.

299

    simpl. rewrite IHX. reflexivity.

300
    Qed.
301
    Lemma nth_error_last_length \{A: Type\} : forall ls (x : A), nth_error (ls ++ [x])
303
     \rightarrow (length ls) = Some x.
    Proof.
304
      intros.
305
```

```
induction ls.
306
      reflexivity.
307
      - simpl. assumption.
308
    Qed.
309
    Lemma Forall2_last_length {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2
311
     \rightarrow R ns ms \rightarrow
      nth_error ns ((length ns) - 1) = Some n \rightarrow
312
      nth error ms ((length ms) - 1) = Some m \rightarrow R n m.
313
    Proof.
314
      intros.
      apply Forall2_idx with ns ms (length ns - 1); try assumption.
316
      - apply Forall2_length in H. rewrite H. assumption.
317
    Qed.
318
319
    Lemma Forall2_last_length_T {A B: Type} : forall (n : A) ns (m : B) ms R ,
320
     \hookrightarrow Forall2_T R ns ms \rightarrow
      nth_error ns ((length ns) - 1) = Some n \rightarrow
321
      nth_error ms ((length ms) - 1) = Some m \rightarrow R n m.
322
    Proof.
323
      intros.
324
      apply Forall2 idx T with ns ms (length ns - 1); try assumption.
325
      apply Forall2_length_T in X. rewrite X. assumption.
    Qed.
327
328
    Lemma Forall2 last {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2 R (ns
329
     \rightarrow ++ [n]) (ms ++ [m]) \rightarrow
      R n m.
330
    Proof.
331
      intros.
332
      apply Forall2_last_length with (ns ++ [n]) (ms ++ [m]).
333
      - assumption.
334
      - rewrite app_length. rewrite plus_comm. simpl. rewrite ← minus_n_0.
335
        apply nth_error_last_length.
      - rewrite app length. rewrite plus comm. simpl. rewrite ← minus n 0.
        apply nth_error_last_length.
338
    Qed.
339
340
    Lemma Forall2 last T {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2 T R
341
     \rightarrow (ns ++ [n]) (ms ++ [m]) \rightarrow
      R n m.
342
    Proof.
343
344
      apply Forall2_last_length_T with (ns ++ [n]) (ms ++ [m]).
345
      - assumption.
346
      - rewrite app_length. rewrite plus_comm. simpl. rewrite ← minus_n_0.
        apply nth_error_last_length.
348
      - rewrite app_length. rewrite plus_comm. simpl. rewrite ← minus_n_0.
349
        apply nth error last length.
350
    Qed.
351
352
    Lemma Forall2 firstn {A B: Type} : forall ns ms (R : A \rightarrow B \rightarrow Prop), Forall2 R
     \hookrightarrow ns ms \rightarrow
      forall n ans ams, firstn n ns = ans \rightarrow firstn n ms = ams \rightarrow Forall2 R ans ams.
354
355
      induction 1.
356
```

```
- intros. destruct n; simpl in *; subst; constructor.
357

    intros. destruct n.

358
        + subst. constructor.
359
        + simpl in *. destruct ans, ams.
360
           { constructor. }
           { inversion H1. }
           { inversion H2. }
363
           { inversion H1. inversion H2. constructor.
364

    subst. assumption.

365
             - eapply IHForall2; try reflexivity. }
366
    Qed.
367
368
    Lemma Forall2_firstn_T {A B: Type} : forall ns ms (R : A \rightarrow B \rightarrow Type),
369
     \rightarrow Forall2 T R ns ms \rightarrow
      forall n ans ams, firstn n ns = ans \rightarrow firstn n ms = ams \rightarrow Forall2_T R ans
370
           ams.
    Proof.
      induction 1.
372

    intros. destruct n; simpl in *; subst; constructor.

373

    intros. destruct n.

374
         + subst. constructor.
375
         + simpl in *. destruct ans, ams.
376
           { constructor. }
           { inversion H. }
378
           { inversion H0. }
379

    inversion H. inversion H0. constructor.

380

    subst. assumption.

381
             - eapply IHX; try reflexivity. }
382
    Qed.
383
384
    Lemma firstn_init_length {A : Type} : forall init (x : A), firstn (length init)
385
     \rightarrow (init ++ [x]) = init.
    Proof.
386
      intros.
      induction init.
      reflexivity.
389

    simpl. rewrite IHinit. reflexivity.

390
    Qed.
391
392
    Lemma Forall2_init_length {A B: Type} : forall (ans : list A) ns (ams : list B)
     \rightarrow ms R , Forall2 R ns ms \rightarrow
      firstn ((length ms) - 1) ms = ams \rightarrow
394
      firstn ((length ns) - 1) ns = ans \rightarrow Forall2 R ans ams.
395
    Proof.
396
      intros.
397
      eapply Forall2_firstn.
      - apply H.
399

    apply H1.

400
      - assert (length ns = length ms).
401
         { eapply Forall2_length. apply H. } rewrite H2. apply H0.
402
    Qed.
403
    Lemma Forall2_init_length_T {A B: Type} : forall (ans : list A) ns (ams : list
405
     \rightarrow B) ms R , Forall2_T R ns ms \rightarrow
      firstn ((length ms) - 1) ms = ams \rightarrow
406
      firstn ((length ns) - 1) ns = ans → Forall2_T R ans ams.
407
```

```
Proof.
408
      intros.
409
      eapply Forall2_firstn_T.
410
      apply X.
411
      apply H0.
      - assert (length ns = length ms).
        { eapply Forall2_length_T. apply X. } rewrite H1. apply H.
414
    Qed.
415
416
    Lemma Forall2 init {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2 R (ns
417
     \rightarrow ++ [n]) (ms ++ [m]) \rightarrow
      Forall2 R ns ms.
418
   Proof.
419
      intros.
420
      eapply Forall2_init_length.
421
      apply H.
422
      - rewrite app_length. rewrite plus_comm. simpl. rewrite ← minus_n_0. apply

→ firstn_init_length.

    rewrite app_length. rewrite plus_comm. simpl. rewrite ← minus_n O. apply

424

    firstn_init_length.

    Qed.
425
426
    Lemma Forall2_init_T {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2_T R
     \rightarrow (ns ++ [n]) (ms ++ [m]) \rightarrow
      Forall2_T R ns ms.
428
    Proof.
429
      intros.
430
      eapply Forall2_init_length_T.
431
      apply X.
      - rewrite app_length. rewrite plus_comm. simpl. rewrite ← minus_n_0. apply
433

→ firstn_init_length.

    rewrite app_length. rewrite plus_comm. simpl. rewrite ← minus_n_0. apply

434

    firstn_init_length.

    Qed.
435
    Lemma Forall2_split_last {A B: Type} : forall (n : A) ns (m : B) ms R , Forall2
437
        R (ns + [n]) (ms + [m]) \rightarrow
      Forall2 R ns ms /\ R n m.
438
    Proof.
439
      intros.
440
      split.
441
      - split.
442
        + eapply Forall2 init. apply H.
443
        + eapply Forall2_last. apply H.
444

    intros. destruct H.

445
        induction H.
        + simpl. constructor.
447
          * assumption.
448
          * constructor.
449
        + simpl. constructor; assumption.
450
    Qed.
451
    Lemma Forall2_split_last_T {A B: Type} : forall (n : A) ns (m : B) ms R ,
453
        Forall2_T R (ns ++ [n]) (ms ++ [m]) \rightarrow
      prod (Forall2_T R ns ms) (R n m).
454
   Proof.
455
```

```
intros.
456
      split.
457
        + eapply Forall2_init_T. apply X.
458
        + eapply Forall2_last_T. apply X.
459
    Qed.
461
    Lemma Forall2_split_last_T_r {A B: Type} : forall (n : A) ns (m : B) ms R , prod
462
     \rightarrow (Forall2_T R ns ms) (R n m) \rightarrow Forall2_T R (ns ++ [n]) (ms ++ [m]).
    Proof.
463
     intros. destruct X.
464
     induction f.
465
        + simpl. constructor.
           * assumption.
467
           * constructor.
468
        + simpl. constructor; assumption.
469
    Qed.
470
    Lemma Forall2_head_to_last {A B : Type} : forall n m ns ms (R : A \rightarrow B \rightarrow Prop),
     \rightarrow Forall2 R (n :: ns) (m :: ms) \rightarrow Forall2 R (ns \leftrightarrow [n]) (ms \leftrightarrow [m]).
   Proof.
473
      intros.
474
      split.
475
      - intros. apply Forall2 split last. inversion H. split; assumption.
      intros. apply Forall2_split_last in H. destruct H. constructor; assumption.
477
    Qed.
478
479
    Lemma Forall2 head to last T \{A B : Type\} : forall n m ns ms <math>\{R : A \rightarrow B \rightarrow B\}
480
     \rightarrow Type), Forall2_T R (n :: ns) (m :: ms) \rightarrow Forall2_T R (ns ++ [n]) (ms ++
     \hookrightarrow [m]).
   Proof.
481
      intros. apply Forall2 split last T r. inversion X. split; assumption.
482
483
484
    Lemma Forall2_head_to_last_T_r {A B : Type} : forall n m ns ms (R : A \rightarrow B \rightarrow
     \rightarrow Type), Forall2 T R (ns ++ [n]) (ms ++ [m]) \rightarrow Forall2 T R (n :: ns) (m ::
     \rightarrow ms).
   Proof.
486
      intros. apply Forall2_split_last_T in X. destruct X. constructor; assumption.
487
    Qed.
488
    Lemma rev_nil_iff_nil \{A : Type\} : forall (ms : list A), [] = rev ms <math>\longrightarrow ms =
490
     → [].
    Proof.
491
      intros.
492
      split.
493

    intros. destruct ms.

        + reflexivity.
495
        + inversion H. apply app_eq_length_eq in H1. rewrite app_length in H1.
496
             simpl in H1.
           rewrite plus comm in H1. inversion H1.
497
      - intros. subst. reflexivity.
498
499
    501
      ms = ns / a = b.
502
```

```
Proof.
503
      induction ms.
504
      induction ns.
505
        + simpl. intros. split. reflexivity. inversion H. reflexivity.
506
        + simpl. intros. inversion H. symmetry in H2. apply app eq nil in H2.
          inversion H2. inversion H3.
      induction ns.
509
        + simpl. intros. inversion H. apply app_eq_nil in H2.
510
          inversion H2. inversion H3.
511
        + simpl in *. intros. split.
512
          { inversion H. apply IHms in H2. inversion H2. subst. reflexivity. }
          { inversion H. apply IHms in H2. inversion H2. assumption. }
   Qed.
515
516
   Lemma rev_eq \{A : Type\} : forall (ms ns: list A), rev ms = rev ns <math>\longrightarrow ms = ns.
517
   Proof.
518
      intros.
      split.
520

    intros. remember (length ms) as lms. generalize dependent ms. generalize

521

→ dependent ns. induction (lms).

        + intros. apply app_eq_length_eq in H. rewrite rev_length in H. rewrite
522

→ rev length in H.

          symmetry in Heqlms. apply length_zero_iff_nil in Heqlms. subst. simpl in
           → H.
          symmetry in H. apply length_zero_iff_nil in H. subst. reflexivity.
524
        + intros. assert (length ms = length ns).
525
          { apply app eq length eq in H. rewrite rev length in H. rewrite rev length
526
           → in H. assumption. }
          assert (1 \leq length ms).
          { rewrite ← Heqlms. firstorder. }
528
          assert (1 \leq length ns). { rewrite \leftarrow H0. assumption. }
529
          apply l1_le_length_split in H1.
530
          apply l1_le_length_split in H2.
531
          inversion H1. inversion H2. inversion H3. inversion H4.
          subst. simpl in H. apply app_last_eq in H.
          destruct H. apply IHn in H.
534
          { subst. reflexivity. }
535
          { inversion Healms. reflexivity. }
536

    intros. subst. reflexivity.

537
   Qed.
538
539
   Lemma rev_cons_iff_last {A : Type} : forall (ms : list A) x l , x :: l = rev ms
540
     \hookrightarrow \longrightarrow
      ms = rev l ++ [x].
541
   Proof.
542
      intros.
      split.
544
      intros.
545
        rewrite ← (rev_involutive ms). rewrite ← rev_head_last.
546
        apply rev_eq. symmetry. assumption.
547

    intros. subst. rewrite ← rev head last. rewrite rev involutive.

548
        reflexivity.
   Qed.
550
551
   Lemma Forall2_is_rev {A B: Type} : forall ns ms {R : A \rightarrow B \rightarrow Prop}, Forall2 R
552
     \rightarrow ns ms \longrightarrow Forall2 R (rev ns) (rev ms).
```

```
Proof.
553
      intros.
554
      split.
555
      intros. induction H.
556
        + simpl. constructor.
        + simpl. apply Forall2_head_to_last. constructor; assumption.
      - intros. remember (rev ns) as rns. remember (rev ms) as rms.
559
        generalize dependent ms. generalize dependent ns.
560
        induction H.
561
        + intros. apply rev nil iff nil in Hegrns.
562
           apply rev_nil_iff_nil in Heqrms. subst. constructor.
        + intros.
           apply rev_cons_iff_last in Heqrns.
565
           apply rev cons iff last in Hegrms.
566
           subst. eapply Forall2_split_last. split.
567
           { apply IHForall2.
568

    symmetry. apply rev_involutive.

             - symmetry. apply rev_involutive. }
570
           assumption.
571
    Qed.
572
573
    Lemma Forall2 T is rev {A B: Type} : forall ns ms {R : A \rightarrow B \rightarrow Type},
     \rightarrow Forall2_T R ns ms \rightarrow Forall2_T R (rev ns) (rev ms).
    Proof.
575
      induction 1.
576
        + simpl. constructor.
577
        + simpl. apply Forall2 head to last T. constructor; assumption.
578
    Qed.
579
    Lemma Forall2_T_is_rev_r {A B: Type} : forall ns ms {R : A \rightarrow B \rightarrow Type},
        Forall2_T R (rev ns) (rev ms) \rightarrow Forall2_T R ns ms.
    Proof.
582
      intros. remember (rev ns) as rns. remember (rev ms) as rms.
583
        generalize dependent ms. generalize dependent ns.
        induction X.
        + intros. apply rev_nil_iff_nil in Heqrns.
           apply rev_nil_iff_nil in Heqrms. subst. constructor.
587
        + intros.
588
           apply rev cons iff last in Hegrns.
589
           apply rev_cons_iff_last in Heqrms.
           subst. eapply Forall2_split_last_T_r. split.
           { apply IHX.
592
             - symmetry. apply rev involutive.
593
             - symmetry. apply rev_involutive. }
594
           assumption.
595
    Qed.
597
598
    Lemma hd_none_impl_nil \{T: Type\} : forall (ms : list T), hd_error ms = None <math>\rightarrow
599
     \rightarrow ms = nil.
    Proof.
600
        intros.
601
        induction ms.
602

    reflexivity.

603

    simpl in H. discriminate H.

604
    Qed.
605
```

```
606
    Lemma list_nonmt_split \{T: Type\} : forall (ms : list T), ms \Leftrightarrow nil \longrightarrow exists
607
        head tail, ms = head :: tail.
    Proof.
608
        intros. split.
        intros. destruct (hd_error ms) eqn:he.
610
           + exists t. exists (tl ms). remember (tl ms) as tail.
611
           assert (hd_error ms = Some t /\ tl ms = tail). auto.
612
           apply hd error tl repr in H0. assumption.
613
           + eapply hd none impl nil in he. contradiction.
614
        intros.
          destruct H as [head [tail Heq]].
           intros F. subst. discriminate F.
617
    Qed.
618
619
    Fixpoint enumerate_aux {T : Type} (ls : list T) (start : nat) : list (nat * T)
620
        match ls with
621
        | [] \Rightarrow []
622
        | x :: xs \Rightarrow (start, x) :: enumerate_aux xs (Datatypes.S start)
623
624
625
    Definition enumerate \{T: Type\} (ls: list T) : list (nat * T) :=
        enumerate_aux ls 0.
627
628
629
    Lemma prooflater : False.
630
    Proof.
631
      Admitted.
632
633
    Ltac prooflater := exfalso; apply prooflater.
634
635
    Lemma list_split_rev {A}: forall (ms : list A), (exists mshead mstail, ms =
636
     → mshead :: mstail) >→ exists msinit mslast, ms = msinit ++ [mslast].
   Proof.
637
      intros. split.
638
      + intros. destruct H as [hd [tl mssplit]]. assert (ms 	⇔ nil).
639
        { intros F. subst. discriminate F. }
640
        apply exists last in H. destruct H as [x [x0 Heq]]. exists x. exists x0.
641

→ assumption.

      + intros. destruct H as [init [tail mssplit]]. destruct ms.
642

    symmetry in mssplit. apply app_eq_nil in mssplit. destruct mssplit. subst.

643

→ discriminate H0.

    exists a. exists ms. reflexivity.

644
    Qed.
645
   Definition eq_ind_T {A : Type } :=
647
    fun (x : A) (P : A \rightarrow Type) (f : P x) (y : A) (e : x = y) \Rightarrow
648
    match e in (_ = y0) return (P y0) with
649
    \mid eq refl \Rightarrow f
650
    end.
651
   Definition list_ind_T {A : Type} :=
653
    fun (P : list A \rightarrow Type) (f : P []) (f0 : forall (a : A) (l : list A), P l \rightarrow P
654
     \rightarrow (a :: 1)) \Rightarrow
    fix F (l : list A) : P l := match l as l0 return (P l0) with
```

```
| [] \Rightarrow f
                                                                                                                                                                  | y :: 10 \Rightarrow f0 y 10 (F 10)
657
                                                                                                                                                                  end.
658
659
                  Definition rev list ind T {A : Type} :=
660
                  fun (P : list A \rightarrow Type) (H : P [])
661
                            (H0 : forall (a : A) (l : list A), P (rev l) \rightarrow P (rev (a :: l))) (l : list A)
662
                            list ind T (fun l0 : list A \Rightarrow P (rev l0)) H (fun (a : A) (l0 : list A) (IHl :
663
                                 \rightarrow P (rev l0)) \Rightarrow H0 a l0 IHl) l.
                  Definition rev_ind_T {A : Type} :=
665
                  fun (P : list A \rightarrow Type) (H : P []) (H0 : forall (x : A) (l : list A), P l \rightarrow P
666
                       \rightarrow (l ++ [x]))
                            (l : list A) \Rightarrow
667
                  (fun E : rev (rev l) = l \Rightarrow
668
                        eq_ind_T (rev (rev l)) (fun l0 : list A \Rightarrow P l0)
                                  (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev l0)) \Rightarrow H0 a (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev l0)) \Rightarrow H0 a (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (l0 : list A) (H1 : P (rev_list_ind_T P H (fun (a : A) (l0 : list A) (l0 : lost A) (l0 
670
                                      \rightarrow l0) H1) (rev l)) l E)
                            (rev involutive l).
671
672
                  Lemma nth error nil \{A: Type\} : forall x, anth error A [] x = None.
673
                  Proof.
                            induction x; reflexivity.
675
                  Qed.
676
677
678
679
                  Lemma nth_error_map \{A B\} : forall x (l : list A) (f: A \rightarrow B),
680
                            nth_error ((map f) l) x = match (nth_error l x) with
681
                                                                                                                                                                                                                                                                                                                                                                                                    l None
682
                                                                                                                                                                                                                                                                                                                                                                                                         → None
                                                                                                                                                                                                                                                                                                                                                                                                     | Some a
                                                                                                                                                                                                                                                                                                                                                                                                                       \Rightarrow
                                                                                                                                                                                                                                                                                                                                                                                                                      Some
                                                                                                                                                                                                                                                                                                                                                                                                                      (f
                                                                                                                                                                                                                                                                                                                                                                                                         → a)
                                                                                                                                                                                                                                                                                                                                                                                                   end.
684
                 Proof.
685
                            induction x.
686
                             intros. destruct l; reflexivity.
687

    intros. destruct l.

688
                                       + reflexivity.
689
                                      + simpl. apply IHx.
690
                  Qed.
691
                  Fixpoint nth_ok \{A\} (l : list A) (n : nat) : (n < length l) \rightarrow A :=
692
                            match n with
693
                             | 0 \Rightarrow match \ l \ with
694
                                                           | [] ⇒ fun proof ⇒ False_rect A (Nat.nlt_0_r _ proof)
695
                                                           | (a :: ) \Rightarrow fun \Rightarrow a
696
                                                          end
                             | S m \Rightarrow match 
698
                                                           | [] ⇒ fun proof ⇒ False_rect A (Nat.nlt_0_r _ proof)
699
                                                           | (a :: aa) \Rightarrow fun proof \Rightarrow nth_ok aa m (Lt.lt_S_n _ proof)
700
                                                           end
701
```

```
end.
702
703
    Lemma nth_ok_nth_error {A} : forall n (l : list A) p a,
704
        nth_ok \ l \ n \ p = a \longrightarrow nth_error \ l \ n = Some \ a.
705
    Proof.
      intros.
707
      split.
708
      revert n l p a.
709
        induction n.
710
        + intros [h l] p a; try solve [inversion p].
711
          simpl. intros. subst. reflexivity.
        + simpl. intros [h l] p a; try solve [inversion p].
713
          simpl in *. apply IHn.
714

    revert n l p a. induction n.

715
        + simpl. intros [h l] p a; intros H; try solve [inversion p]; try
716
         + simpl. intros [h l] p a; intros H; try solve [inversion p]. simpl in *.
         → apply IHn. assumption.
    Qed.
718
719
    Lemma nth_ok_map {A B} : forall n l (f: A \rightarrow B) p, nth_ok (map f l) n p = f
720
     \rightarrow (nth ok l n (rew (map length ) in p)).
    Proof.
721
      intros.
722
      remember (nth_ok (map f l) n p) as Ha.
723
      remember (nth ok l n (rew [lt n] map length f l in p)) as Hb.
724
      symmetry in HeqHa.
725
      symmetry in HeqHb.
726
      rewrite nth_ok_nth_error in HeqHa.
      rewrite nth_ok_nth_error in HeqHb.
728
      rewrite nth error map in HeqHa. destruct (nth error l n).
729
      - inversion HegHb. inversion HegHa. subst. reflexivity.
730
      - discriminate HegHa.
731
    Qed.
732
    Definition tuple dec {A} {eqdec : EqDec A eq} := prod eqdec eqdec eqdec.
734
735
    Fixpoint Inb \{A\} {eqdec : EqDec A eq} (x : A) l {struct l} :=
736
      match l with
737
      | [] \Rightarrow false
      | a :: ms \Rightarrow orb (x = b a) (Inb x ms)
739
740
741
    Fixpoint get_adj \{A\} {eqdec : EqDec A eq} (a : A) (l : list (A * A)) : list A :=
742
      match l with
743
      | [] \Rightarrow []
      | (a', b') :: l' \Rightarrow if (a = a') then b' :: get_adj a l' else get_adj a l'
745
      end.
746
747
    Inductive in_range_dir \{A\} (R : list (A * A)) : nat \rightarrow A \rightarrow A \rightarrow Type :=
748
    | in range base a b : In (a, b) R \rightarrow in range dir R 1 a b
749
    | in range refl a : in range dir R 0 a a
750
    \mid in_range_follow a b n : in_range_dir R n a b \rightarrow in_range_dir R (S n) a b
751
    | in_range_trans a b c n : In (b, c) R \rightarrow in_range_dir R n a b \rightarrow in_range_dir R
752
     \hookrightarrow (S n) a c.
753
```

```
Lemma in_range_dir_le \{A R\} \{a b : A\} : forall m n, n \leq m \rightarrow in_range_dir R n a
        b \rightarrow in range dir R m a b.
    Proof.
755
      intros.
756
      remember (m - n) as diff.
757
      revert n m H X Hegdiff.
      induction diff; intros.
759
      - assert (n = m). { omega. } subst. assumption.
760
      - assert (diff = m - S n).
761
         { omega. } apply (IHdiff (S n) m). omega. apply in range follow. assumption.
762
          → assumption.
    Qed.
763
764
    Lemma in_ex_dec {A} {eqdec : EqDec A eq} (a : A) R : {c & In (a, c) R} + {forall
765
     \rightarrow c :A, In (a, c) R \rightarrow False}.
    Proof.
766
      induction R.
      - right. intros. inversion H.
768

    destruct IHR.

769
        + destruct s. left. exists x. constructor 2. assumption.
770
        + destruct (a = fst a0).
771
           * left. exists (snd a0). constructor. rewrite e. apply surjective pairing.
772
           * right. intros. inversion H.
             -- assert (a = fst a0). subst. reflexivity. contradiction.
774
              -- pose proof f c0. contradiction.
775
    Defined.
776
777
    Inductive ts\_cl\_list \{A\} (R: list (A * A)) : A \rightarrow A \rightarrow Type :=
778
      | ts_R = 1  ts_R_list : forall a b, In (a, b) R \rightarrow ts_cl_list R a b
      | ts_symm_list : forall a b, ts_cl_list R a b → ts_cl_list R b a
780
      | ts_trans_list : forall a b c, ts_cl_list R a b \rightarrow ts_cl_list R b c \rightarrow
781

    ts_cl_list R a c

782
    Inductive eq cl list \{A\} (R: list (A * A)) : A \rightarrow A \rightarrow Type :=
784
        eq_R_list : forall a b, In (a, b) R \rightarrow eq_cl_list R a b
785
        eq_refl_list_l : forall a b, In (a, b) R \rightarrow eq_cl_list R a a
786
      \mid eq\_refl\_list\_r : forall \ a \ b, In (b, a) R \rightarrow eq\_cl\_list \ R \ a \ a
787
      | eq symm list : forall a b, eq cl list R a b \rightarrow eq cl list R b a
788
      \mid eq_trans_list : forall a b c, eq_cl_list R a b \rightarrow eq_cl_list R b c \rightarrow

    eq_cl_list R a c

790
791
    Lemma eq_cl_list_pump \{A\} \{R\} \{a \ b : A\} : eq_cl_list R a b <math>\rightarrow forall r,
792
     → eq_cl_list (r::R) a b.
    Proof.
      intros.
794
      induction X.
795

    constructor. constructor 2. assumption.

796
      - econstructor 2. constructor 2. exact i.
797
      - econstructor 3. constructor 2. exact i.
798
      - constructor 4. assumption.
      - econstructor 5. apply IHX1. apply IHX2.
800
    Qed.
801
802
```

803

```
Inductive sym_hull \{A\} (R: A \rightarrow A \rightarrow Type) : A \rightarrow A \rightarrow Type :=
    | sym_R : forall a b, R a b \rightarrow sym_hull R a b
805
    | sym_sym : forall a b, R a b \rightarrow sym_hull R b a.
806
807
    Definition flip_tuple \{A \ B\} : (A * B) \rightarrow (B * A) := fun tpl <math>\Rightarrow (snd tpl, fst
808
     \rightarrow tpl).
809
    Lemma flip_tuple_invol {A B} (t : A * B) : flip_tuple (flip_tuple t) = t.
810
811
      unfold flip tuple. simpl.
812
      destruct t.
813
      simpl. reflexivity.
    Qed.
815
816
    Definition flipped {A B} (R: list (A \star B)) : list (B \star A) := map flip_tuple R.
817
818
    Definition sym_hull_list \{A\} (R: list (A * A)) : list (A * A) :=
      R ++ flipped R.
820
821
    Lemma In_flipped {A B}: forall (R : list (A \star B)) a b, In (a, b) R \rightarrow In (b, a)
822

→ (flipped R).

    Proof.
823
      induction R.
824
      intros. inversion H.
825

    intros. destruct H.

826
         + simpl. left. subst. reflexivity.
827
         + constructor 2. apply IHR. assumption.
828
    Qed.
829
    Lemma flipped_invol {A B} : forall (R : list (A * B)), flipped (flipped R) = R.
831
832
      induction R.
833
      reflexivity.
834
      - simpl. rewrite flip_tuple_invol. rewrite IHR. reflexivity.
    Qed.
836
837
    Lemma sym_hull_dec \{A\} : forall (R : A \rightarrow A \rightarrow Type), (forall a b , R a b + (R a \rightarrow A \rightarrow Type))
838
     \rightarrow b \rightarrow False)) \rightarrow forall a b, sym_hull R a b + (sym_hull R a b \rightarrow False).
    Proof.
839
      intros.
840
      destruct (X a b).
841
      - left. constructor. assumption.
842
      destruct (X b a).
843
         + left. constructor 2. assumption.
844
         + right. intros F. inversion F; contradiction.
845
    Defined.
847
    Definition diag_dom \{A\} (R: list (A * A)) : list (A * A) :=
848
      map (fun a \Rightarrow (fst a, fst a)) R.
849
850
    Definition diag codom \{A\} (R: list (A * A)) : list (A * A) :=
851
      map (fun a \Rightarrow (snd a, snd a)) R.
852
853
    Definition refl_hull \{A\} (R: list (A * A)) : list (A * A) :=
854
      R ++ diag dom R ++ diag codom R.
855
856
```

```
Lemma diag_codom_eq \{A\} (R : list (A * A)) a b : In (a, b) (diag_codom R) \rightarrow
                                                    a = b.
858
    Proof.
859
      intros. unfold diag_codom in H.
860
      induction R.
861
      inversion H.

    simpl in H. destruct H.

863
        + inversion H. reflexivity.
864
        + apply IHR. apply H.
865
    Qed.
866
867
    Lemma diag_dom_in_dec {A} {eqdec: EqDec A eq} (R: list (A * A))
868
      : forall a b, In (a, b) (diag_dom R) \rightarrow {c & In (a, c) R}.
869
870
      induction R.
871
      intros. inversion H.
872
      - intros. unfold diag_dom in H. simpl map in H. apply In_head_set in H.

→ destruct H.

        + inversion e. subst. exists (snd a). constructor. apply surjective_pairing.
874
        + apply IHR in i. destruct i. eexists. constructor 2. exact i.
875
    Qed.
876
877
    Lemma diag codom in dec {A} {eqdec: EqDec A eq} (R: list (A * A))
      : forall a b, In (a, b) (diag_codom R) \rightarrow {c & In (c, a) R}.
879
    Proof.
880
      induction R.
881

    intros. inversion H.

882

    intros. unfold diag_codom in H. simpl map in H. apply In_head_set in H.

883

→ destruct H.

        + inversion e. exists (fst a). constructor. apply surjective_pairing.
        + apply IHR in i. destruct i. eexists. constructor 2. exact i.
885
    Qed.
886
887
    Lemma diag_dom_eq {A} (R : list (A \star A)) a b : In (a, b) (diag_dom R) \rightarrow
                                                    a = b.
    Proof.
890
      intros. unfold diag_dom in H.
891
      induction R.
892
      inversion H.
893

    simpl in H. destruct H.

        + inversion H. reflexivity.
895
        + apply IHR. apply H.
896
    Qed.
897
898
    Lemma diag_dom_codom_flipped {A} (R : list (A * A)) a:
899
      In a (diag\_dom(R)) \longrightarrow In a (diag\_codom(flipped R)).
   Proof.
901
      revert a.
902
      induction R.
903
      reflexivity.
904
      intros. split.
905
        + simpl. intros. destruct H.
          * left. assumption.
907
          * right. apply IHR. apply H.
908
        + simpl. intros. destruct H.
909
          * left. assumption.
910
```

```
* right. apply IHR. apply H.
911
    Qed.
912
913
    Lemma In_refl {A} {eqdec : EqDec A eq} R : forall (a b : A), In (a, b)
914
     \rightarrow (refl hull R) \rightarrow
                                                                      \{c\& prod (a = b) (In (a,
                                                                       \rightarrow c) R)} +
                                                                      \{c\& prod (a = b) (In (c,
916
                                                                       \rightarrow a) R)} + {In (a, b)
                                                                       \rightarrow R}.
    Proof.
      intros. unfold refl_hull in H. apply In_app_sumbool in H. destruct H.
918
      - right. exact i.
919
      - apply In app sumbool in i. destruct i.
920
        + left. left. pose proof diag_dom_eq _
                                                       _ _ i.
921
           apply diag_dom_in_dec in i. destruct i. eexists. split. assumption. exact
922
            \hookrightarrow i.
        + left. right. pose proof diag_codom_eq _ _ i.
923
           apply diag_codom_in_dec in i. destruct i. eexists. split. assumption.
924
            \rightarrow exact i.
    Qed.
925
926
    Inductive trans hull \{A\} (R : list (A * A)) : A \rightarrow A \rightarrow Type :=
    | trans_R : forall a b, In (a, b) R \rightarrow trans_hull R a b
928
    \mid trans_trans : forall a b c, trans_hull R a b \rightarrow trans_hull R b c \rightarrow trans_hull
929
     → Rac.
930
931
    Inductive trans_refl_hull \{A\} (R: list (A * A)) : A \rightarrow A \rightarrow Type :=
932
    | trans_refl_R : forall a b, In (a, b) R \rightarrow trans_refl_hull R a b
933
    \mid trans_refl_refl_l : forall a b, In (a, b) R \rightarrow trans_refl_hull R a a
934
    | trans_refl_refl_r : forall a b, In (b, a) R \rightarrow trans_refl_hull R a a
935
    \mid trans_refl_trans : forall a b c, trans_refl_hull R a b \rightarrow trans_refl_hull R b
936
     \hookrightarrow c \rightarrow
                                       trans refl hull R a c.
938
    Inductive list_rel \{A\} : list (A * A) \rightarrow A \rightarrow A \rightarrow Type :=
939
    | list_rel_head : forall a b x R, x = (a, b) \rightarrow list_rel (x::R) a b
940
    | list rel tail : forall a b x R, list rel R a b \rightarrow list rel (x::R) a b.
941
    Lemma In_sym_list_dec {A} {eqdec: EqDec A eq} (a b : A) R : In (a, b)
943
     \rightarrow (sym_hull_list R) \rightarrow {In (a, b) R} + {In (b, a) R}.
    Proof.
944
      intros. unfold sym_hull_list in H. apply In_app_sumbool in H. destruct H.
945

    left. auto.

946
      - right. apply In_flipped in i. rewrite flipped_invol in i. auto.
    Defined.
948
949
    Lemma In_sym_list_sym {A} (a b : A) R : In (a, b) (sym_hull_list R) \rightarrow In (b, a)
950
     Proof.
951
      intros.
      unfold sym_hull_list in *. apply in_app_or in H. apply in_or_app.
953
      destruct H. right. apply In_flipped. assumption.
954
      left. apply In flipped in H. rewrite flipped invol in H.
955
      assumption.
956
```

```
Qed.
957
958
    Lemma trans_refll_syml_eq_cl_list {A} {eqdec : EqDec A eq}: forall R (a b : A),
959
        trans_hull (refl_hull (sym_hull_list R)) a b \rightarrow eq_cl_list R a b.
960
    Proof.
961
      intros.
962
      induction X.
963
      apply In_refl in i. destruct i as [[[c [Heq HIn]] | [c [Heq HIn]]] | HIn].
964
        + subst. unfold sym hull list in HIn. apply In app sumbool in HIn. destruct
965
         → HIn.
          * econstructor 2. exact i.
          * econstructor 3. apply In_flipped in i. rewrite flipped_invol in i. exact
        + subst. unfold sym_hull_list in HIn. apply In_app_sumbool in HIn. destruct
968
         → HIn.
          * econstructor 3. exact i.
969
          * econstructor 2. apply In_flipped in i. rewrite flipped_invol in i. exact
           \hookrightarrow i.
        + unfold sym_hull_list in HIn. apply In_app_sumbool in HIn. destruct HIn.
971
          * constructor. assumption.
972
          * constructor 4. constructor. apply In_flipped in i. rewrite flipped_invol
973
           \rightarrow in i. assumption.
      - econstructor 5. exact IHX1. assumption.
    Qed.
975
976
    Lemma trans refl sym eq cl list {A} {eqdec : EqDec A eq}: forall R (a b : A),
977
        trans refl hull (sym hull list R) a b \rightarrow eq cl list R a b.
    Proof.
978
      intros.
979
      induction X.
980
      induction R.
981
        + inversion i.
982
        + unfold sym_hull_list in i. apply In_app_sumbool in i. destruct i.
983
          * apply In_head_set in i. destruct i.
             -- subst. constructor. constructor. reflexivity.
             -- constructor. constructor 2. assumption.
          * unfold flipped in i. simpl map in i. apply In_head_set in i. destruct i.
987
             -- constructor 4. constructor. inversion e. subst. constructor. destruct
988
              → a0. reflexivity.
             -- apply eq_cl_list_pump. apply IHR. unfold sym_hull_list. apply

→ in_or_app. right. assumption.

      - apply In_sym_list_dec in i. destruct i.
990
        + econstructor 2. apply i.
991
        + econstructor 3. apply i.
992
      apply In_sym_list_dec in i. destruct i.
993
        + econstructor 3. apply i.
        + econstructor 2. apply i.
995

    econstructor 5. exact IHX1. assumption.

996
    Qed.
997
998
    Lemma trans refl sym is sym {A} : forall R (a b : A),
999
      trans_refl_hull (sym_hull_list R) b a \rightarrow
      trans_refl_hull (sym_hull_list R) a b.
1001
    Proof.
1002
      intros.
1003
      induction X.
1004
```

```
- constructor. unfold sym_hull_list in *. apply in_app_or in i. apply
1005
           in_or_app.
        destruct i.
1006
         + right. apply In_flipped. assumption.
1007
         + left. apply In flipped in H. rewrite flipped invol in H. assumption.
      - econstructor 2. exact i.

    econstructor 3. exact i.

1010

    econstructor 4. apply IHX2. assumption.

1011
1012
1013
    Lemma refl_skip {A} {eqdec : EqDec A eq} R : forall (a b: A), a \Leftrightarrow b \rightarrow In (a,
        b) (refl hull R) \rightarrow In (a, b) R.
    Proof.
1015
      intros. apply In_refl in H0. destruct H0 as [[[_ [F _]] | [_ [F _]]]]].
1016

    contradiction.

1017

    contradiction.

1018
      - assumption.
    Qed.
1020
1021
    Lemma refl_sym_is_sym {A} : forall R (a b : A),
1022
         In (a, b) (refl_hull (sym_hull_list R)) \rightarrow In (b, a) (refl_hull
1023
            (sym hull list R)).
    Proof.
      intros.
1025
      unfold refl_hull in *.
1026
      apply in app or in H. destruct H.
1027

    apply in or app. left. unfold sym hull list in *. apply in app or in H.

1028

→ destruct H.

         + apply in_or_app. right. apply In_flipped. assumption.
         + apply in_or_app. left. apply In_flipped in H. rewrite flipped_invol in H.
1030
          → assumption.
       - apply in_app_or in H. destruct H.
1031
         + apply in_or_app. right. apply in_or_app. epose proof diag_dom_eq _ _ _ H.
1032

→ subst. left. assumption.

         + apply in_or_app. right. apply in_or_app. epose proof diag_codom_eq _ _ _

→ H. subst. right. assumption.
    Qed.
1034
1035
    Lemma trans sym is sym {A} : forall R (a b : A),
1036
         trans_hull (sym_hull_list R) a b \rightarrow
         trans_hull (sym_hull_list R) b a.
1038
    Proof.
1039
      intros.
1040
      induction X.
1041
      constructor. apply In_sym_list_sym. assumption.
1042

    econstructor 2. exact IHX2. assumption.

    Qed.
1044
1045
    Lemma trans_refl_sym_is_sym2 {A} : forall R (a b : A),
1046
         trans hull (refl hull (sym hull list R)) a b \rightarrow
1047
         trans hull (refl hull (sym hull list R)) b a.
1048
    Proof.
1049
      intros.
1050
      induction X.
1051
      constructor. apply refl_sym_is_sym. assumption.
1052

    econstructor 2. exact IHX2. assumption.

1053
```

```
Qed.
1054
1055
    Lemma ts_cl_list_trans_sym {A} {eqdec: EqDec A eq} : forall R (a b : A),
1056
         ts_cl_list R a b \rightarrow
1057
         trans_hull (sym_hull_list R) a b.
    Proof.
1059
       intros.
1060
       induction X.
1061

    constructor. unfold sym hull list. apply in or app. left. assumption.

1062

    apply trans sym is sym. assumption.

1063

    econstructor 2. exact IHX1. assumption.

    Qed.
1065
1066
    Lemma trans_hull_nil \{A\} : forall (a b : A), trans_hull [] a b \rightarrow False.
1067
    Proof.
1068
       intros.
1069
       induction X.
1070
       - inversion i.
1071
       - auto.
1072
    Qed.
1073
1074
    Lemma trans sym ts cl list {A} {eqdec: EqDec A eq} : forall R (a b : A),
1075
         trans hull (sym hull list R) a b \rightarrow
         ts_cl_list R a b.
1077
    Proof.
1078
       intros.
1079
       induction X.
1080
       - unfold sym_hull_list in i. apply In_app_sumbool in i. destruct i.
1081
         + constructor. assumption.
         + apply In_flipped in i. rewrite flipped_invol in i. constructor 2.
1083

→ constructor. assumption.

    econstructor 3. exact IHX1. exact IHX2.

1084
    Qed.
1085
1087
    Lemma eq_cl_list_trans_refl_sym2 {A} {eqdec: EqDec A eq} : forall R (a b : A),
1088
         eq_cl_list R a b \rightarrow
1089
         trans_hull (refl_hull (sym_hull_list R)) a b.
1090
    Proof.
1091
       intros.
       induction X.
1093

    constructor. unfold refl_hull. apply in_or_app. left. unfold sym_hull_list.

1094
           apply in_or_app. left.
         assumption.
1095

    constructor. unfold refl_hull. apply in_or_app. right. apply in_or_app.

1096
         left. unfold diag_dom. induction R.
         + inversion i.
1098
         + destruct i.
1099
           * simpl. left. subst. reflexivity.
1100
           * right. rewrite map_app. apply in_or_app. left. clear IHR. induction R.
1101
              -- inversion H.
1102
              -- simpl. destruct H.
                 ++ left. destruct a1. injection H. intros. subst. reflexivity.
1104
                 ++ right. apply IHR. assumption.
1105

    constructor. unfold refl_hull. apply in_or_app. right. apply in_or_app.

1106
         right. unfold diag_codom. induction R.
1107
```

```
+ inversion i.
1108
         + destruct i.
1109
            * simpl. left. subst. reflexivity.
1110
            * right. rewrite map_app. apply in_or_app. left. clear IHR. induction R.
              -- inversion H.
              -- simpl. destruct H.
                  ++ left. destruct a1. inversion H. reflexivity.
                  ++ right. apply IHR. assumption.
1115

    apply trans refl sym is sym2. assumption.

1116
       - econstructor 2. apply IHX1. assumption.
1117
    Qed.
1118
1119
1120
    Lemma eq_cl_list_trans_refl_sym {A} {eqdec: EqDec A eq} : forall R (a b : A),
1121
         eq cl list R a b \rightarrow
1122
         trans_refl_hull (sym_hull_list R) a b.
1123
    Proof.
1124
       intros.
1125
       induction X.
1126

    repeat constructor. induction R.

1127
         + inversion i.
1128
         + apply In head set in i. destruct i. constructor. assumption. constructor
1129
          → 2.
            simpl. apply in_or_app. left. assumption.
1130
       econstructor 2. unfold sym_hull_list. apply in_or_app. left. exact i.
1131
       econstructor 3. unfold sym hull list. apply in or app. left. exact i.
1132

    apply trans refl sym is sym. assumption.

1133

    econstructor 4. exact IHX1. assumption.

1134
    Qed.
1135
1136
    Lemma trs_eq_cl_dec {A} {eqdec : EqDec A eq} (R : list (A * A)) : (forall a b,
1137
        trans_refl_hull (sym_hull_list R) a b +
                                   (trans\_refl\_hull (sym\_hull\_list R) a b \rightarrow False)) \rightarrow
1138
                                   (forall a b, eq_cl_list R a b + (eq_cl_list R a b →
                                    → False)).
    Proof.
1140
       intros.
1141
       destruct (X a b).
1142
       - left. apply trans refl sym eq cl list. assumption.
1143
       right. intros. apply f. apply eq_cl_list_trans_refl_sym. assumption.
    Defined.
1145
1146
    Inductive t path \{A\} (R: list (A * A)) : list (A * A) \rightarrow A \rightarrow A \rightarrow Type :=
1147
     | t_path_R : forall \ a \ b, \ In (a, b) \ R \rightarrow t_path \ R [(a, b)] \ a \ b
1148
     \mid t_path_trans : forall a b c p, In (a, b) R \rightarrow t_path R p b c \rightarrow t_path R
1149
     \rightarrow ((a,b) :: p) a c.
1150
    Inductive tr_path \{A\} (R: list (A *A)) : list (A *A) \rightarrow A \rightarrow A \rightarrow Type :=
1151
    | \text{tr_path_R} : \text{forall a b, In (a, b) R} \rightarrow \text{tr_path R [(a, b)] a b}
1152
    | tr_path_refl_l : forall a b, In (a, b) R \rightarrow tr_path R [] a a
1153
    | tr path refl r : forall a b, In (b, a) R \rightarrow tr path R [] a a
1154
    | tr_path_trans : forall a b c p, In (a, b) R \rightarrow tr_path R p b c \rightarrow tr_path R
     \rightarrow ((a,b) :: p) a c.
1156
    Lemma t_path_nil \{A\} : forall (a b : A) p, t_path [] p a b \rightarrow False.
1157
    Proof.
1158
```

```
intros.
1159
       induction X; inversion i.
1160
1161
1162
    Lemma t_path_0 {A} : forall (a b : A) R, t_path R [] a b \rightarrow False.
    Proof.
1164
       intros.
1165
       inversion X.
1166
1167
1168
    Lemma t_path_start {A} R : forall P (a b a' b': A), t_path R ((a, b)::P) a' b'
1169
     \rightarrow a = a'.
    Proof.
1170
       intros.
1171
       inversion X; reflexivity.
1172
    Qed.
1173
1175
    Lemma t_path_{Sn} \{A\} (R : list (A * A)) P : forall (a b x: A),
1176
         t_{path} R ((a, x) :: P) a b \rightarrow prod (In (a, x) R) (t_{path} R P x b) + \{prod \}
1177
          \rightarrow (P = []) (prod (b = x) (In (a, b) R))}.
    Proof.
1178
       intros.
       inversion X.
1180

    subst. right. split. reflexivity. split. reflexivity. assumption.

1181
       subst. left. split; assumption.
1182
    Qed.
1183
1184
    Lemma t_path_pin_R \{A\} (R : list (A * A)) P (a b : A): <math>t_path R P a b \rightarrow forall
     \rightarrow a' b', In (a', b') P \rightarrow In (a', b') R.
    Proof.
1186
       induction 1.
1187

    intros. inversion H. injection H0. intros. subst. assumption. inversion H0.

1188
       - intros. destruct H. injection H. intros. subst. assumption. apply IHX.

→ assumption.

    Qed.
1190
1191
    Lemma t_path_trans2 \{A\} : forall P P' (a b c: A) R , t_path R P a b \rightarrow t_path R
1192
     \rightarrow P' b c \rightarrow t_path R (P ++ P') a c.
    Proof.
       induction P; intros.
1194
       inversion X.
1195
       - destruct a. pose proof t_path_start _ _ _
1196
         pose proof t_path_P_in_R _ _ _ X a0 a1 (in_eq _ _).
1197
         simpl. pose proof t_path_Sn _ _ _ _ X as [[Hin Htr]|[Heq [Hnil Hin]]].
         + constructor 2.
           * assumption.
1200
           * eapply IHP. apply Htr. assumption.
1201
         + subst. simpl. constructor 2. assumption. assumption.
1202
    Qed.
1203
1204
    Lemma t path ex \{A\} R (a b : A) : trans hull R a b \rightarrow
                                              \{P \& t path R P a b\}.
1206
    Proof.
1207
       induction 1.
1208

    eexists. constructor. assumption.

1209
```

```
- destruct IHX1. destruct IHX2. exists (x + x0). eapply t_path_trans2.
1210
         + apply t.
1211
         + assumption.
1212
    Qed.
1213
    Lemma t_path_trh {A} R (a b: A) P : t_path R P a b \rightarrow trans_hull R a b.
1215
    Proof.
1216
       induction 1.
1217

    constructor. assumption.

1218

    econstructor 2. constructor. apply i. exact IHX.

1219
    Qed.
1220
1221
    Lemma t_path_pump \{A\} R (a b : A) P : t_path R P a b \rightarrow forall ab, t_path
1222
     \rightarrow (ab::R) P a b.
    Proof.
1223
       induction 1.
1224

    intros. constructor. constructor 2. assumption.

    intros. constructor 2. constructor 2. assumption. apply IHX.

1226
    Qed.
1227
1228
    Lemma t_path_trans_dec {A} {eqdec : EqDec A eq} R (a b a' b':A)
1229
           (IHR : forall a b : A,
1230
               \{P : list (A * A) \& t_path R P a b\} + \{forall P : list (A * A), t_path\}
                 \rightarrow R P a b \rightarrow False\}):
       {P1 \& \{P2 \& ((t_path R P1 a a') * (t_path R P2 b' b))\%type \}} +
1232
                  {forall P1 P2, ((t path R P1 a a') * (t path R P2 b' b))%type \rightarrow
1233
                   → False}.
    Proof.
1234
       destruct (IHR a a') as [[P Htr] | Htr].
1235
       - destruct (IHR b' b) as [[P' Htr'] | Htr'].
1236
         + left. eexists. eexists. split. apply Htr. apply Htr'.
1237
         + right. intros P1 P2 [t1 t2]. eapply Htr'. apply t2.
1238
       right. intros P1 P2 [t1 t2]. eapply Htr. apply t1.
1239
    Defined.
1240
1241
    Definition ex_{in}rel \{A\} (a : A) R : Type := (\{c \& In (a,c) R\} + \{c \& In (c, a)\} \}
1242
     → R})%type.
1243
    Lemma ex in rel dec {A} {eqdec : EqDec A eq} : forall R (a : A), ex in rel a R +
1244
     \rightarrow (ex in rel a R \rightarrow False).
    Proof.
1245
       induction R.
1246
       - intros. right. intros. unfold ex in rel in X. destruct X as [[c Hin] | [c
1247
        → Hin]]; inversion Hin.
       intros. destruct (IHR a0).
1248
         + left. unfold ex_in_rel. destruct e as [[c H]|[c H]].
           \star left. exists c. constructor 2. assumption.
1250
           * right. exists c. constructor 2. assumption.
         + destruct a. destruct (a0 = a1).
1252
           * left. unfold ex_in_rel. rewrite e. right. exists a. constructor.
1253
            → reflexivity.
           * destruct (a0 = a).
             -- rewrite e. left. left. exists a1. constructor. reflexivity.
              -- right. intros. destruct X as [[c' H]|[c' H]].
1256
                 ++ apply In head set in H. destruct H.
1257
                    ** inversion e. subst. apply c0. reflexivity.
1258
```

```
** apply f. left. exists c'. assumption.
1259
                                   ++ apply In_head_set in H. destruct H.
1260
                                          ** inversion e. subst. apply c. reflexivity.
1261
                                          ** apply f. right. exists c'. assumption.
1262
         Defined.
1263
1264
         Lemma t_path_trans_R_and {A} {eqdec: EqDec A eq} : forall R (a b a' b' : A) P',
1265
                   t_{path} ((a', b')::R) P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b') = (a, b)\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} + \{(a', b')::R\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t_{path} R P a b\} P' a b \rightarrow \{P \& t
1266
                                                                                           {P \& prod (t_path R P a a') (b = b')} +
1267
                                                                                           \{P \ \& \ prod \ (t \ path \ R \ P \ b' \ b) \ (a = a')\} +
1268
                                                                                           \{P1 \ \& \ \{P2 \ \& \ prod \ (t_path \ R \ P1 \ a \ a') \ (t_path \ R
                                                                                             \rightarrow P2 b' b)}}.
         Proof.
1270
              intros.
1271
              induction X.
1272
              apply In_head_set in i. destruct i.
                   + inversion e. subst. left. left. left. right. auto.
                  + left. left. left. eexists. constructor. assumption.

    apply In head set in i. destruct i.

1276
                   + destruct IHX as [[[[[P Htr]| Heq]| [P [Htr Heq]]]| [P [Htr Heq]]]| [P1 [P2
1277
                     → [Htr1 Htr2]]]]; inversion e; subst.
                       * left. right. eexists. split. apply Htr. reflexivity.
1278
                       * left. left. left. right. inversion Heq. subst. assumption.
                       * left. left. left. right. reflexivity.
                       * left. right. eexists. split. apply Htr. reflexivity.
1281
                       * left. right. eexists. split. apply Htr2. reflexivity.
1282
                  + destruct IHX as [[[[[P Htr]| Heq]| [P [Htr Heq]]]| [P [Htr Heq]]]| [P1 [P2
1283
                     → [Htr1 Htr2]]].
                       * left. left. left. eexists. econstructor 2. apply i. apply Htr.
                       * inversion Heq. subst. left. left. right. eexists. split. constructor.
1285

→ assumption. reflexivity.

                       * left. left. right. eexists. split. econstructor 2. apply i. apply Htr.
1286
                         → assumption.
                       * right. eexists. eexists. split. constructor. subst. assumption. apply
                         → Htr.
                       * right. eexists. eexists. split. constructor 2. apply i. apply Htr1.
1288
                         → apply Htr2.
         Qed.
1289
1290
         Lemma prod_dec_ex {A} P Q {pdec : {p : A & P p} + {forall p, P p \rightarrow False}}
                  \{qdec : Q + (Q \rightarrow False)\} :
              {p & (P p * Q)%type} + {forall p, (P p * Q)%type \rightarrow False}.
1292
         Proof.
1293
              destruct pdec as [[p HP]| HP].
1294

    destruct qdec.

1295
                  + left. exists p. split; assumption.
                  + right. intros. apply f. destruct X. assumption.
1297

    right. intros. eapply HP. destruct X. apply p0.

1298
         Defined.
1299
1300
         Lemma t_path_dec {A} {eqdec : EqDec A eq} : forall R (a b : A), {P & t_path R P
1301
           \rightarrow a b} + {forall P, (t_path R P a b \rightarrow False)}.
         Proof.
1302
              induction R; intros.
1303
              - right. apply t_path_nil.
1304

    destruct (IHR a0 b) as [[P IH] | H1].

1305
```

```
{left. eexists. apply t_path_pump. apply IH. }
        destruct a. destruct ((a, a1) = (a0, b)).
1307
         {left. inversion e. subst. eexists. constructor. constructor. reflexivity. }
1308
        destruct (\alphaprod_dec_ex _ (fun P \Rightarrow t_path R P a0 a) (a1 = b) (IHR a0 a)).
1309
         \{ destruct (a1 = b). rewrite e. left. reflexivity. right. intros.
          → contradiction. }
         { destruct s as [p [Htr Heq]]. subst. left. eexists. eapply t_path_trans2.
1311
           apply t_path_pump. apply Htr.
1312
           - constructor. constructor. reflexivity. }
1313
        destruct (@prod_dec_ex _ (fun P \Rightarrow t_path R P a1 b) (a0 = a) (IHR a1 b)).
1314
         { destruct (a0 = a). rewrite e. left. reflexivity. right. intros.
         → contradiction. }
         { destruct s as [p [Htr Heq]]. subst. left. eexists. eapply t_path_trans2.
1316

    constructor. constructor. reflexivity.

1317
           - apply t_path_pump. apply Htr. }
1318
        destruct (t_path_trans_dec _ a0 b a a1 IHR) as [[P1 [P2 [Ht1 Ht2]]]]. left.
1319
         eexists. eapply t_path_trans2. eapply t_path_trans2. apply t_path_pump.
          → exact Ht1. constructor.
         constructor. reflexivity. apply t_path_pump. apply Ht2.
1321
        right. intros.
1322
        pose proof t_path_trans_R_and _
                                                      _ X as [[[[[P' Htr]| Heq]| [P' [Htr
1323
          → Heq]]]| [P' [Htr Heq]]]| [P1 [P2 [Htr1 Htr2]]]].
         + eapply H1. apply Htr.
        + contradiction.
        + eapply f. split. apply Htr. subst. reflexivity.
1326
        + eapply f0. split. apply Htr. assumption.
1327
         + eapply f1. split. apply Htr1. apply Htr2.
1328
    Defined.
1329
    Lemma trans_hull_dec {A} {eqdec: EqDec A eq}: forall R (a b :A), trans_hull R a
1331
        b + (trans\_hull R a b \rightarrow False).
    Proof.
1332
      intros.
1333
      destruct (t_path_dec R a b) as [[P Htr]|Htr].
      - left. eapply t path trh. apply Htr.
      right. intros. apply t_path_ex in X. destruct X. eapply Htr. apply t.
1336
    Defined.
1337
1338
    Lemma ts cl list dec {A} {eqdec: EqDec A eq}: forall R (a b: A), ts cl list R a
1339
     \rightarrow b + (ts_cl_list R a b \rightarrow False).
    Proof.
1340
      intros.
1341
      destruct (trans_hull_dec (sym_hull_list R) a b).
1342
      - left. apply trans sym ts cl list. assumption.
1343
      - right. intros. apply f. apply ts_cl_list_trans_sym. assumption.
1344
    Defined.
    Lemma eq_cl_nil_impl_refl \{A\} : forall (a b :A), eq_cl_list [] a b \rightarrow a = b.
1347
1348
      induction 1; try inversion i; subst; reflexivity.
1349
    Qed.
1350
    Inductive eq_cl_bool \{A\} (R: A \rightarrow A \rightarrow bool) : A \rightarrow A \rightarrow Type :=
1352
      \mid eq_R = 0 eq_R_bool : forall a b, R a b = true \rightarrow eq_Cl_bool R a b
1353
        eq_refl_bool : forall a b, eq_cl_bool R a b \rightarrow eq_cl_bool R a a
1354
      \mid eq_symm_bool : forall a b, eq_cl_bool R a b \rightarrow eq_cl_bool R b a
1355
```

```
\mid eq_trans_bool : forall a b c, eq_cl_bool R a b \rightarrow eq_cl_bool R b c \rightarrow
1356

    eq_cl_bool R a c

1357
1358
     Inductive eq_cl_prop \{A\} (R: A \rightarrow A \rightarrow Prop) : A \rightarrow A \rightarrow Type :=
         eq_R_prop : forall a b, R a b \rightarrow eq_cl_prop R a b
1360
       \mid eq_refl_prop : forall a b, eq_cl_prop R a b \rightarrow eq_cl_prop R a a
1361
       \mid eq_symm_prop : forall a b, eq_cl_prop R a b \rightarrow eq_cl_prop R b a
1362
       | eq trans prop : forall a b c, eq cl prop R a b \rightarrow eq cl prop R b c \rightarrow
1363

→ eq cl prop R a c

1364
     Inductive eq_cl_type \{A\} (R: A \rightarrow A \rightarrow Type) : A \rightarrow A \rightarrow Type :=
1365
       \mid eq_R = forall \ a \ b, R a b \rightarrow eq_cl_type \ R \ a \ b
1366
         eq_refl_type : forall a b, eq_cl_type R a b \rightarrow eq_cl_type R a a
1367
       \mid eq_symm_type : forall a b, eq_cl_type R a b \rightarrow eq_cl_type R b a
1368
       \mid eq_trans_type : forall a b c, eq_cl_type R a b \rightarrow eq_cl_type R b c \rightarrow
1369
        → eq_cl_type R a c
1370
1371
     Definition range (n: nat) : list nat := seq 0 n.
1372
1373
     Lemma app head \{A\} : forall (a : A) aa, a :: aa = [a] ++ aa.
1374
     Proof.
       auto.
1376
     Qed.
1377
1378
1379
     Definition combine_with {A B C} (f: A \rightarrow B \rightarrow C) := fix dummy (l:list A) (l'
1380
          :list B) : list C :=
         match l with
1381
            | [] \Rightarrow []
1382
             | a :: t \Rightarrow match l' with
1383
                          | [] \Rightarrow []
1384
                          | b :: t' \Rightarrow (f a b) :: (dummy t t')
                          end
          end.
1387
1388
1389
     Lemma nth error combine \{A B\}: forall n a b (x : A) (y : B), nth error (combine
1390
         a b) n = Some (x, y) \rightarrow
            nth\_error a n = Some x / nth\_error b n = Some y.
1391
       induction n; intros.
1392
       - inversion H. destruct a; destruct b; try discriminate H1.
1393
          simpl in *. inversion H. subst. split; reflexivity.
1394

    inversion H. destruct a; destruct b; try discriminate H1.

1395
          simpl in \star. apply IHn in H1. assumption.
     Qed.
1397
1398
1399
     Fixpoint all some {A} (l : list (option A)) : option (list A) :=
1400
       match l with
1401
       | [] \Rightarrow Some []
       | Some a :: xs ⇒ match all_some xs with
1403
                             | Some xxs \Rightarrow Some (a :: xxs)
1404
                             | None ⇒ None
1405
                             end
1406
```

```
| None :: xs \Rightarrow None
1407
       end.
1408
1409
    Definition option_concat {A} (l: list (option (list A))) : option (list A) :=
1410
       match all some l with
       | Some xs \Rightarrow Some (concat xs)
1412
       | None \Rightarrow None
1413
       end.
1414
1415
    Instance option egdec {A: Type} {equiv0 : Equivalence eq} {inneregdec: EqDec A
1416
         eq} : EqDec (option A) eq.
    Proof.
1417
       intros.
1418
       unfold EqDec.
1419
       intros.
1420
       destruct x; destruct y; try destruct (a = a0) eqn:Haa.
       - left. inversion e. reflexivity.
       right. intros F. apply some_eq in F. contradiction.

    right. intros F. discriminate F.

1424

    right. intros F. discriminate F.

1425

    left. reflexivity.

1426
    Defined.
1427
    Lemma in_not_first \{A\} : forall b a (x : A), In x (a :: b) \rightarrow a \Leftrightarrow x \rightarrow In x b.
1429
    Proof.
1430
       intros.
1431
       inversion H.
1432
       - contradiction.
1433

    assumption.

    Qed.
1435
1436
    Lemma all_some_some {A} {eq_dec : EqDec A eq} : forall ls l (x : option (A)),
1437
         all_some ls = Some l \rightarrow In \times ls \rightarrow
                                                                             \{ y \& x = Some y / \setminus
                                                                              \rightarrow In y l \}.
    Proof.
1439
       induction ls.
1440

    intros. inversion H0.

1441

    intros. simpl in H.

1442
         destruct a eqn:Ha; try discriminate H.
         destruct (all_some ls) eqn:Hls; try discriminate H.
         apply some_eq in H.
1445
         destruct (option_eqdec x a).
1446
         + subst. exists a0. split. rewrite e. reflexivity. constructor. reflexivity.
1447
         + apply in_not_first in H0; try (rewrite ← Ha; intros F; symmetry in F;
1448
          \rightarrow contradiction).
           pose proof (IHls l0 x eq_refl H0). subst. destruct X as [y [Hx Hin]].
1449
           exists y. split. assumption. constructor 2. assumption.
1450
    Qed.
1451
1452
    Lemma all_some_none_head \{A\} (a : A) ls : all_some (Some a :: ls) = None \rightarrow
1453
     → all some ls = None.
    Proof.
1454
       simpl.
1455
       intros.
1456
       destruct (all_some ls); try discriminate H. reflexivity.
1457
```

```
Qed.
1458
1459
    Lemma all_some_none_last \{A\} (a : A) ls : all_some (ls ++ [Some a]) = None \rightarrow
1460
     → all some ls = None.
    Proof.
      intros.
1462
      induction ls.
1463

    simpl in H. discriminate H.

1464

    simpl in H. destruct a0.

1465
         destruct (all some (ls ++ [Some a])).
1466
         + discriminate H.
         + apply IHls in H.
           simpl. rewrite H. reflexivity.
1469
         + reflexivity.
1470
    Qed.
1471
1472
    Lemma all_some_some_app_l {A} (l1 l2 : list (option A)) l3 : all_some (l1 ++ l2)
1473
        = Some 13 \rightarrow
                                                                               exists 14,
1474
                                                                                → all_some l1
                                                                                \rightarrow = Some 14.
    Proof.
1475
      revert 13.
      induction l1.
1477
      simpl. exists []. reflexivity.
1478

    intros. simpl. rewrite ← app comm cons in H. destruct a eqn:Ha.

1479
         + simpl in H.
1480
           destruct (all_some (_ ++ _)); try discriminate H. pose proof (IHl1 l
1481
            → eq_refl). destruct H0.
           rewrite H0. eexists. reflexivity.
1482
         + simpl in H. discriminate H.
1483
    Qed.
1484
1485
    Lemma all_some_forall_eq {A} {eqdec : EqDec A eq} (l : list (option A)) sl :
1486
        all some l = Some \ sl \rightarrow
                                                                         forall i, In i l \rightarrow
1487
                                                                          \rightarrow { s & i = Some s
                                                                          \rightarrow /\ In s sl}.
    Proof.
1488
      revert sl.
      induction l.
1490
      - intros. inversion H0.
1491

    intros. unfold all_some in H at 1. destruct a; try discriminate H. fold

1492
           (@all some A) in H.
         destruct (all_some l) eqn:Hall; try discriminate H. apply some_eq in H.
1493
         rewrite app_head in H0. apply In_app_sumbool in H0. destruct H0.
         + exists a. subst. split. symmetry. inversion i0. assumption. exfalso.
1495

→ inversion H. constructor. reflexivity.

         + subst. pose proof (IHl _ eq_refl i i0) as [s [Hisome Hin]].
1496
           exists s. split. assumption. constructor 2. assumption.
1497
    Qed.
1498
    Lemma le_tail {A n ms} \{n0 : A\}: S n < length (n0 :: ms) \rightarrow n < length ms.
1500
1501
      intros. simpl in H. apply lt_S_n in H. assumption.
1502
    Qed.
1503
```

```
1504
    Lemma nth_ok_skip {A}: forall (n0 :A) n ms p1, nth_ok (n0 :: ms) (S n) p1 =
1505
         nth_ok ms n (le_tail p1).
    Proof.
1506
       intros.
       simpl.
1508
       remember (nth_ok ms n _) as n1.
1509
       symmetry in Heqn1. symmetry. apply nth_ok_nth_error. apply nth_ok_nth_error in
1510

→ Heqn1. assumption.

    Qed.
1511
1512
    Lemma in_concat \{A\} : forall \{x : A\}, In [x] \} \to \{x \in A\} in \{x \in A\}.
1513
    Proof.
1514
     induction 1.
1515
      - intros. inversion H.
1516
      simpl. intros x [Heq | Hin].
1517
        + subst. constructor. reflexivity.
        + apply in_or_app. right. apply IHl. assumption.
1519
    Qed.
1520
1521
    Lemma combine_with_map {A B C}: forall ms ns (f : A \rightarrow B \rightarrow C),
1522
                       combine with f ms ns = map (fun mi \Rightarrow f (fst mi) (snd mi))
1523
                        Proof.
1524
       induction ms.
1525

    reflexivity.

1526
       - intros. destruct ns. { reflexivity. }
1527
         simpl. rewrite IHms. reflexivity.
1528
    Qed.
1529
1530
    Lemma length_combine_with {A} {B}: forall (ms : list A) (f : A \rightarrow nat \rightarrow B),
1531
     → length (combine_with f ms (range (length ms))) = length ms.
    Proof.
1532
       intros.
       rewrite combine with map.
       rewrite map_length.
1535
       rewrite combine_length.
1536
       unfold range.
1537
       rewrite seq length. apply Nat.min id.
1538
    Qed.
1539
1540
    Lemma lt_comb {A B n} {ms : list A} : forall (f : A \rightarrow nat \rightarrow B), n < length ms
1541
      \rightarrow n <
       length ( combine_with f ms (range (length ms))).
1542
    Proof.
1543
       intros.
       rewrite (length_combine_with ms).
1545
       assumption.
1546
    Qed.
1547
1548
1549
    Lemma seq nth error : forall len start n,
1550
         n < len \rightarrow nth\_error (seq start len) n = Some (start+n).
1551
       Proof.
1552
         induction len; intros.
1553
         inversion H.
1554
```

```
simpl seq.
1555
                        destruct n; simpl.
1556
                        auto with arith.
1557
                        rewrite IHlen; simpl; auto with arith.
1558
            Qed.
1559
1560
            Lemma pair_eq {A B} : forall (a: A) (b: B) a' b', a = a' / b = b' \rightarrow (a, b) = (a, 
1561
               \rightarrow (a', b').
            Proof.
1562
                  intros.
1563
                  split.
                  intros [HA HB]. subst. reflexivity.

    intros. inversion H. subst. split; reflexivity.

1566
1567
1568
            Lemma equivb_prop {A} {eqdec : EqDec A eq} : forall (a b : A), a =b b = true
1569
               \rightarrow \rightarrow a = b.
           Proof.
1570
                  intros. split; intros.
1571

    - unfold equiv_decb in H. destruct (a = b). inversion e. subst. reflexivity.

1572

→ discriminate H.

    subst. unfold equiv decb. destruct (equiv dec b b).

1573
                        + reflexivity.
                        + unfold complement in c. exfalso. apply c. reflexivity.
1575
            Qed.
1576
1577
            Lemma nequivb_prop {A} {eqdec : EqDec A eq} : forall (a b : A), a ⇔b b = true
1578
               \rightarrow \rightarrow a \diamondsuit b.
            Proof.
                  intros. split; intros; unfold nequiv_decb in *.
1580

    apply Bool.negb_true_iff in H. unfold equiv_decb in H. destruct (a = b).

1581

    inversion H. assumption.

                  - apply Bool.negb_true_iff. unfold equiv_decb. destruct (a = b).
1582
                        + contradiction.
                        + reflexivity.
            Qed.
1585
1586
            Definition Forall2_T_map \{A \ B\} \{P : A \rightarrow B \rightarrow Type\} \{Q \ aa \ bb\} (forall2t : A \rightarrow B \rightarrow Type)
1587
               \rightarrow Forall2 T (fun a b \Rightarrow {p : P a b & Q a b p}) aa bb) : Forall2 T P aa bb :=
                  Forall2_T_rect _ _ (fun l l' _ \Rightarrow Forall2_T P l l') (Forall2_T_nil _) (
1588
                                                                      fun a b aa bb head tail result ⇒ Forall2_T_cons _ _ _ _
1589
                                                                         → (projT1 head) result
                                                                ) aa bb forall2t.
1590
1591
            Definition Forall2_T_pair {A B} (P : A \rightarrow B \rightarrow Type) Q F ms pis f3 :=
1592
                                                                  _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} _{-} 
                  Forall2_T_rect
1593
                     → pi p } ) l l') (Forall2_T_nil _) (
                                                                      fun a b aa bb head tail result ⇒
1594
                                                                            Forall2_T_cons _ a b aa bb (existT (Q a b) head (F a b
1595
                                                                               → head)) result
                                                                ) ms pis f3.
1596
            Lemma Forall2_T_map_inv {A B} (P : A \rightarrow B \rightarrow Type) Q aa bb f3 : forall F,
1598
                        Forall2_T_map (Forall2_T_pair P Q F aa bb f3) = f3.
           Proof.
1599
                  intro F.
1600
```

```
induction f3.
1601

    constructor.

1602

    simpl. apply f_equal. assumption.

1603
1604
    Definition Rsub {A} (R R' : A\rightarrow A\rightarrow Type) := forall (pi pi' : A), R pi pi' \rightarrow R'
1606
         pi pi'.
1607
    Definition Rsub_list {A} R R' := forall (pi pi' : A), In (pi, pi') R \rightarrow In (pi,
1608
      → pi') R'.
1609
    Lemma Rsub_list_ts \{A\} R R' : \emptysetRsub_list A R R' \rightarrow \emptysetRsub A (ts_cl_list R)
1610

    (ts_cl_list R').

    Proof.
1611
       unfold Rsub_list.
1612
       unfold Rsub.
1613
       intros.
       induction X.
1615

    constructor. apply H. assumption.

1616

    constructor 2. assumption.

1617
       econstructor 3.
1618
         + apply IHX1.
1619
          + apply IHX2.
1620
    Qed.
1621
1622
    Lemma seq head : forall b a, seq a (S b) = a :: seq (S a) b.
1623
    Proof.
1624
       induction b.
1625

    reflexivity.

       - intros. rewrite IHb.
1627
         + reflexivity.
1628
    Qed.
1629
1630
    Lemma list_cons_eq \{A\} : forall (x : A) l1 l2, l1 = l2 \longrightarrow x :: l1 = x :: l2.
1631
    Proof.
1632
       intros.
1633
       split; intros; subst. reflexivity. inversion H. reflexivity.
1634
    Qed.
1635
1636
    Lemma cons_app \{A\} : forall (a : A) l, a :: l = [a] ++ l.
1637
    Proof.
1638
       intros.
1639
       auto.
1640
    Qed.
1641
1642
    Lemma forallb_existsb \{A\} : forall P (ls : list A), forallb P ls = false \rightarrow
         existsb (fun x \Rightarrow \text{negb}(P x)) ls = true.
    Proof.
1644
       intros.
1645
       induction ls.
1646
       inversion H.
1647
       simpl. apply Bool.orb_true_intro.
          simpl in H. apply Bool.andb_false_iff in H as [H1 | H2].
1649
         + left. rewrite H1. reflexivity.
1650
         + right. apply IHls. apply H2.
1651
    Qed.
1652
```

```
1653
    Fixpoint ntimes \{A\} (n : nat) (f :forall m, A) : list A :=
1654
       match n with
1655
       | 0 \Rightarrow []
1656
       | S n \Rightarrow f n :: ntimes n f
       end.
1658
1659
    Definition ntimes_proof \{A\} (n: nat) : (forall m, (m < n) \rightarrow A) \rightarrow list A :=
1660
       (fix ntimes proof inner (n': nat) (proof : n' < S n) (f : forall m, (m < n) \rightarrow
1661
        → A) {struct n'} : list A :=
          (match n' as n'' return (n' = n'') \rightarrow list A with
          \mid 0 \Rightarrow \mathsf{fun} \perp \Rightarrow []
1663
          | S n'' \Rightarrow fun Heq \Rightarrow let newproof := (Lt.lt_S_n n'' n (rew [fun \_ \Rightarrow \_] Heq
1664
          → in proof)) in
                       f n'' newproof :: ntimes_proof_inner n'' (Nat.lt_lt_succ_r _ _
1665
                        → newproof) f
         end) eq_refl) n (Nat.lt_succ_diag_r n).
1667
    Definition Forall_with_proof (n: nat) : (forall m, (m < n) \rightarrow Prop :=
1668
       (fix ntimes_proof_inner (n': nat) (proof : n' < S n) (f : forall m, (m < n) \rightarrow
1669
        \rightarrow Prop) {struct n'} : Prop :=
         (match n' as n'' return (n' = n'') \rightarrow Prop with
1670
          | 0 \Rightarrow \mathbf{fun} \Rightarrow \mathsf{True}
          | S n'' \Rightarrow fun Heq \Rightarrow let newproof := (Lt.lt_S_n n'' n (rew [fun \_ \Rightarrow \_] Heq
1672
           → in proof)) in
                       f n'' newproof /\ ntimes_proof_inner n'' (Nat.lt_lt_succ_r _ _
1673
                        → newproof) f
         end) eq_refl) n (Nat.lt_succ_diag_r n).
1674
    Ltac bsplit := apply andb_true_intro; split.
1676
1677
    Lemma seq_skip : forall n len, seq n (S len) = n :: seq (S n) len.
1678
    Proof.
1679
       reflexivity.
    Qed.
1681
1682
    Lemma nth_error_map_range \{A\} : forall n (f : nat \rightarrow A) lms, n < lms \rightarrow
1683
          nth_error (map f (range lms)) n = Some (f n).
1684
    Proof.
1685
       induction n.
       - intros. unfold range. destruct lms. {inversion H. } reflexivity.
1687
       - intros. unfold range. simpl. destruct lms. {inversion H. } simpl.
1688
         rewrite ← seg shift. rewrite map map. rewrite IHn.
1689
          + reflexivity.
1690
         + auto with arith.
1691
    Qed.
1692
1693
    Fixpoint map_i {A B} (f : nat \rightarrow A \rightarrow B) (start : nat) (ls : list A) : list B :=
1694
       match ls with
1695
       | [] \Rightarrow []
1696
       | (hd :: tl) \Rightarrow f \text{ start hd } :: map i f (S \text{ start}) tl
1697
1699
    Lemma nth_ok_proof_irel {A} : forall n (ms : list A) p1 p2, nth_ok ms n p1 =
1700
      \rightarrow nth ok ms n p2.
    Proof.
1701
```

```
intros.
1702
       remember (nth_ok _ _ _) as n1.
1703
       symmetry. symmetry in Heqn1. apply nth_ok_nth_error. apply nth_ok_nth_error in
1704
       rewrite ← Hegn1. reflexivity.
    Qed.
1706
1707
    Lemma fold_left_max_acc : forall ls i j, fold_left Nat.max ls i < j \rightarrow i < j.
1708
1709
       induction ls; intros.
1710

    assumption.

1711
       - simpl in H. apply IHls in H. apply Nat.max_lub_lt_iff in H. destruct H.
1712
        → assumption.
    Qed.
1713
1714
    Lemma fold_left_max_in : forall ls i j, fold_left Nat.max ls i \langle j \rightarrow forall k,
1715
      \rightarrow In k ls \rightarrow k < j.
    Proof.
1716
       induction ls; intros.
1717

    inversion H0.

1718

    simpl in H0. destruct H0.

1719
         + simpl in H. apply fold left max acc in H. apply Nat.max lub lt iff in H.
1720

→ destruct H. subst. assumption.

         + eapply IHls. simpl in H. apply H. assumption.
1721
    Qed.
1722
1723
    Lemma in_fold_left_max : forall ls j i, (forall k, In k ls \rightarrow k < j) \rightarrow i < j \rightarrow
1724
     → fold_left Nat.max ls i < j.</pre>
    Proof.
       induction ls; intros.
1726
       - simpl. assumption.
1727
       simpl.
1728
         destruct (dec_le i a).
1729
         + rewrite (Nat.max_r _ H1).
           eapply IHls.
           * intros. apply H. constructor 2. assumption.
1732
           * apply H. constructor. reflexivity.
1733
         + apply not_le in H1. assert (a \leq i). {
1734
              omega.
1735
           } rewrite (Nat.max l H2).
           apply IHls. intros. apply H. constructor 2. assumption. assumption.
1737
    Qed.
1738
1739
    Lemma all_some_none_exists {A} : forall (ls : list (option A)), all_some ls =
1740
     \rightarrow None \rightarrow In None ls.
    Proof.
1741
       induction ls.
1742
       - intros F. inversion F.
1743
       intros. destruct a.
1744
         + apply all_some_none_head in H. apply IHls in H. constructor 2. assumption.
1745
         + constructor. reflexivity.
1746
    Qed.
1747
1748
    Lemma nth_error_Some2 \{A\} : forall (ms : list A) n, n < length ms \rightarrow \{x \in A\}
1749
     \rightarrow nth error ms n = Some x}.
    Proof.
1750
```

```
intros.
1751
             apply nth_error_Some in H.
1752
             destruct (nth error ms n).
1753
             exists a. reflexivity.
1754
             - exfalso. apply H. reflexivity.
         Qed.
1756
1757
         Lemma nth_error_nth_ok \{A\} : forall ms (x : A) n, nth_error ms n = Some x \rightarrow \{
1758
           \rightarrow lp & nth ok ms n lp = x }.
         Proof.
1759
             intros.
             assert (nth_error ms n \diamond None).
1761
1762
                  intros F. rewrite H in F. discriminate F.
1763
1764
             rewrite nth_error_Some in H0. exists H0.
1765
             revert ms H H0.
             induction n; intros.
1767

    destruct ms. inversion H0. simpl in H0.

1768
                  simpl. simpl in H. apply some_eq in H. assumption.
1769
              destruct ms. inversion H0. simpl in H. pose proof (IHn ms H (lt_S_n _ _
1770
                \rightarrow H0)).
                  simpl. assumption.
1771
         Qed.
1772
1773
         Lemma In nth error set \{A\} {eqdec : EqDec A eq} \{A\} (x : A) : In x \{A\} + \{A\}
1774
           \rightarrow nth error l n = Some x}.
        Proof.
1775
             induction l.
             - intros. inversion H.
1777

    intros. rewrite app_head in H. apply In_app_sumbool in H. destruct H.

1778
                  + exists 0. simpl. inversion i. subst. reflexivity. inversion H.
1779
                  + apply IHl in i. destruct i. exists (S \times0). rewrite \leftarrow e. reflexivity.
1780
         Qed.
1781
1782
         Lemma in_map_set {A B} {eqdec : EqDec B eq} {f : A \rightarrow B} : forall (l : list A)
1783
           \rightarrow y, In y (map f l) \rightarrow { x & f x = y /\ In x l}.
         Proof.
1784
             induction 1.
1785
             - intros. inversion H.
             intros. rewrite app_head in H. rewrite map_app in H. apply In_app_sumbool in
                → H. destruct H.
                  + exists a. split. inversion i. assumption. inversion H. constructor.
1788

→ reflexivity.

                  + apply IHl in i. destruct i as [x [Hf Hin]]. exists x. split. assumption.
1789
                    \rightarrow constructor 2. assumption.
         Qed.
1790
1791
         Lemma forall_length_in \{A\} : forall ms (Pr : A \rightarrow Prop), (forall n pr, Pr
1792
           \rightarrow (nth_ok ms n pr)) \rightarrow (forall m, In m ms \rightarrow Pr m).
        Proof.
1793
             induction ms; intros.
             inversion H0.
1795

    destruct H0.

1796
                  + subst. pose proof (H 0 (Nat.lt 0 succ )).
1797
                       simpl in H0. assumption.
1798
```

```
+ apply IHms.
1799
           * intros. pose proof (H (S n) (lt_n_S _ _ pr)).
1800
             rewrite nth_ok_skip in H1.
1801
             erewrite nth_ok_proof_irel in H1.
1802
             apply H1.
           * assumption.
    Qed.
1805
1806
    Lemma option concat head {A}: forall (m: list (option (list A))) a oms,
1807
     \rightarrow option concat (a :: m) = Some oms \rightarrow
                                                                                    exists
                                                                                        omms,
                                                                                        option_conca
                                                                                        m =
                                                                                        Some
                                                                                        omms.
    Proof.
1809
      unfold option_concat.
1810
      simpl.
1811
      intros.
1812
      destruct a; try discriminate H.
1813
      destruct (all some m); try discriminate H.
1814
      exists (concat 10).
      reflexivity.
1816
    Qed.
1817
1818
    Lemma all_some_app_l {A} : forall (m1 : list (option A)) m2 ams, all_some (m1 ++
1819
     \rightarrow m2) = Some ams \rightarrow
                                                           { amms & all_some m1 = Some
                                                            \rightarrow amms \}.
    Proof.
1821
      induction m1.
1822
      intros. exists []. reflexivity.
1823
      - intros. simpl in H. destruct a eqn:Hl; try discriminate H.
         simpl. destruct (all some (m1 ++ m2)) eqn:Hb; try discriminate H.
         apply IHm1 in Hb. destruct Hb. rewrite e. eexists. reflexivity.
1826
    Qed.
1827
1828
    Lemma all some app l sub \{A\}: forall (m1 : list (option A)) m2 ams, all some
1829
     \rightarrow (m1 ++ m2) = Some ams \rightarrow
                                                           { amms & all_some m1 = Some amms
1830
                                                            \rightarrow /\ forall i, In i amms \rightarrow In
                                                            → i ams}.
    Proof.
1831
      induction m1.
1832
       intros. exists []. split. reflexivity. intros. inversion H0.

    intros. simpl in H. destruct a eqn:Hl; try discriminate H.

1834
         simpl. destruct (all_some (m1 ++ m2)) eqn:Hb; try discriminate H.
1835
         apply IHm1 in Hb. destruct Hb as [amms [Hall Hin]]. destruct (all_some m1);
1836

→ try discriminate Hall. eexists. split.

         + reflexivity.
1837
         + intros. apply some_eq in H. apply some_eq in Hall. rewrite ← H. destruct
          → H0.
           * constructor. assumption.
1839
           * constructor 2. apply Hin. subst. assumption.
1840
    Qed.
1841
```

```
1842
    Lemma all_some_app {A} : forall (m1 : list (option A)) m2 ams,
1843
         all some (m1 + m2) = Some ams
1844
         \rightarrow { ams1 & { ams2 & all_some m1 = Some ams1 /\ all_some m2 = Some ams2 /\
1845
          \rightarrow ams = ams1 \leftrightarrow ams2}}.
    Proof.
1846
       induction m1.
1847
       intros. exists []. exists ams. split.
1848
         + reflexivity.
1849
         + split.
1850
           * assumption.
           * reflexivity.
1852
       intros. rewrite ← app_comm_cons in H. destruct a; try (simpl in H;
1853

→ discriminate H). simpl in H.

         destruct (all_some (m1 ++ m2)) eqn:Hamm; try (simpl in H; discriminate H).
1854
         apply IHm1 in Hamm. destruct Hamm as [ams1 [ams2 [H1 [H2 H3]]]].
1855
         eexists. eexists.
         split.
1857
         + simpl. rewrite H1. reflexivity.
1858
         + split.
1859
           * apply H2.
1860
           * apply some eq in H. subst. reflexivity.
1861
    Qed.
1862
1863
1864
1865
    Lemma option_concat_app_l {A} : forall (m1 : list (option (list A))) m2 oms,
1866
         option_concat (m1 \leftrightarrow m2) = Some oms \rightarrow
                                                     { omms & option_concat m1 = Some omms}.
    Proof.
1868
       unfold option_concat.
1869
       intros.
1870
       destruct (all_some (m1 ++ m2)) eqn:Hb; try discriminate H.
1871
       apply all_some_app_l in Hb. destruct Hb. rewrite e. eexists. reflexivity.
    Qed.
1873
1874
    Lemma in_l_in_concat \{A\} : forall (x : list A) l, In x l \rightarrow (forall i, In i x \rightarrow l)
1875
     \rightarrow In i (concat l)).
    Proof.
1876
       induction 1.
       - intros F. inversion F.
1878
       - intros. simpl in *. destruct H.
1879
         + apply in_or_app. left. subst. assumption.
1880
         + apply in_or_app. right. apply IHl.
1881
           * assumption.
1882
           * assumption.
    Qed.
1884
1885
    Lemma concat_in \{A\} : forall l1 l2 (x : A), (forall i, In i l1 \rightarrow In i l2) \rightarrow In
1886
     \rightarrow x (concat l1) \rightarrow In x (concat l2).
    Proof.
1887
       induction l1.

    intros. inversion H0.

1889

    intros.

1890
         simpl in H0.
1891
         apply in_app_or in H0.
1892
```

```
destruct H0.
1893
         + eapply in_l_in_concat in H0.
1894
           * apply H0.
1895
           * apply H. constructor. reflexivity.
1896
         + apply IHl1.
           intros. apply H. constructor 2. assumption. assumption.
    Qed.
1899
1900
    Lemma option concat app l sub {A}: forall (m1: list (option (list A))) m2 oms,
1901
         option concat (m1 + m2) = Some \ oms \rightarrow \{ omms \& option \ concat \ m1 = Some \ omms \}
      \rightarrow /\ forall x, In x omms \rightarrow In x oms}.
    Proof.
1902
       unfold option_concat.
1903
1904
       destruct (all some (m1 ++ m2)) eqn:Hb; try discriminate H.
1905
       apply all_some_app_l_sub in Hb. destruct Hb as [amms [Hall Hin]].
1906
       rewrite Hall. exists (concat amms). split. reflexivity. intros.
       apply some_eq in H. subst.
1908
       eapply concat_in.
1909
       - intros. apply Hin. apply H.
1910
       - assumption.
1911
1912
    Qed.
    Lemma option_concat_app {A} : forall (m1 : list (option (list A))) m2 oms,
1914
         option_concat (m1 \leftrightarrow m2) = Some oms \rightarrow { oms1 & {oms2 & option_concat m1 =
         Some oms1 /\ option concat m2 = Some oms2 /\ oms = oms1 \leftrightarrow oms2}\.
    Proof.
1915
       unfold option_concat.
1916
       intros.
1917
       destruct (all_some (m1 ++ m2)) eqn:Hb; try discriminate H.
1918
       apply all some app in Hb. destruct Hb as [ams1 [ams2 [H1 [H2 Heq]]]].
1919
       rewrite H1. rewrite H2. rewrite ← some_eq in H. rewrite Heq in H.
1920
       rewrite concat_app in H. exists (concat ams1). exists (concat ams2).
1921
        → firstorder.
    Qed.
1922
1923
    Lemma ts_cl_list_sub \{A\} : forall (R: list (A \star A)) R', Rsub_list R R' \rightarrow Rsub
1924
     → (ts cl list R) (ts cl list R').
    Proof.
1925
       unfold Rsub.
       intros.
1927
       induction X.
1928
       - constructor. apply H. assumption.
1929

    constructor 2. assumption.

1930

    econstructor 3. apply IHX1. assumption.

1931
    Qed.
1932
1933
    Lemma Rsub_trans \{t\} : forall (A B C : t \rightarrow t \rightarrow Type) , Rsub A B \rightarrow Rsub B C \rightarrow
1934
         Rsub A C.
    Proof.
1935
       firstorder.
1936
    Qed.
1937
1938
    Lemma Rsublist_app \{A\} : forall ls1 ls2 (ms : list (A * A)), Rsub_list ls1 ms \rightarrow
1939
     → Rsub_list ls2 ms → Rsub_list (ls1 ++ ls2) ms.
    Proof.
1940
```

```
unfold Rsub list.
1941
              intros.
1942
              apply in_app_or in H1.
1943
              destruct H1.
1944
              + apply H. assumption.
              + apply H0. assumption.
1946
         Qed.
1947
1948
         Lemma Rsub in app \{A: Type\} {eqdec : EqDec A eq} : forall (R : A \rightarrow A \rightarrow Type)
1949
                   oms1 oms2 , (Rsub (fun p p' \Rightarrow In (p, p') oms1) R) \rightarrow Rsub (fun p p' \Rightarrow In
            \rightarrow (p, p') oms2) R \rightarrow Rsub (fun p p' \Rightarrow In (p, p') (oms1 \leftrightarrow oms2)) R.
         Proof.
1950
              unfold Rsub. intros.
1951
              apply In_app_sumbool in H.
1952
              destruct H.
1953
              + apply X. assumption.
1954
              + apply X0. assumption.
         Qed.
1956
1957
         Lemma Rsub_in_concat \{A: Type\} {eqdec : EqDec A eq} : forall (R : A \rightarrow A \rightarrow Type)
1958
                   (forall m, In m l \rightarrow (Rsub (fun p p' \Rightarrow In (p, p') m) R)) \rightarrow Rsub (fun p
1959
                      \rightarrow p' \Rightarrow In (p, p') (concat l)) R.
         Proof.
1960
              induction l.
1961
              - intros. unfold Rsub. intros. inversion H.
1962
               - intros. simpl. apply Rsub in app.
1963
                   + apply X. constructor. reflexivity.
1964
                   + apply IHl. intros. apply X. constructor 2. assumption.
         Qed.
1966
1967
         Lemma all_some_map2 {A B} {eqdec : EqDec A eq} : forall ms (f: A \rightarrow option B) l,
1968
                   all_some (map f ms) = Some l \rightarrow \mathbf{forall} \ m, In m ms \rightarrow \{x \ \delta \ f \ m = Some \ x\}.
         Proof.
1969
              induction ms.
1970

    intros. inversion H0.

1971
              - intros. simpl in H. destruct (f a) eqn:Hfa; try discriminate H.
1972
                   destruct (all_some _) eqn:Hall; try discriminate H. apply In_head_set in H0.
1973
                      → destruct H0.
                   + subst. exists b. apply Hfa.
                   + eapply IHms.
1975
                        * apply Hall.
1976
                        * assumption.
1977
         Qed.
1978
1979
         Lemma all_some_map {A B} {eqdec : EqDec B eq} : forall ms (f : A \rightarrow option B) l,
1980
                   all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n ms / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n ms / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n ms / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n ms / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n ms / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n ms / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n m s / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n m s / f n = all_some (map f ms) = Some l \rightarrow (forall m, In m l \rightarrow \{n \& In n m s / f n \&
1981
                      → Some m}).
         Proof.
1982
              induction ms.
1983
               - intros. inversion H. subst. exfalso. inversion H0.
1984

    intros. simpl in *. destruct (f a) eqn:Hfa; try discriminate H.

                   destruct (all_some _) eqn:Hall; try discriminate H. apply some_eq in H.
                   pose proof (IHms f _ Hall m).
1987
                   rewrite \leftarrow H in H0. rewrite app head in H0. apply In app sumbool in H0.
1988

→ destruct H0.
```

```
+ assert (m = b). subst. inversion i. subst. reflexivity. inversion H.

→ subst. exists a. split.

           * left. reflexivity.
1990
           * assumption.
1991
         + apply X in i. destruct i. destruct a0. exists x.
           split.
           * right. assumption.
           * assumption.
1995
    Qed.
1996
1997
    Lemma in_combine_range {A} : forall (ls : list A) a n start pr, (In (a, start +
1998
     \rightarrow n) (combine ls (seq start (length ls)))) \rightarrow nth_ok ls n pr = a.
    Proof.
1999
       intros. rewrite nth_ok_nth_error. apply In_nth_error in H. destruct H as [n0
2000
        → H]. apply nth_error_combine in H. destruct H.
       assert (nth_error ls n0 \Leftrightarrow None). \{ intros F. rewrite H in F. discriminate F.
2001
        → }
       apply nth_error_Some in H1.
2002
       pose proof (seq_nth_error (length ls) start n0 H1). rewrite H0 in H2. apply
2003

→ some_eq in H2. assert (n = n0). omega.
       subst. assumption.
2004
    Qed.
2005
2007
    Inductive Exists_T \{A : Type\} (P : A \rightarrow Type) : list A \rightarrow Type :=
2008
    | Exists T cons hd : forall (x : A) (l : list A), P x \rightarrow Exists T P (x :: l)
2009
    | Exists_T_cons_tl : forall (x : A) (l : list A), Exists_T P l \rightarrow Exists_T P (x
2010
     2011
    Lemma nth_error_Some3 \{A\} : forall ms n (a :A), nth_error ms n = Some a \rightarrow n <
2012
     → length ms.
    Proof.
2013
       intros.
2014
       apply nth_error_Some.
       intros F. rewrite H in F. discriminate F.
    Qed.
2017
2018
    Lemma all_some_length \{A\} : forall ls (ls' : list A), all_some ls = Some ls' \rightarrow
2019
     → length ls = length ls'.
    Proof.
       induction ls.
2021
       - intros. inversion H. reflexivity.
2022
       - intros. destruct a; try discriminate H. simpl in *. destruct (all_some ls)
2023

→ eqn:Hall; try discriminate H.
         destruct ls'; try discriminate H. apply some_eq in H. rewrite \leftarrow H. simpl.
2024

→ apply eq_S.

         apply IHls. reflexivity.
2025
    Qed.
2026
2027
    Lemma map_in {A B}: forall ls (f: A \rightarrow B) x, In x ls \rightarrow In (f x) (map f ls).
2028
    Proof.
2029
       induction ls.
       intros. inversion H.
2031

    intros. destruct H.

2032
         + simpl. subst. left. reflexivity.
2033
         + simpl. right. apply IHls. assumption.
2034
```

```
Qed.
2035
2036
    Lemma nth_ok_in {A} : forall (ls : list A) x pr, In (nth_ok ls x pr) ls.
2037
    Proof.
2038
      intros.
      remember (nth_ok ls x pr). symmetry in Heqa. apply nth_ok_nth_error in Heqa.
2040
           eapply nth_error_In. apply Heqa.
    Qed.
2041
2042
    Lemma all some some head {A} : forall a (a0: A) ls ls0,
2043
         all_some (a :: ls) = Some (a0 :: ls0) \rightarrow a = Some a0 / all_some ls = Some
          → ls0.
    Proof.
2045
      intros.
2046
      simpl in H. destruct a eqn:Ha; try discriminate H. destruct (all_some ls) eqn:
2047

→ Hall; try discriminate H.

      apply some_eq in H. inversion H. split; reflexivity.
    Qed.
2049
2050
    Lemma all_some_nth {A} : forall (ls : list (option A)) ls' (Hall : all_some ls =
2051
        Some ls') x (pr : x < (length ls)),
        nth_ok ls x pr = Some (nth_ok ls' x (rew (all_some_length _ _ Hall) in pr)).
2052
    Proof.
      induction ls.
2054

    intros. inversion pr.

2055
      intros. destruct ls'.
2056
        + pose proof (all_some_length _ _ Hall). inversion H.
2057
        + pose proof (all_some_some_head _ _ _ _ Hall) as [Ha Hall2]. destruct x.
2058
           * simpl. assumption.
           * simpl.
2060
             erewrite nth_ok_proof_irel.
2061
             erewrite (nth_ok_proof_irel _ ls').
2062
             apply IHls. Unshelve.
2063
             ** firstorder.
             ** assumption.
    Qed.
2066
2067
    Lemma repeat_rev {A} : forall (a :A) n, repeat a (S n) = repeat a n ++ [a].
2068
    Proof.
2069
      induction n.
2070
      reflexivity.
2071

    simpl. rewrite ← IHn. reflexivity.
```

2072

2073

Qed.