CSC488S Source Language Semantic Analysis

This document describes the semantic analysis that should be performed on the project source language. The semantic analysis actions are described in terms of a set of semantic analysis actions **S**??

Semantic Analysis Rules

program:	**S**00 scope **S**01
statement:	variable ':' '=' expression **S**34 , '**if**' expression **S**30 '**then**' statement '**fi**' , '**if**' expression **S**30 '**then**' statement '**else**' statement '**fi**' , '**loop**' statement '**pool**' '**exit**' **S**50 , '**exit**' integer **S**50 **S**53 , '**exit**' '**when**' expression **S**30 **S**50 , '**exit**' integer '**when**' expression **S**30 **S**50 **S**53 , '**result**' expression **S**51 **S**35 , '**return**' **S**52 , '**put**' output , '**get**' input , procedurename '(' **S**44 argumentList ')' **S**43 , **S**06 scope **S**07 , statement statement
scope	'**begin**' declaration **S**02 statement '**end**' , '**begin**' statement '**end**'
declaration:	type variablenames **S**47 , functionHead **S**49 **S**04 **S**54 scope **S**05 **S**13 , procedureHead **S**49 **S**08 scope **S**09 **S**13 , '**forward**' functionHead **S**11 '**forward**' procedureHead **S**17 declaration declaration
functionHead:	type functionname '(' **S**14 parameterList ')' **S**12
procedureHead:	'**proc**' procedurename '(' **S**14 parameterList ')' **S**18
variablenames:	variablename **S**10 , variablename '[' integer ']' **S**48 , variablename '[' bound '.' '.' bound **S**46 ']' **S**19 , variablenames ',' variablenames
bound	integer , '-' integer

type: '**integer**' **S**21 ,
'**boolean**' **S**20

output: expression **S**31 ,
text ,
'**skip**' ,
output ',' output

input: variable **S**31 ,
input ',' input

argumentList: arguments ,
% EMPTY

arguments: expression **S**45 **S**36 ,
arguments ',' arguments

parameterList: parameters ,
% EMPTY

parameters: type parametername **S**16 **S**15 ,
parameters ',' parameters

variable: variablename **S**26 ,
arrayname '[' expression **S**31 ']' **S**27

expression: integer **S**21 ,
'-' expression **S**31 **S**21 ,
expression **S**31 '+' expression **S**31 **S**21 ,
expression **S**31 '-' expression **S**31 **S**21 ,
expression **S**31 '*' expression **S**31 **S**21 ,
expression **S**31 '/' expression **S**31 **S**21 ,
'**true**' **S**20 ,
'**false**' **S**20 ,
'!' expression **S**30 **S**20
expression **S**30 '&' expression **S**30 **S**20 ,
expression **S**30 '|' expression **S**30 **S**20 ,
expression '=' expression **S**32 **S**20 ,
expression '!' '=' expression **S**32 **S**20 ,
expression **S**31 '<' expression **S**31 **S**20 ,
expression **S**31 '<' '=' expression **S**31 **S**20 ,
expression **S**31 '>' expression **S**31 **S**20 ,
expression **S**31 '>' '=' expression **S**31 **S**20 ,
'(' expression ')' **S**23 ,
variable ,
functionname '(' **S**44 argumentList ')' **S**43 **S**28 ,
parametername **S**25 ,

variablename: identifier **S**37

arrayname: identifier **S**38

functionname: identifier **S**40

procedurename: identifier **S**41

parametername: identifier **S**39

Semantic Analysis Operators

Scopes and Program

These semantic operators are used to keep track of scopes in the program being compiled.

S00 Start program scope.
S01 End program scope.
S02 Associate declaration(s) with scope.
S04 Start function scope.
S05 End function scope.
S06 Start statement scope.
S07 End statement scope.
S08 Start procedure scope.
S09 End procedure scope.

Declarations

These semantic operators make entries in the symbol table for the current scope. All of the *Declare...* operators should check that the identifier being declared has not already been declared in the current scope.

S10 Declare scalar variable.
S11 Declare forward function .
S12 Declare function with parameters (if any) and specified type.
S13 Associate scope with function/procedure.
S14 Set parameter count to zero.
S15 Declare parameter with specified type.
S16 Increment parameter count by one.
S17 Declare forward procedure.
S18 Declare procedure with parameters (if any).
S19 Declare array variable with specified lower and upper bounds.
S46 Check that lower bound is <= upper bound.
S47 Associate type with variables.
S48 Declare array variable with specified upper bound.
S49 If function/procedure was declared forward, verify forward declaration matches.
S54 Associate parameters if any with scope.

Statement Checking

These semantic operators check various correctness conditions for statements.

S50 Check that **exit** statement is inside a loop.
S51 Check that **result** is inside a function
S52 Check that **return** statement is inside a procedure.
S53 Check that integer is > 0 and <= number of containing loops.

Expressions Types

These semantic operators are used to keep track of the type of expressions. In the model used in this document, a type (integer or boolean) is a associated with the left hand side of each rule in the expression part of the grammar. The *Set result type to ...* semantic operators (somehow) associate a type with the left hand side. This same mechanism is used to keep track of types in declarations.

S20 Set result type to boolean.
S21 Set result type to integer.
S23 Set result type to type of expression.
S25 Set result type to type of parametername.
S26 Set result type to type of variablename.
S27 Set result type to type of array element.
S28 Set result type to result type of function.

Expression Type Checking

These semantic operators check that the type of an expression is correct for the use that is being made of the expression.

S30 Check that type of expression is boolean.
S31 Check that type of expression or variable is integer.
S32 Check that left and right operand expressions are the same type.
S34 Check that variable and expression in assignment are the same type.
S35 Check that expression type matches the return type of enclosing function.
S36 Check that type of argument expression matches type of corresponding formal parameter.
S37 Check that identifier has been declared as a scalar variable.
S38 Check that identifier has been declared as an array.
S39 Check that identifier has been declared as a parameter.

Functions, procedures and arguments

These semantic operators are used the check that procedures and functions are being used correctly.

S40 Check that identifier has been declared as a function.
S41 Check that identifier has been declared as a procedure.
S43 Check that the number of arguments is equal to the number of formal parameters.
S44 Set the argument count to zero.
S45 Increment the argument count by one.

Addressing for Variables

The pseudo machine that you will be generating code for uses a simple form of addressing for variables and parameters called *lexic level , order number* addressing.

An address is a pair of numbers: `( lexic level , order number )` where:

lexic level is the *static* depth of nesting of scopes in the program. The main program is lexic level zero, scopes directly inside the main program are lexic level one, etc. It is your design choice whether the lexic level gets incremented for all scopes or only for *major scopes* (program, functions and procedures). Only for major scopes is strongly recommended.

order number is an integer that uniquely identifies each variable declared in a scope. Conventionally the first variable declared in a scope is assigned order number zero, the second order number one, etc. This is particularly convenient for addressing on the pseudo machine that you will be generating code for.

It is really convenient to setup this basis for addressing variables during semantic analysis when the scopes in a program are being identified and the declarations in each scope are being processed. Every variable and parameter needs to be associated with a scope (its *lexic level*) and its location in the scope (its *order number*.

Functions and procedures are addressed using the memory address of their first instruction. Addressing for functions and procedures will be discussed in the forthcoming code generation document.

No Change to Project Language Syntax

Nothing in this document is intended to change the syntax of the course project language. If you discover a case where the syntax of the language in this document differs from the Source Language Reference Grammar, it is an error in this document, not an intentional change. Please notify the instructor if you think there is an error in this document.