## 1. Moment Generating Function

a. We define the MGF  $\phi_X(t)$  of a random variable X as the following function of a real variable t:

$$\phi_X(t) = \mathbb{E}ig[e^{tX}ig] = \mathbb{E}igg[1 + tX + rac{(tX)^2}{2!} + rac{(tX)^3}{3!} + \ldotsig]$$

- b. When X is Bernoulli (p),  $\phi_X(t) = E[e^{t}] = (1-p)e^{t} + p^*e^{t} = (1-p)e^{t}$
- c. Moments derive when t > 0

$$\phi_X(t) = \mathbb{E}ig[e^{tX}ig] = \sum_x e^{tx} p_X(x)$$

i. If X is discrete, then

$$\phi_X(t) = \mathbb{E}ig[e^{tX}ig] = \int_{-\infty}^{+\infty} e^{tx} f_X(x) dx$$

- ii. If X is continuous, then
- d. Why do we call it moment generating function?
  - i D

$$_{ ext{e. Claim:}}\,M_k(X)=\mathbb{E}ig[X^kig]=\phi_X^{(k)}(0)$$

- f. In other words, taking the k-th derivative, and setting t = 0 yields the k-th moment
- g. Proof sketch (induction)
  - i. Base case:

$$\phi'(0) = rac{d}{dt}\phi(t)ig|_{t=0} = \mathbb{E}ig[Xe^{0X}ig] = \mathbb{E}[X]$$

ii. IH: 
$$\phi^{(k)}(t) = \mathbb{E}\left[X^k e^{tX}\right]_{ ext{ so we obtain}}$$
  $\phi^{(k+1)}(t) = rac{d}{dt}\phi^{(k)}(t) = rac{d}{dt}\mathbb{E}\left[X^k e^{tX}
ight] = \mathbb{E}\left[X^{k+1} e^{tX}
ight]$ 

Substitute t = 0

- 2. Exercises on mgfs
  - a. Let  $\varphi_X(t)$  be the mgf of X. Let Y=aX+b. What is the mgf of Y?

$$\phi_Y(t) = \mathbb{E}\Big[e^{t(aX+b)}\Big] = e^{tb}\mathbb{E}\Big[e^{(at)X}\Big] = e^{tb}\phi_X(at)$$

b. Let X, Y be independent RVs with mgfs  $\phi_X(t)$ ,  $\phi_Y(t)$ . What is the mgf  $\phi_{X+Y}(t)$  of the sum X+Y? What if we have the sum of n ind. variables?

$$\phi_{X+Y}(t) = \mathbb{E}\Big[e^{t(X+Y)}\Big] = \mathbb{E}ig[e^{tX}e^{tY}ig] = \mathbb{E}ig[e^{tX}ig]\mathbb{E}ig[e^{tY}ig] = \phi_X(t)\phi_Y(t)$$

c. What is the MGF of the normal N(0,1)?

$$\phi_{\mathcal{N}(0,t)}(X) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} e^{tx} e^{-\frac{x^2}{2}} dx = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} e^{-\frac{1}{2}(x^2 - 2tx)} dx = \frac{e^{-\frac{t^2}{2}}}{\sqrt{2\pi}} \int_{-\infty}^{\infty} e^{-\frac{1}{2}(x^2 - 2tx + t^2)} dx = e^{-t^2/2}$$

PDF of e^(-0.5\*(x-t)^2)) is sqrt( $2\pi$ ) since it is N ~ (t,1)

d. What is the mgf of  $X \sim Bernoulli$  (p)?

$$\phi_X(t) = pe^t + (1-p)$$

e. What is the mgf of  $X \sim Bin(n,p)$ ?

$$\sum k = 0 \text{ to } n \to {}_{n}C_{k} * p^{k} * (1-p)^{n}(n-k) * e^{t}(k)$$

$$= \sum k = 0 \text{ to } n \to {}_{n}C_{k} * (p^{*}e^{t})^{k} * (1-p)^{n}(n-k)$$

$$\phi_X(t) = \prod_{i=1}^n ig(pe^t + (1-p)ig) = ig(1 + ig(e^t - 1)pig)ig)^n$$

f. What is the mgf of  $X \sim Poisson(\lambda)$ ?

$$\phi_X(t)=\sum_{n=0}^{+\infty}\Pr(X=n)e^{tn}=\sum_{n=0}^{+\infty}\,rac{e^{-\lambda}\lambda^n}{n!}e^{tn}=\sum_{n=0}^{+\infty}\,rac{e^{-\lambda}ig(\lambda e^tig)^n}{n!}=e^{\lambdaig(e^t-1ig)}$$

- 3. General form of Chernoff bound
  - a. Let X be a random variable. Then for any  $\epsilon \geq 0$  , we have that

$$\Pr(X \geq \epsilon) \leq e^{-\phi(\epsilon)}$$
 where

$$\phi(\epsilon) = \max_{s>0}(s\epsilon - \log M_X(s)), \ ext{ where } M_X(s) = \mathbb{E}ig[e^{sX}ig]$$

- b. Proof (sketch)
  - i. We apply Markov's inequality as follows, and then we minimize LHS over all s:

$$\Pr(X \geq \epsilon) = \Prig(e^{sX} \geq e^{s\epsilon}ig) \leq rac{M_X(s)}{e^{s\epsilon}} = e^{-s\epsilon + \log M_X(s)}$$

- 4. Chernoff bound
  - a. Let  $X_i = 1$  with probability  $p_i$ , 0 with prob. 1- $p_i$

$$X = \sum_{i=1}^{n} X_i$$
, and  $\mu = \mathbb{E}[X] = \sum_{i=1}^{n} p_i$ 

- b. Define
- c. Then, the following probability inequality holds

$$\Pr(|X - \mu| \geq \delta \mu) \leq 2e^{-rac{\delta^2}{3}\mu} \, ext{for all} \, 0 < \delta < 1$$

(much sharper bound by  $\mu$  since it goes down exponentially)  $\rightarrow$  when we know more information about the distribution unlike chebyshev (we do not specifically know the distribution and is more general)

- 5. Application: flipping a fair coin
  - a. Suppose we flip a fair coin n times. Let X be the number of Heads
    - i. Chebyshev's inequality:

$$\Pr\bigg(\bigg|\frac{X}{n} - \frac{1}{2}\bigg| \ge \epsilon\bigg) = \Pr\bigg(\bigg|X - \frac{n}{2}\bigg| \ge \epsilon n\bigg) \le \frac{\mathbb{V}\mathrm{ar}[X]}{n^2\epsilon^2} = \frac{1}{4n\epsilon^2}$$

ii. Chernoff bound (for which  $\varepsilon$  is this valid? Always check the assumptions before applying a theorem)

$$\Pr\bigg(\bigg|\frac{X}{n} - \frac{1}{2}\bigg| \geq \epsilon\bigg) = \Pr\bigg(\bigg|X - \frac{n}{2}\bigg| \geq \epsilon n\bigg) \leq 2e^{-(2\epsilon)^2\frac{n}{6}} = 2e^{-\frac{2n\epsilon^2}{3}}$$

- 6. Sampling theorem
  - a. Chernoff bound comes in many variations
  - b. A beautiful corollary of Chernoff bound is the following
    - i. Sampling theorem: Given n independent 0-1 RVs Xi such that Pr(Xi=1) = p(i=1...n), where  $n \ge \frac{3}{\epsilon^2} \ln \left(\frac{2}{\delta}\right)$  then the following holds:

$$\left| \Pr \left( \left| rac{\sum_{i=1}^n X_i}{n} - p 
ight| \leq \epsilon 
ight) \geq 1 - \delta$$

c. Let's assume we want to be  $\varepsilon$ -accurate with a certain confidence 1- $\delta$  about the fraction of population that will vote for Biden:

$$n \geq rac{3}{\epsilon^2} \mathrm{ln}\left(rac{2}{\delta}
ight)$$

The poll size does not depend on the size of the total population

- 7. The Thumbtack problem
  - a. A billionaire from the suburbs of Boston asks you a question:
    - i. He says: I have thumbtack, if I flip it, what's the probability it will fall with the nail up?
    - ii. You say: Please flip it a few times
  - b. Assumptions
    - i. Pr(Heads) = p, Pr(Tails) = 1-p
    - ii. Flips are iid
      - 1. Consider the sequence HTHHTHT  $\rightarrow$  p<sup>4</sup>(1-p)<sup>3</sup>