

1. Control hijacking attacks

- a. Attacker's goal: take over target machine (e.g. web server)
- b. Execute arbitrary code on target by hijacking application control flow
- c. Examples:
 - Buffer overflow and integer overflow attacks
 - Format string vulnerabilities
 - Use after free

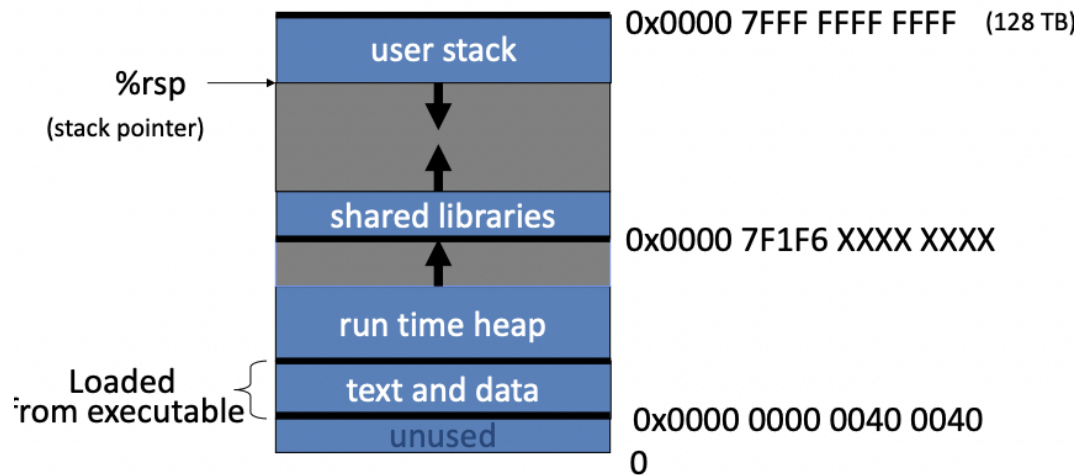
2. First example: buffer overflow

- a. Extremely common bug in C/C++ programs
- b. First major exploit: 1988 Internet Worm, Fingerd

3. What is needed

- a. Understanding C functions, the stack, and the heap
- b. Know how system calls are made
- c. The `exec()` system call
- d. Attacker needs to know which CPU and OS used on the target machine:
 - Our examples are for x86 running Linux or Windows
 - Details vary slightly between CPUs and OSs

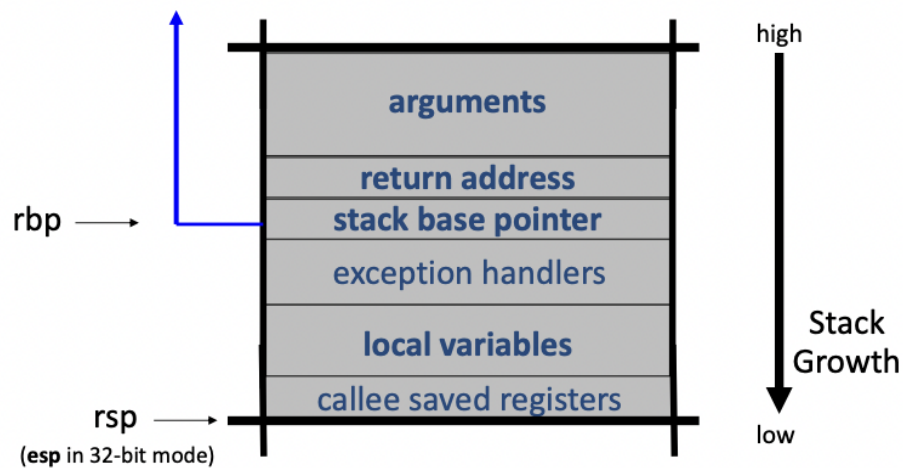
4. Linux process memory layout



a.

Don't Break

5. Stack Frame



a.

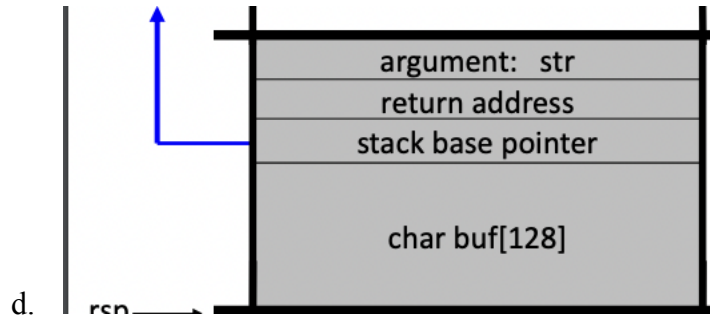
6. What are buffer overflows?

a. Suppose a web server contains a function:

```
void func(char *str) {
    char buf[128];
    strcpy(buf, str);
    do-something(buf);
}
```

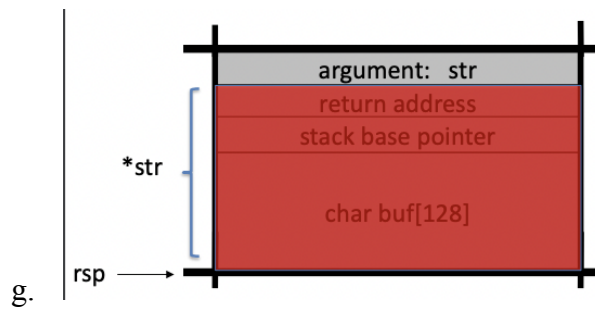
b.

c. After `func()` is called stack looks like:



e. What if *url is 144 bytes long?

f. After strcpy:

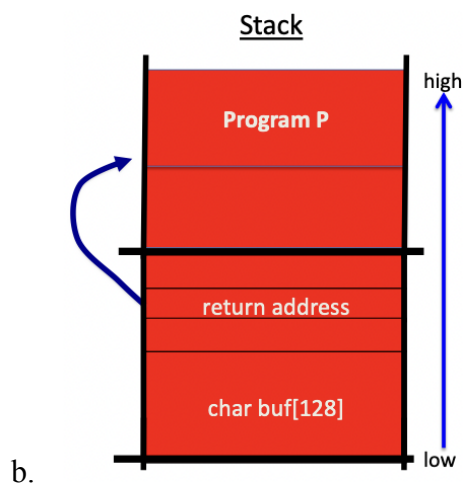


h. Poisoned return address!

i. Problem: no bounds checking in strcpy()

7. Basic stack exploit

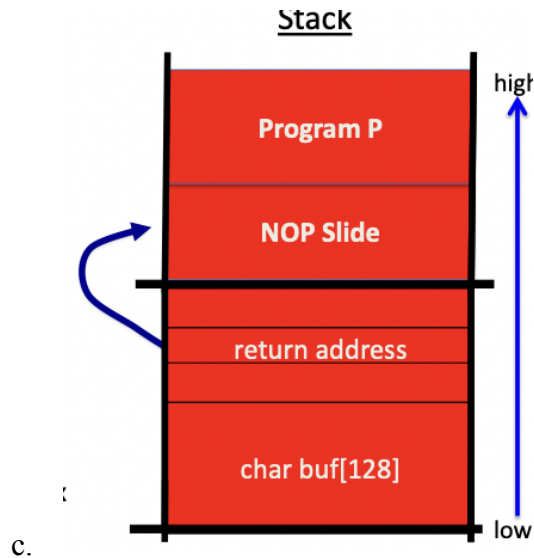
a. Suppose *url is such that after strcpy stack looks like:



c. Program P: `exec("/bin/sh")`

8. The NOP slide

- a. Problem: how does attacker determine ret-address?
- b. Solution: NOP slide
 - Guess approximate stack state when func() is called
 - Insert many NOPs before program P: nop, xor eax, eax, inc ax



9. Details and examples

- a. Some complications:
 - Program P should not contain the '\0' character
 - Overflow should not crash program before func() exits
- b. Famous remote stack smashing overflows:
 - Overflow in Windows animated cursors (ANI)
 - Buffer overflow in Symantec virus detection

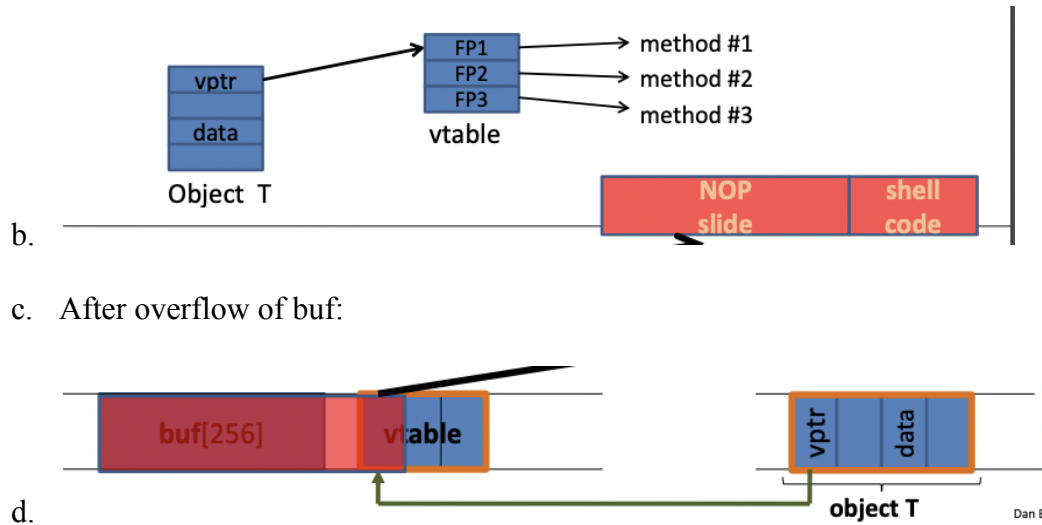
10. Many unsafe libc functions

- a. strcpy (char *dest, const char *src)
- b. strcat (char * dest, const char *src)

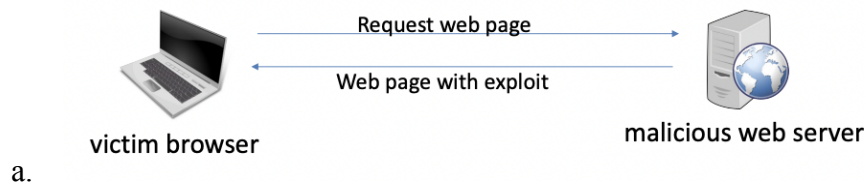
- c. Gets (char *s)
- d. Scanf (const char *format,...) and many more
- e. “Safe” libc versions strncpy(), strncat() are misleading

11. Heap exploits: corrupting virtual tables

- a. Compiler generated function pointers (e.g. C++ code)



12. Exploiting the browser heap



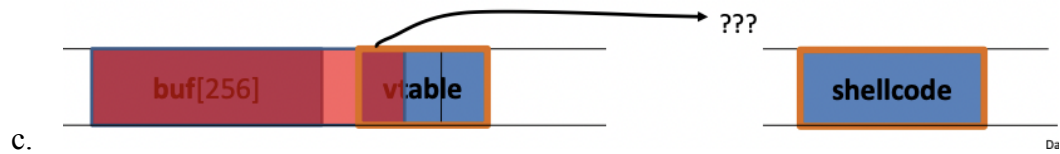
- b. Attacker’s goal is to infect browsers visiting the web site
- c. How: send javascript to browser that exploits a heap overflow

13. A reliable exploit?

- a.

```
<SCRIPT language="text/javascript">
shellcode = unescape("%u4343%u4343%..."); // allocate in heap
overflow-string = unescape("%u2332%u4276%...");
cause-overflow(overflow-string); // overflow buf[ ]
</SCRIPT>
```

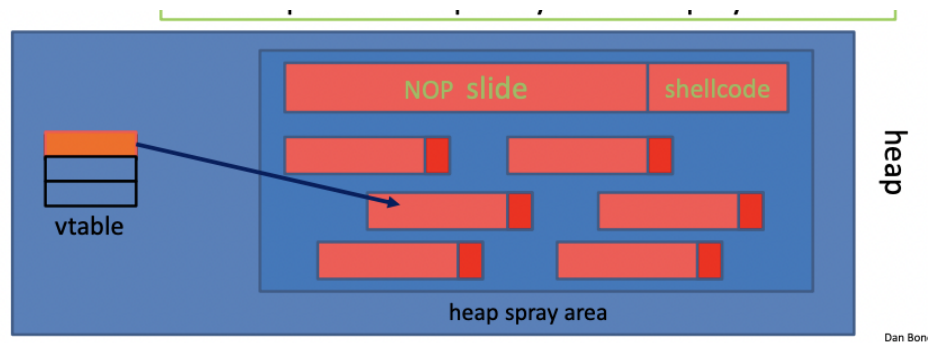
- b. Problem: attacker does not know where browser places shellcode on the heap



14. Heap spraying

- a. Idea:

- Use javascript to spray heap with shellcode (and NOP slides)
- Then point vtable ptr anywhere in spray area



15. Javascript Heap spraying

```
var nop = unescape("%u9090%u9090")
while (nop.length < 0x100000) nop += nop;

var shellcode = unescape("%u4343%u4343%...");

var x = new Array ()
for (i=0; i<1000; i++) {
    x[i] = nop + shellcode;
}
```

- a.

- b. Pointing function-ptr almost anywhere in heap will cause shellcode to execute

16. More Hijacking opportunities

- Integer overflows
- Double free

- c. Use after free
- d. Format string vulnerabilities

17. Integer Overflows

- a. Problem: what happens when int exceeds max value?
- b. $\text{int } m \rightarrow 32 \text{ bits}$, $\text{short } s \rightarrow 16 \text{ bits}$, $\text{char } c \rightarrow 8 \text{ bits}$

$c = 0x80 + 0x80 = 128 + 128$	$\Rightarrow c = 0$
$s = 0xff80 + 0x80$	$\Rightarrow s = 0$
$m = 0xffffffff80 + 0x80$	$\Rightarrow m = 0$

- c.
- d. Example:

```
void func( char *buf1, *buf2, unsigned int len1, len2) {
    char temp[256];
    if (len1 + len2 > 256) {return -1}           // length check
    memcpy(temp, buf1, len1);                   // cat buffers
    memcpy(temp+len1, buf2, len2);
    do-something(temp);                         // do stuff
}
```

- e.

What if $\text{len1} = 0x80$, $\text{len2} = 0xffffffff80$?
 $\Rightarrow \text{len1} + \text{len2} = 0$

Second `memcpy()` will overflow heap !!

- f.

- g. Example: better length check

```
void func( char *buf1, *buf2, unsigned int len1, len2) {
    char temp[256];
    // length check
    if (len1 > 256) || (len2 > 256) || (len1 + len2 > 256)
        return -1;
    memcpy(temp, buf1, len1);                   // cat buffers
    memcpy(temp+len1, buf2, len2);
    do-something(temp);                         // do stuff
}
```

- h.

18. Format string bugs

- a. Format string problem

```
int func(char *user) {  
    fprintf(stderr, user);  
}
```

- b.

- c. Problem: what if *user = “%s%s%s%s%s%s%s%s”??

- Most likely program will crash: DoS
- If not, program will print memory contents.
- Full exploit using user = “%n”

- d. Correct form: fprintf(stdout, “%s”, user);

19. Vulnerable functions

- a. Any function using a format string

- b. Printing:

- Printf, fprintf, sprintf
- vprintf, vfprintf, vsprintf

- c. logging:

- Syslog, err, warn

20. Exploit

- a. Dumping arbitrary memory:

- b. Writing to arbitrary memory

21. Use after free exploits

- a. 1E11 Example: CVE-2014-0282

(IE11 written in C++)

```

<form id="form">
  <textarea id="c1" name="a1" ></textarea>
  <input id="c2" type="text" name="a2" value="val">
</form>

<script>
  function changer() {
    document.getElementById("form").innerHTML = "";
    CollectGarbage();    // erase c1 and c2 fields
  }

  document.getElementById("c1").onpropertychange = changer;
  document.getElementById("form").reset();
</script>

```

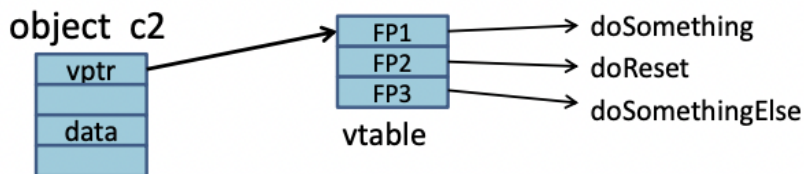
Loop on form elements:
c1.DoReset()
c2.DoReset()

b.

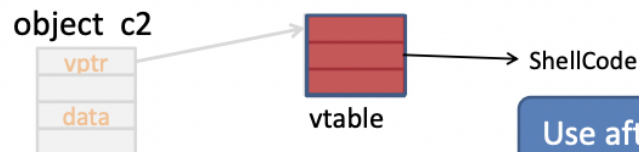
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22. What just happened?

- a. c1.doReset() causes changer() to be called and free object c2



b.



Use after free !

Suppose attacker allocates a string of same size as vtable

When c2.DoReset() is called, attacker gets shell

c.