Spark (cont.)

1. Overview

- a. RDD → represents dataset (collection of records that is physically distributed in class of machines)
- b. $RDDs \rightarrow support operations$
 - i. Passed into data operator (as inputs and outputs)
 - ii. RDDs are mutable data structures → generates fault tolerance
- c. Spark & MapReduce
 - i. Spark supports multiple inputs
 - ii. Spark gives more control on intermediate data (specify as the user to keep the data on disk or memory unlike MapReduce), MapReduce stores data in local disk on mapper
 - iii. Programming model is different
 - 1. MapReduce → define map and reduce functions (singular) and may have to implement many different versions
 - 2. Spark \rightarrow more intuitive



d.

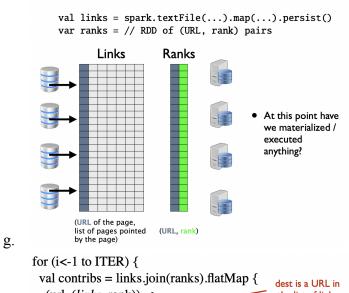
- e. All the nodes are data file operators
- f. Data file operator \rightarrow each instance applies the logic to the joint partitions
- g. Task parallelism \rightarrow applying different task on the same data
 - i. ex) Multi Tread \rightarrow different tasks on the same data (concurrent task)
- h. Data parallelism \rightarrow same function, different data
- i. In reality, system breaks down into series of tasks that look like Map and reduce functions
- j. Workers always communicate with the driver/master (so that it can give the output of one worker to another worker as input)
- k. Transformations (map, join, filter, reduce by key) → take as inputs as RDD and produce RDDs and pass down the output to the next operator in the graph
- 1. Pipeline allows data to flow and make data progress in parallel
 - i. Makes better use of resources
- m. Do we have pipeline in map and reduce?
 - i. No, we wait for the mappers to finish and start the reduce phase

n. Transactions/actions → transformations pipeline the output to the next operator and actions output/materialize

- 0.
- i. Point \rightarrow RDD
- ii. $map \rightarrow transformation$
- iii. Gradient \rightarrow RDD
- iv. For loop \rightarrow Anonymous function that defines logic of the map
- v. Reduce → action that materializes gradient (because we need to subtract gradient)
- 2. PageRank

a.

- b. Links, ranks \rightarrow RDDs
- c. links.join(ranks) → graph with two inputs (links and ranks) that go as input in join(transformation) that generates a new RDD that goes input into flatMap, which also produces an RDD
- d. flatMap will apply the logic inside and will generate the records
- e. Input format of flatMap \rightarrow (url, and tuple of (links, rank)), and therefore, the links.join(ranks) should return such format
- f. For each record in the collection, apply logic (desk \rightarrow (dest, rank/links.size))



for (i<-1 to ITER) {
 val contribs = links.join(ranks).flatMap {
 (url, (links, rank)) => the list of links
 links.map(dest => (dest, rank/links.size))
 }
 Links Ranks
 normalized contribution

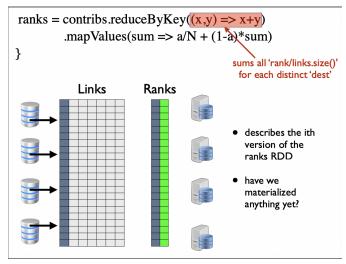
• what does this say?

• have we materialized anything yet?

h.

j.

i. Join will join the two tables on the common attribute (URL) and will produce a RDD

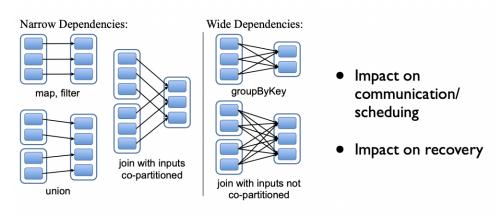


k. Output of the flatmap is inserted as input to reduceByKey (transformation)

3. What an RDD is composed of

Operation	Meaning
partitions()	Return a list of Partition objects
preferredLocations(p)	List nodes where partition <i>p</i> can be accessed faster due to data locality
dependencies()	Return a list of dependencies
iterator(p, parentIters)	Compute the elements of partition <i>p</i> given iterators for its parent partitions
partitioner()	Return metadata specifying whether the RDD is hash/range partitioned

- a.
- b. RDD dependencies → which RDD depends on which RDD
- c. This means that if something fails, we know the RDDs that depend on, so we can handle fault tolerance
- d. Dependencies store parent RDDs, and data in case of failure in order to reconstruct the partition



- e.
- f. Dependencies look like the figure above
- g. Wide dependencies in mapReduce
- 4. Why does an RDD carry metadata on its partitioning?
 - a. Each RDD also has some information about the way it is partitioned
 - b. Transformations that depend on multiple RDDs know whether they need to shuffle data (wide dependency) or not (narrow) and allows users control locality and reduce shuffles
- 5. Handling Faults
 - a. No replication by default even when persisted (because it is expensive)
 - i. Ram or disk node failure could mean permanent loss of data
 - b. Assume heartbeats used to detect lost worker

- i. If worker lost need a way to recompute it's partitions will need all dependent partitions
- ii. Recompute if they can't be found in RAM or on disk
- iii. If wide dependency will need all partitions of dependent RDD if narrow then only one partition
- c. So two mechanisms enable recovery from faults: lineage and policy of what partitions to persist (driver and actual replication and partition strategy)