ECE 350 Instruction Set Architecture										
Instruction	Opcode (ALU Op)	Туре	Usage	Operation						
add	00000 (00000)	R	add \$rd, \$rs, \$rt	\$rd = \$rs + \$rt						
addi	00101	I	addi \$rd, \$rs, N	\$rd = \$rs + N						
sub	00000 (00001)	R	sub\$rd, \$rs, \$rt	\$rd = \$rs - \$rt						
and	00000 (00010)	R	and \$rd, \$rs, \$rt	\$rd = \$rs AND \$rt						
or	00000 (00011)	R	or \$rd, \$rs, \$rt	\$rd = \$rs OR \$rt						
sll	00000 (00100)	R	sll \$rd, \$rs, shamt	\$rd = \$rs << shamt						
sra	00000 (00101)	R	sra \$rd, \$rs, shamt	\$rd = \$rs / 2^shamt						
mul	00000 (00110)	R	mul \$rd, \$rs, \$rt	\$rd = \$rs * \$rt (32b X 32b); STATUS=1 if overflow						
div	00000 (00111)	R	div \$rd, \$rs, \$rt	\$rd = \$rs / \$rt (32b ÷ 32b); STATUS=1 if div0						
CUSTOM R (in ALU)	00000 (01000) 00000 (11111)	R	CUSTOM_R# \$rd, \$rs, \$rt	\$rd = CUSTOM#(\$rs, \$rt)						
a51start (multicyc xor loadkeystream	le) 11111 00000 (01000 00000 (01001)	,	n a51start xor \$rd, \$rs, \$rt stom loadkeysteam \$rd	, load #						
j	00001	JI	j N	PC = N						
bne	00010	I	bne \$rd, \$rs, N	if(\$rd != \$rs) PC = PC+1+N						
jal	00011	JI	jal N	\$r31 = PC+1; PC=N						
jr	00100	JII	jr \$rd	PC = \$rd						
blt	00110	I	blt \$rd, \$rs, N	if(\$rd < \$rs) PC=PC+1+N						
bex	10110	JI	bex N	if(STATUS > 0) PC=PC+1+N						
setx	10101	JI	setx N	STATUS = N						
SW	00111	I	sw \$rd, N(\$rs)	MEM[\$rs + N] = \$rd						

Instruction Type	Instruction Format								
R	Opcode [31:27]	RD [26:22]	RS [21:17]	RT [16:12]	shiftamt [11:7]	ALUop [6:2]	Zeros [1:0]		
1	Opcode [31:27]	RD [26:22]		RS [21:1	7]	Immediate [16:0]			
<u>'</u>									
	Opcode [31:27]	Target [26:0]							
J (I/II)	Opcode [31:27]	RD [26:22]							

lw \$rd, N(\$rs)

rd = MEM[rs + N]

I-type immediate field [16:0] is signed 2's complement and sign-extended to the full 32-bit word size.

J-type target field [26:0] is extended to the full 32-bit PC size using the upper bits from the current PC+1.

Register fields that are undefined are filled with zeroes by the assembler.

01000

lw

Register \$r0 always equals zero. Registers \$r1 through \$r30 are general purpose. Register \$r31 stores the link address of a jump-and-link instruction.

Instructions that change control flow (beq, blt, j, jal, jr) do not have a delay slot.

Memory is word-addressed. The instruction and data memory address spaces are separate. Static data begins at data memory address zero. Stack data begins at the end of the data memory and grows downwards. There is no preset boundary between the end of static data and the start of the upwardsgrowing heap; this is a property of the assembly program.

After a reset, all register values are zero and program execution begins from instruction memory address zero. The memory's contents are not reset.

Useful hints for the Assembulator:

This assembly fragment (paste it into, and save- always save the file before assembling- the main window):

.text
main: lw \$r3, wow(\$r0)
lw \$r4, wow(\$r0)
mul \$r5, \$r3, \$r4
bex dead
addi \$r7, \$r0, 0x0FEEDF00
j quit

dead: addi \$r7, \$r0, 0x0DEADF00

quit: halt .data

wow: .word 0x0000B504

mystring: .string ASDASDASDASDASD

var: .char Z label: .char A

heapsize: .word 0x00000000 myheap: .word 0x00000000

A simple program fragment, illustrating various assembly syntax for the Assembulator, which calculates the square of the value stored in location

"wow" and sets the value of R7 to either 0x0FEEDF00 if the product doesn't

overflow, otherwise 0x0DEADF00.

Every program needs a ".text" and ".data" region, and a "main:" label. Program execution starts at "main:" Data variables can be allocated in the .data segment and referenced by their labels. For example, "mystring" is shorthand for the "location in memory of a .data labelled 'mystring'".

To assemble this fragment, enter the code (or open a file), save the code (to a file which INCLUDES a "." in it's name, like "mything.asm"), verify it (button with a check), and then you can simulate it (button with a play symbol). Next, you'll want to output some machine code and data in .mif files. Use the Assemble (or Assemble To) buttons (down arrows) to specify an output name (like "mything_output") which will be used to create two files (dmem.mif and imem.mif). You can inspect these files from within the Assembulator, or add them to your project in Quartus (ultimately).

The assembulator also expands certain "pseudo opcodes" like this:

```
Idia = $rd, Iabel + Idi $rd, N
Idi = $rd, N + addi $rd, $r0, N
ret = ir $ra
```

Comments

 $\label{eq:bgt} \begin{array}{l} \text{bgt} = \text{\$rd}, \, \text{\$rs}, \, \text{N} + \text{blt} \, \text{\$rs}, \, \text{\$rd}, \, \text{N} \\ \text{nop} = \text{add} \, \text{\$r0}, \, \text{\$r0}, \, \text{\$r0} \\ \text{halt} = \text{j PC} \end{array}$

It will also expand certain constants like:

\$zero = \$R0 \$Ra = \$R31

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CD.