Benchmarks

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1 Benchmarks

The benchmarks we used in our evaluation cover a broad range of well-known applications of SAT solvers. Table 1 provides an overview; for each benchmark set, we list the number of total instances ('#inst'), the number of training instances ('#train') we used for automated configuration, the number of testing instances ('#test'), as well as basic statistics for the number of variables ('#var') and clauses ('#clause'), including the average ('avg'), median ('median'), 0.25-and 0.75-quantiles ('q25' and 'q75', respectively). Further information on each set is given in the following. The total number of instances used in this work is 26 932, and to our best knowledge is much more than those used in previous work on SAT.

FCC-SAT + FCC-UNKNOWN Recently, SAT solvers have been prominently used by the US Federal Communication Commission (FCC) for revenue-optimising spectrum repacking in the context of a bandwidth auction that resulted in a 7 billion dollar revenue for the US government [4, 9]. We used the SAT-encoded benchmark instances made available by Newman et al. [9], which contains 10 000 SAT instances, 9482 of which are known to be satisfiable and 121 unsatisfiable, while the satisfiability of the remaining 397 instances is unknown. Since SLS solvers such as PbO-CCSAT are unable to prove unsatisfiability, we discarded all unsatisfiable instances. Furthermore, we selected 100 of the 9482 satisfiable instances as the training set used for automated solver configuration; we refer to this set as FCC-SAT [Train]. The remaining 9382 satisfiable instances make up our test set, FCC-SAT [Test]; we also considered an additional test set, FCC-UNKNOWN, consisting of the 397 instances whose satisfiability was unknown prior to our experiments.

⁴ https://www.cs.ubc.ca/labs/beta/www-projects/SATFC/cacm_cnfs.tar.gz

⁵ Since we need training sets in which most instances can be solved within a given reasonable cutoff time, we selected 80 instances uniformly at random from those instances which could be solved by the default version of *PbO-CCSAT* within a cutoff time of 5000 CPU seconds and 20 instances uniformly at random from those instances which could not be solved by the default version of *PbO-CCSAT* within a cutoff time of 5000 CPU seconds.

SC17-mp1-9

20

Benchmark Reference #inst #train/#test #var #clause median median avg q25avg q25q75 $\frac{q75}{27993.0}$ FCC-SAT 9482100/938227633.3 25712647 061.5 [9] $653\,302.0$ 30 477 762 108 553185FCC-UNKNOWN 397 0/39732 562.9 32.562.0 867 708.0 [9] 867 845.4 30 564 $34\,593$ 763157PTN 23 11/123721.0[8] 14321.514 160.0 3696 3733 13802 $14\,455$ ${ t PTN-More}^b$ 556 0/5567413.77416.0 $17\,800.1$ 17799.0 [8] 7208 7620 17226 $18\,366$ 100/16334 $24\,906.0$ SMT-QF-BV $71\,287.0\ 305\,428$ [12] $16\,434$ 174748.3480 460.9 843 $113\,975$ 17092 200.0 9 086.0 Community 20 5/152200.09086.0[6] 2 200 9 086 22009086

Table 1: Bechmarks for training and testing in our experiments.

1458.0

 $280\,139.1$

280139.5

280 140

[2]

1458.0

5/15

PTN + PTN-More SAT techniques have been recently applied to (and play prominent roles in) tackling a long-standing open problem in mathematics known as Boolean Pythagorean Triples (PTN) [8, 7]. We used 23 publically available SATencoded PTN instances, 2 of which - plain7824.cnf⁶ and bce7824.cnf⁷ - were published by Heule et al. [8], while the 21 others were taken from the crafted benchmark of the 2016 SAT Competition. 8 Of these 23 SAT-encoded instances, we selected 11 uniformly random for use as a training set; we refer to this set as PTN [Train]. The remaining 12 instances make up our test set, PTN [Test]. Moreover, in order further explore the efficiency of PbO-CCSAT on PTN instances, we used the PTN encoder⁹ by Heule et al. [8] to generate 556 additional satisfiable instances, according to the suggestions by the authors, to obtain an additional test set, PTN-More.

SMT-QF-BV SLS solvers for SAT have recently achieved promising results on solving the Satisfiability Modulo Theories (SMT) problem in the theory of quantifierfree bit-vectors (QF-BV) [5]. We first downloaded 16436 SMT-encoded QF-BV instances made available by Niemetz et al. [11, 12]; 10 then, we utilised Boolector 11 [10] and aigtocn^{f12} [3] to translate those SMT-encoded instances into SAT, according to suggestions by Niemetz et al. [11, 12]. Since the translation process of 2 instances (bench_3774.smt2 and bench_3843.smt2) could not be completed within 24 CPU hours on our reference machines, we discarded those 2 instances and used the resulting 16434 SAT-encoded instances. Of these, we selected 100

¹⁴⁵⁸ 1458 $280\,136$ a For PbO-CCSAT and its competitors, the configurations trained on FCC-SAT [Train] are the ones tested on FCC-UNKNOWN.

b For PbO-CCSAT and its competitors, the configurations trained on PTN [Train] are the ones tested on PTN-More.

⁶ http://www.cs.utexas.edu/~marijn/ptn/plain7824.cnf

⁷ http://www.cs.utexas.edu/~marijn/ptn/bce7824.cnf

⁸ https://baldur.iti.kit.edu/sat-competition-2016/downloads/crafted16.zip

⁹ http://www.cs.utexas.edu/~marijn/ptn/ptn-encode.c

¹⁰ http://fmv.jku.at/fmsd16/fmsd16-benchmarks-sat.7z

¹¹ http://fmv.jku.at/boolector/boolector-2.4.1-with-lingeling-bbc.tar.bz2

¹² http://fmv.jku.at/aiger/aiger-1.9.9.tar.gz

instances as training set; we refer to this set as SMT-QF-BV [Train]. The remaining 16 334 instances make up our test set, SMT-QF-BV [Test].

Community This benchmark consists of instances with community structure which are generated by an industrial model called *Community Attachment* [6]. Also, this benchmark is included in the application track of the 2016 SAT Competition [1]. We used all 20 Community instances from the application track of the 2016 SAT Competition. ¹⁴ We selected 5 instances uniformly at random to obtain a training set Community [Train]; the remaining 15 instances form the testing set Community [Test].

SC17-mp1-9 Finally, we considered a set of satisfiable instances from the main track of the 2017 SAT Competition [2]. 15 Starting with the complete set of application instances from the competition, we filtered out all instances known to be unsatisfiable, resulting in a reduced set, SC17. We note that the application benchmarks used in SAT competitions are heterogeneous in that they include various types of structured instances, with most instance families containing relatively few instances. Since automated algorithm configuration is primarily intended for performance optimisation on relatively homogeneous families of problem instances, we selected from SC17 those instance families containing at least 20 instances. This resulted in three instance families, mp1-9, mp1-ps and g2-ak128, from SC17. In preliminary experiments, we found that all solvers we considered (including SLS, CDCL and hybrid solvers) were unable to solve any of the instances in mp1-ps and may well be unsatisfiable, as are all instances from this family that were solved in the competition. Furthermore, in preliminary experiments, 5 of the instances in g2-ak128 could be solved by SLS solvers within the cutoff time of 5000 CPU seconds, but these 5 instances were all trivial for SLS solvers such as PbO-CCSAT if all variables were set to false in the initial assignment. In light of this, since our interest is in solving non-trivial, structured SAT instances, we also eliminated instance family g2-ak128. This left us with one benchmark: SC17-mp1-9, consisting of all 20 instances in mp1-9. We selected 5 instances uniformly at random from each of the set to obtain the training set, SC17-mp1-9 [Train]; the remaining 15 instances form the testing set, SC17-mp1-9 [Test].

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 $^{^{13}}$ We selected 80 instances uniformly at random from those instances that were solved by the default configuration of PbO-CCSAT within a cutoff time of 5000 CPU seconds, and 20 instances uniformly at random from those instances that remained unsolved within this cutoff time.

 $^{^{14}\ \}mathtt{https://baldur.iti.kit.edu/sat-competition-2016/downloads/app16.zip}$

¹⁵ https://baldur.iti.kit.edu/sat-competition-2017/benchmarks/Main.zip

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