Statistical mechanical models for quantum codes subject to correlated noise

Christopher Chubb (USyd), Steve Flammia (USyd/Yale)

(me@)christopherchubb.com

arXiv:1809.10704

AIP 2018-12





Two important questions about quantum codes:

- How do I decode?
- What is the threshold?

¹Dennis, Kitaev, Landahl, Preskill, JMP 2002, doi:10/cs2mtf, arXiv:quant-ph/0110143

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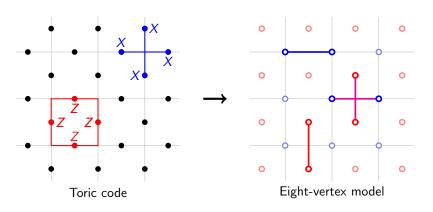
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Statistical mechanical mapping

The idea here it to construct a family of statistical mechanical models, whose thermodynamic properties reflect the error correction properties of the code.



This will allow us to use the analytic and numerical tools developed to study stat mech systems to study quantum codes.

Statistical mechanical mapping

Statistical mechanical mapping



- General stat mech mapping construction for arbitrary codes and correlated noise
- Numerically show that mild correlations can lower the threshold of the toric code considerably
- Extend to spatio-temporal correlations, such as circuit noise
- Give efficient TN (approximate) maximum likelhood decoders, generalising BSV

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Structure of the (independent) mapping

The stat mech mapping gives a bond-disordered Ising/Potts-type model with:

For a stabiliser code generated by $\{S_k\}_k$,

$$H_E(\vec{s}) := -\sum_i \sum_{\sigma \in \mathcal{P}_i} \underbrace{J_i(\sigma)}_{Coupling \ Disorder} \underbrace{\prod_{k: [\![\sigma,S_k]\!]=-1}^{DoF} s_k}_{k: [\![\sigma,S_k]\!]=-1},$$

where $s_k=\pm 1$, and $\llbracket A,B'
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 $\begin{array}{cccc} \mathsf{States} & \longrightarrow & \mathsf{Stabilisers} \\ \mathsf{Disorder} & \longrightarrow & \mathsf{Error} \\ \mathsf{Couplings} & \longrightarrow & \mathsf{Single\text{-}site} \; \mathsf{Paulis} \\ \mathsf{Coupling} \; \mathsf{strengths} & \longrightarrow & \mathsf{Fourier} \; \mathsf{transform} \; \mathsf{of} \\ & & \mathsf{error} \; \mathsf{model} \end{array}$

For a stabiliser code generated by $\{S_k\}_k$,

$$H_{E}(\vec{s}) := -\sum_{i} \sum_{\sigma \in \mathcal{P}_{i}} \underbrace{J_{i}(\sigma)}_{\text{Coupling Disorder}} \underbrace{\prod_{k: \llbracket \sigma, S_{k} \rrbracket = -1}^{\text{DoF}} s_{k}}_{\text{K}: \llbracket \sigma, S_{k} \rrbracket = -1}$$

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Independent case: Properties

The stabiliser group is encoded as a gauge symmetry of the form

$$s_k \rightarrow -s_k$$
 and $E \rightarrow ES_k$

which implies that $Z_E = Z_{ES_k}$.

If the coupling strengths satisfy the Nishimori condition, then

$$e^{-\beta H_E(\vec{0})} = \Pr(E)$$
 \Longrightarrow $Z_E = \Pr(\overline{E}) := \sum_S \Pr(ES).$

Nishimori condition:
$$\beta J_i(\sigma) = rac{1}{4} \sum_{ au \in \mathcal{P}} \log p_i(au) \left[\!\left[\sigma, au\right]\!\right].$$

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Step 0: Code and noise model

Toric code with iid bit-flips

$$s_{\!\scriptscriptstyle V}=\pm 1$$
 on each vertex v

Step 2: Interactions

$$H_I = -\sum_{v \sim v'} J \, s_v s_{v'}$$

$$H_E = -\sum_{v \sim v'} Je_{vv'} \, s_v s_{v'}$$

$$\text{where } e_{vv'} = \begin{cases} +1 & E_{vv'} = I, \\ -1 & E_{vv'} = X. \end{cases}$$



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Step 1: Degrees of freedom

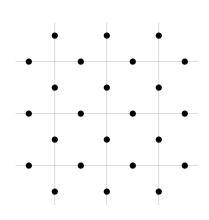
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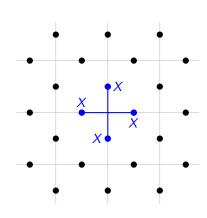
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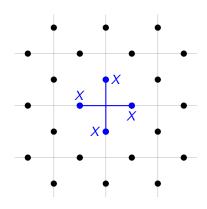
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$$\Pr(X_e) = p, \quad \Pr(I_e) = 1 - p.$$



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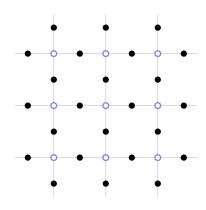
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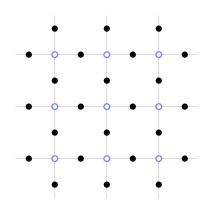
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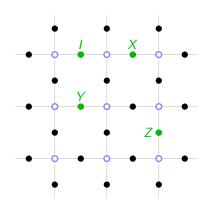
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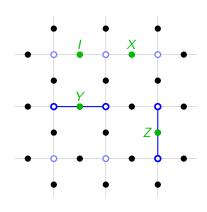
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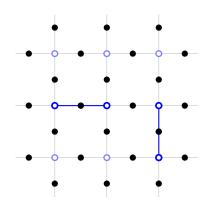
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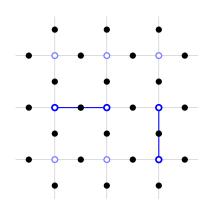
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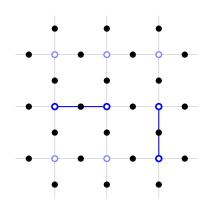
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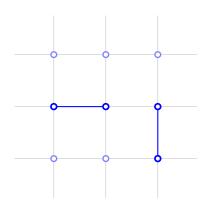
$$H_I = -\sum_{v \sim v'} J s_v s_{v'}$$

Step 3: Disorder

$$H_E = -\sum_{v \sim v'} Je_{vv'} \, s_v s_{v'}$$

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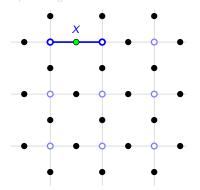
$$Pr(+J) = p$$
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 $\pm J$ Random-bond Ising Model

Toric code

Bit-flip \rightarrow Random-bond Ising¹ Indep. $X\&Z \rightarrow 2\times$ Random-bond Ising Depolarising \rightarrow Random 8-vertex model



Colour code

Bit-flip \rightarrow Random 3-spin Ising Indep. $X\&Z \rightarrow 2 \times \text{Random 3-spin Ising}$ Depolarising \rightarrow Random interacting 8-vertex²

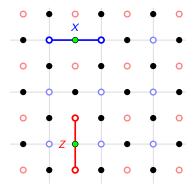


Dennis et.al., JMP 2002, doi:10/cs2mtf, arXiv:quant-ph/0110143

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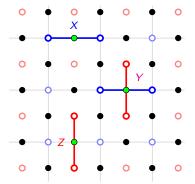


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 $\begin{array}{l} \text{Bit-flip} \rightarrow \text{Random-bond Ising}^1 \\ \text{Indep. } X\&Z \rightarrow 2\times \text{Random-bond Ising} \\ \text{Depolarising} \rightarrow \text{Random 8-vertex model}^2 \end{array}$



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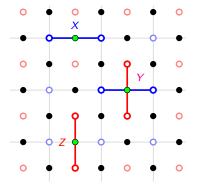


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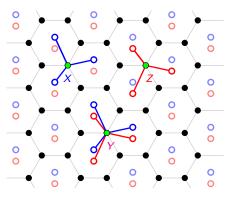
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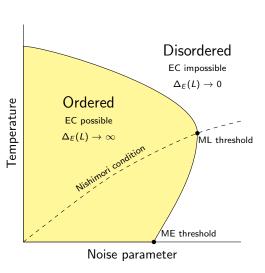
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Phase diagram sketch



Consider the free energy cost of a logical error L (generalised domain wall),

$$\Delta_E(L) := F_{EL} - F_E$$

$$= \frac{\log Z_E - \log Z_{EL}}{\beta}.$$

Along the Nishimori line we have

$$\beta \Delta_E(L) = \log \frac{\Pr(\overline{E})}{\Pr(\overline{EL})},$$

which means

Below th.: $\Delta_{\it E}(\it L)
ightarrow \infty$ (in mean)

Above th. : $\Delta_E(L) \rightarrow 0$ (in prob.)

The key point independence gave us was the ability to factor our noise model

$$\Pr(E) = \prod_i p_i(E_i).$$

We can also consider the following correlated models:

Factored distribution

An error model factors over regions $\{R_j\}_j$ if there exist $\phi_j:\mathcal{P}_{R_j}\to\mathbb{R}$ such that

$$\Pr(E) = \prod_{j} \phi_{j} \left(E_{R_{j}} \right)$$

This includes many probabilistic graphical models, such as Bayesian Networks and Markov/Gibbs Random Fields.

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By construction, we can extend to the correlated case by changing $\sigma \in \mathcal{P}_i$ to $\sigma \in \mathcal{P}_{R_i}$:

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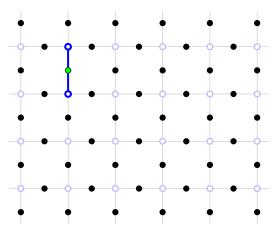
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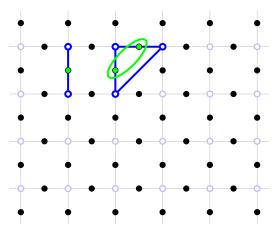




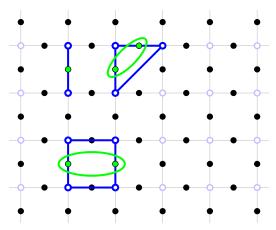
Toric code with correlated bit-flipsCorrelations induce longer-range interactions



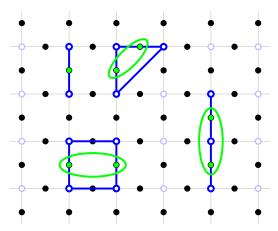
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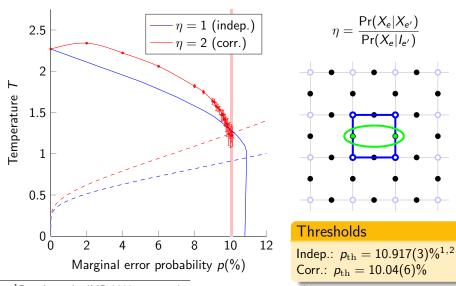
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Monte Carlo simulations



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²Toldin et.al., JSP 2009, doi:10/c3r2kc, arXiv:0811.2101

Other results

TN approximations of partition function give efficient approximate ML decoders (generalises BSV).

Other codes:

- Qudit subsystem codes
- Abelian quantum doubles (finite and infinite)
- General quantum doubles with pure-magnetic or pure-electric noise

Other noise models:

- Faulty measurements
- Leakage errors
- Spatio-temporal correlations, e.g. circuit noise

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Conclusions and further work

- Extended the stat. mech. mapping to correlated models
- Can apply stat. mech. mapping to circuit noise via the history code
- Numerically evaluated the threshold of correlated bit-flips in the toric code
- Stat. mech. mapping gives tensor network maximum likelihood decoders

- Experimentally relevant correlated models?
- Non-stochastic noise? Non-Pauli noise? Coherent noise?
- Can we use the decoders to understand to better understand the connection between correlation and the threshold (ongoing work with David Tuckett and Benjamin Brown).

Thank you!

ArXiv:1809.10704

(me@)christopherchubb.com