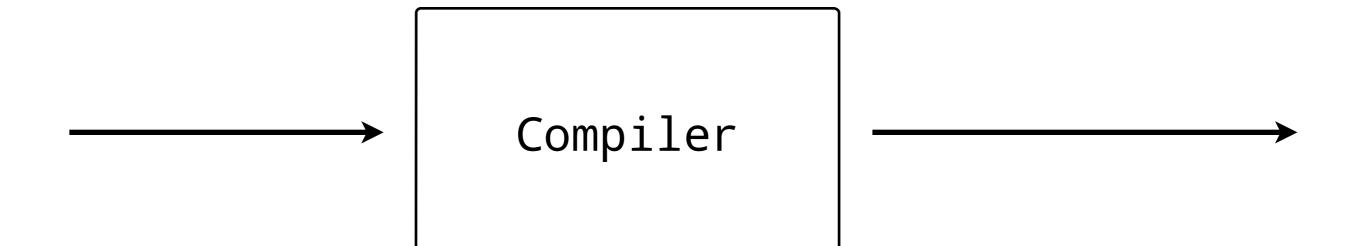
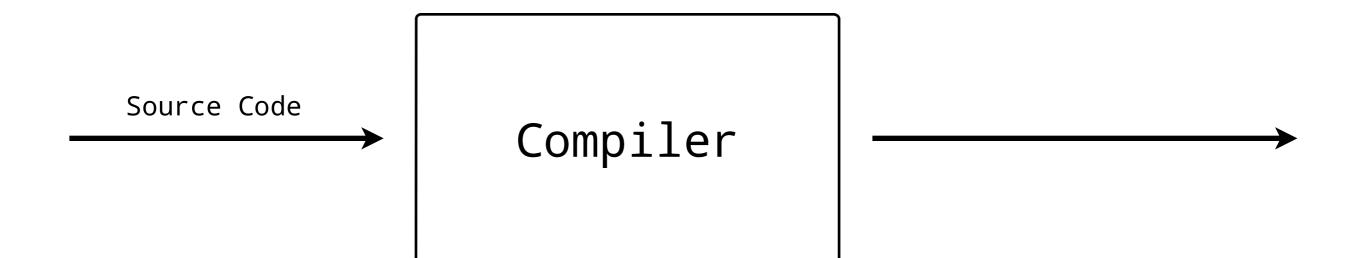
Let's Write a Type Checker

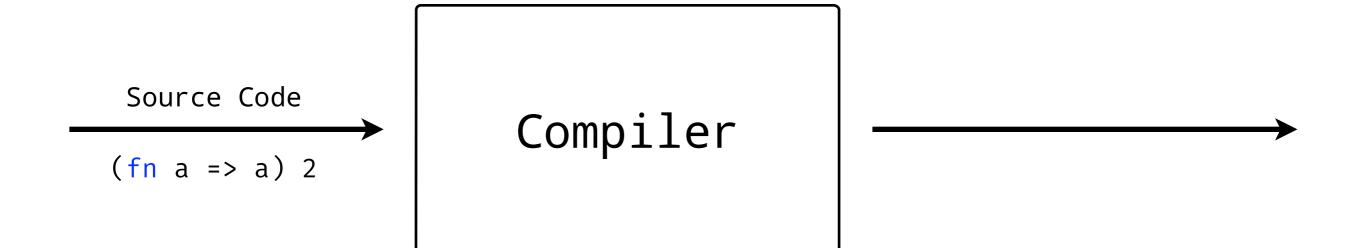
Ionuţ G. Stan — I T.A.K.E. — May 2015

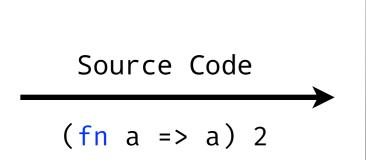
The Plan

- Part 1
 - Compilers Overview
 - Type Checking vs Type Inference
 - Vehicle Language
 - Wand's Type Inference Algorithm
- Part 2
 - Live Demonstration



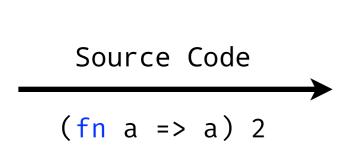






Compiler

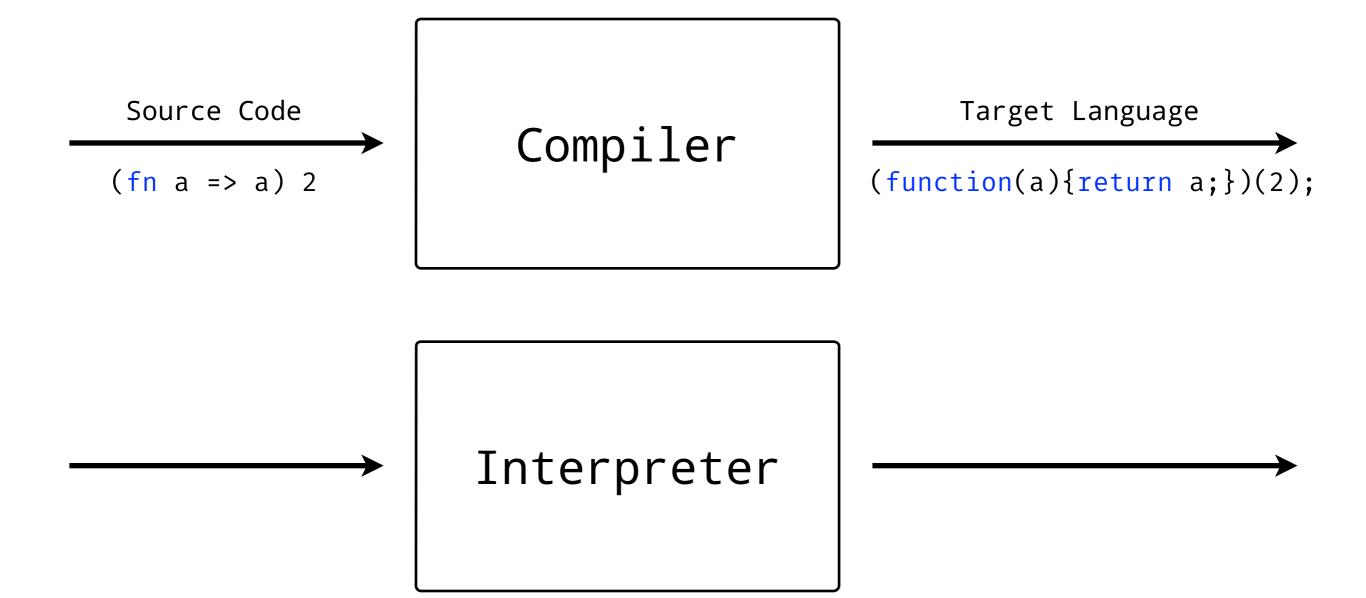
Target Language

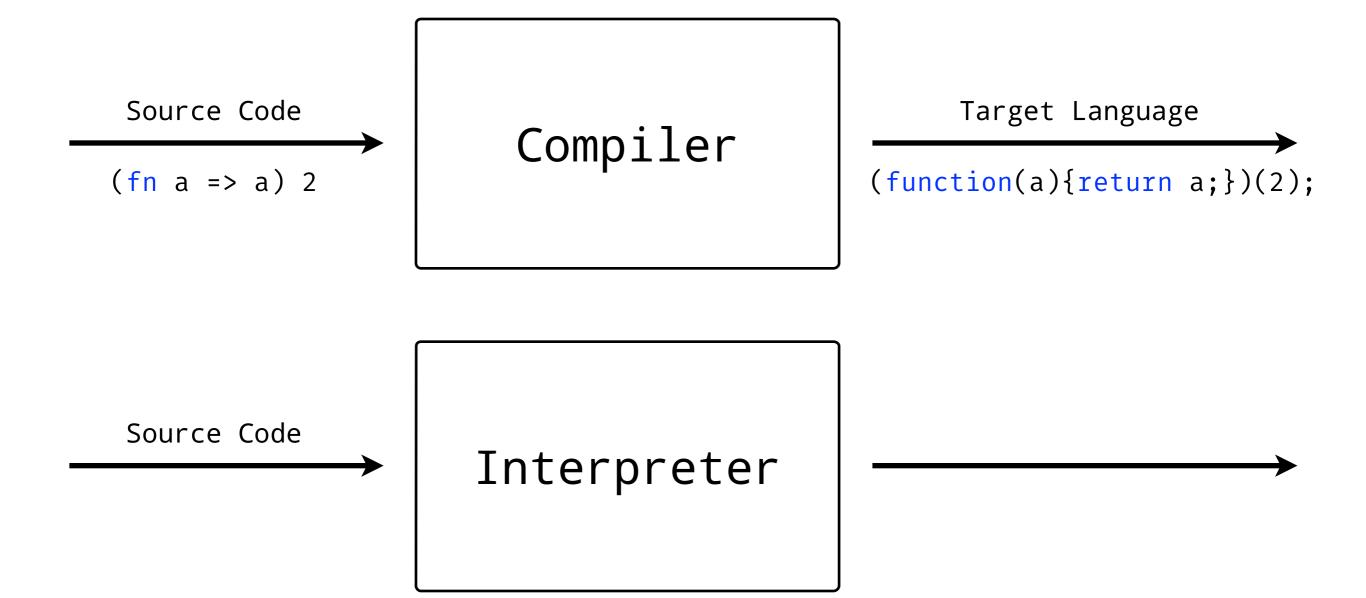


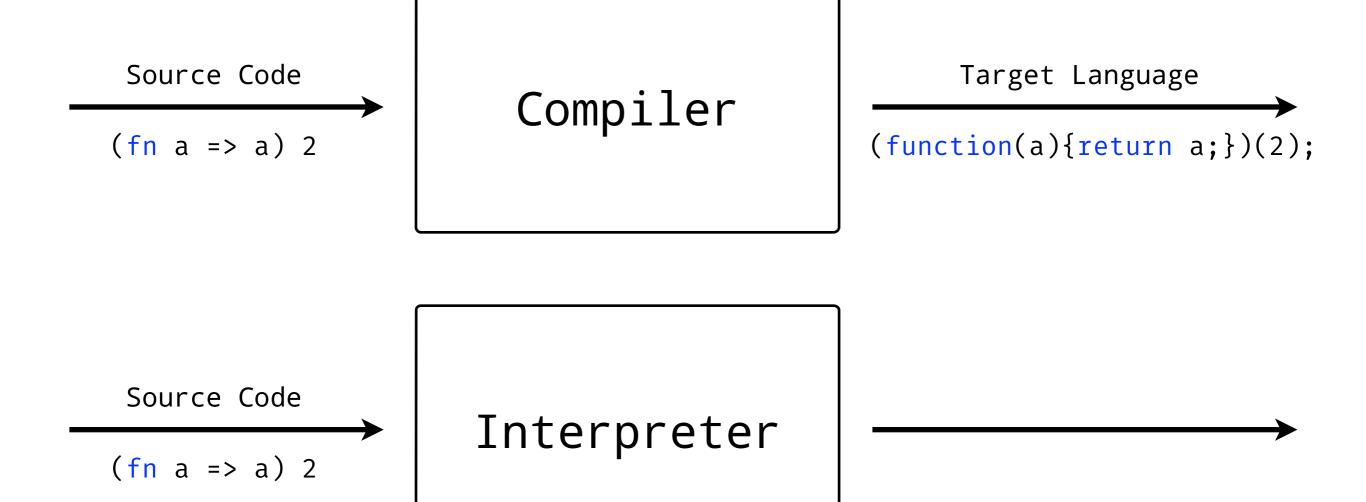
Compiler

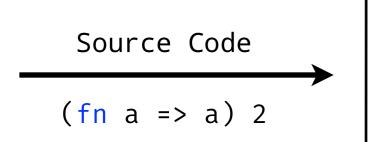
```
Target Language

(function(a){return a;})(2);
```









Compiler

Target Language

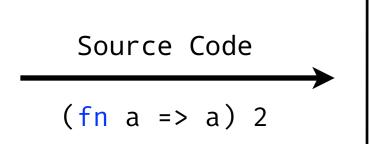
(function(a){return a;})(2);

```
Source Code

(fn a => a) 2
```

Interpreter

Evaluation Result



Compiler

Target Language

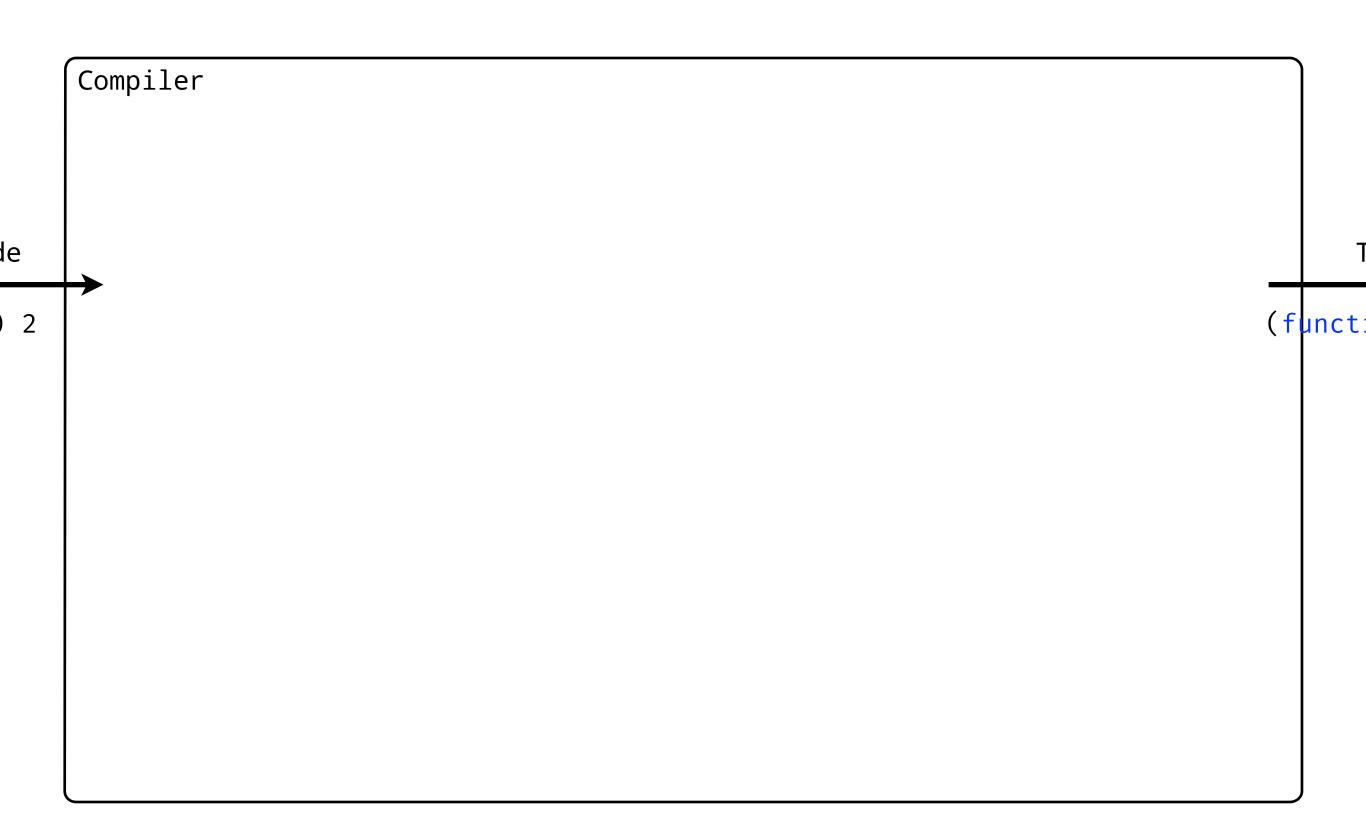
(function(a){return a;})(2);

```
Source Code

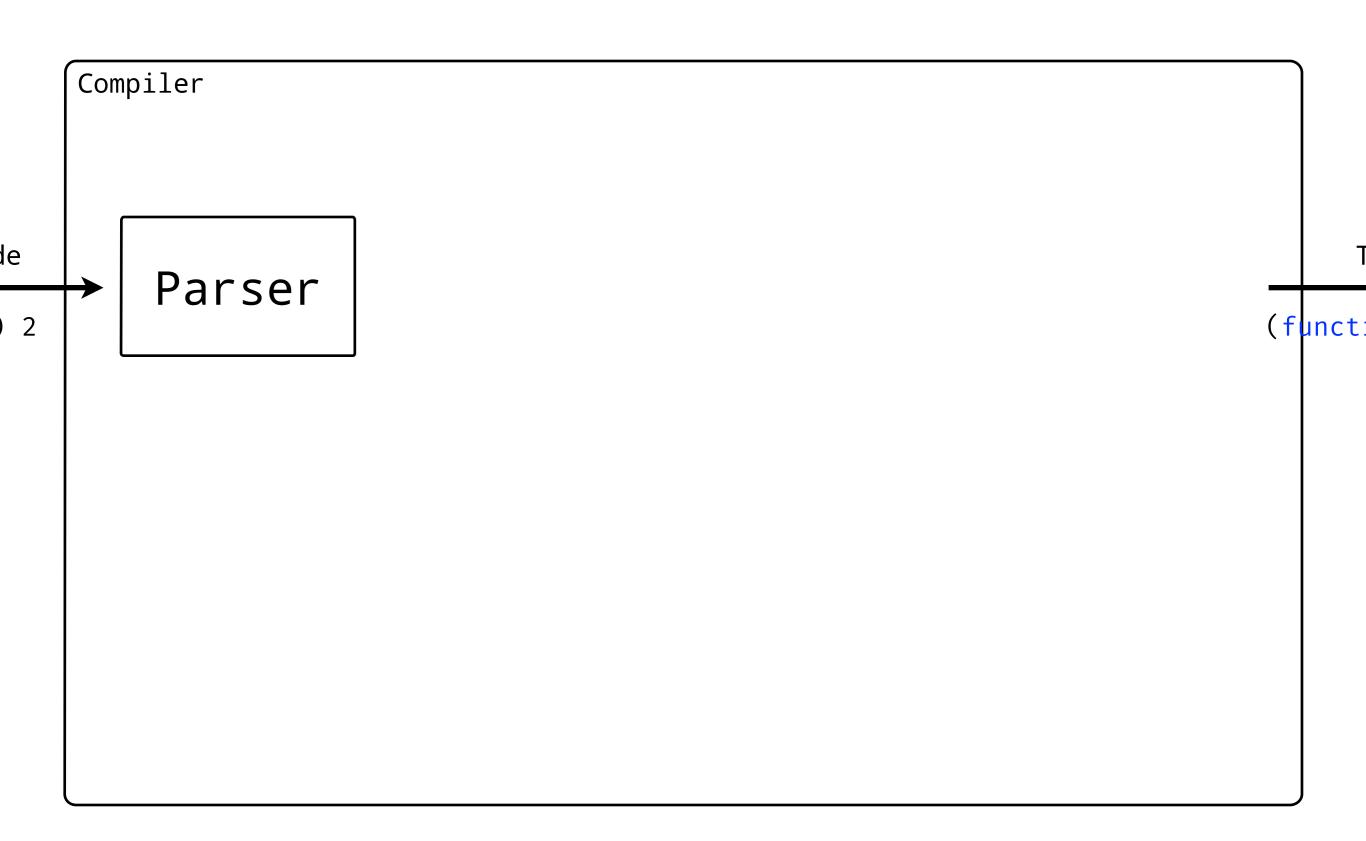
(fn a => a) 2
```

Interpreter

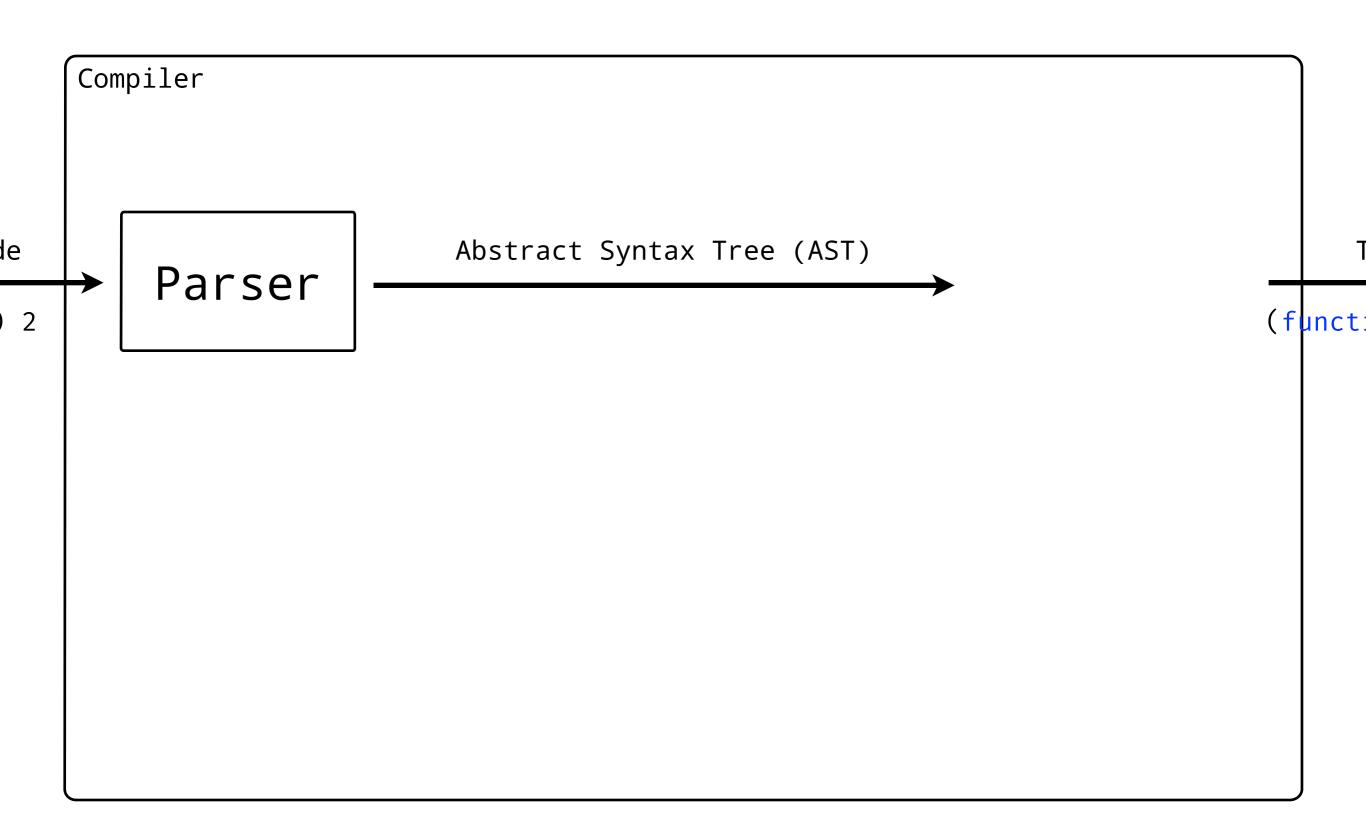
Evaluation Result



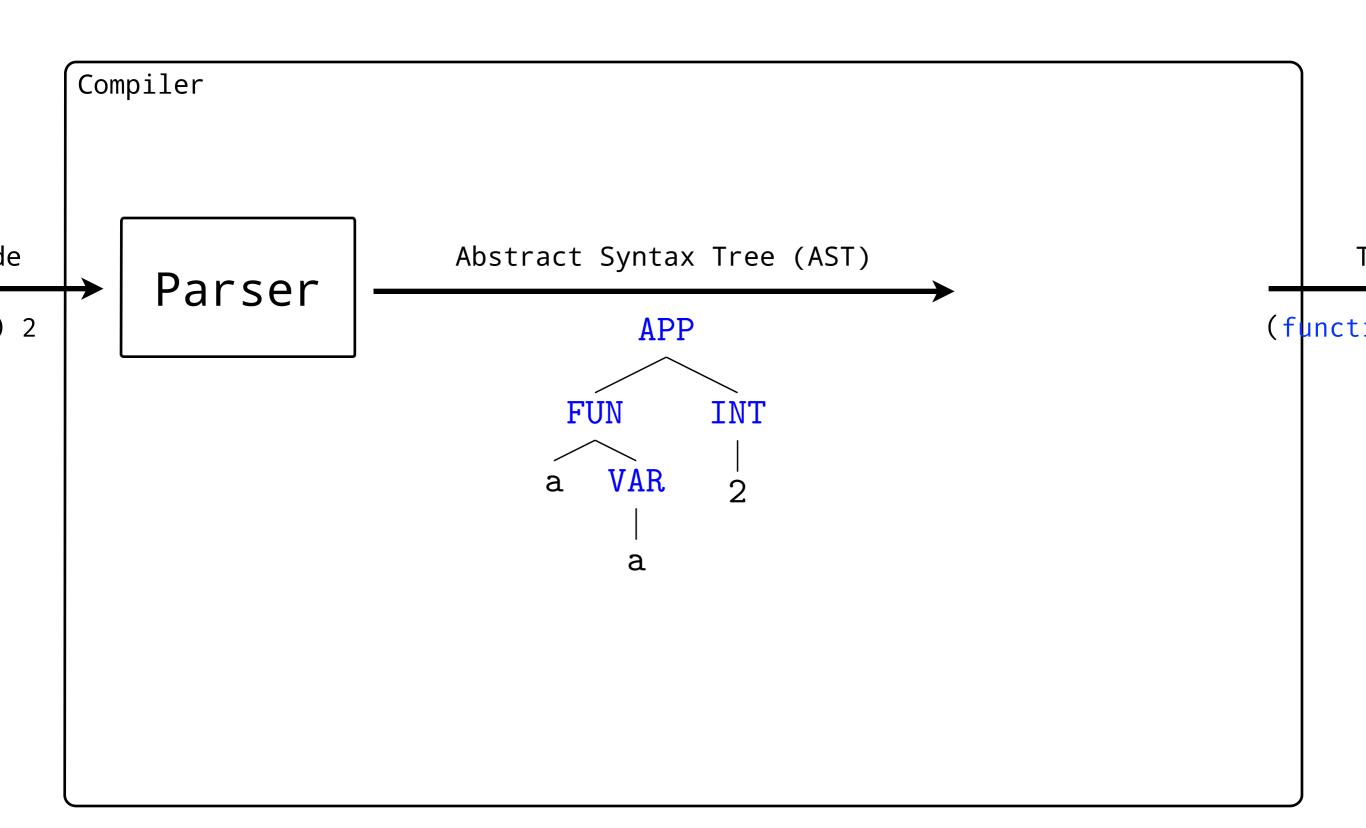
Parsing



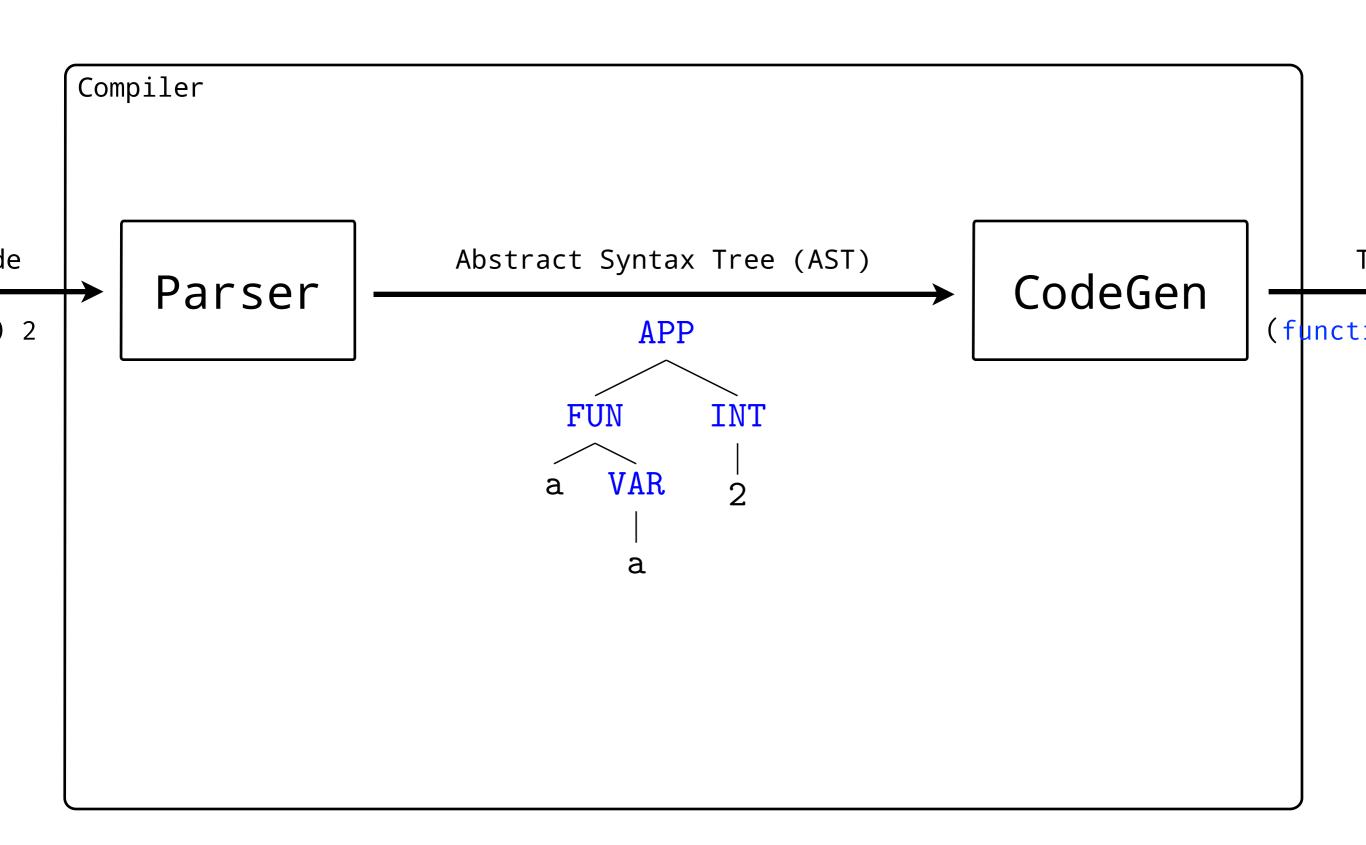
Abstract Syntax Tree



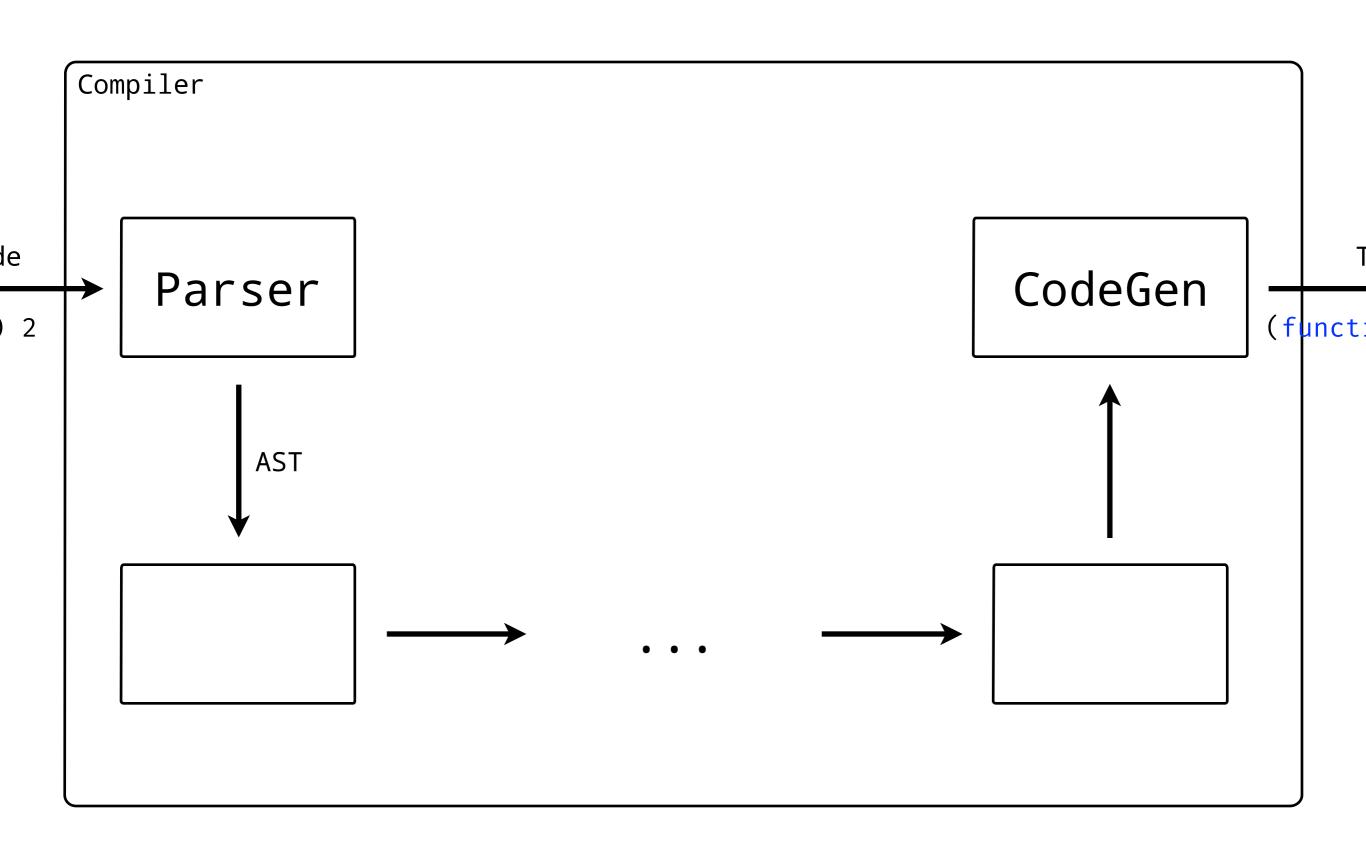
Abstract Syntax Tree



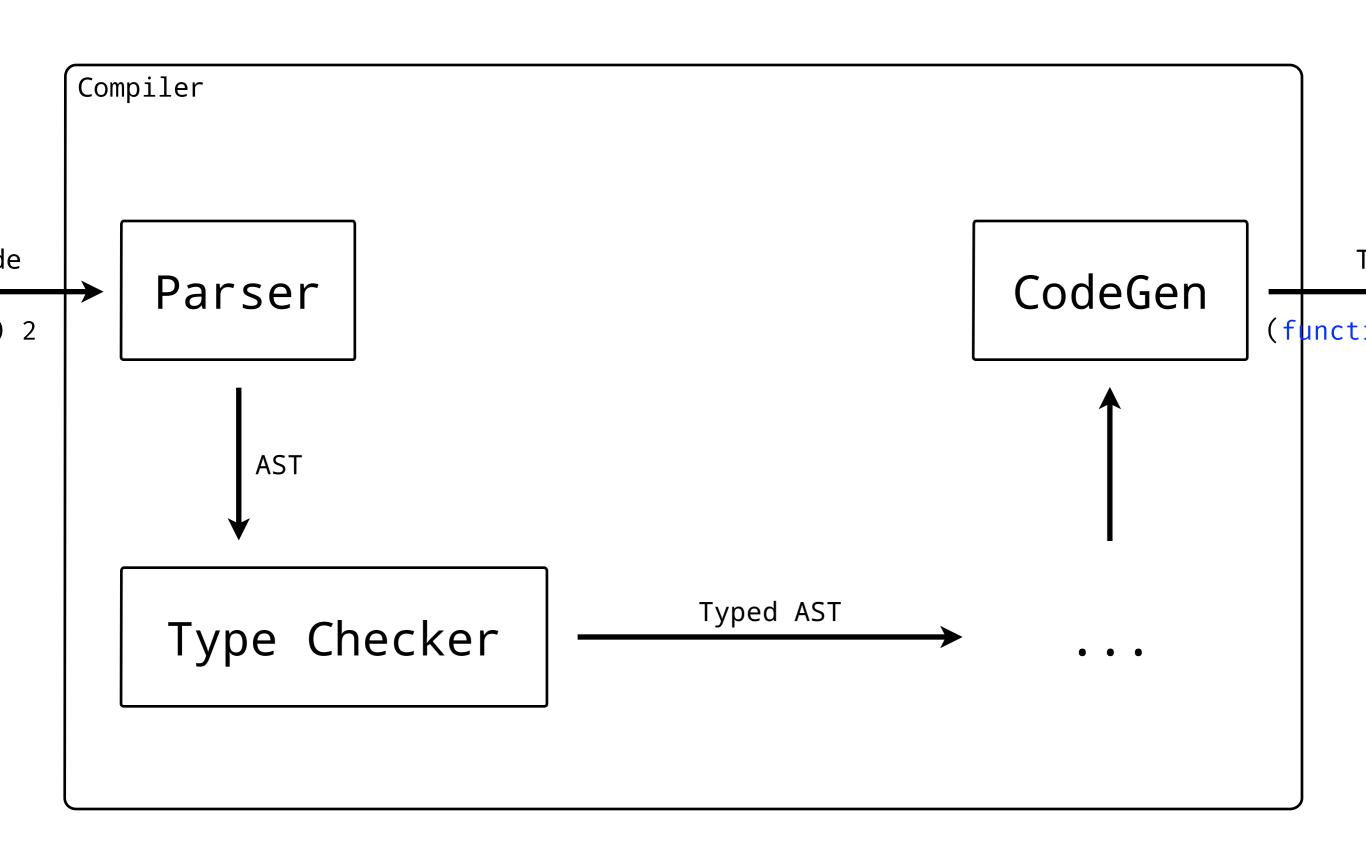
Code Generation



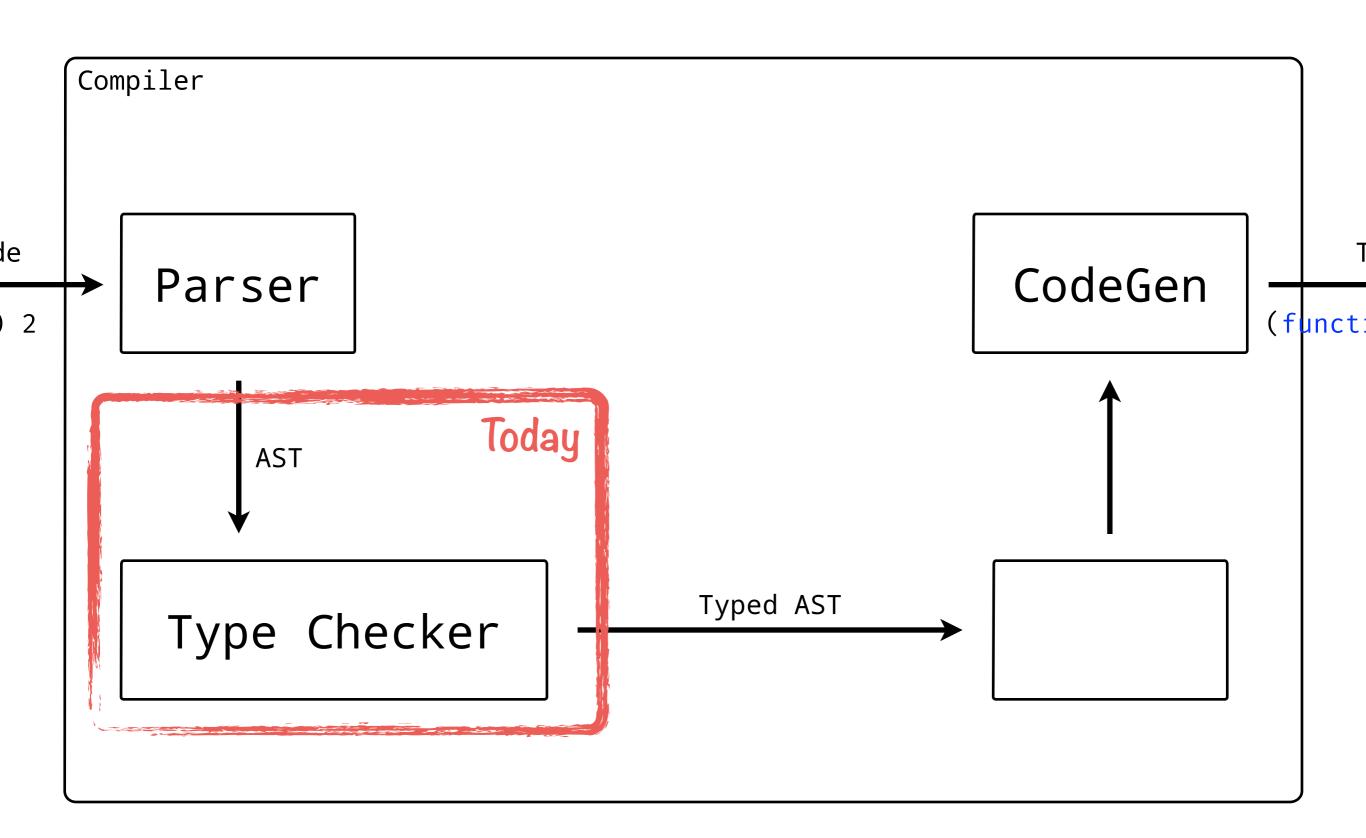
Many Intermediate Phases



Type Checking



Today's Talk



Type Checking

VS

Type Inference

Type Checking vs Inference

- Type Checking
 - Ensures declared types are used consistently
 - All types must be declared
 - Traverse AST and compare def site with use site
- Type Inference
 - Ensures consistency as well
 - Types need not be declared, though; are deduced
 - Two main classes of algorithms
 - We'll see one instance today

Vehicle Language

Vehicle Language

- Surface Syntax
 - What's the concrete syntax of the language
- Type System
 - What types are supported by the language

1. Numbers: 1, 2, 3, ...

- 1. Numbers: 1, 2, 3, ...
- 2. Booleans: true and false

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- 3. Anonymous functions (lambdas): fn a => a

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- 2. Booleans: true and false
- 3. Anonymous functions (lambdas): fn a => a
- 4. Function application: inc 42

- 1. Numbers: 1, 2, 3, ...
- 2. Booleans: true and false
- 3. Anonymous functions (lambdas): fn a => a
- 4. Function application: inc 42
- 5. If expressions: if cond then t else f

6. Let blocks/expressions:

```
let
  val name = ...
in
  name
end
```

Small Example

```
let
   val inc = fn a => a + 1
in
   inc 42
end
```

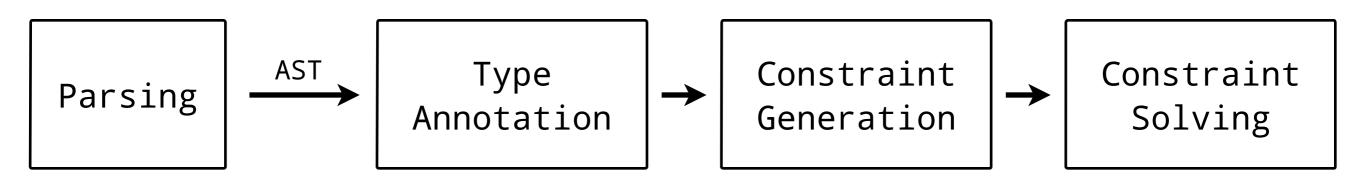
1. Integer type: int

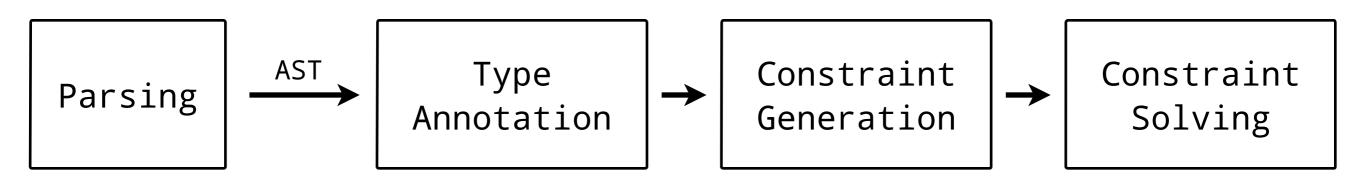
- 1. Integer type: int
- 2. Boolean type: bool

- 1. Integer type: int
- 2. Boolean type: bool
- 3. Function type: int -> bool

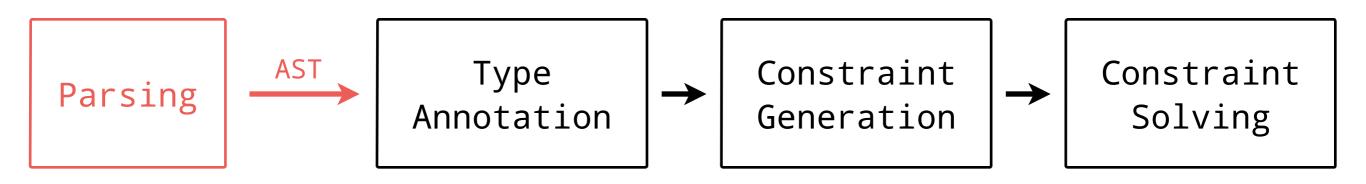
- 1. Integer type: int
- 2. Boolean type: bool
- 3. Function type: int -> bool
- 4. Generic type variables: 'a, 'b, 'c, etc.

Today's Algorithm Overview





```
fn isZero =>
  if isZero 1
  then 2
  else 3
```



```
fn isZero =>

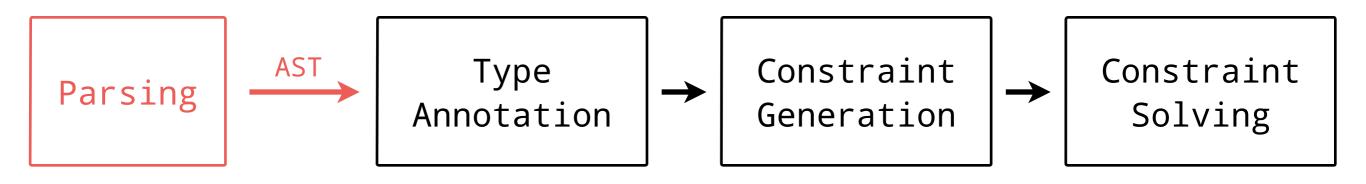
if isZero 1

then 2

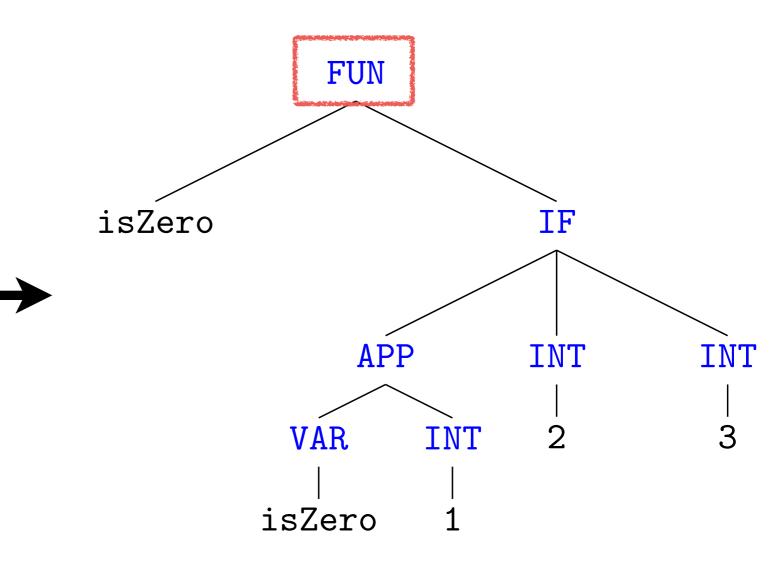
else 3

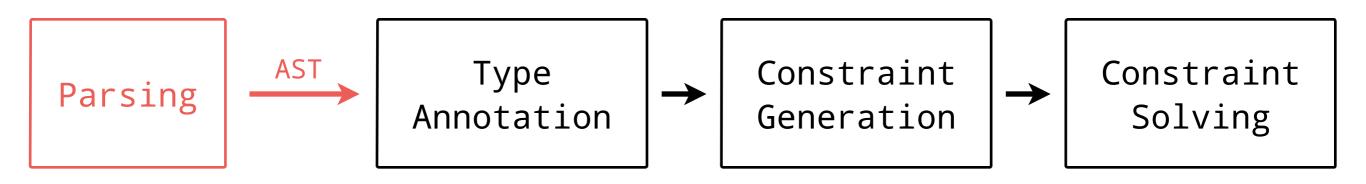
VAR INT 2

isZero 1
```

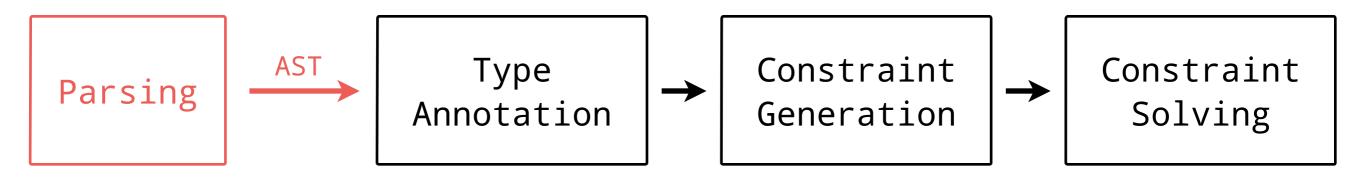


```
fn isZero =>
  if isZero 1
  then 2
  else 3
```

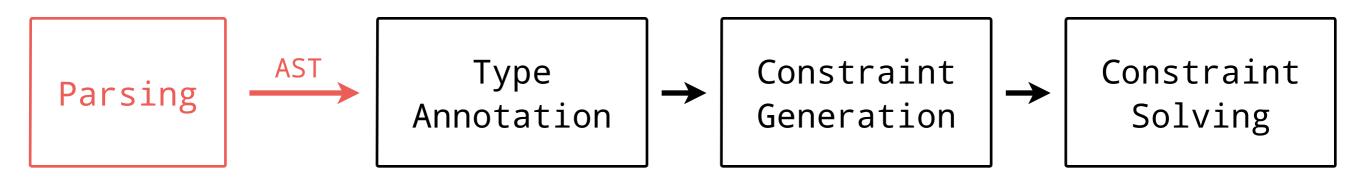




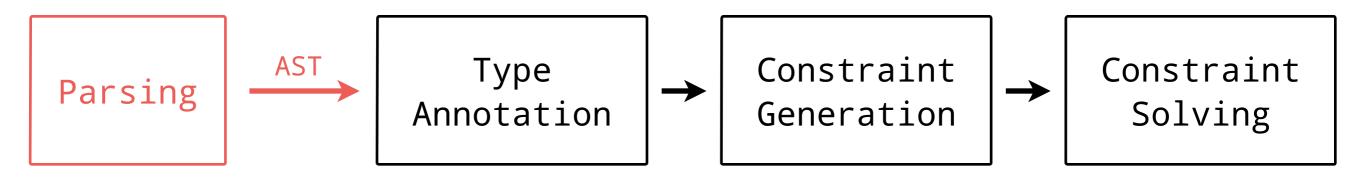
```
FUN
fn isZero =>
  if isZero 1
                          isZero
                                              IF
  then 2
                                      APP
                                              INT
                                                     INT
  else 3
                                   VAR
                                         INT
                                  isZero
```



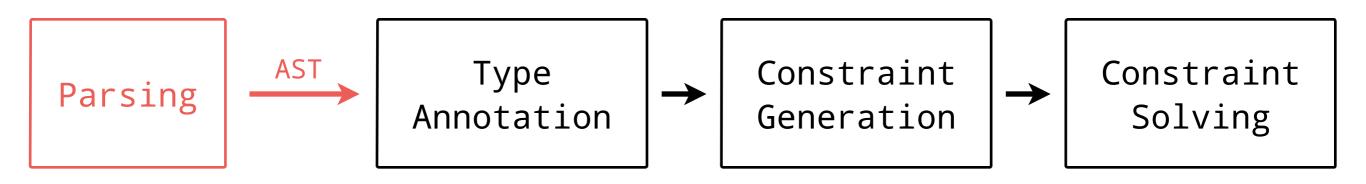
```
FUN
fn isZero =>
  if isZero 1
                                              IF
                          isZero
  then 2
                                      APP
                                             INT
                                                     INT
  else 3
                                         INT
                                   VAR
                                 isZero
```



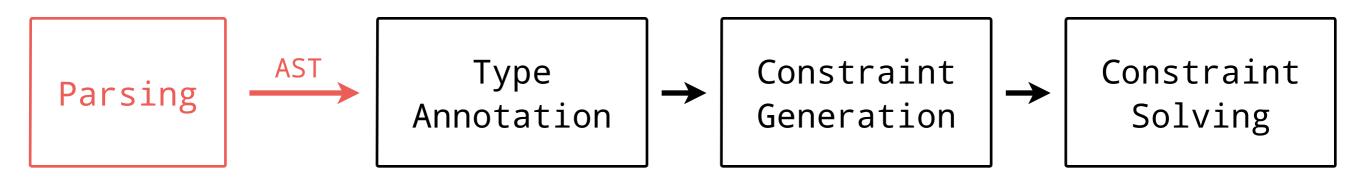
```
FUN
fn isZero =>
  if isZero 1
                          isZero
                                              IF
  then 2
                                      APP
                                             INT
                                                     INT
  else 3
                                         INT
                                   VAR
                                  isZero
```



```
FUN
fn isZero =>
  if isZero 1
                          isZero
                                              IF
  then 2
                                      APP
                                             INT
                                                     INT
  else 3
                                   VAR
                                         INT
                                  isZero
```



```
FUN
fn isZero =>
  if isZero 1
                          isZero
                                              IF
  then 2
                                      APP
                                             INT
                                                     INT
  else 3
                                         INT
                                   VAR
                                 isZero
```



```
fn isZero =>

if isZero 1

then 2

else 3

FUN

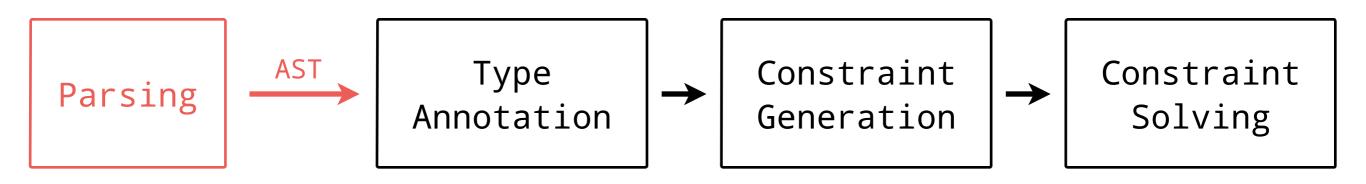
isZero

IF

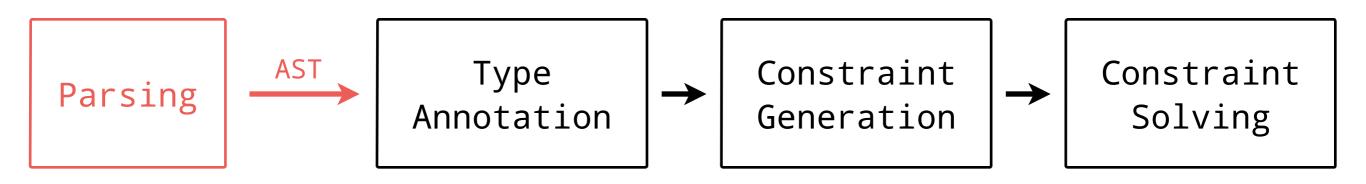
APP
INT
INT
2

3
```

isZero



isZero



```
fn isZero =>

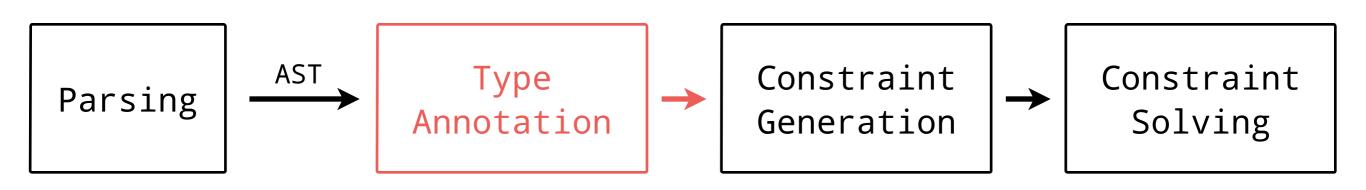
if isZero 1

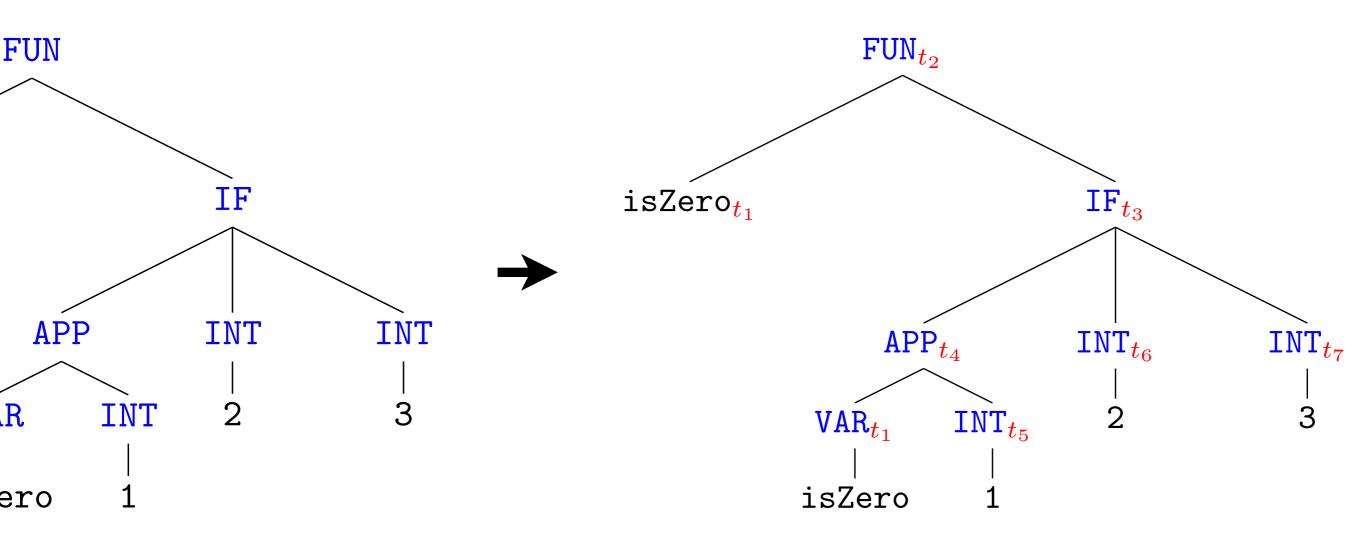
then 2

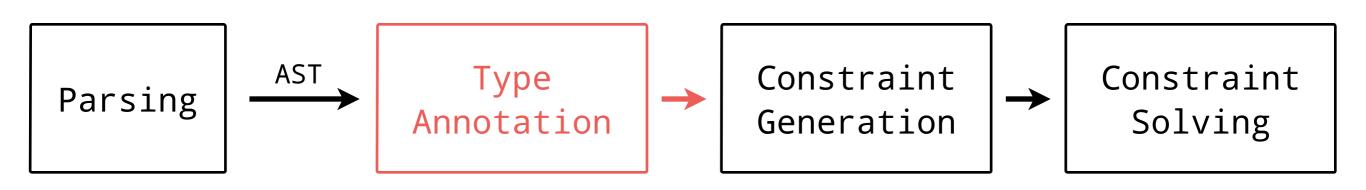
else 3

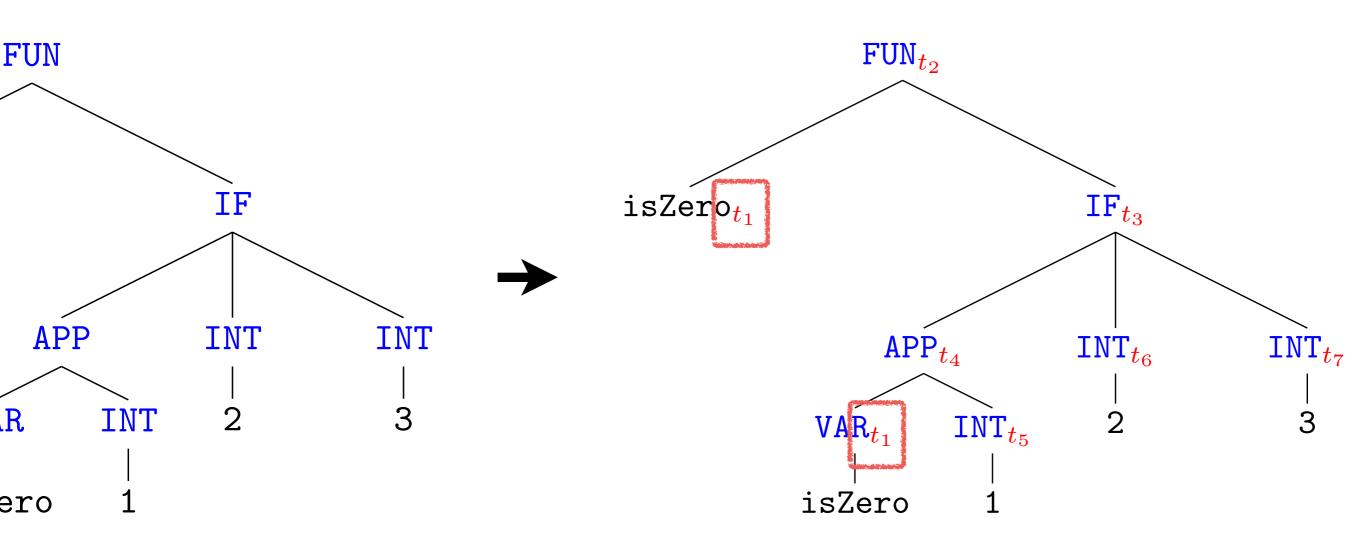
VAR INT 2

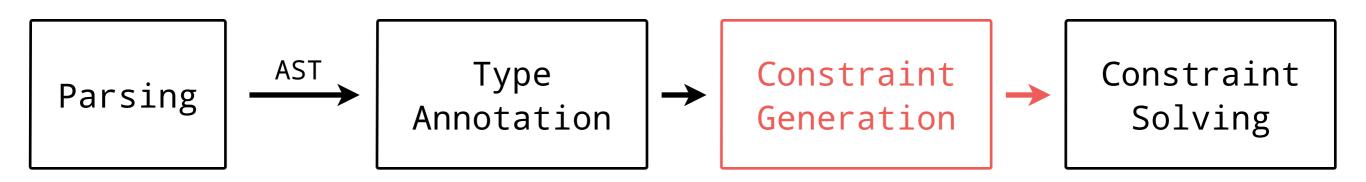
isZero 1
```

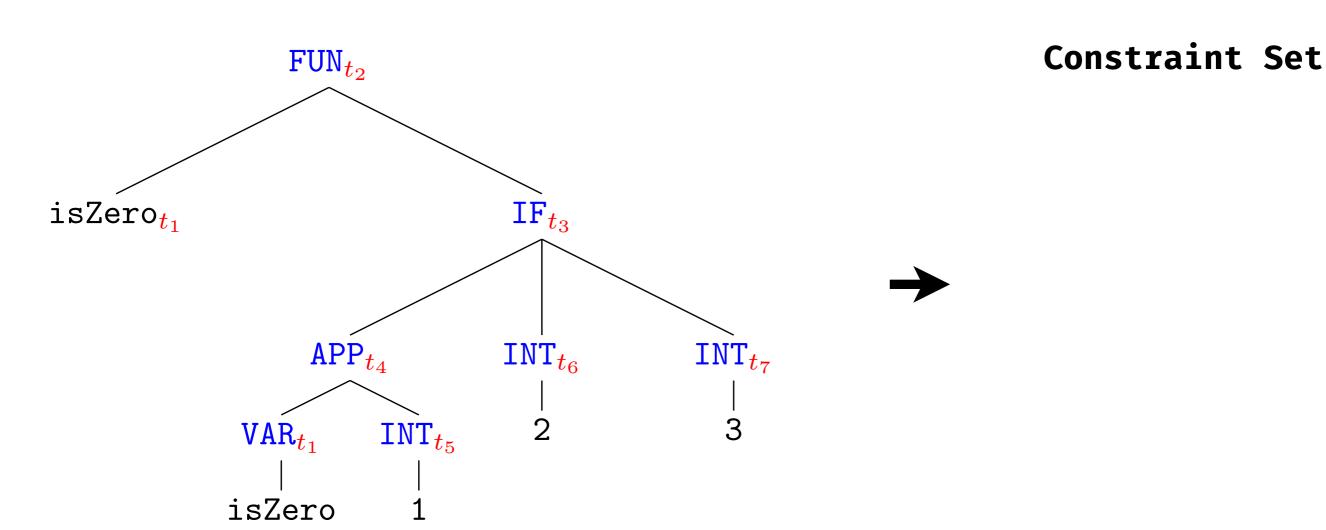


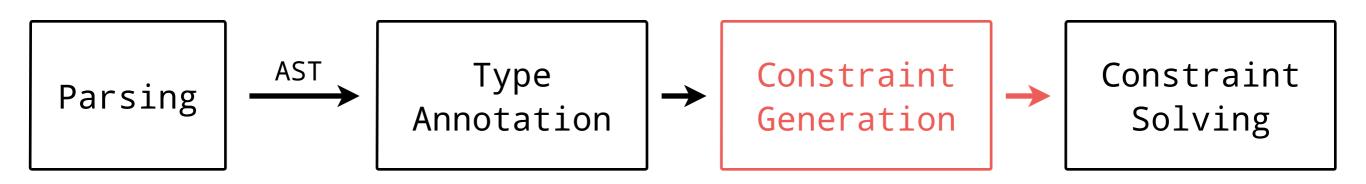


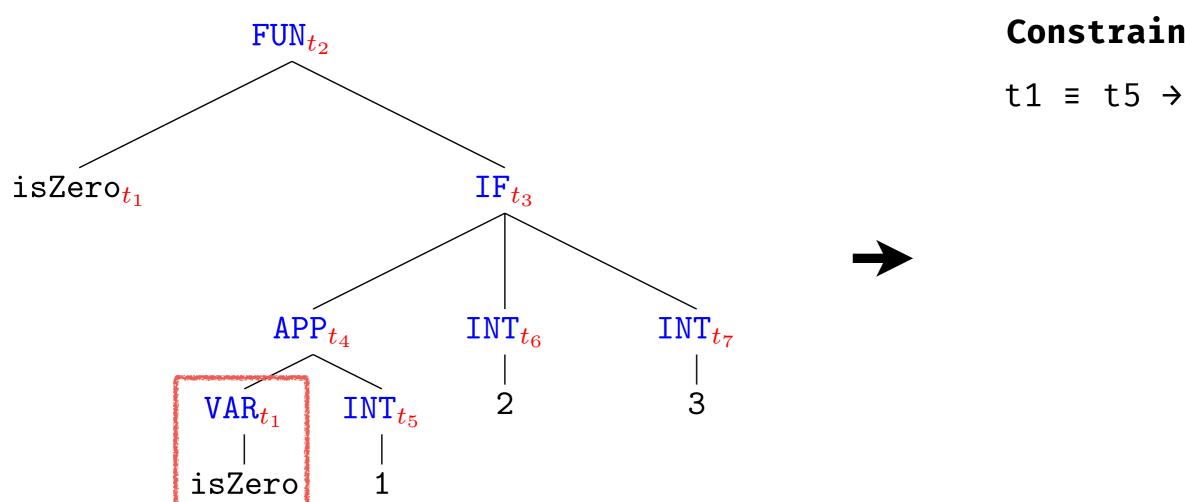




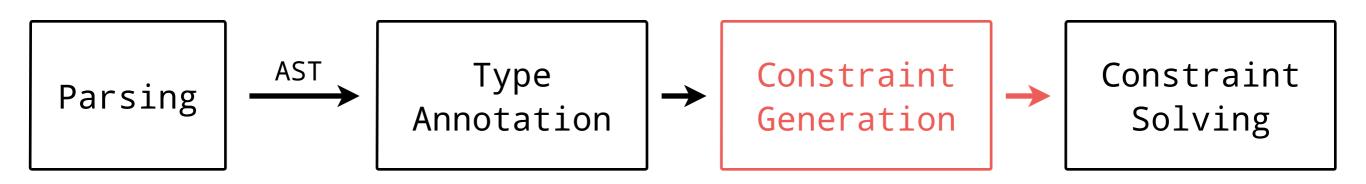


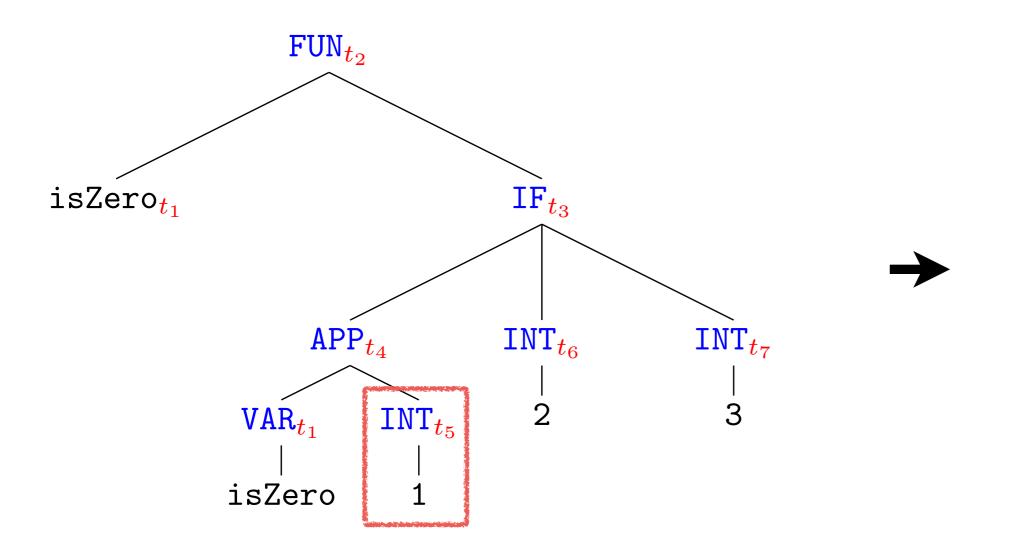






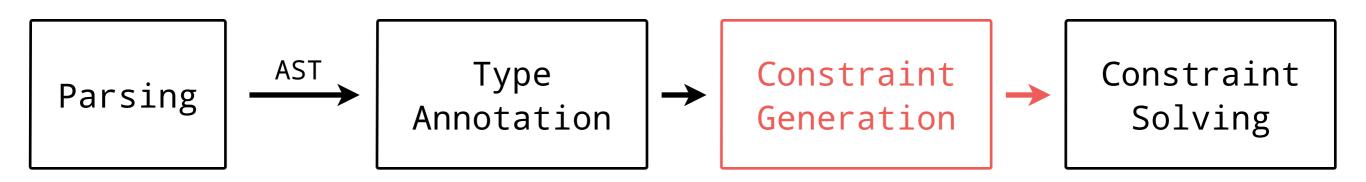
$$t1 \equiv t5 \rightarrow t4$$

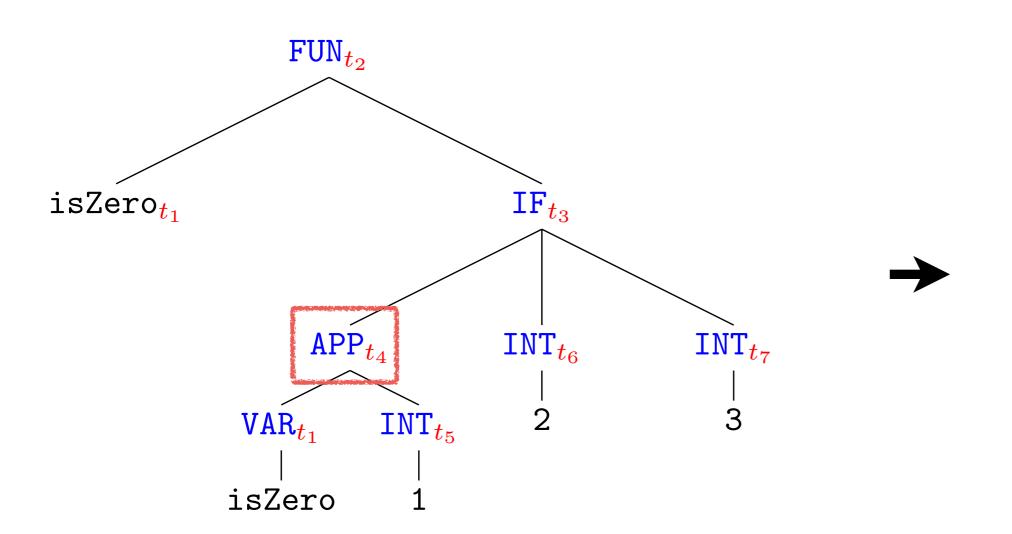




$$t1 \equiv t5 \rightarrow t4$$

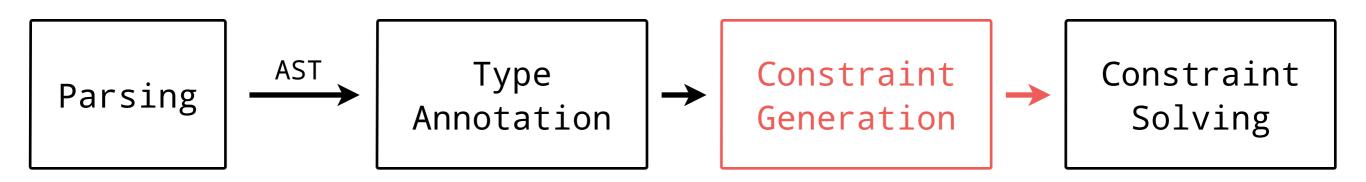
$$t5 \equiv int$$

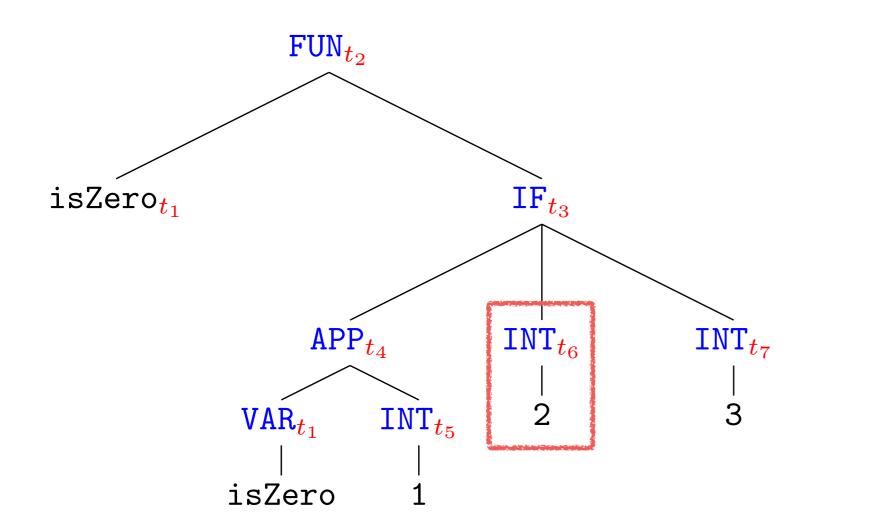




$$t1 \equiv t5 \rightarrow t4$$

$$t5 \equiv int$$





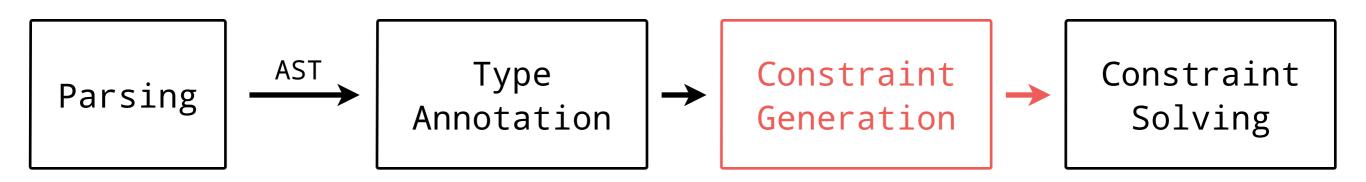
Constraint Set

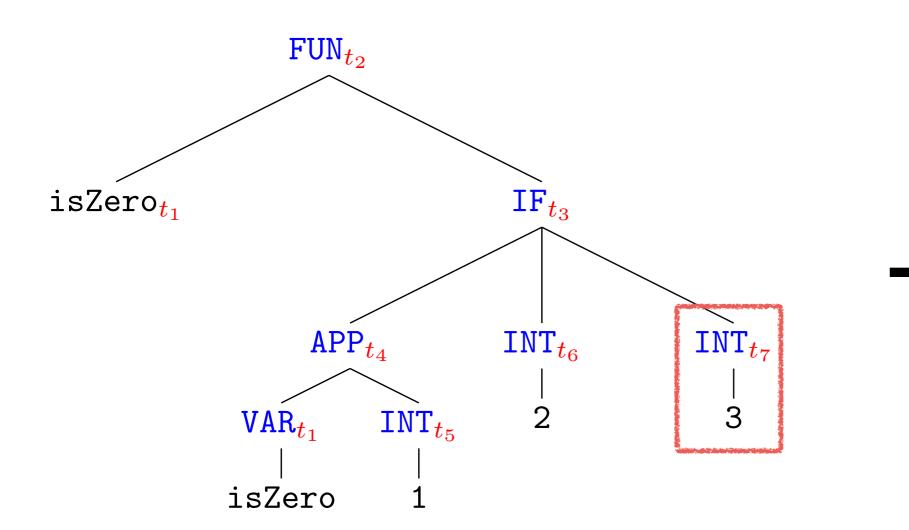
 $t1 \equiv t5 \rightarrow t4$

 $t5 \equiv int$

t4 ≡ bool

t6 ≡ int





Constraint Set

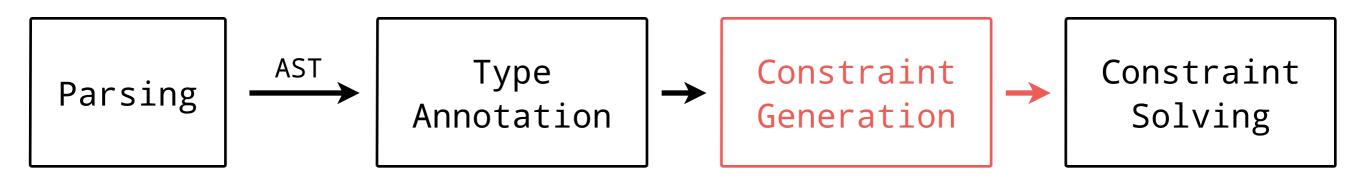
 $t1 \equiv t5 \rightarrow t4$

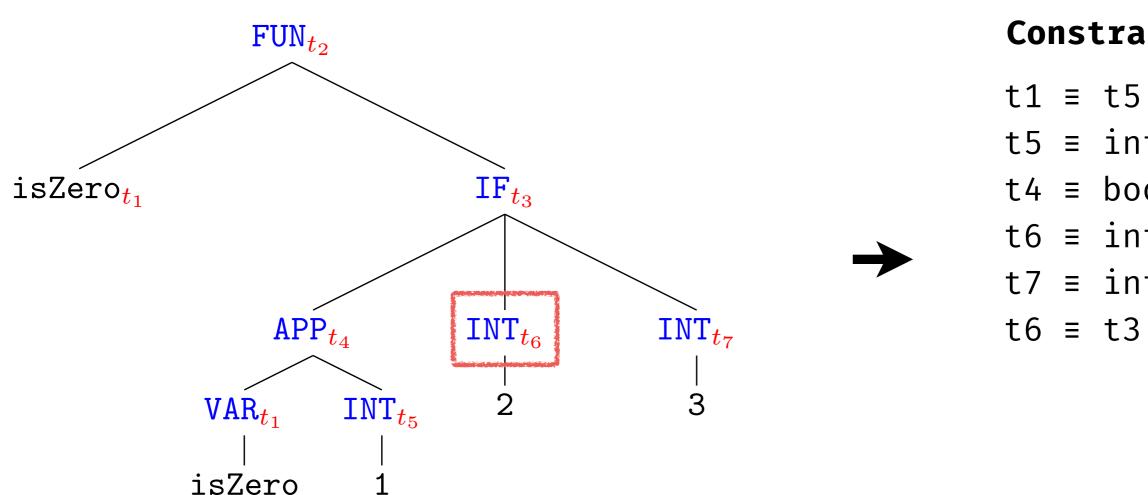
 $t5 \equiv int$

t4 ≡ bool

t6 ≡ int

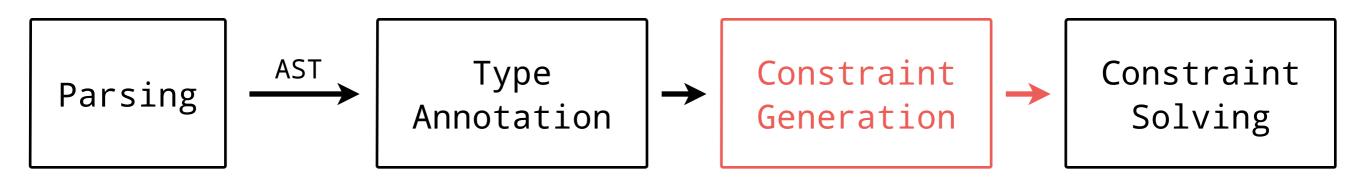
t7 ≡ int

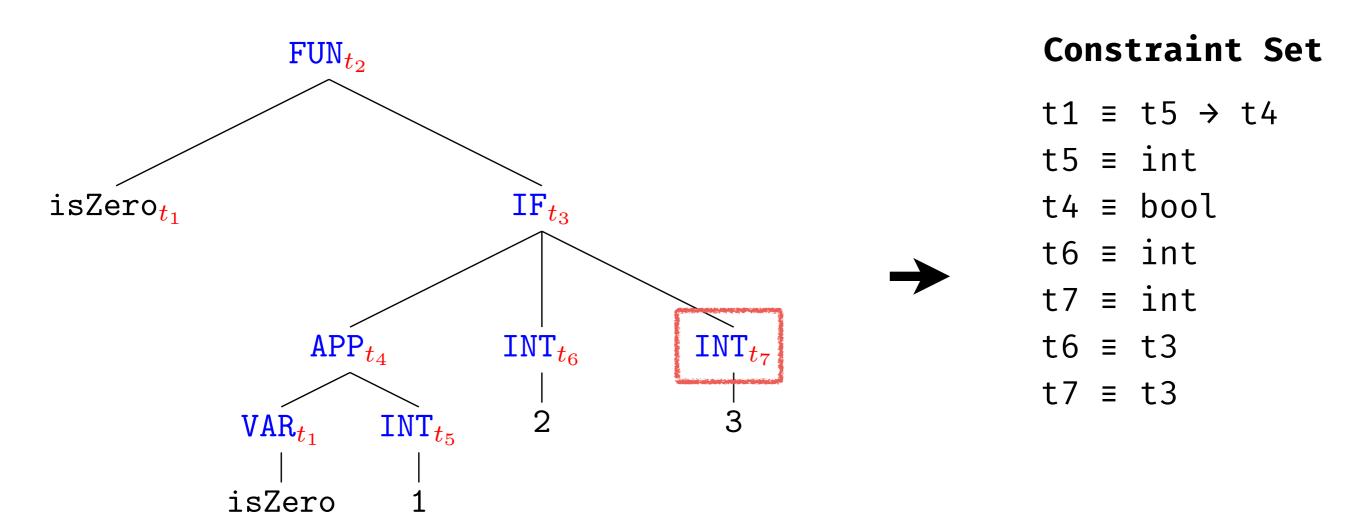


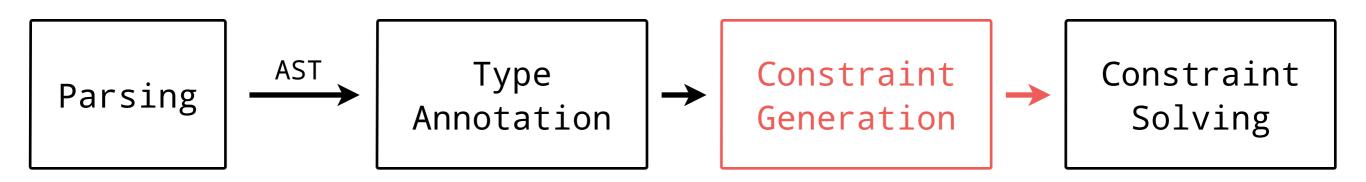


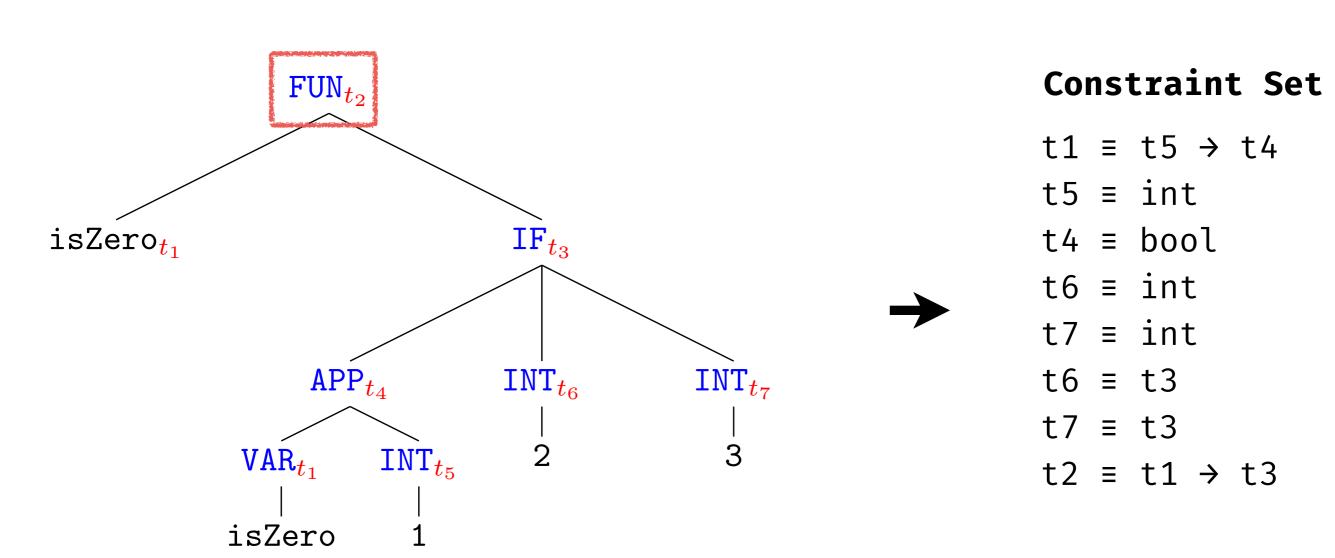
Constraint Set

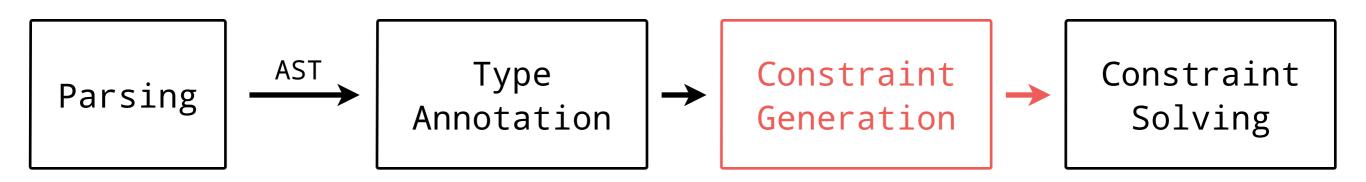
 $t1 \equiv t5 \rightarrow t4$ $t5 \equiv int$ t4 ≡ bool t6 ≡ int t7 ≡ int

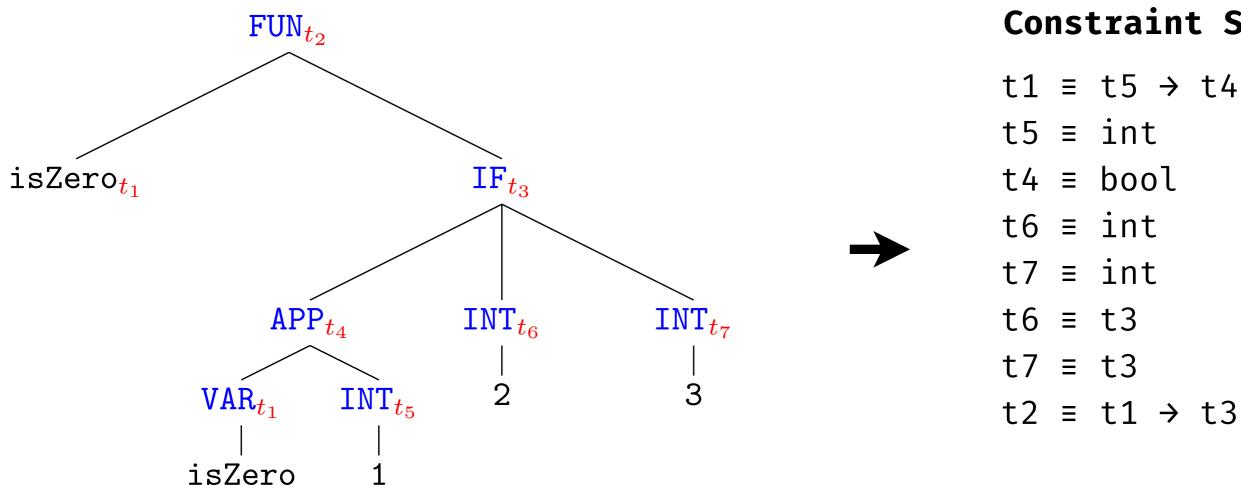


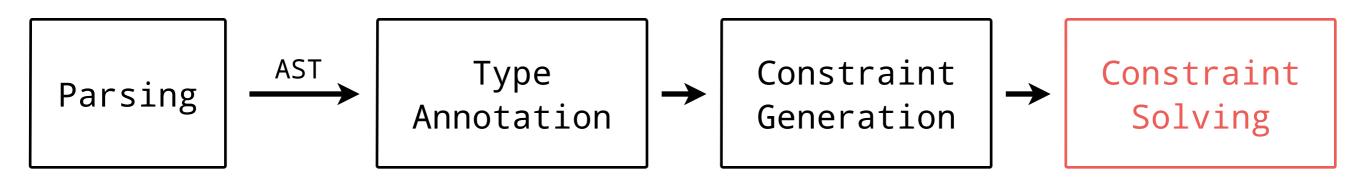








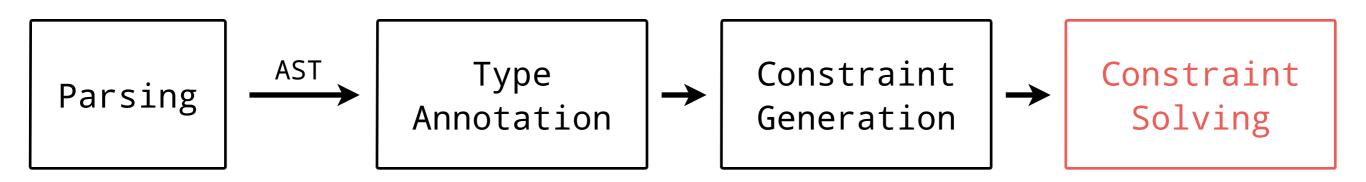




Constraint Set

```
t1 \equiv t5 \rightarrow t4
t5 \equiv int
t4 \equiv bool
t6 \equiv int
t7 \equiv int
t6 \equiv t3
t7 \equiv t3
t2 \equiv t1 \rightarrow t3
```

Solution Map



Constraint Set

```
t5 ≡ int

t4 ≡ bool

t6 ≡ int

t7 ≡ int

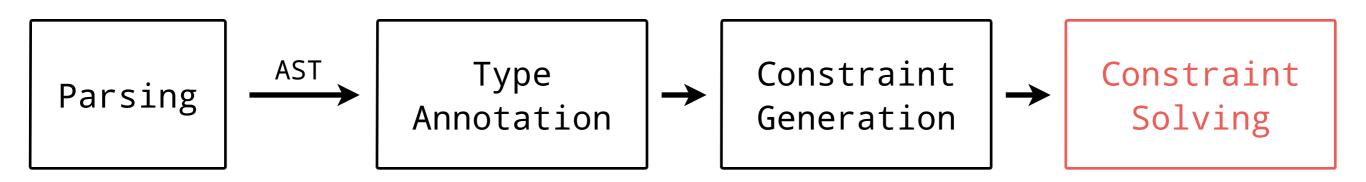
t6 ≡ t3

t7 ≡ t3

t2 ≡ t1 → t3
```

Solution Map

t1: $t5 \rightarrow t4$

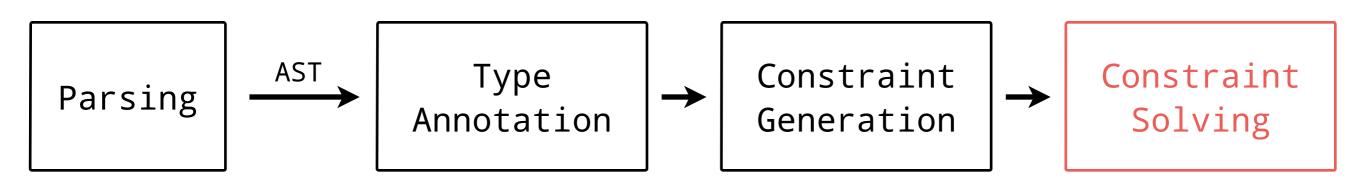


Constraint Set

```
t5 ≡ int
t4 ≡ bool
t6 ≡ int
t7 ≡ int
t6 ≡ t3
t7 ≡ t3
t2 ≡ (t5 → t4) → t3
```

Solution Map

t1: $t5 \rightarrow t4$



Constraint Set

```
t4 \equiv bool

t6 \equiv int

t7 \equiv int

t6 \equiv t3

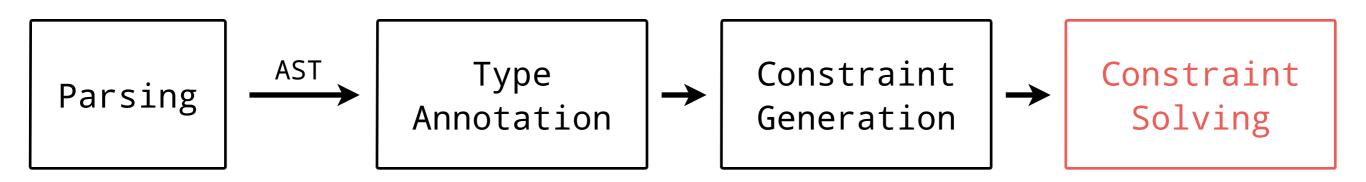
t7 \equiv t3

t2 \equiv (t5 \Rightarrow t4) \Rightarrow t3
```

Solution Map

t1: $t5 \rightarrow t4$

t5: int



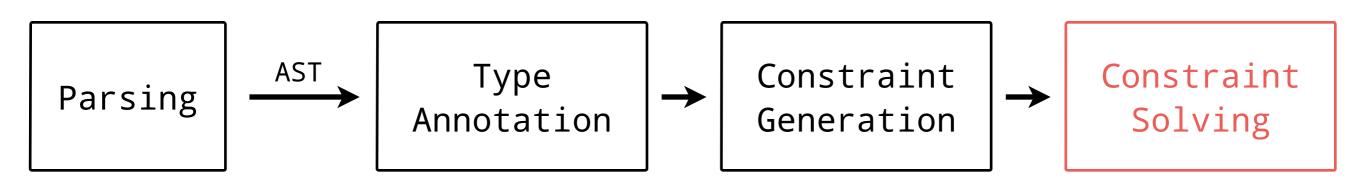
Constraint Set

```
t4 ≡ bool
t6 ≡ int
t7 ≡ int
t6 ≡ t3
t7 ≡ t3
t2 ≡ (int → t4) → t3
```

Solution Map

t1: int \rightarrow t4

t5: int



Constraint Set

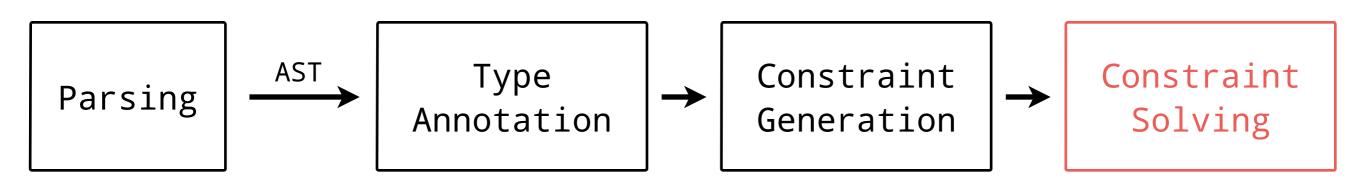
```
t6 ≡ int
t7 ≡ int
t6 ≡ t3
t7 ≡ t3
t2 ≡ (int → t4) → t3
```

Solution Map

t1: int \rightarrow t4

t5: int

t4: bool



Constraint Set

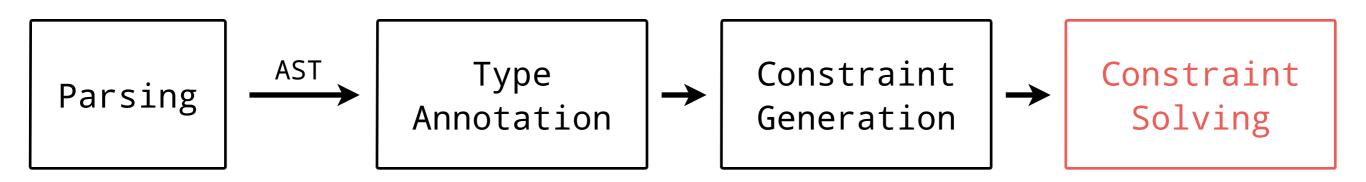
```
t6 ≡ int
t7 ≡ int
t6 ≡ t3
t7 ≡ t3
t2 ≡ (int → bool) → t3
```

Solution Map

t1: int → bool

t5: int

t4: bool



Constraint Set

```
t7 \equiv int
t6 \equiv t3
t7 \equiv t3
t2 \equiv (int \Rightarrow bool) \Rightarrow t3
```

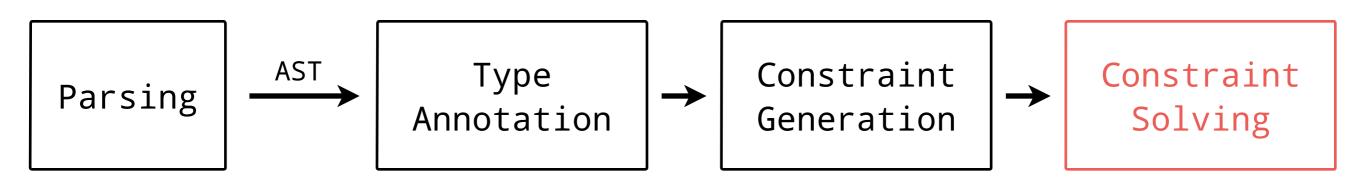
Solution Map

t1: int → bool

t5: int

t4: bool

t6: int



Constraint Set

```
t7 ≡ int
int ≡ t3
t7 ≡ t3
t2 ≡ (int → bool) → t3
```

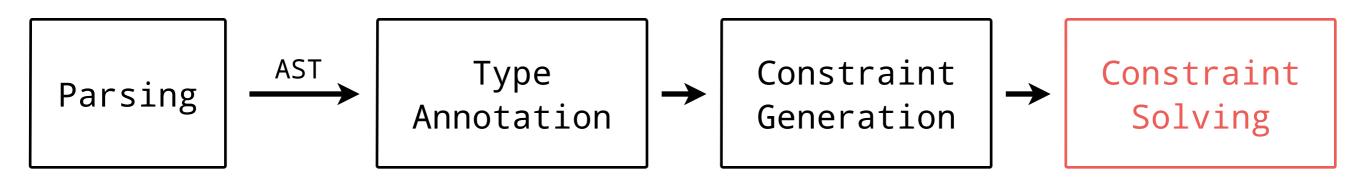
Solution Map

t1: int → bool

t5: int

t4: bool

t6: int



Constraint Set

```
int \equiv t3

t7 \equiv t3

t2 \equiv (int \Rightarrow bool) \Rightarrow t3
```

Solution Map

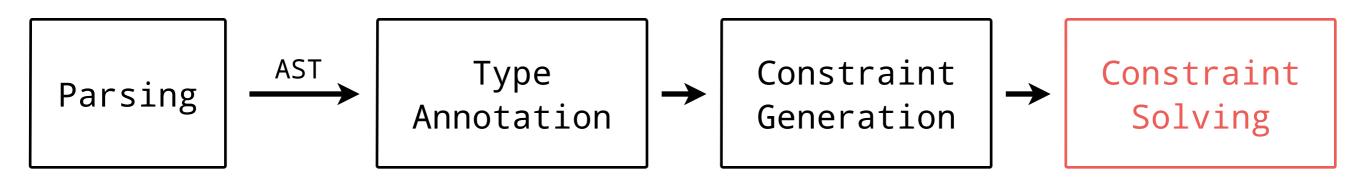
t1: int → bool

t5: int

t4: bool

t6: int

t7: int



Constraint Set

```
int ≡ t3
int ≡ t3
t2 ≡ (int → bool) → t3
```

Solution Map

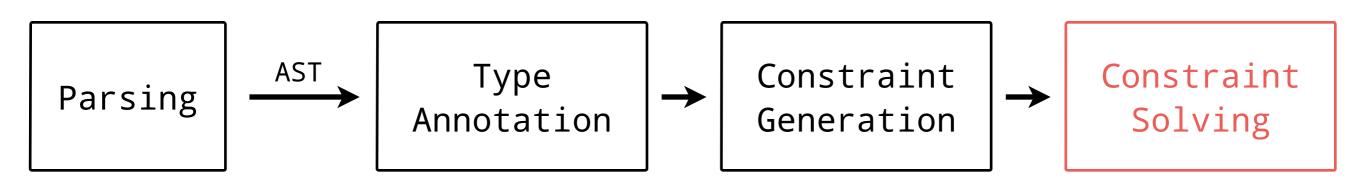
t1: int → bool

t5: int

t4: bool

t6: int

t7: int



Constraint Set

```
int \equiv t3
t2 \equiv (int \rightarrow bool) \rightarrow t3
```

Solution Map

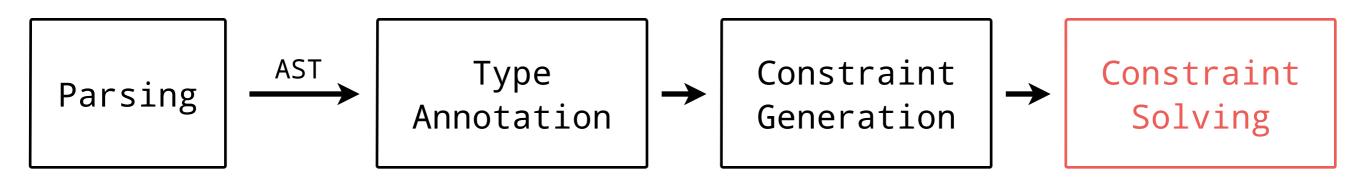
t1: int → bool

t5: int

t4: bool

t6: int

t7: int



Constraint Set

```
int ≡ int
t2 ≡ (int → bool) → int
```

Solution Map

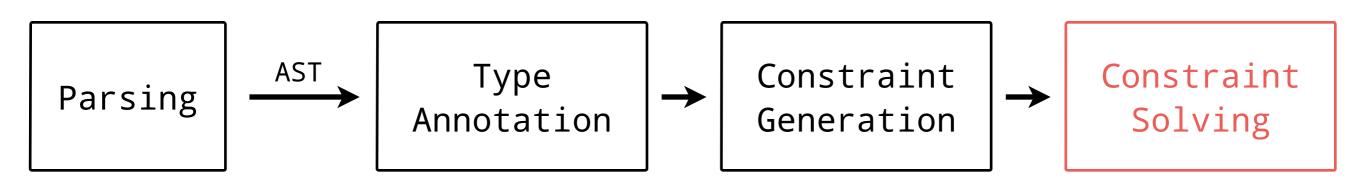
t1: int → bool

t5: int

t4: bool

t6: int

t7: int



Constraint Set

```
int ≡ int
t2 ≡ (int → bool) → int
```

Solution Map

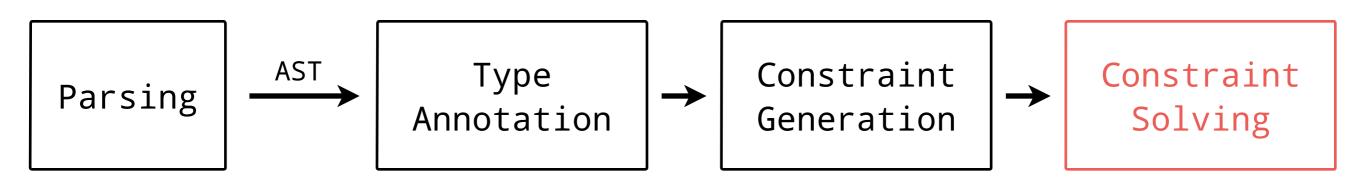
t1: int → bool

t5: int

t4: bool

t6: int

t7: int



Constraint Set

```
t2 \equiv (int \rightarrow bool) \rightarrow int
```

Solution Map

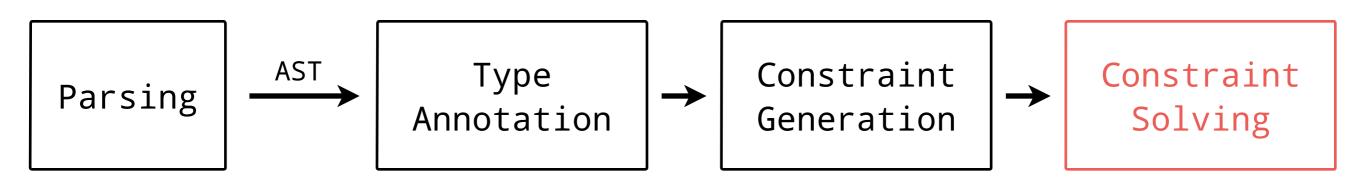
t1: int → bool

t5: int

t4: bool

t6: int

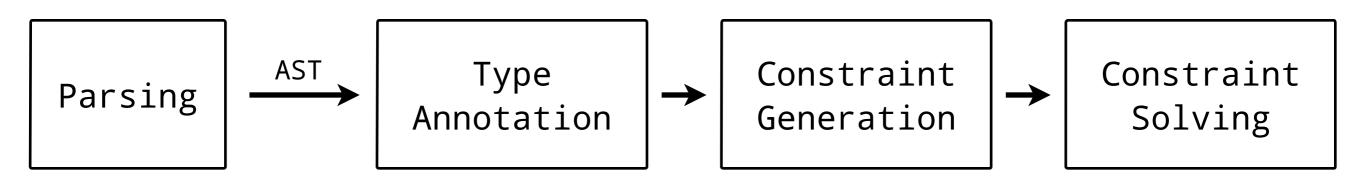
t7: int



Constraint Set

Solution Map

```
t1: int → bool
t5: int
t4: bool
t6: int
t7: int
t3: int
t2: (int → bool) → int
```



Solution Map

t1: int → bool

t5: int

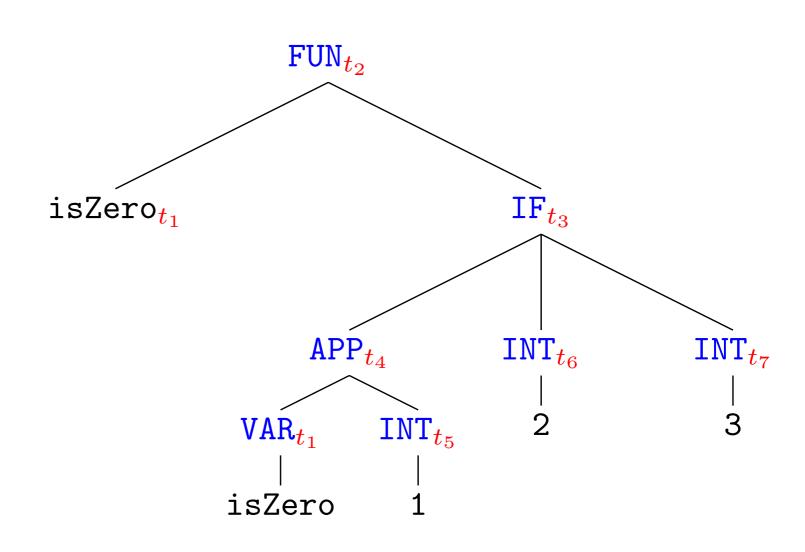
t4: bool

t6: int

t7: int

t3: int

t2: $(int \rightarrow bool) \rightarrow int$



```
Parsing → Type Annotation → Constraint Generation → Constraint Solving
```

Type Checking Algorithms

- As I said, there are two main classes
- Constraint-based ones. We've just seen one example — Wand's algorithm
- Substitution-based ones, where constraint generation and solving are not two separate processes, they are interleaved. Example: the classic Hindley-Milner Algorithm W.

Live Demonstration

https://github.com/igstan/itake-2015

Thank You!

Resources

- Sunil Kothari and James L. Caldwell. Type Reconstruction Algorithms A Survey
- Mitchell Wand. A simple algorithm and proof for type inference.
- Bastiaan Heeren, Jurriaan Hage and Doaitse Swierstra. Generalizing Hindley-Milner Type Inference Algorithms
- Oleg Kiselyov and Chung-chieh Shan. Interpreting Types as Abstract Values
- Shriram Krishnamurthi. Programming Languages: Application and Interpretation, chapter 15
- Shriram Krishnamurthi. Programming Languages: Application and Interpretation, lecture 24
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- **Bastiaan Heeren**. Top Quality Type Error Messages
- Stephen Diehl. Write You a Haskell, chapter 6
- Andrew Appel. Modern Compiler Implementation in ML, chapter 16
- Benjamin Pierce. Types and Programming Languages, chapter 22
- Martin Odersky. Scala by Example, chapter 16
- Danny Gratzer. https://github.com/jozefg/hm
- Arlen Cox. ML Type Inference and Unification
- Radu Rugină. CS 312, Type Inference