

## 1

### 1.1 a

any  $x \leq 50$  will result in a speedup  $\geq 1$

### 1.2 b

$$E = S/N$$

$$N = 2$$

$$E = 0.7$$

$$S = E * N$$

$$S = 2 * 0.7$$

$$S = 1.4$$

### 1.3 c

$$S = T_{old}/T_{new} \quad S = 1000/(500 + (100 + 10 * (10))) \quad S = 1000/(500 + (100 + 100))$$

$$S = 1000/700 \quad S = 10/7$$

### 1.4 d

## 2

### 2.1 a

myrank =1

lock M2

CONTEXT SWITCH

myrank =2

lock M1

wait for M2

CONTEXT SWITCH

wait for M1

both threads would now be waiting for a lock the other holds, with no way to unlock.

## **2.2 b**

yes! each would still be waiting for the other to finish what it is doing, without the other being able to finish.

## **2.3 c**

A simple solution would be whenever both MUTEX may be needed, always try to acquire the same one first (like M1) before trying the other (M2)

## **3**

### **3.1 d**

The times decreased with the number of threads, but busy waiting still timed out. (works better with a smaller n)

### **3.2 e**

needing to protected the approx value slows execution