# CMPT 830 - Bioinformatics and Computational Biology

Chapter 6: RNA Structure

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October 17, 2019

#### **RNA**

- In this section, we are going to try to examine the sequence of RNAs, and predict its structure.
- In RNA, thymine (T) is substituted with uracil (U).
- There are many types of RNA including mRNA, tRNA and rRNA.

#### **RNA**

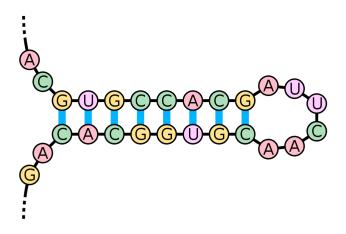
- mRNA is transcribed from DNA, then translated to proteins.
- They are called coding RNA. They do code for proteins.
- RNA can also have catalytic and structural functions, like proteins.
- These are called RNA genes, or non-coding RNA. The most important examples are tRNA and rRNA.
- RNA can have regulatory properties; e.g. microRNA.

#### RNA

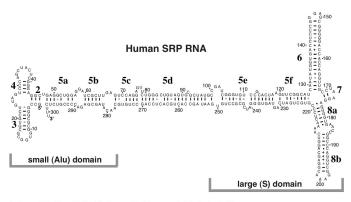
- RNA can also come in double-stranded form.
- The RNA world hypothesis proposes that the common ancestor to all life used RNA to carry genetic information and to catalyze biochemical reactions.

#### **RNA Stucture**

Some primary features of RNA secondary structure include single stranded regions, stems, bulges and loops.



#### **RNA Structure**

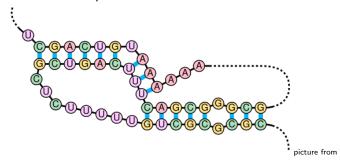


Andersen, E.S., Rosenblad, M.A., Larsen, N., Westergaard, J.C., Burks, J., Wower, I.K., Wower, J., Gorodkin, J., Samuelsson, T., and Zwieb, C. (2006) The tmRDB and SRPDB resources. Nucleic Acids Res. 34, D163-D168.

picture from http://de.wikipedia.org/wiki/Datei:RF00017\_hsap.jpg

#### **Pseudoknots**

- There is also a secondary structure called the pseudoknot.
- This results from a region of a loop pairing with a complementary region outside the loop.



http://en.wikipedia.org/wiki/File:Pseudoknot.svg

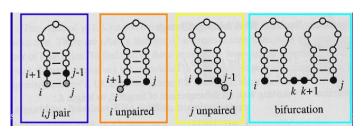
- Often times, we ignore this structure with secondary prediction, because it is very hard.
- However, without pseudoknots, we've got an  $O(n^3)$  algorithm, while allowing pseudoknots makes the problem exponential.

- After we've found a possible conformation of the RNA structure, we would like to evaluate it, to see if it is good.
- We can assign a score, much like we would with an alignment.
- As a first try, it seems the more basepairs we obtain the better, so why don't we let the number of basepairs be our score?

- Now we must *find* the structure with the highest score.
- One method to do this is called *Nussinov's Algorithm* which uses dynammic programming.
- A little bit like a sequence alignment, although we are aligning it with itself.

- Let  $w = a_1 a_2 \cdots a_n$  be our input sequence, where  $a_i \in \{A, U, G, C\}$  for  $i, 1 \le i \le n$ .
- We move from the "inside" of the sequence towards the "outside".
- Suppose we have already figured out every optimal "substructure" strictly between a<sub>i</sub> and a<sub>i</sub>.
- So, for example, we already have figured out the optimal structure of the sequence  $a_{i+1}\cdots a_j$ , and  $a_{i+2}\cdots a_j$ , and also  $a_i\cdots a_{j-1}$  amongst many others.

- There are four ways to extend that information to a best structure of a<sub>i</sub>···a<sub>i</sub>.
  - 1. add an unpaired base i to the best structure for the subsequence  $a_{i+1} \cdots a_j$ ,
  - 2. add an unpaired base j to the best structure for the subsequence  $a_i \cdots a_{j-1}$ ,
  - 3. add paired bases i, j to the best structure for the subsequence  $a_{i+1} \cdots a_{j-1}$ ,
  - 4. combine two optimal substructures  $a_i \cdots a_k$  and  $a_{k+1} \cdots a_j$ , for some k.



- Cases 1 and 2 (orange and yellow outline) should not result in an increase in score.
- Case 3 (blue outline) should result in an increase of score.
- Case 4 (teal outline) does not increase the score, however, it requires we try every possible value of *k* to see which one gives you the structure with the best score.

- A recursive procedure would not be a very good way to program the problem.
- Indeed, by starting with the full string and then calculating its
  optimal structure by calculating smaller ones recursively inwards, we
  are calculating the optimal structure of many of the sequences many
  times.

- We only want to calculate each one once.
- This is exactly what dynamic programming can do.
- Start with small instances, calculate the optimal structure once and then use those to build larger instances.

We start by building an  $n \times n$  matrix, and plot w against itself.

#### **Example**

For this example, let w = UGAU.

	U	G	$\overline{}$	U
U				
U G				
Α				
U				

#### Initialization

- We need to initialize it somehow, much as we did with sequence alignment.
- Instead of going from the beginning of the sequences towards the end, as with sequence alignment, we are trying to go from "small subsequences" outwards towards "bigger subsequences".

#### note

We know that the situation in aligning everything between  $a_i$  and  $a_i$  is boring. There are no basepairs. So, we should initialize all the entries on the diagonal to 0.

#### Initialization

- Also, notice that we are mainly concerned with the situation where i ≤ j.
- Thus, we can ignore much of the table.
- It turns out that it is a good idea to initialize the main diagonal to be zeros, but also the diagonal just below as well.

#### **Example**

For our example, this gives us:

	U	G	Α	U
U	0			
Ğ	0	0		
Α	Х	0	0	
U	Х	Х	0	0

#### Initialization Pseudocode

#### we initialize the 0's as follows:

```
\begin{array}{l} \text{for } i \leftarrow 0 \text{ to } n-1 \\ M(i,i) \leftarrow 0 \\ \text{for } i \leftarrow 1 \text{ to } n-1 \\ M(i,i-1) \leftarrow 0 \end{array}
```

#### Recurrence

## Recall the matrix:

	U	G	Α	U
U	0			
U G A U	0	0		
Α	×	0	0	
U	×	X	0	0

## We then fill in the rest of the matrix with the following recurrence:

$$\begin{split} M(i,j) &= \max\{M(i+1,j),\\ & M(i,j-1),\\ & M(i+1,j-1) + d(i,j),\\ & \max\{M(i,k) + M(k+1,j) \mid \text{ for } i \leq k < j\},\\ & \}, \end{split}$$

where 
$$d(i,j) = \begin{cases} 1 & \text{if } w(i) \text{ and } w(j) \text{ are complementary,} \\ 0 & \text{otherwise} \end{cases}$$

## Finding the Score

#### **Example**

We obtain the following matrix:

	U	G	Α	U
U	0	0	1	1
G	0	0	0	1
Α	Х	0	0	1
U	Х	Х	0	0

- Entry M(0, n-1) is the optimal score.
- Of course, now we'd actually like to build one such optimal structure.

#### Create the Structure

- As with sequence alignment, we trace back from the optimal score.
- Problem: the traceback is complicated by the bifurcation defined by the fourth term in the recurrence relation.

### Stacks — a Tangent

- A stack is a computational data structure, similar to an array.
- Intuitively, it is similar to a stack of books.
- We can look at the top book on the stack, we can "pop" a book from the top, or we can "push" a new book onto the top of the stack.

### Stacks — a Segue

- When we push a new book onto the top, everything below is buried, but they are still there.
- We can get at them, but only if we keep on popping books until we get to the one we want.

#### **Example**

2 6 : 5 II

- We can have a stack of numbers, or anything else for that matter.
- We can also have a stack of ordered pairs of integers (like our next example). So each cell will contain (i, j) for some i, j.

#### Create the Structure

#### the algorithm

```
PUSH((0, n-1)) //adds the ordered pair (1, n) to the stack
while READ()\neq II //until we hit the bottom of the stack, do the following
    (i, i) \leftarrow POP()
    if i > i //if i is greater than or equal to i, do nothing
    else if M(i + 1, j) == M(i, j) PUSH ((i + 1, j))
    else if M(i, j-1) == M(i, j) PUSH ((i, j-1))
    else if M(i, j) == M(i + 1, j - 1) + d(i, j)
          output "i pairs with j" //this outputs to the screen
          PUSH((i + 1, i - 1))
    else
          for k \leftarrow i + 1 to i - 1
              if M(i, k) + M(k + 1, i) == M(i, i)
                   PUSH((k+1,i))
                   PUSH((i,k))
                   break //exits the for loop
```

- Nussinov's algorithm is from 1978.
- Most algorithms now try to minimize free energy instead of counting basepairs.

 The most prominent program is the mfold package of Michael Zucker.

#### website

```
http://mfold.rna.albany.edu/?q=mfold
```

- mfold does not take pseudoknots into account.
- This was merged with DINAMelt to create a newer package called UNAFold.

#### website

```
http://unafold.rna.albany.edu/?q=DINAMelt
```

- There are also algorithms which take pseudoknots into account.
- They typically sacrifice accuracy for speed.
- Indeed, when pseudoknots are included, using minimization of free energy as criteria, the problem has been shown to be NP-complete (see chapter 1).
- Thus, exponential algorithms are likely the best we will be able to do, if we want to be completely accurate.