# Advanced OpenCL Programming

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# Multiple Devices

- We have been using one device for computation.
- Now we will use multiple devices to solve a problem.
- We will use matrix multiplication as an example.

### **Partition**

- Since we are very poor and cannot afford four GPUs, we will partition the data by rows (not by block) among devices.
- We will partition matrix by rows, so that the memory assigned to a device is contiguous.
- Each kernel will run on a device, and compute part of the answers.

### **Partition**

- We will partition A and C by rows, and do *not* partition B.
- For example, if we use two devices, then the first kernel will compute the top half of *C*, and another kernel will compute the bottom half of *C*.
- The reason for not partitioning *B* is that all kernels need the entire *B*.

### Constants

- MAXLOG is the maximum number of bytes for storing compilation log. More on this later.
- The number of device is DEVICENUM. We will use 2 in our humble installation.
- ITEMPERDEVICE is the number of work item in NDRange of a kernel.

### Constants

Example 1: (matrixMul-time-copy-local-multidevice.c)

```
#include <stdio.h>
4
   #include <assert.h>
5
   #include <CL/cl.h>
6
   #define N 1024
   #define Blk 64
9
   #define BSIDE (N / Blk)
10
   #define MAXGPU 10
11
   #define MAXK 1024
12
   #define MAXI.OG 4096
13
   #define DEVICENUM 2
14
   #define ITEMPERDEVICE
                           (N * N / DEVICENUM)
15
   #define NAND2SECOND 1000000000.0
16
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   cl_uint A[N][N], B[N][N], C[N][N];
```

### Discussion

• If *N* is 1024, and the number of device is 2, then how many works items are there in a kernel?

### Select Devices

- We will select the first DEVICENUM GPU device from our humble installation.
- The code is similar to those earlier version, except that now we need to make sure that the number of GPUs found is at least DEVICENUM.

### **Devices**

#### Example 2: (matrixMul-time-copy-local-multidevice.c)

```
30
     cl_device_id GPU[MAXGPU];
31
     cl_uint GPU_id_got;
32
     status =
33
        clGetDeviceIDs(platform_id, CL_DEVICE_TYPE_GPU,
34
                       MAXGPU, GPU, &GPU_id_got);
35
     assert(status == CL_SUCCESS &&
36
             GPU_id_got >= DEVICENUM);
37
     printf("There are %d GPU devices\n", GPU_id_got);
```

### Context

- Since a context can consist of multiple devices, we need only one context.
- Here we include the first DEVICENUM GPUs in the context.

## Context

#### Example 3: (matrixMul-time-copy-local-multidevice.c)

# Multiple Queue

- Recall that a command queue connects to a device.
- Since we have multiple devices, we need multiple command queues.
- We put the command queues in the commandQueue array.
- Note that we set the CL\_QUEUE\_PROFILING\_ENABLE to enable profiling.

# Command Queue

#### Example 4: (matrixMul-time-copy-local-multidevice.c)

```
44
     const cl_queue_properties properties[] =
45
        {CL_QUEUE_PROPERTIES, CL_QUEUE_PROFILING_ENABLE, O};
46
     cl_command_queue commandQueue[DEVICENUM];
47
     for (int device = 0; device < DEVICENUM; device++) {</pre>
48
        commandQueue[device] =
49
          clCreateCommandQueueWithProperties(context, GPU[devic
50
                                               properties, &statu
51
        assert(status == CL SUCCESS):
52
     }
```

# One Context, Multiple Devices

- Up to this point, we can image there is only one context, which has multiple devices.
- Within this context we will have one command queue to connect to each device.
- Later we will send commands to these devices. The command sent into a command queue will run on the device this command queue connects to.

## Discussion

• What will be sent into these command queues as commands?

## Program

- The kernel will be sent into the command queue as commands.
- Before this can happen we need to compile the kernel.
- Recall that the "kernel" at this point is only a set of strings, which need to be compiled.

# Build Program

- We call clBuildProgram to build executable.
- Note that we pass all devices we want to use as parameters, so
   OpenCL will build executable for all our DEVICENUM devices.

# Compilation Log

- The kernel source is compiled when an OpenCL program runs.
- It is very inconvenient because when we run an OpenCL program the execution will not display compilation errors of the kernel.
- We will use clGetProgramBuildInfo to show the error message if the compilation of the kernel fails.

#### Prototype 5: clGetProgramBuildInfo.h

### **Parameters**

```
program The compiled program.
  device The device you want to query.
param_name The information you want to query.
param_value_size The length of param_value in bytes.
param_value The location for the query answer.
param_value_size_ret The number of bytes returned from the query.
```

### Error

- When there is an error in clBuildProgram, we will call clGetProgramBuildInfo to find out.
- Since we are not sure about which device causes the compilation error, we list the information from all of them.
- We query with CL\_PROGRAM\_BUILD\_LOG for the compilation log, and place it into the log buffer we prepared.

# **Build Program**

#### Example 6: (matrixMul-time-copy-local-multidevice.c)

```
66
      status =
67
        clBuildProgram (program, DEVICENUM, GPU, NULL,
68
                        NULL. NULL):
69
      if (status != CL_SUCCESS) {
70
        char log[MAXLOG];
71
        size_t logLength;
72
        for (int device = 0; device < DEVICENUM; device++)</pre>
73
          clGetProgramBuildInfo(program, GPU[device],
74
                                  CL_PROGRAM_BUILD_LOG,
75
                                  MAXLOG, log, &logLength);
76
          puts(log);
77
78
        exit(-1);
79
80
      printf("Build program completes\n");
```

## Discussion

- What is the type of the compilation log?
- What does puts do?

### **Partition**

- Now we are ready to partition host buffers A, B and C.
- We will create DEVICENUM buffers for A, and each of which will be given to a kernel as a parameter for computations.
- We partition A into sub-matrix of N / DEVICENUM rows of A each, and assign each partition to an OpenCL buffer.

### Buffer A

- ITEMPERDEVICE means the number of items per device.
- The index device determines the starting position of a buffer, i.e., where the matrix A will be partitioned.

## Buffer for A

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#### Example 7: (matrixMul-time-copy-local-multidevice.c)

```
cl_mem bufferA[DEVICENUM];
for (int device = 0; device < DEVICENUM; device++) {
  bufferA[device] =
    clCreateBuffer(context,
        CL_MEM_READ_ONLY | CL_MEM_COPY_HOST_PTR,
        ITEMPERDEVICE * sizeof(cl_uint),
        ((cl_uint *)A) + device * ITEMPERDEVICE,
        &status);
  assert(status == CL_SUCCESS);
}</pre>
```

## Buffer B

• We do not need to partition *B* because it will be used by all kernels.

## Buffer for B

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#### Example 8: (matrixMul-time-copy-local-multidevice.c)

```
cl_mem bufferB =
    clCreateBuffer(context,
        CL_MEM_READ_ONLY | CL_MEM_COPY_HOST_PTR,
        N * N * sizeof(cl_uint), B, &status);
assert(status == CL_SUCCESS);
```

## Buffer C

- Buffers for C are similarly built as for A.
- Unlike A and B that will copy host buffers to device buffers,
   C will use host memory directly.

## Buffer for C

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#### Example 9: (matrixMul-time-copy-local-multidevice.c)

```
cl_mem bufferC[DEVICENUM];
for (int device = 0; device < DEVICENUM; device++) {
  bufferC[device] =
    clCreateBuffer(context,
        CL_MEM_WRITE_ONLY | CL_MEM_USE_HOST_PTR,
        ITEMPERDEVICE * sizeof(cl_uint),
        ((cl_uint *) C) + device * ITEMPERDEVICE,
        &status);
  assert(status == CL_SUCCESS);
}
printf("Build buffers completes\n");</pre>
```

## Discussion

 Make sure that you understand the offset calculation in A and C.

## **NDRange**

- According to our partition we declare NDRange as a two dimension array with N / DEVICENUM rows and N columns.
- This is consistent with C, where each work item will compute an element of C.
- The work group is still BSIDE by BSIDE since we will use the previous local memory algorithm.
- We will have DEVICENUM kernels, so we will need DEVICENUM events for synchronization and profiling.

# **NDRange**

#### Example 10: (matrixMul-time-copy-local-multidevice.c)

```
size_t globalThreads[] =
{(size_t)(N / DEVICENUM), (size_t)N};
size_t localThreads[] = {BSIDE, BSIDE};
cl_event events[DEVICENUM];
```

# Set Argument

- We will launch DEVICENUM kernels. Each of them will compute N / DEVICENUM rows and N columns of C.
- We loop through all devices, and assign the argument order as A, B, and C.
- Note that we need to supply the correct A and C since there
  are DEVICENUM of them.

# Set Arguments

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#### Example 11: (matrixMul-time-copy-local-multidevice.c)

```
for (int device = 0; device < DEVICENUM; device++) {</pre>
        status = clSetKernelArg(kernel, 0, sizeof(cl_mem),
                                  (void*)(&bufferA[device]));
        assert(status == CL SUCCESS):
        status = clSetKernelArg(kernel, 1, sizeof(cl_mem),
                                 (void*)&bufferB):
        assert(status == CL SUCCESS):
133
        status = clSetKernelArg(kernel, 2, sizeof(cl_mem),
                                  (void*)(&bufferC[device])):
135
        assert(status == CL_SUCCESS);
        printf("Set kernel arguments completes\n");
```

### Kernel

- Now we are ready to launch kernel.
- Note that we need to supply the command queue to the device as a parameters.
- The same kernel is launched again and again for DEVICENUM times. The only difference is the buffers A and C used in launching the kernel.
- We associate an event with each kernel launching.

# Start Kernel

### Example 12: (matrixMul-time-copy-local-multidevice.c)

```
status =

clEnqueueNDRangeKernel(commandQueue[device],

kernel, 2, NULL,

globalThreads, localThreads,

0, NULL, &(events[device]));

assert(status == CL_SUCCESS);

}
```

# Wait for Completion

- Since kernel launching is non-blocking. we need to wait for them to complete.
- We simply use clWaitForEvents, and wait for all events to complete.

# Wait for Kernels

### Example 13: (matrixMul-time-copy-local-multidevice.c)

# Discussion

How many events do we need to wait for the whole thing to complete?

### Kernel

- The kernel function is very similar to the previous local memory version.
- We also use the constant DEVICENUM to denote the number of devices.

# Kernel

### Example 14: (mul-local-multidevice-kernel.cl)

```
2 #define N 1024
3 #define Blk 64
4 #define DEVICENUM 2
5 #define BSIDE (N / Blk)
```

# Kernel

#### Example 15: (mul-local-multidevice-kernel.cl)

```
7
   __kernel void mul(__global int A[N/DEVICENUM][N],
8
                      __global int B[N][N],
9
                      __global int C[N/DEVICENUM][N])
   {
10
11
     int globalRow = get_global_id(0);
12
     int globalCol = get_global_id(1);
13
     int localRow = get_local_id(0);
14
     int localCol = get_local_id(1);
15
16
     __local int ALocal[BSIDE][BSIDE];
17
     __local int BLocal[BSIDE][BSIDE];
```

# **NDRange**

- We declare A and C as N / DEVICENUM by N arrays, because the domain is partitioned among devices.
- The rest of the code is not changed at all!

### Reason

- The reason for this "no need to change" is that a kernel function is acting from a local view, i.e., it is a work item within a work group, so all the operations are still the same.
- The kernel still thinks from the point of view of blocks, and the related operations are done in local indices.
- ullet We only need the global index while accessing A, B, and C.
- Since the kernel is given the corresponding parts of A and B, so the operation is correct.

# Discussion

• Convince yourself that the code is correct.



# **Timing**

- We would like to know the detailed timing of the kernels.
- In particular we would like to prove the two kernel run *in* parallel.
- To do so we need the absolute time of the events from all devices.

### Base

- The time returned by the event is difficult to interpret.
- We would like to establish a relative time for inspection.
- We choose the time the first kernel entering the queue as the base time, and report the relative time to it for ease understanding.

### Example 16: (matrixMul-time-copy-local-multidevice.c)

- We loop through all devices and get timing information from all of them.
- This part is the same as before.

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### Example 17: (matrixMul-time-copy-local-multidevice.c)

```
for (int device = 0; device < DEVICENUM; device++) {</pre>
  cl_ulong timeEnterQueue, timeSubmit, timeStart,
    timeEnd;
  status =
    clGetEventProfilingInfo(events[device],
      CL_PROFILING_COMMAND_QUEUED,
      sizeof(cl_ulong), &timeEnterQueue, NULL);
  assert(status == CL_SUCCESS);
  status =
    clGetEventProfilingInfo(events[device],
      CL_PROFILING_COMMAND_SUBMIT,
      sizeof(cl_ulong), &timeSubmit, NULL);
  assert(status == CL_SUCCESS);
```

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### Example 18: (matrixMul-time-copy-local-multidevice.c)

```
status =
176
           clGetEventProfilingInfo(events[device],
             CL_PROFILING_COMMAND_START,
178
             sizeof(cl_ulong), &timeStart, NULL);
179
        assert(status == CL_SUCCESS);
        status =
           clGetEventProfilingInfo(events[device],
182
             CL_PROFILING_COMMAND_END,
183
              sizeof(cl_ulong), &timeEnd, NULL);
        assert(status == CL_SUCCESS);
184
```

## Relative Time

- In addition to queuing time, submission time, and execution, we also report the relative time when the four events happened.
- The relative time is in nanosecond.

### Print Time

### Example 19: (matrixMul-time-copy-local-multidevice.c)

```
186
        printf("\nkernel entered queue at %f\n",
187
                (timeEnterQueue - base) / NANO2SECOND);
188
        printf("kernel submitted to device at %f\n",
189
                (timeSubmit - base) / NANO2SECOND);
190
        printf("kernel started at %f\n",
191
                (timeStart - base) / NANO2SECOND):
192
        printf("kernel ended at %f\n",
193
                (timeEnd - base) / NANO2SECOND);
194
        printf("kernel queued time %f seconds\n",
195
                (timeSubmit - timeEnterQueue) / NANO2SECOND);
196
        printf("kernel submission time %f seconds\n",
197
                (timeStart - timeSubmit) / NANO2SECOND);
198
        printf("kernel execution time %f seconds\n",
199
              (timeEnd - timeStart) / NANO2SECOND);
      }
200
```

# Demonstration

 Run the matrixMul-time-copy-local-multidevice-cl program.



# Discussion

• Does the kernels run in parallel? Observe the results and give your answer.

# Dependency

- We may launch multiple kernels into multiple devices.
- Kernels may have dependency, and we need to ensure dependency among kernels.

# Example

- In the following example we compute G = (A + B) + (D + E), where all vector has length N.
- We compute C = A + B, then F = D + E, then we compute G = C + F. The last computation must wait for the first two to complete.
- We will use one device for each of the three additions.

# Constants

• We first declare the constants and variables.

### Example 20: (vectorAdd-dependency.c)

```
#define CL_USE_DEPRECATED_OPENCL_2_O_APIS
3
   #include <stdio.h>
4
   #include <assert.h>
5
   #include <CL/cl.h>
6
   #define N (65536 * 4)
   #define MAXGPU 10
8
   #define MAXK 1024
   #define MAXI.OG 4096
10
   #define DEVICENUM 3
11
   #define NANO2SECOND 100000000.0
12
13
   cl_uint A[N], B[N], C[N], D[N], E[N], F[N], G[N];
```

# Initialization

• We initialize the A, B, D and E vectors.

### Example 21: (vectorAdd-dependency.c)

# **Buffers**

- We then create buffers for kernel parameters.
- A and B are read only, and need to be copied into device.
- C is both read and write enable.

#### Example 22: (vectorAdd-dependency.c)

```
cl mem bufferA =
 clCreateBuffer(context,
                 CL MEM READ ONLY | CL MEM COPY HOST PTR.
                 N * sizeof(cl_uint), A, &status);
assert(status == CL_SUCCESS);
cl mem bufferB =
 clCreateBuffer(context.
                 CL MEM READ ONLY | CL MEM COPY HOST PTR.
                 N * sizeof(cl uint). B. &status):
assert(status == CL_SUCCESS);
cl mem bufferC =
 clCreateBuffer(context.
                 CL_MEM_READ_WRITE,
                 N * sizeof(cl uint), NULL, &status):
assert(status == CL_SUCCESS);
```

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# **Buffers**

• Similarly we create buffers for D, E and F.

#### Example 23: (vectorAdd-dependency.c)

```
cl mem bufferD =
 clCreateBuffer(context,
                 CL MEM READ ONLY | CL MEM COPY HOST PTR.
                 N * sizeof(cl_uint), D, &status);
assert(status == CL_SUCCESS);
cl mem bufferE =
 clCreateBuffer(context.
                 CL MEM READ ONLY | CL MEM COPY HOST PTR.
                 N * sizeof(cl uint), E, &status):
assert(status == CL_SUCCESS);
cl mem bufferF =
 clCreateBuffer(context,
                 CL_MEM_READ_WRITE,
                 N * sizeof(cl uint), NULL, &status):
assert(status == CL_SUCCESS);
```

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# **Buffers**

- We create buffers for G.
- This buffer is write only and use the host memory directly, so the host can print it directly.

#### Example 24: (vectorAdd-dependency.c)

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## Discussion

• What are the characteristics for read only, read and write, and write only buffers respectively?

# **Buffers**

- Both NDRange and the work group are one dimensional.
- The size of NDRange is N.
- The size of a work group is 256.

#### Example 25: (vectorAdd-dependency.c)

```
size_t globalThreads[] = {(size_t)N};
size_t localThreads[] = {256};
```

$$C = A + B$$

- Now we launch the first kernel for C = A + B.
- We have already built three command queues for three devices.
- We use commandQueue[0] for the first addition.
- We also declared three events to denote the completion of the three kernels that we will launch.

#### Example 26: (vectorAdd-dependency.c)

```
131
       cl event events[3]:
132
       status = clSetKernelArg(kernel, 0, sizeof(cl_mem),
133
                                (void*)&bufferA):
134
       assert(status == CL_SUCCESS);
135
       status = clSetKernelArg(kernel, 1, sizeof(cl_mem),
136
                                (void*)&bufferB):
137
       assert(status == CL_SUCCESS);
138
       status = clSetKernelArg(kernel, 2, sizeof(cl_mem),
139
                                (void*)&bufferC):
140
       assert(status == CL_SUCCESS);
141
142
       status =
143
         clEnqueueNDRangeKernel(commandQueue[0], kernel, 1, NULL,
144
                                 globalThreads, localThreads,
145
                                 0, NULL, &(events[0]));
146
       assert(status == CL SUCCESS):
```

$$F = D + E$$

- Similarly we launch the second kernel for F = D + E.
- Note that we used commandQueue[1] and events[1].

#### Example 27: (vectorAdd-dependency.c)

```
148
       status = clSetKernelArg(kernel, 0, sizeof(cl_mem),
149
                                (void*)&bufferD):
150
       assert(status == CL SUCCESS):
151
       status = clSetKernelArg(kernel, 1, sizeof(cl_mem),
152
                                (void*)&bufferE):
153
       assert(status == CL SUCCESS):
154
       status = clSetKernelArg(kernel, 2, sizeof(cl_mem),
155
                                (void*)&bufferF):
156
       assert(status == CL SUCCESS):
157
       status =
158
         clEnqueueNDRangeKernel(commandQueue[1], kernel, 1, NULL,
159
                                 globalThreads, localThreads,
160
                                 0, NULL, &(events[1]));
161
       assert(status == CL SUCCESS):
```

$$G = C + F$$

- Finally we launch the third kernel for G = C + F.
- Note that we used commandQueue[2] and events[2].

#### Example 28: (vectorAdd-dependency.c)

```
163
       status = clSetKernelArg(kernel, 0, sizeof(cl_mem),
164
                                (void*)&bufferC);
165
       assert(status == CL SUCCESS):
166
       status = clSetKernelArg(kernel, 1, sizeof(cl_mem),
167
                                (void*)&bufferF):
168
       assert(status == CL SUCCESS):
169
       status = clSetKernelArg(kernel, 2, sizeof(cl_mem),
170
                                (void*)&bufferG):
171
       assert(status == CL SUCCESS):
172
       status =
173
         clEnqueueNDRangeKernel(commandQueue[2], kernel, 1, NULL,
174
                                 globalThreads, localThreads,
175
                                 0, NULL, &(events[2]));
176
       assert(status == CL SUCCESS):
```

# Wait for Events

• We wait for all three kernel to finish.

#### Example 29: (vectorAdd-dependency.c)

```
178
       status = clWaitForEvents(DEVICENUM, events);
179
       assert(status == CL SUCCESS):
180
    #ifdef COPYG
181
       clEnqueueReadBuffer(commandQueue[0], bufferG, CL_TRUE,
182
                            0, N * sizeof(cl_uint), G,
183
                            O, NULL, NULL);
184
    #endif
185
       printf("All three kernels complete.\n");
```

### Wait for Events

- The rest of the code just prints the timing information, checks for correctness, and releases resources.
- Please refer to previous discussion.

## Demonstration

• Run the vectorAdd-dependency-cl.

## Discussion

- Does the program produce the correct answer?
- If not, what is the possible reason?

### **Problem**

- The previous program waits for all three kernels to complete, but the third kernel did *not* wait for the first two.
- We will use the wait for events mechanism to launch the third kernel.

#### Prototype 30: clEnqueueNDRangeKernel.h

```
cl int
   clEnqueueNDRangeKernel (cl_command_queue command_queue,
3
                             cl_kernel kernel,
4
                             cl uint work dim.
5
                             const size_t *global_work_offset,
6
                             const size_t *global_work_size,
7
                             const size t *local work size.
8
                             cl_uint num_events_in_wait_list,
9
                             const cl_event *event_wait_list,
10
                             cl_event *event);
```

$$G = C + F$$

- We launch the third kernel and wait for the first two events in the events array.
- Set num\_events\_in\_wait\_list to 2 and event\_wait\_list to events.

#### Example 31: (vectorAdd-dependency-correct.c)

```
163
       status = clSetKernelArg(kernel, 0, sizeof(cl_mem),
164
                                (void*)&bufferC):
165
       assert(status == CL SUCCESS):
166
       status = clSetKernelArg(kernel, 1, sizeof(cl_mem),
167
                                (void*)&bufferF):
168
       assert(status == CL SUCCESS):
169
       status = clSetKernelArg(kernel, 2, sizeof(cl_mem),
170
                                (void*)&bufferG):
171
       assert(status == CL SUCCESS):
172
       status =
173
         clEnqueueNDRangeKernel(commandQueue[2], kernel, 1, NULL,
174
                                 globalThreads, localThreads,
175
                                 2, events, &(events[2]));
176
       assert(status == CL SUCCESS):
```

## Demonstration

• Run the vectorAdd-dependency-correct-cl.

### Discussion

- Does the program produce the correct answer?
- Observe the timing and make sure the first two kernels run in parallel, and the third kernel will wait for the first two.

### clFinish

- The host can also explicitly wait for the first two kernels to finish before launching the third kernel by clFinish.
- We just need to wait for the command queues to the first two devices to become empty.

#### Example 32: (vectorAdd-dependency-clfinish.c)

```
clFinish(commandQueue[0]);
168
       clFinish(commandQueue[1]);
       status = clSetKernelArg(kernel, 0, sizeof(cl_mem),
170
                                (void*)&bufferC):
       assert(status == CL_SUCCESS);
       status = clSetKernelArg(kernel, 1, sizeof(cl_mem),
173
                                (void*)&bufferF):
174
       assert(status == CL SUCCESS):
       status = clSetKernelArg(kernel, 2, sizeof(cl_mem),
                                (void*)&bufferG):
177
       assert(status == CL SUCCESS):
178
       status =
         clEnqueueNDRangeKernel(commandQueue[0], kernel, 1, NULL,
180
                                 globalThreads, localThreads,
                                 0, NULL, &(events[2]));
       assert(status == CL SUCCESS):
```

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## Demonstration

• Run the vectorAdd-dependency-clfinish-cl program.

### Discussion

- Does the program produce the correct answer?
- Observe the timing and make sure the first two kernels run in parallel, and the third kernel will wait for the first two.
- Compare the timing results between vectorAdd-dependency-correct-cl and vectorAdd-dependency-clfinish-cl, especially in when the last kernel joined the command queue.

# **Group Size**

- The previous matrix multiplication program has extremely good performance due to two reasons.
  - It uses local memory to speed up data access.
  - It has a larger group size (256).
- Now we would like to study the effects of group size on performance.

# Compute Units

- A device has only a limited number of compute units.
- The work items of a work group will occupy a compute unit.
- For a given number of work items, if the work group size is small, then the number of work groups becomes large, and some work groups will wait for compute unit.
- In other words, a small work group size limits the the number of work items that we can process in parallel.

### Classroom

- A school has only a limited number of classrooms.
- The students are divided into group, and one group will occupy a classroom.
- For a given number of students, if the group size is small, then the number of groups becomes large, and some groups will wait for classroom.
- In other words, a small group size limits the the number of students that can study in parallel.

## Discussion

• Give your own example of this phenomenon.

# Compute Unit Number

- We need only one GPU device.
- We call clGetDeviceInfo with CL\_DEVICE\_MAX\_COMPUTE\_UNITS to get the number of compute units on GPU.
- We need the number of compute units to determine the work group size.

### Example 33: (vectorAdd-groupsize.c)

```
24
     cl_device_id GPU[MAXGPU];
25
     cl_uint GPU_id_got;
26
     status =
27
        clGetDeviceIDs(platform_id, CL_DEVICE_TYPE_GPU,
28
                       MAXGPU, GPU, &GPU_id_got);
29
     assert(status == CL_SUCCESS && GPU_id_got >= 1);
30
     printf("There are %d GPU devices\n", GPU_id_got);
31
     cl_uint unit;
32
     status =
33
       clGetDeviceInfo(GPU[0], CL_DEVICE_MAX_COMPUTE_UNITS,
34
                        sizeof(cl_uint), &unit, NULL);
35
     assert(status == CL_SUCCESS);
36
     printf("# of compute units is %d\n", unit);
```

# **Group Size**

- We fix the number of work items to *N*, and vary the size of the work group from 1 to 256.
- We wait for the event before launching the next kernel.
- We will record the timing information into a file vectorAdd-groupsize.dat.

### Example 34: (vectorAdd-groupsize.c)

```
110
       size_t workItem[] = {(size_t)N};
111
       FILE *timefp = fopen("vectorAdd-groupsize.dat", "w");
112
       assert(timefp != NULL);
113
       for (int groupSize = 1; groupSize <= 256; groupSize *= 2)</pre>
114
         cl event event:
115
         size t localSize[1]:
116
         localSize[0] = groupSize;
117
         status =
118
           clEnqueueNDRangeKernel (commandQueue, kernel, 1, NULL,
119
                                   workItem, localSize,
120
                                   O, NULL, &event);
121
         assert(status == CL_SUCCESS);
122
         /* waitforevent */
123
         status = clWaitForEvents(1, &event);
124
         assert(status == CL_SUCCESS);
125
         printf("The kernel with group size %d completes.\n",
126
                groupSize);
```

# Get Timing Information

- The rest of the code will retrieve timing information from the event.
- We output the group size and the execution time into a file vectorAdd-groupsize.dat.

### Example 35: (vectorAdd-groupsize.c)

```
printf("\nkernel entered queue at %f\n",
         (timeEnterQueue - base) / NANO2SECOND);
  printf("kernel submitted to device at %f\n".
         (timeSubmit - base) / NANO2SECOND);
  printf("kernel started at %f\n".
         (timeStart - base) / NANO2SECOND):
  printf("kernel ended at %f\n",
         (timeEnd - base) / NANO2SECOND):
  printf("kernel queued time %f seconds\n",
         (timeSubmit - timeEnterQueue) / NANO2SECOND);
  printf("kernel submission time %f seconds\n".
         (timeStart - timeSubmit) / NANO2SECOND);
  printf("kernel execution time %f seconds\n",
         (timeEnd - timeStart) / NANO2SECOND):
  fprintf(timefp, "%d %f\n", groupSize,
          (timeEnd - timeStart) / NANO2SECOND);
}
```

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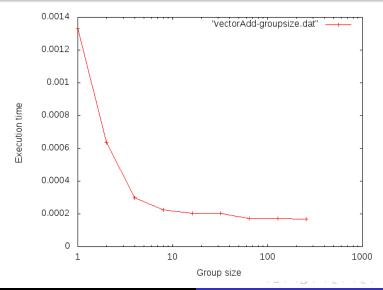
173

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### Demonstration

- We plot the kernel execution time for different group sizes.
- The x-xis is in log scale.

# **Execution Time**



### Discussion

• Observe the execution time and draw you conclusion about the effects of group size on performance.