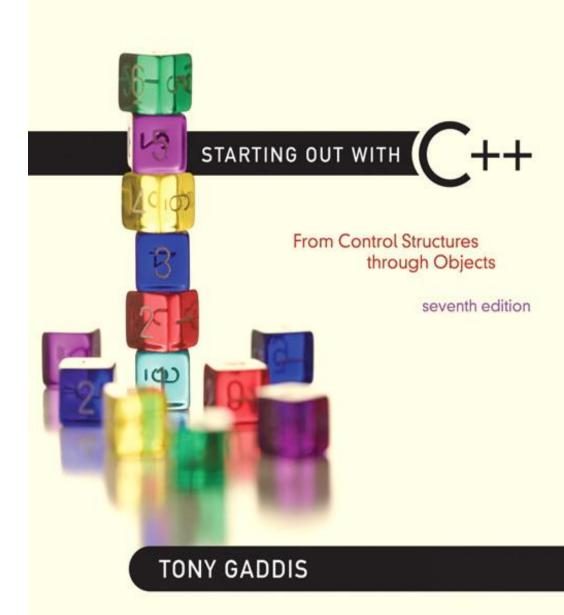
#### **Chapter 20:**

**Binary Trees** 



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20.1

## Definition and Application of Binary Trees

#### Definition and Application of Binary Trees

- Binary tree:
  - a nonlinear linked list in which each node may point to 0, 1, or two other nodes

 Each node contains one or more data fields and two pointers Tree pointer

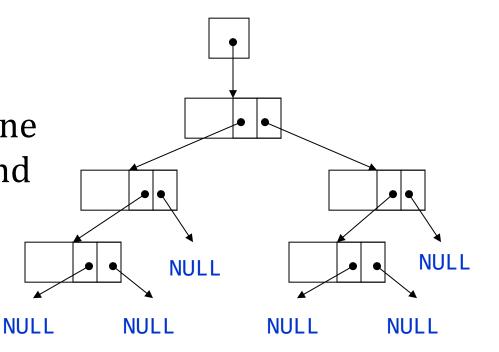
Root node

• Parent

• Child (children)

• Leaf node

• Subtree (left/right)

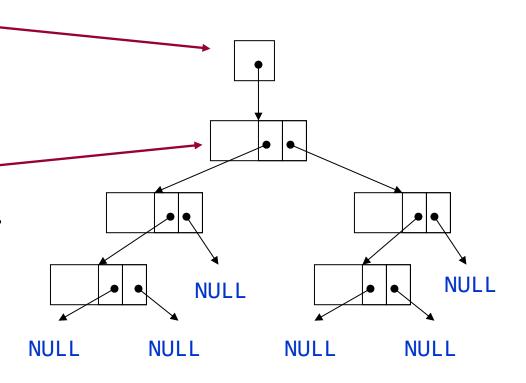


#### Tree pointer:

 like a head pointer for a linked list, it points to the first node in the binary tree

#### Root node:

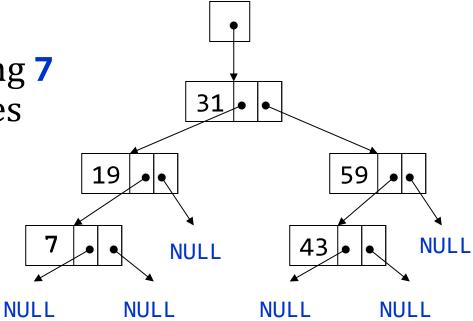
 the node at the top of the tree



#### Leaf nodes:

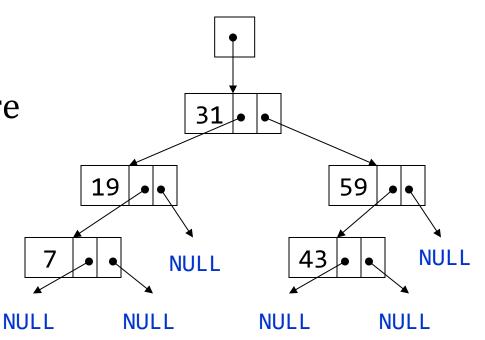
 nodes that have no children

The nodes containing 7 and 43 are leaf nodes



- Child nodes (children):
  - nodes below a given node

The children of the node containing 31 are the nodes containing 19 and 59



- Parent node:
  - node above a given node

The parent of the node containing 43 is the node containing 59

NULL

**NULL** 

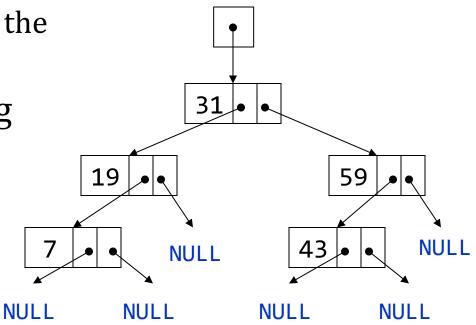
NULL

**NULL** 

#### • Subtree:

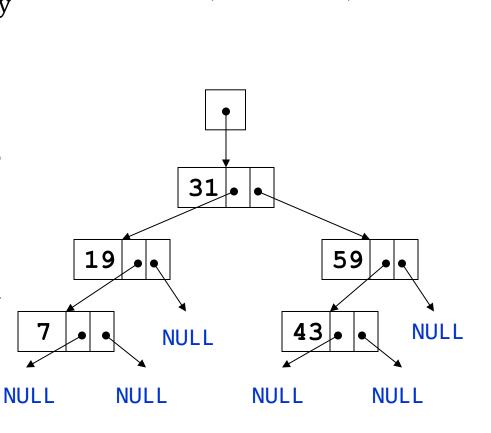
 the portion of a tree from a node down to the leaves

The nodes containing 19 and 7 are the left subtree of the node containing 31



## Uses of Binary Trees

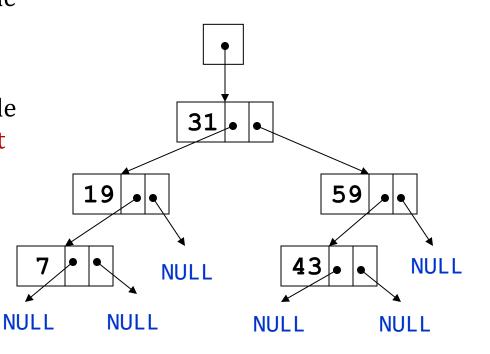
- Binary search tree:
  - data is organized in a binary tree to simplify (speed up ) searches
- Left subtree of a node contains data values < the data in the node
- Right subtree of a node contains values > the data in the node



## Searching in a Binary Tree

- 1) Start at root node
- 2) Examine node data:
  - a) Is it the desired value (=)? Done
  - b) Else, is the desired data < node data? Repeat step 2 with left subtree
  - c) Else, is the desired data > node data? Repeat step 2 with right subtree
- 3) Continue until the desired value is found or a NULL pointer is reached

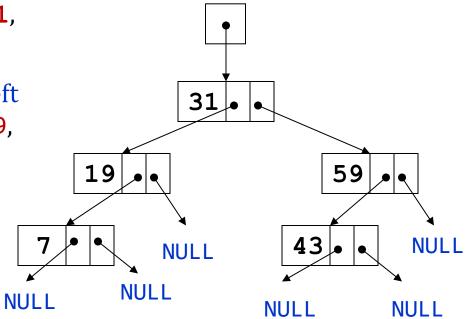
data stored in ascending sequence



## Searching in a Binary Tree

#### To locate the node containing 43,

- Examine the root node (31) first
- Since 43 > 31, examine the right child of the node containing 31,
   (59)
- Since 43 < 59, examine the left child of the node containing 59,</li>
   (43)
- The node containing
   43 has been found



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20.2

#### Binary Search Tree Operations

#### Binary Search Tree Operations

- Create a binary search tree
  - organize data into a binary search tree
- Insert a node into a binary tree
  - put node into tree in its correct position to maintain order
- Search for a node in a binary tree
  - locate a node with particular data value
- Delete a node from a binary tree
  - remove a node and adjust links to maintain binary tree

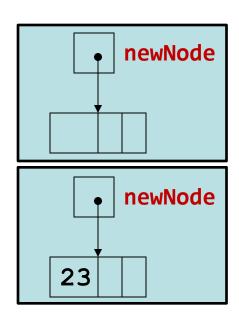
#### Binary Search Tree Node

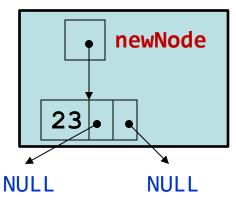
 A node in a binary tree is like a node in a linked list with two node pointer fields:

```
struct TreeNode
{
  int value;
  TreeNode *left;
  TreeNode *right;
}
```

#### Creating a New Node

- Allocate memory for new node: TreeNode \*newNode = new TreeNode;
- 2. Initialize the contents of the node:
  newNode->value = num;
- 3. Set the pointers to NULL:
  - newNode->Left
  - = newNode->Right
  - = NULL;

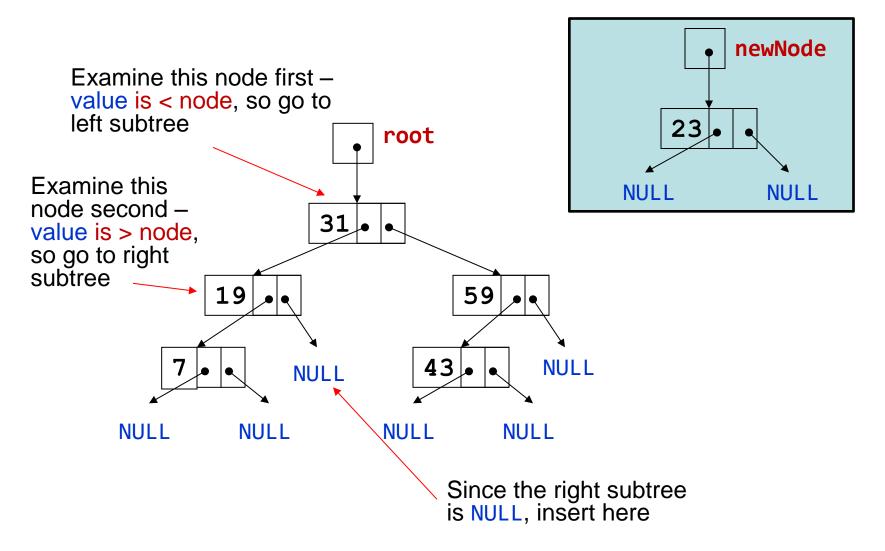




### Inserting a Node in aBinary Search Tree

- 1) If the tree is empty, insert the new node as the root node
- Else, compare new node against left or right child, depending on whether data value of new node is
   or > root node
- 3) Continue comparing and choosing left or right subtree until the NULL pointer is found
- 4) Set this **NULL** pointer to point to the new node

### Inserting a Node in a Binary Search Tree

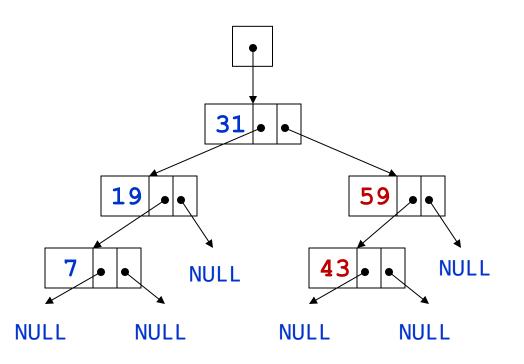


#### Traversing a Binary Tree

#### Three traversal methods:

- 1) <u>Inorder</u>: left / data / right
  - a) Traverse left subtree of node
  - b) Process data in node
  - c) Traverse right subtree of node
- 2) <u>Preorder</u>: data / left / right
  - a) Process data in node
  - b) Traverse left subtree of node
  - c) Traverse right subtree of node
- 3) <u>Postorder</u>: left / right / data
  - a) Traverse left subtree of node
  - b) Traverse right subtree of node
  - c) Process data in node

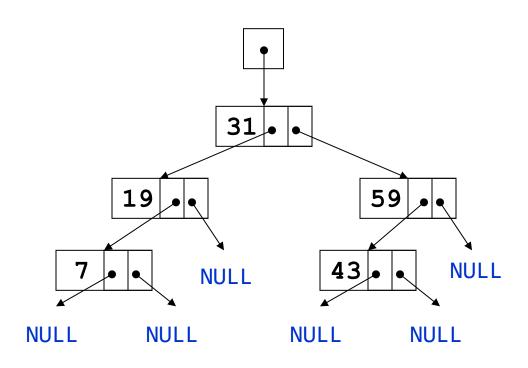
#### Traversing a Binary Tree



TRAVERSAL METHOD	NODES VISITED IN ORDER
Inorder	7, 19, 31,
L/D/R	43, 59
Preorder	31, 19, 7,
D/L/R	59, 43
Postorder	7, 19, 43,
L/R/D	59, 31

## Searching in a Binary Tree

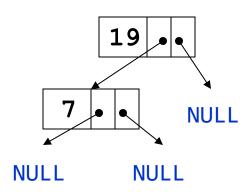
- Start at the root node, traverse the tree looking for a value
- Stop when the value is found or a NULL pointer is detected
- Can be implemented as a bool function



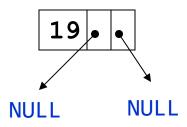
```
Search for 43 -> returns true
Search for 17 -> returns false
```

### Deleting a Node from a Binary Tree - Leaf Node

 If the node to be deleted is a leaf node, replace the parent node's pointer to it with a NULL pointer, then delete the node



Deleting node with 7 – before deletion

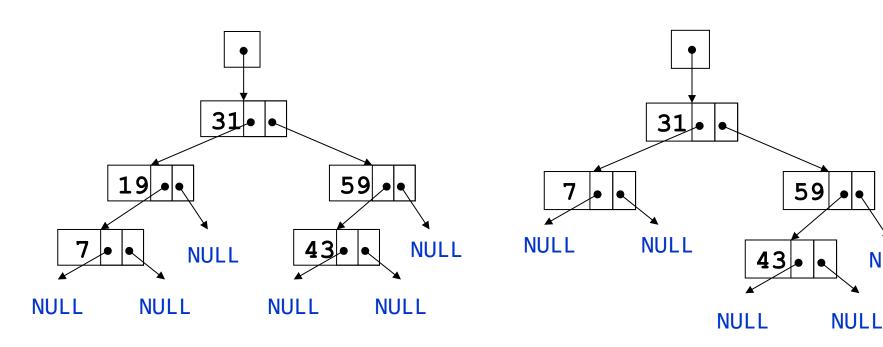


Deleting node with 7 – after deletion

#### Deleting a Node from a Binary Tree - One Child

 If the node to be deleted has one child node, adjust the pointers so that the parent of the node to be deleted points to the child of node to be deleted, then delete the node

#### Deleting a Node from a Binary Tree - One Child



Deleting node with 19 – before deletion

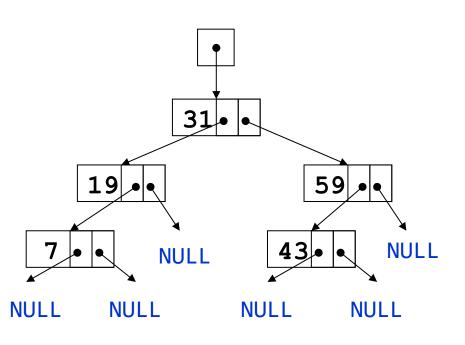
Deleting node with 19 – after deletion

**NULL** 

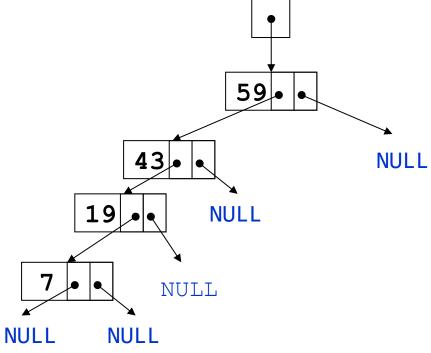
# Deleting a Node from a Binary Tree - Two Children

- If the node to be deleted has both left and right children,
  - 'Promote' one child to take the place of the deleted node
  - Locate correct position for other child in subtree of promoted child
- The textbook convention: promote the right child, then position the left subtree underneath

# Deleting a Node from a Binary Tree - Two Children



Deleting node with 31 – before deletion



Deleting node with 31 – after deletion

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## Template Considerations for Binary Search Trees

#### Template Considerations for Binary Search Trees

 A binary tree can be implemented as a template allowing flexibility in determining the type of data stored

 The implementation must support relational operators >, <, and == to allow comparison of nodes