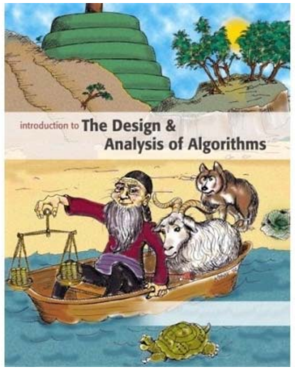
Introduction to

Algorithm Design and Analysis

[3] Recursion



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In the Last Class …

- **Asymptotic growth rate**
  - O, Ω, Θ
  - o, ω
- **Brute force algorithms**
  - By iteration
  - By recursion

Recursion

- **Recursion in algorithm design**
  - The divide and conquer strategy
  - Proving the correctness of recursive procedures
- **Solving recurrence equations**
  - Some elementary techniques
  - Master theorem

Recursion in Algorithm Design

- **Computing n! with Fac(n)**
  - o if n=1 then return 1 else return Fac(n-1)*n

M(1)=0 and M(n)=M(n-1)+1 for n>0

(critical operation: multiplication)

- **Hanoi Tower**
  - o if n=1 then move d(1) to peg3 else

    Hanoi(n-1, peg1, peg2); move d(n) to peg3; Hanoi(n-1, peg2, peg3)

M(1)=1 and M(n)=2M(n-1)+1 for n>1

(critical operation: move)

Recursion in Algorithm Design

- **Counting the Number of Bits**
  - Input: a positive decimal integer *n*
  - Output: the number of binary digits in *n*'s binary representation

```
Int BitCounting (int n)
If(n==1) return 1;
Else
  return BitCounting(n div 2) +1;
```

$$T(n) = \begin{cases} 0 & n = 1 \\ T(\lfloor n/2 \rfloor) + 1 & n > 1 \end{cases}$$

Divide and Conquer

- **Divide**
  - Divide the “big” problem to smaller ones
- **Conquer**
  - Solve the “small” problems by recursion
- **Combine**
  - Combine results of small problems, and solve the original problem

Divide and Conquer

The general pattern

```
solve($I$)
    $n$=size($I$);
    if ($n$≤smallSize)
        solution=directlySolve($I$);
    else
        divide $I$ into $I_1$,… $I_k$;
        for each $i$∈{1,…,k}
            $S_i$=solve($I_i$);
        solution=combine($S_1$ ,… ,$S_k$);
    return solution
```

$T(n)=B(n)$ for $n \le smallSize$

$$T(n)=D(n)+\sum_{i=1}^{k} T(size(I_i))+C(n)$$

for $n>smallSize$

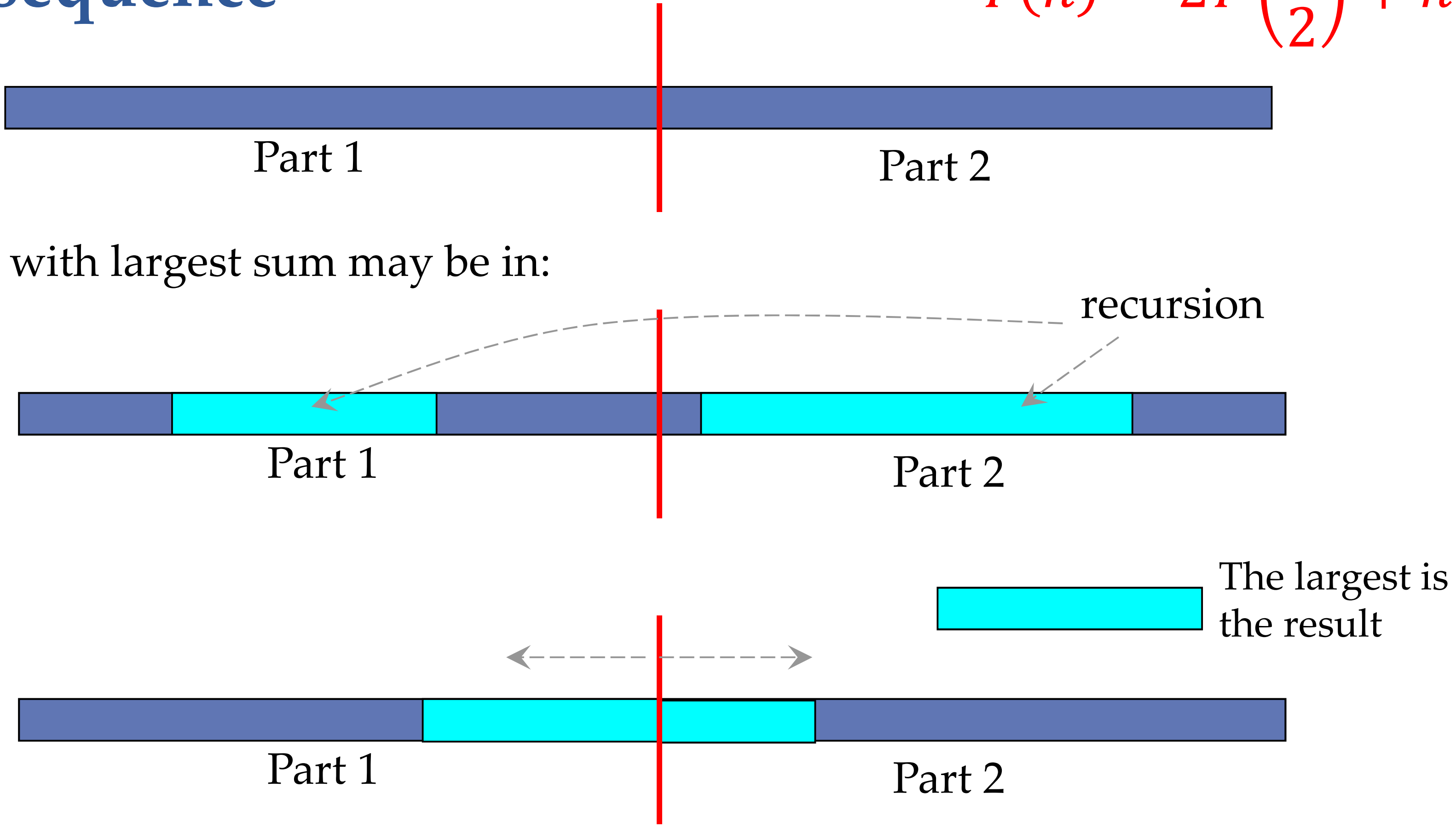
Divide Conquer

- **The BF recursion**
  - Problem size: often decreases linearly
    - "n, n-1, n-2, …"
- **The D&C recursion**
  - Problem size: often decrease exponentially
    - "n, n/2, n/4, n/8, …"

Examples

Max sum subsequence

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$

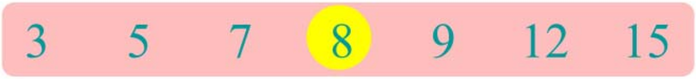

Part 1
Part 2
the sub with largest sum may be in:
recursion
Part 1
Part 2
or:
The largest is the result
Part 1
Part 2

Examples

- **Maxima**
- **Frequent element**
- **Multiplication**
  - Integer
  - Matrix
- **Nearest point pair**

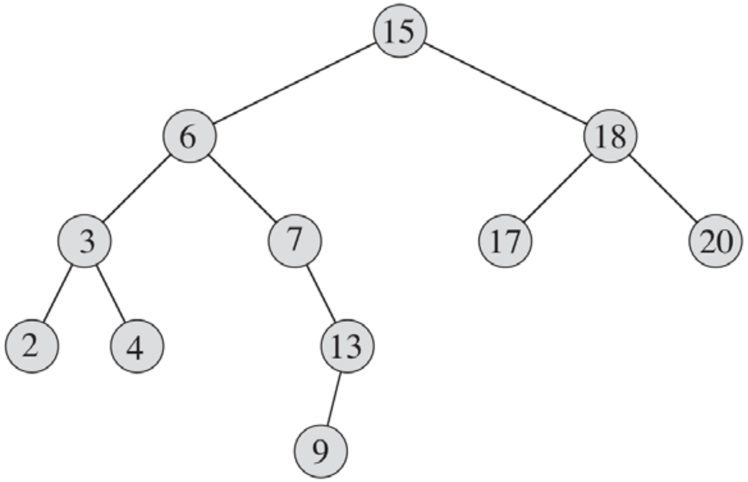
Examples

- Arrays



3 5 7 8 9 12 15

- Trees



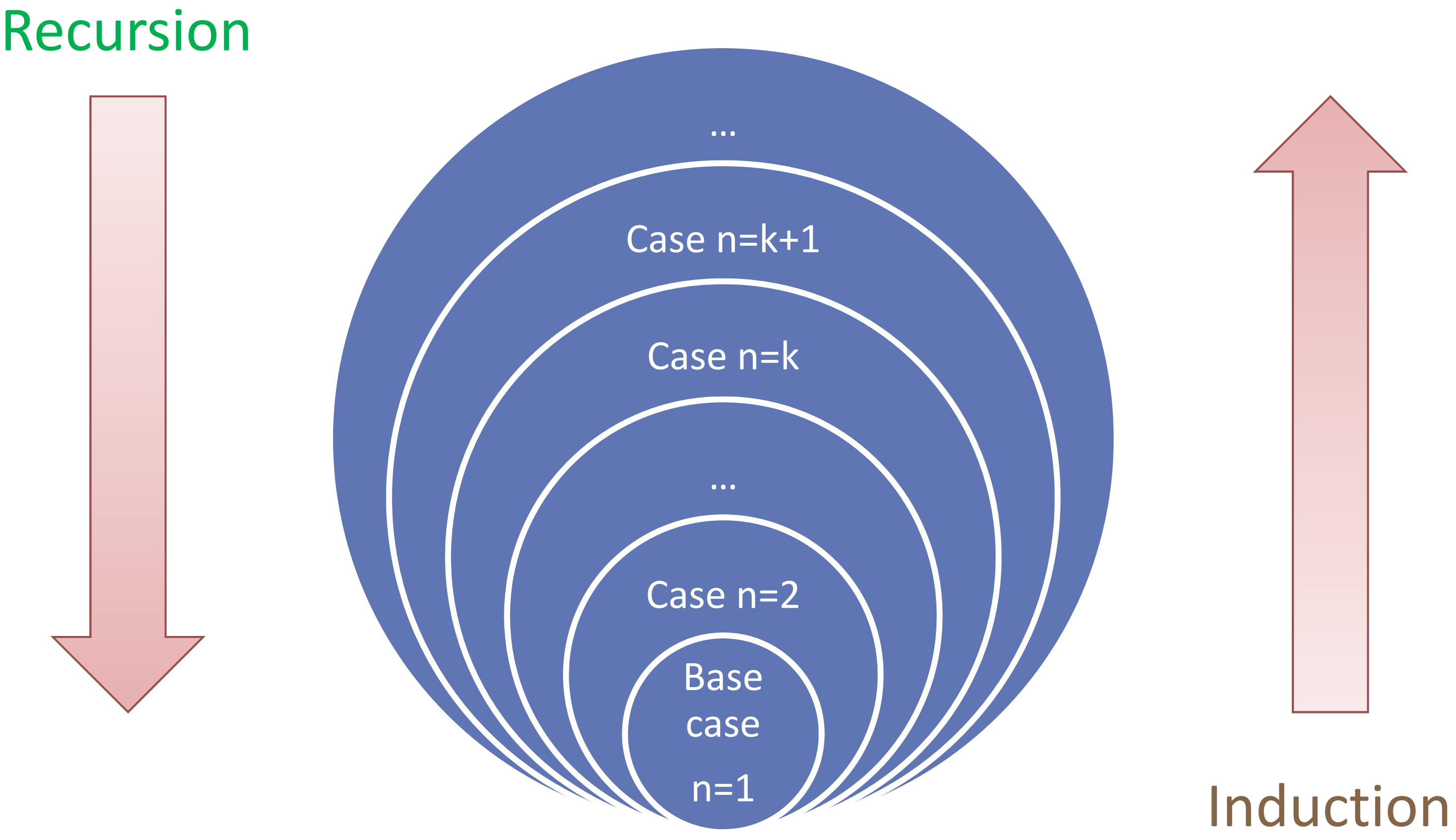
Workhorse

- “Hard division, easy combination”

- “Easy division, hard combination”

Usually, the “real work” is in one part.

Correctness of Recursion



Analysis of Recursion

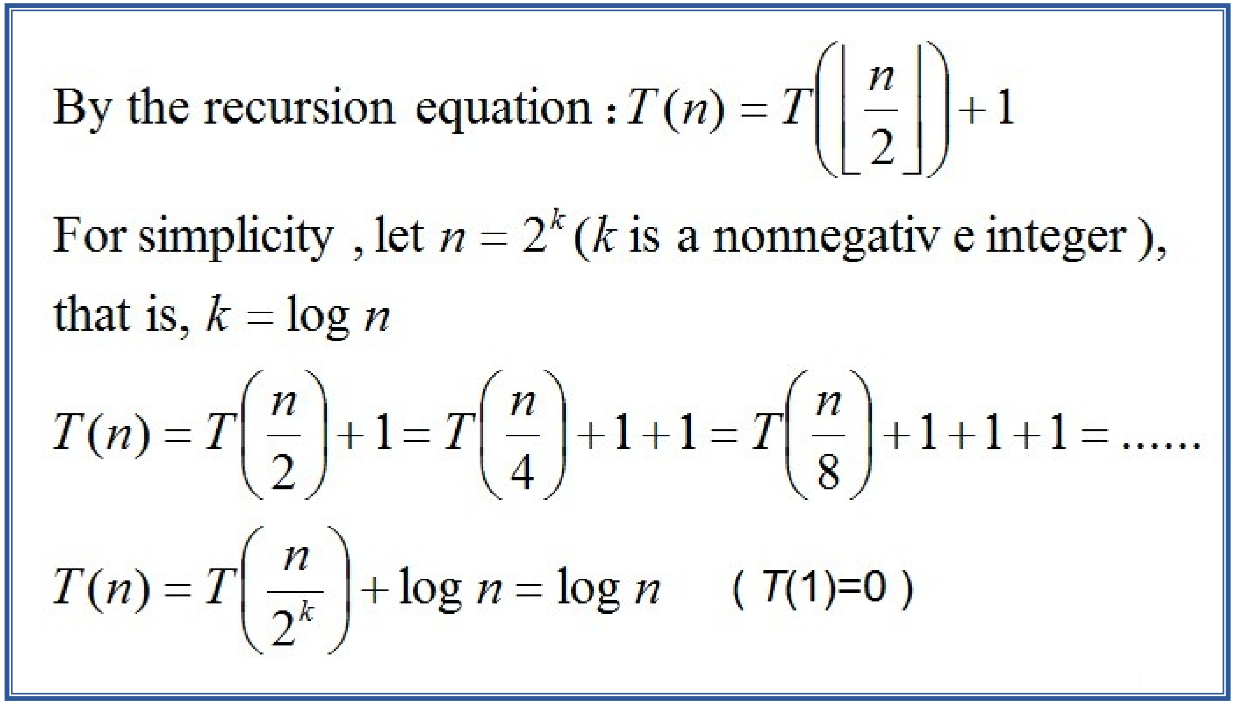
- **Solving recurrence equations**

- **E.g., Bit counting**
  - Critical operation: add
  - The recurrence relation

$$T(n) = \begin{cases} 0 & n = 1 \\ T(\lfloor n/2 \rfloor) + 1 & n > 1 \end{cases}$$

Analysis of Recursion

- **Backward substitutions**



By the recursion equation : $T(n) = T\left(\left\lfloor \frac{n}{2} \right\rfloor\right) + 1$

For simplicity , let $n = 2^k$ (k is a nonnegativ e integer), that is, $k = \log n$

$$T(n) = T\left(\frac{n}{2}\right) + 1 = T\left(\frac{n}{4}\right) + 1 + 1 = T\left(\frac{n}{8}\right) + 1 + 1 + 1 = \ldots\ldots$$

$$T(n) = T\left(\frac{n}{2^k}\right) + \log n = \log n \qquad (\ T(1)=0\ )$$

Smooth Functions

- $f(n)$
  - Nonnegative eventually non-decreasing function defined on the set of natural numbers
- **$f(n)$ is called smooth**
  - If $f(2n) \in \Theta(f(n))$.

- **Examples of smooth functions**
  - $\log n$, n, $n\log n$ and n^{α} ($\alpha \geq 0$)
  - E.g., $2n\log 2n = 2n(\log n + \log 2) \in \Theta(n\log n)$

Even Smoother

- **Let $f(n)$ be a smooth function, then, for any fixed integer $b \geq 2$, $f(bn) \in \Theta(f(n))$.**
  - That is, there exist positive constants c_b and d_b and a nonnegative integer n_0 such that

$$d_b f(n) \leq f(bn) \leq c_b f(n) \quad \text{for } n \geq n_0.$$

It is easy to prove that the result holds for $b = 2^k$,

for the second inequality :

$f(2^k n) \leqslant c_2^k f(n)$ *for* $k = 1,2,3\ldots$ *and* $n \geqslant n_0$.

For an arbitrary integer $b \geqslant 2, 2^{k-1} \leqslant b \leqslant 2^k$

Then, $f(bn) \leqslant f(2^k n) \leqslant c_2^k f(n)$, *we can use* c_2^k *as* c_b.

Smoothness Rule

- **Let $T(n)$ be an eventually non-decreasing function and $f(n)$ be a smooth function.**
  - If $T(n)\in\Theta(f(n))$ for values of n that are powers of $b(b\geq 2)$, then $T(n)\in\Theta(f(n))$.

Just proving the big-Oh part:

By the hypothsis: $T(b^k)\leqslant cf(b^k)$ *for* $b^k\geqslant n_0$.

By the prior result: $f(bn)\leqslant c_b f(n)$ *for* $n\geqslant n_0$.

Let $n_0\leqslant b^k\leqslant n\leqslant b^{k+1}$,

$$T(n)\underset{\text{Non-decreasing}}{\leqslant} T(b^{k+1})\underset{\text{hypothesis}}{\leqslant} cf(b^{k+1})=cf(bb^k)\underset{\text{Prior result}}{\leqslant} cc_b f(b^k)\underset{\text{Non-decreasing}}{\leqslant} cc_b f(n)$$

Computing the Fibonacci Number

$T(0)=0$

$T(1)=1$

$T(n)= T(n-1)+T(n-2)$

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...

$$a_n = r_1 a_{n-1} + r_2 a_{n-2} + \cdots + r_k a_{n-k}$$

is called linear homogeneous relation of degree *k*.

For the special case of Fibonacci: $a_n=a_{n-1}+a_{n-2}$, $r_1=r_2=1$

Characteristic Equation

- For a linear homogeneous recurrence relation of degree *k*

$$a_n = r_1 a_{n-1} + r_2 a_{n-2} + \cdots + r_k a_{n-k}$$

the polynomial of degree *k*

$$x^k = r_1 x^{k-1} + r_2 x^{k-2} + \cdots + r_k$$

is called its characteristic equation.

- The characteristic equation of linear homogeneous recurrence relation of degree 2 is:

$$\boxed{x^2 - r_1 x - r_2 = 0}$$

Solution of Recurrence Relation

- If the characteristic equation $x^2 - r_1 x - r_2 = 0$ of the recurrence relation $a_n = r_1 a_{n-1} + r_2 a_{n-2}$ has two distinct roots s_1 and s_2, then

$$a_n = u s_1^n + v s_2^n$$

where u and v depend on the initial conditions, is the explicit formula for the sequence.

- If the equation has a single root s, then, both s_1 and s_2 in the formula above are replaced by s

Back to Fibonacci Sequence

$f_0=0$

$f_1=1$

$f_n= f_{n-1}+ f_{n-2}$

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ……

Explicit formula for Fibonacci Sequence

The characteristic equation is x^2-x-1=0, which has roots:

$$s_1 = \frac{1+\sqrt{5}}{2} \quad and \quad s_2 = \frac{1-\sqrt{5}}{2}$$

Note: (by initial conditions) $f_1 = us_1 + vs_2 = 1 \;\; and \;\; f_2 = us_1^{\,2} + vs_2^{\,2} = 1$

which means:

$$f_n = \frac{1}{\sqrt{5}}\left(\frac{1+\sqrt{5}}{2}\right)^n - \frac{1}{\sqrt{5}}\left(\frac{1-\sqrt{5}}{2}\right)^n$$

Guess and Prove

- **Example: T(n)=2T(⌊n/2⌋) +n**
- **Guess**
  - o $T(n) \in O(n)$?
    - $T(n) \leq cn$, to be pr
  - o $T(n) \in O(n^2)$?
    - $T(n) \leq cn^2$, to be prove
  - o **Or maybe**, $T(n) \in O(n\log$
    - $T(n) \leq cn\log n$, to be prove
- **Prove**
  - o by substitution

Try to prove $T(n) \leq cn$:

However:

$T($

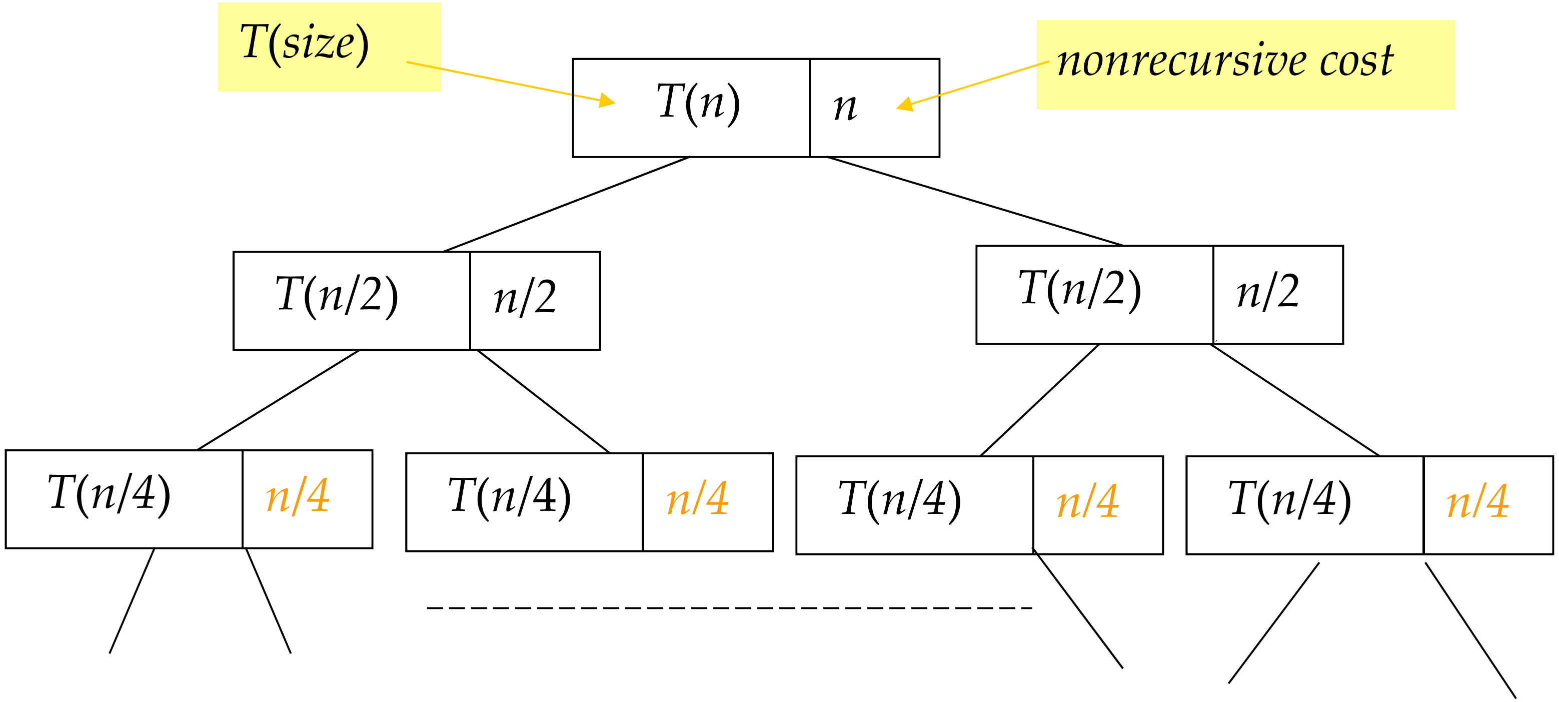
$$
\begin{aligned}
T(n) &= 2T(\lfloor n/2 \rfloor) + n \\
&\leq 2(c\lfloor n/2 \rfloor \log(\lfloor n/2 \rfloor)) + n \\
&\leq cn\log(n/2) + n \\
&= cn\log n - cn\log 2 + n \\
&= cn\log n - cn + n \\
&\leq cn\log n \quad \text{for } c \geq 1
\end{aligned}
$$

Divide and Conquer Recursions

- **Divide and conquer**
  - Divide the "big" problem to smaller ones
  - Solve the "small" problems by recursion
  - Combine results of small problems, and solve the original problem
- **Divide and conquer recursion**

$$T(n) = \underbrace{b}_{\text{divide}}\ \underbrace{T}_{\text{conquer}}(n/\underbrace{c}_{\text{divide}}) + \underbrace{f(n)}_{\text{combine}}$$

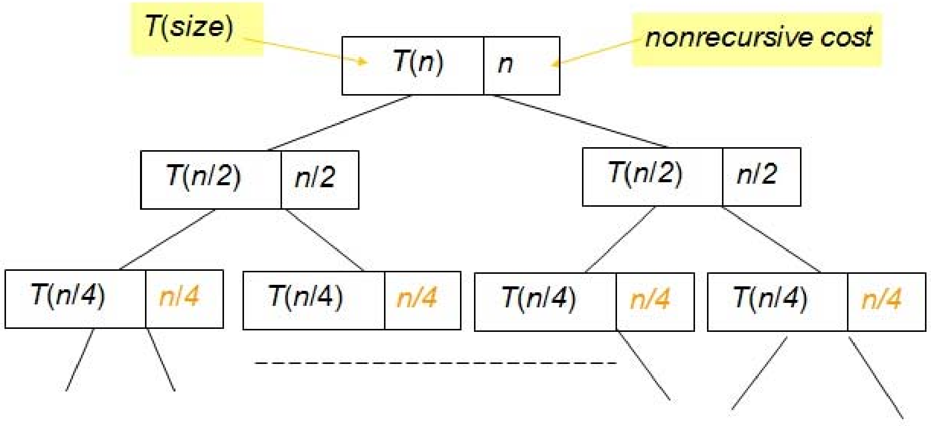
Recursion Tree



The recursion tree for T(n) = 2T(n/2) + n

Recursion Tree

- **Node**
  - Non-leaf
    - Non-recursive cost
    - Recursive cost
  - Leaf
    - Base case
- **Edge**
  - Recursion



The recursion tree for T(n)=T(n/2)+T(n/2)+n

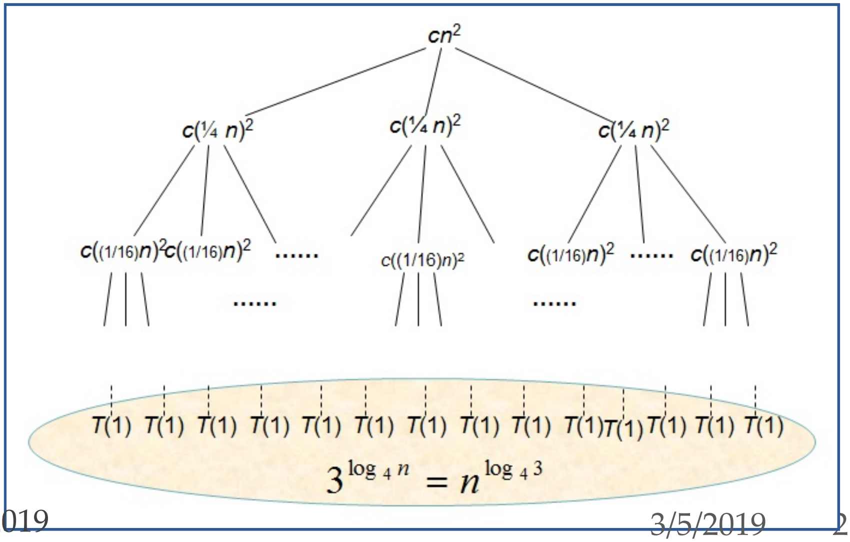
Recursion Tree

- $T(n) = \overbrace{\underbrace{3}_{\text{\# of sub-problems}}T(n/\underbrace{4}_{\text{size of sub-problems}})}^{\text{Recursive cost}} + \overbrace{\Theta\,(n^2)}^{\text{Non-recursive cost}}$

- **Total cost**

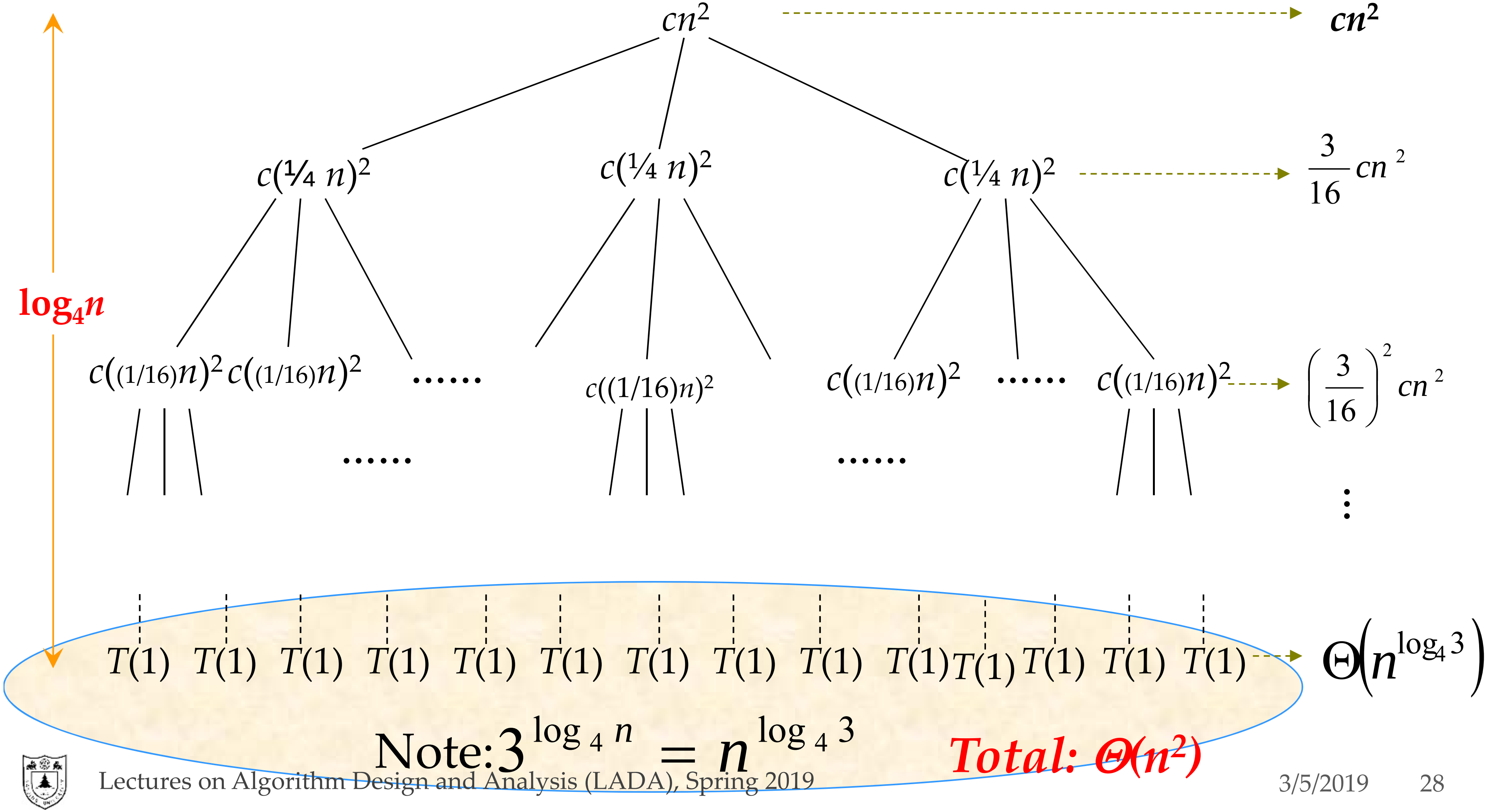
$$\Sigma$$

Sum of row sums



Sum of Row-sums

$T(n)=3T(\lfloor n/4 \rfloor)+\Theta(n^2)$



$\log_4 n$

cn^2 → cn^2

$c(\frac{1}{4}n)^2$ $c(\frac{1}{4}n)^2$ $c(\frac{1}{4}n)^2$ → $\frac{3}{16}cn^2$

$c((1/16)n)^2$ $c((1/16)n)^2$ …… $c((1/16)n)^2$ $c((1/16)n)^2$ …… $c((1/16)n)^2$ → $\left(\frac{3}{16}\right)^2 cn^2$

…… ……

⋮

$T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ → $\Theta\left(n^{\log_4 3}\right)$

Note: $3^{\log_4 n} = n^{\log_4 3}$

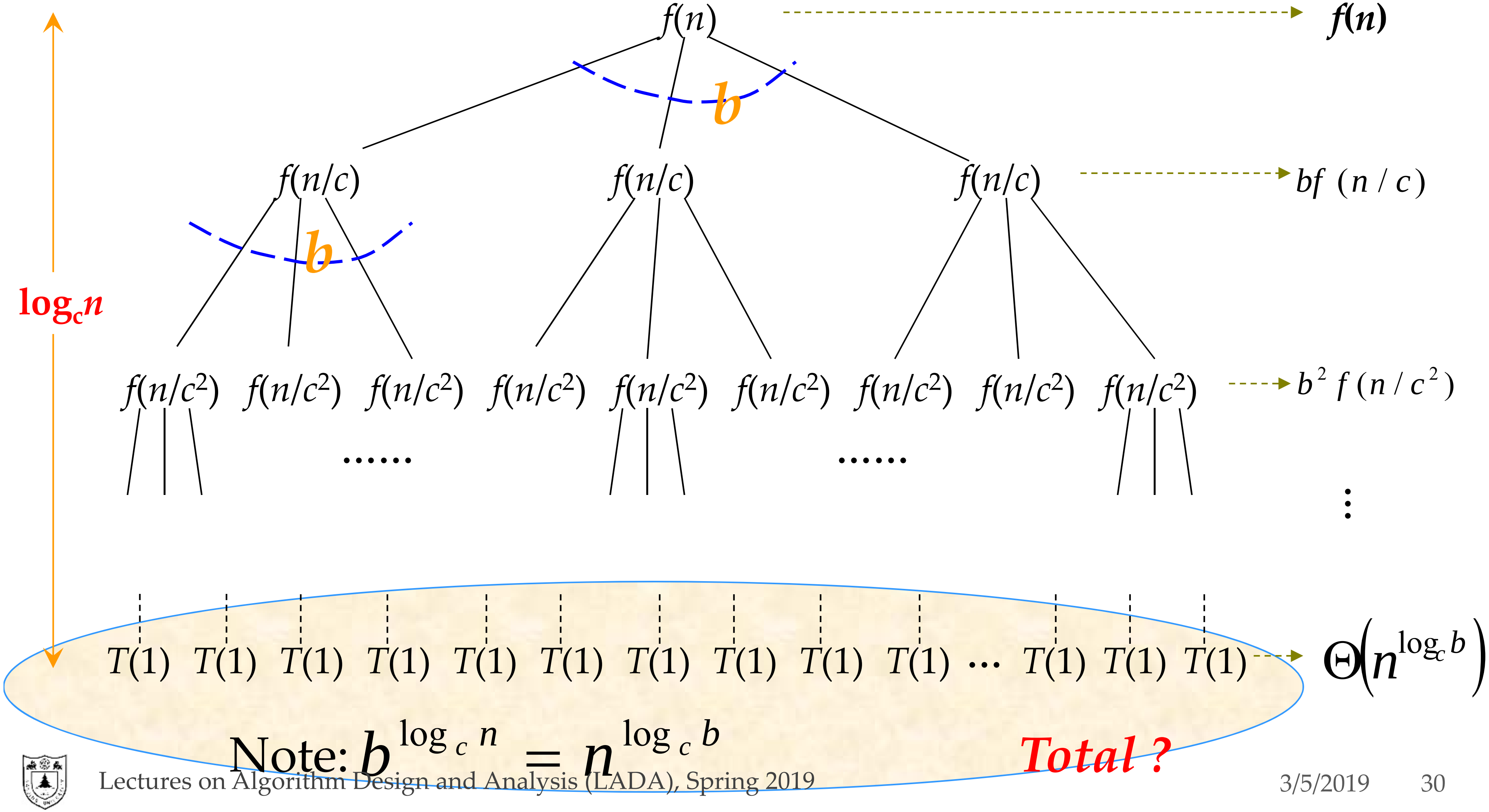
Total: $\Theta(n^2)$

Solving the Divide-and-Conquer Recurrence

- **The recursion equation for divide-and-conquer, the general case: $T(n)=bT(n/c)+f(n)$**
- **Observations:**
  - Let base-cases occur at depth D(leaf), then $n/c^{D}=1$, that is $D=\log(n)/\log(c)$
  - Let the number of leaves of the tree be L, then $L=b^{D}$, that is $L=b^{(\log(n)/\log(c))}$.
  - By a little algebra: $L=n^{E}$, where $E=\log(b)/\log(c)$, called *critical exponent*.

Recursion Tree for

$T(n)=bT(n/c)+f(n)$



$f(n)$
b
$f(n/c)$ $f(n/c)$ $f(n/c)$
b
$\log_c n$
$f(n/c^2)$ $f(n/c^2)$ $f(n/c^2)$ $f(n/c^2)$ $f(n/c^2)$ $f(n/c^2)$ $f(n/c^2)$ $f(n/c^2)$ $f(n/c^2)$
......
......
$T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ $T(1)$ ... $T(1)$ $T(1)$ $T(1)$
$f(n)$
$bf(n/c)$
$b^2 f(n/c^2)$
⋮
$\Theta\left(n^{\log_c b}\right)$

Note: $b^{\log_c n} = n^{\log_c b}$

Total ?

Divide-and-Conquer - the Solution

- **The solution of divide-and-conquer equation is the non-recursive costs of all nodes in the tree, which is the sum of the row-sums**
  - The recursion tree has depth D=log(n)/ log(c), so there are about that many row-sums.

- **The 0th row-sum**
  - is $f(n)$, the nonrecursive cost of the root.
- **The D^{th} row-sum**
  - is n^E, assuming base cases cost 1, or $\Theta(n^E)$ in any event.

Solution by Row-sums

- **[Little Master Theorem] Row-sums decide the solution of the equation for divide-and-conquer:**
  - Increasing geometric series: $T(n) \in \Theta(n^E)$
  - Constant: $T(n) \in \Theta(f(n) \log n)$
  - Decreasing geometric series: $T(n) \in \Theta(f(n))$

This can be generalized to get a result not using explicitly row-sums.

Master Theorem

- **Loosening the restrictions on *f*(*n*)**
  - Case 1: $f(n) \in O(n^{E-\varepsilon})$, ($\varepsilon>0$), then:
    $$T(n) \in \Theta(n^{E})$$
  - Case 2: $f(n) \in \Theta(n^{E})$, as all node depth contribute about equally:
    $$T(n) \in \Theta(f(n)\log(n))$$
  - case 3: $f(n) \in \Omega(n^{E+\varepsilon})$, ($\varepsilon>0$), and if $bf(n/c) \leq \theta f(n)$ for some constant $\theta < 1$ and all sufficiently large n, then:
    $$T(n) \in \Theta(f(n))$$

The positive ε is critical, resulting gaps between cases as well

Using Master Theorem

- Example 1: $T(n) = 9T(\frac{n}{3}) + n$

  $b = 9, c = 3, E = 2, f(n) = n = O(n^{E-1})$

  Case 1 applies: $T(n) = \Theta(n^2)$

- Example 2: $T(n) = T(\frac{2}{3}n) + 1$

  $b = 1, c = \frac{3}{2}, E = 0, f(n) = 1 = \Theta(n^E)$

  Case 2 applies: $T(n) = \Theta(\log n)$

- Example 3: $T(n) = 3T(\frac{n}{4}) + n\log n$

  $b = 3, c = 4, E = \log_4 3, f(n) = \Omega(n^{E+\epsilon})$

  $bf(\frac{n}{4}) = \frac{3}{4}n\log n - \frac{3}{2}n$

  Case 3 applies: $T(n) = \Theta(n\log n)$

Using Master Theorem

- **T(n) = 2T(n/2) + nlogn**
  - o Does Case 3 apply? Why?

- **T(n)=$\sqrt{n}$ T($\sqrt{n}$) + n**

- **The gap between the 3 cases**
  - o Often, none of the 3 cases apply
  - o Your task: design more non-solvable recursions

Thank you!

Q & A

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