

Codebook

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1 Setup

1.1 Template

```
1 #include <bits/stdc++.h>
2 #include <bits/extc++.h>
3 #define F first
4 #define S second
5 #define pb push_back
6 #define pob pop_back
7 #define pf push_front
8 #define pof pop_front
9 #define mp make_pair
10 #define mt make_tuple
11 #define all(x) (x).begin(),(x).end()
12 using namespace std;
13 //using namespace __gnu_pbds;
14 using pii = pair<long long,long long>;
```

```
15 using ld = long double;
16 using ll = long long;
17 const int mod = 1000000007;
18 const int mod2 = 998244353;
19 const ld PI = acos(-1);
20 #define Bint __int128
21 #define int long long
```

1.2 vimrc

```
1 syntax on
2 set mouse=a
3 set nu
4 set ts=4
5 set sw=4
6 set smartindent
7 set cursorline
8 set hlsearch
9 set incsearch
10 set t_Co=256
11 nnoremap y ggyG
12 colorscheme afterglow
13 au BufNewFile *.cpp Or ~/default_code/default.cpp |
   ↪ let IndentStyle = "cpp"
```

2 Data-structure

2.1 PBDS

```
1 gp_hash_table<T, T> h;
2 tree<T, null_type, less<T>, rb_tree_tag,
   ↪ tree_order_statistics_node_update> tr;
3 tr.order_of_key(x); // find x's ranking
4 tr.find_by_order(k); // find k-th minimum, return
   ↪ iterator
```

2.2 LazyTagSegtree

```
1 struct segment_tree{
2     int seg[N << 2];
3     int tag1[N << 2], tag2[N << 2];
4     void down(int l, int r, int idx, int pidx){
5         int v = tag1[pidx], vv = tag2[pidx];
6         if(v)
7             tag1[idx] = v, seg[idx] = v * (r - l + 1),
           ↪ tag2[idx] = 0;
8         if(vv)
```

```

9     tag2[idx] += vv, seg[idx] += vv * (r - l +
↪ 1);
10 }
11 void Set(int l, int r, int ql, int qr, int v, int
↪ idx = 1){
12     if(ql == l && qr == r){
13         tag1[idx] = v;
14         tag2[idx] = 0;
15         seg[idx] = v * (r - l + 1);
16         return;
17     }
18     int mid = (l + r) >> 1;
19     down(l, mid, idx << 1, idx);
20     down(mid + 1, r, idx << 1 | 1, idx);
21     tag1[idx] = tag2[idx] = 0;
22     if(qr <= mid)
23         Set(l, mid, ql, qr, v, idx << 1);
24     else if(ql > mid)
25         Set(mid + 1, r, ql, qr, v, idx << 1 | 1);
26     else{
27         Set(l, mid, ql, mid, v, idx << 1);
28         Set(mid + 1, r, mid + 1, qr, v, idx << 1 |
↪ 1);
29     }
30     seg[idx] = seg[idx << 1] + seg[idx << 1 | 1];
31 }
32 void Increase(int l, int r, int ql, int qr, int
↪ v, int idx = 1){
33     if(ql == l && qr == r){
34         tag2[idx] += v;
35         seg[idx] += v * (r - l + 1);
36         return;
37     }
38     int mid = (l + r) >> 1;
39     down(l, mid, idx << 1, idx);
40     down(mid + 1, r, idx << 1 | 1, idx);
41     tag1[idx] = tag2[idx] = 0;
42     if(qr <= mid)
43         Increase(l, mid, ql, qr, v, idx << 1);
44     else if(ql > mid)
45         Increase(mid + 1, r, ql, qr, v, idx << 1 |
↪ 1);
46     else{
47         Increase(l, mid, ql, mid, v, idx << 1);
48         Increase(mid + 1, r, mid + 1, qr, v, idx << 1 |
↪ 1);
49     }
50     seg[idx] = seg[idx << 1] + seg[idx << 1 | 1];
51 }
52 int query(int l, int r, int ql, int qr, int idx =
↪ 1){
53     if(ql == l && qr == r)
54         return seg[idx];
55     int mid = (l + r) >> 1;
56     down(l, mid, idx << 1, idx);
57     down(mid + 1, r, idx << 1 | 1, idx);
58     tag1[idx] = tag2[idx] = 0;
59     if(qr <= mid)
60         return query(l, mid, ql, qr, idx << 1);
61     else if(ql > mid)
62         return query(mid + 1, r, ql, qr, idx << 1 |
↪ 1);
63     return query(l, mid, ql, mid, idx << 1) +
↪ query(mid + 1, r, mid + 1, qr, idx << 1 | 1);
64 }

```

```

65 void modify(int l, int r, int ql, int qr, int v,
↪ int type){
66     // type 1: increasement, type 2: set
67     if(type == 2)
68         Set(l, r, ql, qr, v);
69     else
70         Increase(l, r, ql, qr, v);
71 }

```

2.3 LiChaoTree

```

1 struct line{
2     int m, c;
3     int val(int x){
4         return m * x + c;
5     }
6     line(){}
7     line(int _m, int _c){
8         m = _m, c = _c;
9     }
10 };
11 struct Li_Chao_Tree{
12     line seg[N << 2];
13     void ins(int l, int r, int idx, line x){
14         if(l == r){
15             if(x.val(l) > seg[idx].val(l))
16                 seg[idx] = x;
17             return;
18         }
19         int mid = (l + r) >> 1;
20         if(x.m < seg[idx].m)
21             swap(x, seg[idx]);
22         // ensure x.m > seg[idx].m
23         if(seg[idx].val(mid) <= x.val(mid)){
24             swap(x, seg[idx]);
25             ins(l, mid, idx << 1, x);
26         }
27         else
28             ins(mid + 1, r, idx << 1 | 1, x);
29     }
30     int query(int l, int r, int p, int idx){
31         if(l == r)
32             return seg[idx].val(l);
33         int mid = (l + r) >> 1;
34         if(p <= mid)
35             return max(seg[idx].val(p), query(l, mid, p,
↪ idx << 1));
36         else
37             return max(seg[idx].val(p), query(mid + 1, r,
↪ p, idx << 1 | 1));
38     }

```

2.4 Treap

```

1 mt19937
↪ mtrd(chrono::steady_clock::now().time_since_epoch())
2 struct Treap{
3     Treap *l, *r;
4     int pri, key, sz;
5     Treap(){}
6     Treap(int _v){

```

```

7     l = r = NULL;
8     pri = mtrd();
9     key = _v;
10    sz = 1;
11 }
12 ~Treap(){
13     if ( l )
14         delete l;
15     if ( r )
16         delete r;
17 }
18 void push(){
19     for(auto ch : {l, r}){
20         if(ch){
21             // do something
22         }
23     }
24 }
25 };
26 int getSize(Treap *t){
27     return t ? t->sz : 0;
28 }
29 void pull(Treap *t){
30     t->sz = getSize(t->l) + getSize(t->r) + 1;
31 }
32 Treap* merge(Treap* a, Treap* b){
33     if(!a || !b)
34         return a ? a : b;
35     if(a->pri > b->pri){
36         a->push();
37         a->r = merge(a->r, b);
38         pull(a);
39         return a;
40     }
41     else{
42         b->push();
43         b->l = merge(a, b->l);
44         pull(b);
45         return b;
46     }
47 }
48 void splitBySize(Treap *t, Treap *&a, Treap *&b,
49     ↪ int k){
50     if(!t)
51         a = b = NULL;
52     else if(getSize(t->l) + 1 <= k){
53         a = t;
54         a->push();
55         splitBySize(t->r, a->r, b, k - getSize(t->l) -
56     ↪ 1);
57         pull(a);
58     }
59     else{
60         b = t;
61         b->push();
62         splitBySize(t->l, a, b->l, k);
63         pull(b);
64     }
65 }
66 void splitByKey(Treap *t, Treap *&a, Treap *&b, int
67     ↪ k){
68     if(!t)
69         a = b = NULL;
70     else if(t->key <= k){
71         a = t;

```

```

69         a->push();
70         splitByKey(t->r, a->r, b, k);
71         pull(a);
72     }
73     else{
74         b = t;
75         b->push();
76         splitByKey(t->l, a, b->l, k);
77         pull(b);
78     }
79 }
80 // O(n) build treap with sorted key nodes
81 void traverse(Treap *t){
82     if(t->l)
83         traverse(t->l);
84     if(t->r)
85         traverse(t->r);
86     pull(t);
87 }
88 Treap *build(int n){
89     vector<Treap*>st(n);
90     int tp = 0;
91     for(int i = 0, x; i < n; i++){
92         cin >> x;
93         Treap *nd = new Treap(x);
94         while(tp && st[tp - 1]->pri < nd->pri)
95             nd->l = st[tp - 1], tp--;
96         if(tp)
97             st[tp - 1]->r = nd;
98         st[tp++] = nd;
99     }
100     if(!tp){
101         st[0] = NULL;
102         return st[0];
103     }
104     traverse(st[0]);
105     return st[0];
106 }

```

3 Graph

3.1 RoundSquareTree

```

1 int cnt;
2 int dep[N], low[N]; // dep == -1 -> unvisited
3 vector<int>G[N], rstree[2 * N]; // 1 ~ n: round, n
4     ↪ + 1 ~ 2n: square
5 vector<int>stk;
6 void init(){
7     cnt = n;
8     for(int i = 1; i <= n; i++){
9         G[i].clear();
10        rstree[i].clear();
11        rstree[i + n].clear();
12        dep[i] = low[i] = -1;
13    }
14    dep[1] = low[1] = 0;
15 }
16 void tarjan(int x, int px){
17     stk.push_back(x);
18     for(auto i : G[x]){
19         if(dep[i] == -1){

```

```

19     dep[i] = low[i] = dep[x] + 1;
20     tarjan(i, x);
21     low[x] = min(low[x], low[i]);
22     if(dep[x] <= low[i]){
23         int z;
24         cnt++;
25         do{
26             z = stk.back();
27             rstree[cnt].push_back(z);
28             rstree[z].push_back(cnt);
29             stk.pop_back();
30         }while(z != i);
31         rstree[cnt].push_back(x);
32         rstree[x].push_back(cnt);
33     }
34 }
35 else if(i != px)
36     low[x] = min(low[x], dep[i]);
37 }
38 }

```

3.2 SCC

```

1 struct SCC{
2     int n;
3     int cnt;
4     vector<vector<int>>>G, revG;
5     vector<int>stk, sccid;
6     vector<bool>vis;
7     SCC(): SCC(0) {}
8     SCC(int _n): n(_n), G(_n + 1), revG(_n + 1),
    ↪ sccid(_n + 1), vis(_n + 1), cnt(0) {}
9     void addEdge(int u, int v){
10         // u -> v
11         assert(u > 0 && u <= n);
12         assert(v > 0 && v <= n);
13         G[u].push_back(v);
14         revG[v].push_back(u);
15     }
16     void dfs1(int u){
17         vis[u] = 1;
18         for(int v : G[u]){
19             if(!vis[v])
20                 dfs1(v);
21         }
22         stk.push_back(u);
23     }
24     void dfs2(int u, int k){
25         vis[u] = 1;
26         sccid[u] = k;
27         for(int v : revG[u]){
28             if(!vis[v])
29                 dfs2(v, k);
30         }
31     }
32     void Kosaraju(){
33         for(int i = 1; i <= n; i++)
34             if(!vis[i])
35                 dfs1(i);
36         fill(vis.begin(), vis.end(), 0);
37         while(!stk.empty()){
38             if(!vis[stk.back()])
39                 dfs2(stk.back(), ++cnt);

```

```

40         stk.pop_back();
41     }
42 }
43 };

```

3.3 2SAT

```

1 struct two_sat{
2     int n;
3     SCC G; // u: u, u + n: ~u
4     vector<int>ans;
5     two_sat(): two_sat(0) {}
6     two_sat(int _n): n(_n), G(2 * _n), ans(_n + 1) {}
7     void disjunction(int a, int b){
8         G.addEdge((a > n ? a - n : a + n), b);
9         G.addEdge((b > n ? b - n : b + n), a);
10    }
11    bool solve(){
12        G.Kosaraju();
13        for(int i = 1; i <= n; i++){
14            if(G.sccid[i] == G.sccid[i + n])
15                return false;
16            ans[i] = (G.sccid[i] > G.sccid[i + n]);
17        }
18        return true;
19    }
20 };

```

3.4 bridge

```

1 int dep[N], low[N];
2 vector<int>G[N];
3 vector<pair<int, int>>bridge;
4 void init(){
5     for(int i = 1; i <= n; i++){
6         G[i].clear();
7         dep[i] = low[i] = -1;
8     }
9     dep[1] = low[1] = 0;
10 }
11 void tarjan(int x, int px){
12     for(auto i : G[x]){
13         if(dep[i] == -1){
14             dep[i] = low[i] = dep[x] + 1;
15             tarjan(i, x);
16             low[x] = min(low[x], low[i]);
17             if(low[i] > dep[x])
18                 bridge.push_back(make_pair(i, x));
19         }
20         else if(i != px)
21             low[x] = min(low[x], dep[i]);
22     }
23 }

```

3.5 BronKerbosch_a algorithm

```

1 vector<vector<int>>>maximal_clique;
2 int cnt, G[N][N], all[N][N], some[N][N],
    ↪ none[N][N];

```

```

3 void dfs(int d, int an, int sn, int nn)
4 {
5     if(sn == 0 && nn == 0){
6         vector<int>v;
7         for(int i = 0; i < an; i++)
8             v.push_back(all[d][i]);
9         maximal_clique.push_back(v);
10        cnt++;
11    }
12    int u = sn > 0 ? some[d][0] : none[d][0];
13    for(int i = 0; i < sn; i++)
14    {
15        int v = some[d][i];
16        if(G[u][v])
17            continue;
18        int tsu = 0, tnn = 0;
19        for(int j = 0; j < an; j++)
20            all[d + 1][j] = all[d][j];
21        all[d + 1][an] = v;
22        for(int j = 0; j < sn; j++)
23            if(g[v][some[d][j]])
24                some[d + 1][tsu++] = some[d][j];
25        for(int j = 0; j < nn; j++)
26            if(g[v][none[d][j]])
27                none[d + 1][tnn++] = none[d][j];
28        dfs(d + 1, an + 1, tsu, tnn);
29        some[d][i] = 0, none[d][nn++] = v;
30    }
31 }
32 void process(){
33     cnt = 0;
34     for(int i = 0; i < n; i++)
35         some[0][i] = i + 1;
36     dfs(0, 0, n, 0);
37 }

```

3.6 Theorem

- Kosaraju's algorithm visit the strong connected components in topological order at second dfs.
- Euler's formula on planar graph: $V - E + F = C + 1$
- Kuratowski's theorem: A simple graph G is a planar graph iff G doesn't has a subgraph H such that H is homeomorphic to K_5 or $K_{3,3}$
- A complement set of every vertex cover correspond to a independent set. \Rightarrow Number of vertex of maximum independent set + Number of vertex of minimum vertex cover = V
- Maximum independent set of G = Maximum clique of the complement graph of G .
- A planar graph G colored with three colors iff there exist a maximal clique I such that $G - I$ is a bipartite.

4 Flow

4.1 Dinic

```

1 struct Max_Flow{
2     struct Edge{

```

```

3         int cap, to, rev;
4         Edge(){}
5         Edge(int _to, int _cap, int _rev){
6             to = _to, cap = _cap, rev = _rev;
7         }
8     };
9     const int inf = 1e18+10;
10    int s, t; // start node and end node
11    vector<vector<Edge>>G;
12    vector<int>dep;
13    vector<int>iter;
14    void addE(int u, int v, int cap){
15        G[u].pb(Edge(v, cap, G[v].size()));
16        // direct graph
17        G[v].pb(Edge(u, 0, G[u].size() - 1));
18        // undirect graph
19        // G[v].pb(Edge(u, cap, G[u].size() - 1));
20    }
21    void bfs(){
22        queue<int>q;
23        q.push(s);
24        dep[s] = 0;
25        while(!q.empty()){
26            int cur = q.front();
27            q.pop();
28            for(auto i : G[cur]){
29                if(i.cap > 0 && dep[i.to] == -1){
30                    dep[i.to] = dep[cur] + 1;
31                    q.push(i.to);
32                }
33            }
34        }
35    }
36    int dfs(int x, int fl){
37        if(x == t)
38            return fl;
39        for(int _ = iter[x] ; _ < G[x].size() ; _++){
40            auto &i = G[x][_];
41            if(i.cap > 0 && dep[i.to] == dep[x] + 1){
42                int res = dfs(i.to, min(fl, i.cap));
43                if(res <= 0)
44                    continue;
45                i.cap -= res;
46                G[i.to][i.rev].cap += res;
47                return res;
48            }
49            iter[x]++;
50        }
51        return 0;
52    }
53    int Dinic(){
54        int res = 0;
55        while(true){
56            fill(all(dep), -1);
57            fill(all(iter), 0);
58            bfs();
59            if(dep[t] == -1)
60                break;
61            int cur;
62            while((cur = dfs(s, INF)) > 0)
63                res += cur;
64        }
65        return res;
66    }
67    void init(int _n, int _s, int _t){

```

```

68     s = _s, t = _t;
69     G.resize(_n + 5);
70     dep.resize(_n + 5);
71     iter.resize(_n + 5);
72 }
73 };

```

5 Math

5.1 FastPow

```

1 long long qpow(long long x, long long powcnt, long
  ↳ long tomod){
2     long long res = 1;
3     for(; powcnt ; powcnt >>= 1 , x = (x * x) %
  ↳ tomod)
4         if(1 & powcnt)
5             res = (res * x) % tomod;
6     return (res % tomod);

```

5.2 EXGCD

```

1 // ax + by = c
2 // return (gcd(a, b), x, y)
3 tuple<long long, long long, long long>exgcd(long
  ↳ long a, long long b){
4     if(b == 0)
5         return make_tuple(a, 1, 0);
6     auto[g, x, y] = exgcd(b, a % b);
7     return make_tuple(g, y, x - (a / b) * y);

```

5.3 EXCRT

```

1 long long inv(long long x){ return qpow(x, mod - 2,
  ↳ mod); }
2 long long mul(long long x, long long y, long long
  ↳ m){
3     x = ((x % m) + m) % m, y = ((y % m) + m) % m;
4     long long ans = 0;
5     while(y){
6         if(y & 1)
7             ans = (ans + x) % m;
8         x = x * 2 % m;
9         y >>= 1;
10    }
11    return ans;
12 }
13 pii ExCRT(long long r1, long long m1, long long r2,
  ↳ long long m2){
14     long long g, x, y;
15     tie(g, x, y) = exgcd(m1, m2);
16     if((r1 - r2) % g)
17         return {-1, -1};
18     long long lcm = (m1 / g) * m2;
19     long long res = (mul(mul(m1, x, lcm), ((r2 - r1)
  ↳ / g), lcm) + r1) % lcm;
20     res = (res + lcm) % lcm;
21     return {res, lcm};

```

```

22 }
23 void solve(){
24     long long n, r, m;
25     cin >> n;
26     cin >> m >> r; // x == r (mod m)
27     for(long long i = 1 ; i < n ; i++){
28         long long r1, m1;
29         cin >> m1 >> r1;
30         if(r != -1 && m != -1)
31             tie(r, m) = ExCRT(r, m, r1, m1);
32     }
33     if(r == -1 && m == -1)
34         cout << "no solution\n";
35     else
36         cout << r << '\n';
37 }

```

5.4 Generating Functions

- Ordinary Generating Function $A(x) = \sum_{i \geq 0} a_i x^i$

$$\begin{aligned}
 - A(rx) &\Rightarrow r^n a_n \\
 - A(x) + B(x) &\Rightarrow a_n + b_n \\
 - A(x)B(x) &\Rightarrow \sum_{i=0}^n a_i b_{n-i} \\
 - A(x)^k &\Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k} \\
 - xA(x)' &\Rightarrow n a_n \\
 - \frac{A(x)}{1-x} &\Rightarrow \sum_{i=0}^n a_i
 \end{aligned}$$

- Exponential Generating Function $A(x) = \sum_{i \geq 0} \frac{a_i}{i!} x^i$

$$\begin{aligned}
 - A(x) + B(x) &\Rightarrow a_n + b_n \\
 - A(x)B(x) &\Rightarrow \sum_{i=0}^n n i a_i b_{n-i} \\
 - A(x)^k &\Rightarrow \sum_{i_1+i_2+\dots+i_k=n} n i_1 i_2 \dots i_k a_{i_1} a_{i_2} \dots a_{i_k} \\
 - xA(x) &\Rightarrow n a_n
 \end{aligned}$$

- Special Generating Function

$$\begin{aligned}
 - \frac{(1+x)^n}{1-x} &= \sum_{i \geq 0} n i x^i \\
 - \frac{1}{(1-x)^n} &= \sum_{i \geq 0} \binom{n+i-1}{i} x^i
 \end{aligned}$$

5.5 Numbers

- Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups. $S(n, k) = S(n-1, k-1) + kS(n-1, k)$, $S(n, 1) = S(n, n) = 1$ $S(n, k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n$ $x^n = \sum_{i=0}^n S(n, i) (x)_i$
- Catalan numbers $C_n = \frac{1}{n+1} 2nn = 2nn - 2nn + 1$, $\forall n \geq 0$ $C_{n+1} = \sum_{i=0}^n C_i C_{n-i} = \frac{2(2n+1)}{n+2} C_n$, $C_0 = 1$

5.6 Theorem

- Cayley's Formula

- Given a degree sequence d_1, d_2, \dots, d_n for each *labeled* vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\dots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of *labeled* forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

- Erdős–Gallai theorem A sequence of nonnegative integers $d_1 \geq \dots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + \dots + d_n$ is even and $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$ holds for every $1 \leq k \leq n$.

- Gale–Ryser theorem A pair of sequences of nonnegative integers $a_1 \geq \dots \geq a_n$ and b_1, \dots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i, k)$ holds for every $1 \leq k \leq n$.

- Flooring and Ceiling function identity

$$\begin{aligned} - \lfloor \frac{\lfloor \frac{a}{b} \rfloor}{c} \rfloor &= \lfloor \frac{a}{bc} \rfloor \\ - \lceil \frac{\lceil \frac{a}{b} \rceil}{c} \rceil &= \lceil \frac{a}{bc} \rceil \\ - \lceil \frac{a}{b} \rceil &\leq \frac{a+b-1}{b} \\ - \lfloor \frac{a}{b} \rfloor &\leq \frac{a-b+1}{b} \end{aligned}$$

- Möbius inversion formula

$$\begin{aligned} - f(n) &= \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d}) \\ - f(n) &= \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d) \\ - \sum_{d|n}^{n=1} \mu(d) &= 1 \\ - \sum_{d|n}^{n \neq 1} \mu(d) &= 0 \end{aligned}$$

- Spherical cap

$$\begin{aligned} - &\text{A portion of a sphere cut off by a plane.} \\ - &r: \text{ sphere radius, } a: \text{ radius of the base of the cap, } h: \\ &\text{height of the cap, } \theta: \arcsin(a/r). \\ - &\text{Volume} = \pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2 + \\ &\cos \theta)(1 - \cos \theta)^2/3. \\ - &\text{Area} = 2\pi r h = \pi(a^2 + h^2) = 2\pi r^2(1 - \cos \theta). \end{aligned}$$