

Codebook

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1 Setup

1.1 Template

```

1 #include <bits/stdc++.h>
2 #include <bits/extc++.h>
3 #define F first
4 #define S second
5 #define pb push_back
6 #define pob pop_back
7 #define pf push_front
8 #define pof pop_front

```

```

9 #define mp make_pair
10 #define mt make_tuple
11 #define all(x) (x).begin(),(x).end()
12 using namespace std;
13 //using namespace __gnu_pbds;
14 using pii = pair<long long,long long>;
15 using ld = long double;
16 using ll = long long;
17 const int mod = 1000000007;
18 const int mod2 = 998244353;
19 const ld PI = acos(-1);
20 #define Bint __int128
21 #define int long long

```

1.2 vimrc

```

1 syntax on
2 set mouse=a
3 set nu
4 set ts=4
5 set sw=4
6 set smartindent
7 set cursorline
8 set hlsearch
9 set incsearch
10 set t_Co=256
11 nnoremap y ggyG
12 colorscheme afterglow
13 au BufNewFile *.cpp Or ~/default_code/default.cpp |
   ↪ let IndentStyle = "cpp"

```

2 Data-structure

2.1 PBDS

```

1 gp_hash_table<T, T> h;
2 tree<T, null_type, less<T>, rb_tree_tag,
   ↪ tree_order_statistics_node_update> tr;
3 tr.order_of_key(x); // find x's ranking
4 tr.find_by_order(k); // find k-th minimum, return
   ↪ iterator

```

2.2 LazyTagSegtree

```

1 struct segment_tree{
2     int seg[N << 2];
3     int tag1[N << 2], tag2[N << 2];
4     void down(int l, int r, int idx, int pidx){

```

```

5   int v = tag1[pidx], vv = tag2[pidx];
6   if(v)
7       tag1[idx] = v, seg[idx] = v * (r - l + 1),
↪   tag2[idx] = 0;
8   if(vv)
9       tag2[idx] += vv, seg[idx] += vv * (r - l +
↪   1);
10  }
11  void Set(int l, int r, int ql, int qr, int v, int
↪   idx = 1){
12      if(ql == l && qr == r){
13          tag1[idx] = v;
14          tag2[idx] = 0;
15          seg[idx] = v * (r - l + 1);
16          return;
17      }
18      int mid = (l + r) >> 1;
19      down(l, mid, idx << 1, idx);
20      down(mid + 1, r, idx << 1 | 1, idx);
21      tag1[idx] = tag2[idx] = 0;
22      if(qr <= mid)
23          Set(l, mid, ql, qr, v, idx << 1);
24      else if(ql > mid)
25          Set(mid + 1, r, ql, qr, v, idx << 1 | 1);
26      else{
27          Set(l, mid, ql, mid, v, idx << 1);
28          Set(mid + 1, r, mid + 1, qr, v, idx << 1 |
↪   1);
29      }
30      seg[idx] = seg[idx << 1] + seg[idx << 1 | 1];
31  }
32  void Increase(int l, int r, int ql, int qr, int
↪   v, int idx = 1){
33      if(ql == l && qr == r){
34          tag2[idx] += v;
35          seg[idx] += v * (r - l + 1);
36          return;
37      }
38      int mid = (l + r) >> 1;
39      down(l, mid, idx << 1, idx);
40      down(mid + 1, r, idx << 1 | 1, idx);
41      tag1[idx] = tag2[idx] = 0;
42      if(qr <= mid)
43          Increase(l, mid, ql, qr, v, idx << 1);
44      else if(ql > mid)
45          Increase(mid + 1, r, ql, qr, v, idx << 1 |
↪   1);
46      else{
47          Increase(l, mid, ql, mid, v, idx << 1);
48          Increase(mid + 1, r, mid + 1, qr, v, idx << 1
↪   | 1);
49      }
50      seg[idx] = seg[idx << 1] + seg[idx << 1 | 1];
51  }
52  int query(int l, int r, int ql, int qr, int idx =
↪   1){
53      if(ql == l && qr == r)
54          return seg[idx];
55      int mid = (l + r) >> 1;
56      down(l, mid, idx << 1, idx);
57      down(mid + 1, r, idx << 1 | 1, idx);
58      tag1[idx] = tag2[idx] = 0;
59      if(qr <= mid)
60          return query(l, mid, ql, qr, idx << 1);
61      else if(ql > mid)

```

```

62      return query(mid + 1, r, ql, qr, idx << 1 |
↪   1);
63      return query(l, mid, ql, mid, idx << 1) +
↪   query(mid + 1, r, mid + 1, qr, idx << 1 | 1);
64  }
65  void modify(int l, int r, int ql, int qr, int v,
↪   int type){
66      // type 1: increasement, type 2: set
67      if(type == 2)
68          Set(l, r, ql, qr, v);
69      else
70          Increase(l, r, ql, qr, v);
71  }

```

2.3 LiChaoTree

```

1  struct line{
2      int m, c;
3      int val(int x){
4          return m * x + c;
5      }
6      line(){}
7      line(int _m, int _c){
8          m = _m, c = _c;
9      }
10 }
11 struct Li_Chao_Tree{
12     line seg[N << 2];
13     void ins(int l, int r, int idx, line x){
14         if(l == r){
15             if(x.val(l) > seg[idx].val(l))
16                 seg[idx] = x;
17             return;
18         }
19         int mid = (l + r) >> 1;
20         if(x.m < seg[idx].m)
21             swap(x, seg[idx]);
22         // ensure x.m > seg[idx].m
23         if(seg[idx].val(mid) <= x.val(mid)){
24             swap(x, seg[idx]);
25             ins(l, mid, idx << 1, x);
26         }
27         else
28             ins(mid + 1, r, idx << 1 | 1, x);
29     }
30     int query(int l, int r, int p, int idx){
31         if(l == r)
32             return seg[idx].val(l);
33         int mid = (l + r) >> 1;
34         if(p <= mid)
35             return max(seg[idx].val(p), query(l, mid, p,
↪   idx << 1));
36         else
37             return max(seg[idx].val(p), query(mid + 1, r,
↪   p, idx << 1 | 1));
38     }

```

2.4 Treap

```

1 mt19937
   ↪ mtrd(chrono::steady_clock::now().time_since_epoch().count());
2 struct Treap{
3     Treap *l, *r;
4     int pri, key, sz;
5     Treap(){
6         Treap(int _v){
7             l = r = NULL;
8             pri = mtrd();
9             key = _v;
10            sz = 1;
11        }
12        ~Treap(){
13            if ( l )
14                delete l;
15            if ( r )
16                delete r;
17        }
18        void push(){
19            for(auto ch : {l, r}){
20                if(ch){
21                    // do something
22                }
23            }
24        }
25    };
26    int getSize(Treap *t){
27        return t ? t->sz : 0;
28    }
29    void pull(Treap *t){
30        t->sz = getSize(t->l) + getSize(t->r) + 1;
31    }
32    Treap* merge(Treap* a, Treap* b){
33        if(!a || !b)
34            return a ? a : b;
35        if(a->pri > b->pri){
36            a->push();
37            a->r = merge(a->r, b);
38            pull(a);
39            return a;
40        }
41        else{
42            b->push();
43            b->l = merge(a, b->l);
44            pull(b);
45            return b;
46        }
47    }
48    void splitBySize(Treap *t, Treap *&a, Treap *&b,
   ↪ int k){
49        if(!t)
50            a = b = NULL;
51        else if(getSize(t->l) + 1 <= k){
52            a = t;
53            a->push();
54            splitBySize(t->r, a->r, b, k - getSize(t->l) -
   ↪ 1);
55            pull(a);
56        }
57        else{
58            b = t;
59            b->push();
60            splitBySize(t->l, a, b->l, k);
61            pull(b);
62        }
63    }
64    void splitByKey(Treap *t, Treap *&a, Treap *&b, int
   ↪ k){
65        if(!t)
66            a = b = NULL;
67        else if(t->key <= k){
68            a = t;
69            a->push();
70            splitByKey(t->r, a->r, b, k);
71            pull(a);
72        }
73        else{
74            b = t;
75            b->push();
76            splitByKey(t->l, a, b->l, k);
77            pull(b);
78        }
79    }
80    // O(n) build treap with sorted key nodes
81    void traverse(Treap *t){
82        if(t->l)
83            traverse(t->l);
84        if(t->r)
85            traverse(t->r);
86        pull(t);
87    }
88    Treap *build(int n){
89        vector<Treap*> st(n);
90        int tp = 0;
91        for(int i = 0, x; i < n; i++){
92            cin >> x;
93            Treap *nd = new Treap(x);
94            while(tp && st[tp - 1]->pri < nd->pri)
95                nd->l = st[tp - 1], tp--;
96            if(tp)
97                st[tp - 1]->r = nd;
98            st[tp++] = nd;
99        }
100        if(!tp){
101            st[0] = NULL;
102            return st[0];
103        }
104        traverse(st[0]);
105        return st[0];
106    }

```

3 Graph

3.1 RoundSquareTree

```

1 int cnt;
2 int dep[N], low[N]; // dep == -1 -> unvisited
3 vector<int> G[N], rstree[2 * N]; // 1 ~ n: round, n
   ↪ + 1 ~ 2n: square
4 vector<int> stk;
5 void init(){
6     cnt = n;
7     for(int i = 1; i <= n; i++){
8         G[i].clear();

```

```

9     rstree[i].clear();
10    rstree[i + n].clear();
11    dep[i] = low[i] = -1;
12 }
13 dep[1] = low[1] = 0;
14 }
15 void tarjan(int x, int px){
16     stk.push_back(x);
17     for(auto i : G[x]){
18         if(dep[i] == -1){
19             dep[i] = low[i] = dep[x] + 1;
20             tarjan(i, x);
21             low[x] = min(low[x], low[i]);
22             if(dep[x] <= low[i]){
23                 int z;
24                 cnt++;
25                 do{
26                     z = stk.back();
27                     rstree[cnt].push_back(z);
28                     rstree[z].push_back(cnt);
29                     stk.pop_back();
30                 }while(z != i);
31                 rstree[cnt].push_back(x);
32                 rstree[x].push_back(cnt);
33             }
34         }
35         else if(i != px)
36             low[x] = min(low[x], dep[i]);
37     }
38 }

```

3.2 SCC

```

1 struct SCC{
2     int n;
3     int cnt;
4     vector<vector<int>>>G, revG;
5     vector<int>stk, sccid;
6     vector<bool>vis;
7     SCC(): SCC(0) {}
8     SCC(int _n): n(_n), G(_n + 1), revG(_n + 1),
9     ↪ sccid(_n + 1), vis(_n + 1), cnt(0) {}
10    void addEdge(int u, int v){
11        // u -> v
12        assert(u > 0 && u <= n);
13        assert(v > 0 && v <= n);
14        G[u].push_back(v);
15        revG[v].push_back(u);
16    }
17    void dfs1(int u){
18        vis[u] = 1;
19        for(int v : G[u]){
20            if(!vis[v])
21                dfs1(v);
22        }
23        stk.push_back(u);
24    }
25    void dfs2(int u, int k){
26        vis[u] = 1;
27        sccid[u] = k;
28        for(int v : revG[u]){
29            if(!vis[v])
30                dfs2(v, k);
31        }
32    }
33 }

```

```

30     }
31 }
32 void Kosaraju(){
33     for(int i = 1; i <= n; i++){
34         if(!vis[i])
35             dfs1(i);
36     }
37     fill(vis.begin(), vis.end(), 0);
38     while(!stk.empty()){
39         if(!vis[stk.back()])
40             dfs2(stk.back(), ++cnt);
41         stk.pop_back();
42     }
43 };

```

3.3 2SAT

```

1 struct two_sat{
2     int n;
3     SCC G; // u: u, u + n: ~u
4     vector<int>ans;
5     two_sat(): two_sat(0) {}
6     two_sat(int _n): n(_n), G(2 * _n), ans(_n + 1) {}
7     void disjunction(int a, int b){
8         G.addEdge((a > n ? a - n : a + n), b);
9         G.addEdge((b > n ? b - n : b + n), a);
10    }
11    bool solve(){
12        G.Kosaraju();
13        for(int i = 1; i <= n; i++){
14            if(G.sccid[i] == G.sccid[i + n])
15                return false;
16            ans[i] = (G.sccid[i] > G.sccid[i + n]);
17        }
18        return true;
19    }
20 };

```

3.4 bridge

```

1 int dep[N], low[N];
2 vector<int>G[N];
3 vector<pair<int, int>>bridge;
4 void init(){
5     for(int i = 1; i <= n; i++){
6         G[i].clear();
7         dep[i] = low[i] = -1;
8     }
9     dep[1] = low[1] = 0;
10 }
11 void tarjan(int x, int px){
12     for(auto i : G[x]){
13         if(dep[i] == -1){
14             dep[i] = low[i] = dep[x] + 1;
15             tarjan(i, x);
16             low[x] = min(low[x], low[i]);
17             if(low[i] > dep[x])
18                 bridge.push_back(make_pair(i, x));
19         }
20         else if(i != px)
21             low[x] = min(low[x], dep[i]);
22     }
23 }

```

```

22 }
23 }

```

3.5 BronKerbosch_a algorithm

```

1 vector<vector<int>>>maximal_clique;
2 int cnt, G[N][N], all[N][N], some[N][N],
  ↪ none[N][N];
3 void dfs(int d, int an, int sn, int nn)
4 {
5     if(sn == 0 && nn == 0){
6         vector<int>v;
7         for(int i = 0; i < an; i++)
8             v.push_back(all[d][i]);
9         maximal_clique.push_back(v);
10        cnt++;
11    }
12    int u = sn > 0 ? some[d][0] : none[d][0];
13    for(int i = 0; i < sn; i++)
14    {
15        int v = some[d][i];
16        if(G[u][v])
17            continue;
18        int tsn = 0, tnn = 0;
19        for(int j = 0; j < an; j++)
20            all[d + 1][j] = all[d][j];
21        all[d + 1][an] = v;
22        for(int j = 0; j < sn; j++)
23            if(g[v][some[d][j]])
24                some[d + 1][tsn++] = some[d][j];
25        for(int j = 0; j < nn; j++)
26            if(g[v][none[d][j]])
27                none[d + 1][tnn++] = none[d][j];
28        dfs(d + 1, an + 1, tsn, tnn);
29        some[d][i] = 0, none[d][nn++] = v;
30    }
31 }
32 void process(){
33     cnt = 0;
34     for(int i = 0; i < n; i++)
35         some[0][i] = i + 1;
36     dfs(0, 0, n, 0);
37 }

```

3.6 Theorem

- Kosaraju's algorithm visit the strong connected components in topological order at second dfs.
- Euler's formula on planar graph: $V - E + F = C + 1$
- Kuratowski's theorem: A simple graph G is a planar graph iff G doesn't has a subgraph H such that H is homeomorphic to K_5 or $K_{3,3}$
- A complement set of every vertex cover correspond to a independent set. \Rightarrow Number of vertex of maximum independent set + Number of vertex of minimum vertex cover = V
- Maximum independent set of G = Maximum clique of the complement graph of G .

- A planar graph G colored with three colors iff there exist a maximal clique I such that $G - I$ is a bipartite.

4 String

4.1 RollingHash

```

1 struct Rolling_Hash{
2     int n;
3     const int P[5] = {146672737, 204924373,
  ↪ 585761567, 484547929, 116508269};
4     const int M[5] = {922722049, 952311013,
  ↪ 955873937, 901981687, 993179543};
5     vector<int>PW[5], pre[5], suf[5];
6     Rolling_Hash(): Rolling_Hash("") {}
7     Rolling_Hash(string s): n(s.size()){
8         for(int i = 0; i < 5; i++){
9             PW[i].resize(n), pre[i].resize(n),
  ↪ suf[i].resize(n);
10            PW[i][0] = 1, pre[i][0] = s[0] - 'a';
11            suf[i][n - 1] = s[n - 1] - 'a';
12        }
13        for(int i = 1; i < n; i++){
14            for(int j = 0; j < 5; j++){
15                PW[j][i] = PW[j][i - 1] * P[j] % M[j];
16                pre[j][i] = (pre[j][i - 1] * P[j] + s[i] -
  ↪ 'a') % M[j];
17            }
18        }
19        for(int i = n - 2; i >= 0; i--){
20            for(int j = 0; j < 5; j++){
21                suf[j][i] = (suf[j][i + 1] * P[j] + s[i] -
  ↪ 'a') % M[j];
22            }
23        }
24        int _substr(int k, int l, int r) {
25            int res = pre[k][r];
26            if(l > 0)
27                res -= 1LL * pre[k][l - 1] * PW[k][r - l + 1]
  ↪ % M[k];
28            if(res < 0)
29                res += M[k];
30            return res;
31        }
32        vector<int>substr(int l, int r){
33            vector<int>res(5);
34            for(int i = 0; i < 5; ++i)
35                res[i] = _substr(i, l, r);
36            return res;
37        }
38 };

```

4.2 SuffixArray

```

1 struct Suffix_Array{
2     int n, m; // m is the range of s
3     string s;
4     vector<int>sa, rk, lcp;
5     Suffix_Array(): Suffix_Array(0, 0, "") {}
6     Suffix_Array(int _n, int _m, string _s): n(_n),
  ↪ m(_m), sa(_n), rk(_n), lcp(_n), s(_s) {}

```

```

7 void Sort(int k, vector<int>&bucket,
→ vector<int>&idx, vector<int>&lst){
8     for(int i = 0; i < m; i++)
9         bucket[i] = 0;
10    for(int i = 0; i < n; i++)
11        bucket[lst[i]]++;
12    for(int i = 1; i < m; i++)
13        bucket[i] += bucket[i-1];
14    int p = 0;
15    // update index
16    for(int i = n - k; i < n; i++)
17        idx[p++] = i;
18    for(int i = 0; i < n; i++)
19        if(sa[i] >= k)
20            idx[p++] = sa[i] - k;
21    for(int i = n - 1; i >= 0; i--)
22        sa[--bucket[lst[idx[i]]]] = idx[i];
23 }
24 void build(){
25     vector<int>idx(n), lst(n), bucket(max(n, m));
26     for(int i = 0; i < n; i++)
27         bucket[lst[i] = (s[i] - 'a')]++;
28     for(int i = 1; i < m; i++)
29         bucket[i] += bucket[i - 1];
30     for(int i = n - 1; i >= 0; i--)
31         sa[--bucket[lst[i]]] = i;
32     for(int k = 1; k < n; k <= 1){
33         Sort(k, bucket, idx, lst);
34         // update rank
35         int p = 0;
36         idx[sa[0]] = 0;
37         for(int i = 1; i < n; i++){
38             int a = sa[i], b = sa[i - 1];
39             if(lst[a] == lst[b] && a + k < n && b + k <
→ n && lst[a + k] == lst[b + k]);
40             else
41                 p++;
42             idx[sa[i]] = p;
43         }
44         if(p == n - 1)
45             break;
46         for(int i = 0; i < n; i++)
47             lst[i] = idx[i];
48         m = p + 1;
49     }
50     for(int i = 0; i < n; i++)
51         rk[sa[i]] = i;
52     buildLCP();
53 }
54 void buildLCP(){
55     // lcp[rk[i]] >= lcp[rk[i - 1]] - 1
56     int v = 0;
57     for(int i = 0; i < n; i++){
58         if(!rk[i])
59             lcp[rk[i]] = 0;
60         else{
61             if(v)
62                 v--;
63             int p = sa[rk[i] - 1];
64             while(i + v < n && p + v < n && s[i + v] ==
→ s[p + v])
65                 v++;
66             lcp[rk[i]] = v;
67         }
68     }

```

```

69     }
70 };

```

5 Flow

5.1 Dinic

```

1 struct Max_Flow{
2     struct Edge{
3         int cap, to, rev;
4         Edge(){}
5         Edge(int _to, int _cap, int _rev){
6             to = _to, cap = _cap, rev = _rev;
7         }
8     };
9     const int inf = 1e18+10;
10    int s, t; // start node and end node
11    vector<vector<Edge>>G;
12    vector<int>dep;
13    vector<int>iter;
14    void addE(int u, int v, int cap){
15        G[u].pb(Edge(v, cap, G[v].size()));
16        // direct graph
17        G[v].pb(Edge(u, 0, G[u].size() - 1));
18        // undirect graph
19        // G[v].pb(Edge(u, cap, G[u].size() - 1));
20    }
21    void bfs(){
22        queue<int>q;
23        q.push(s);
24        dep[s] = 0;
25        while(!q.empty()){
26            int cur = q.front();
27            q.pop();
28            for(auto i : G[cur]){
29                if(i.cap > 0 && dep[i.to] == -1){
30                    dep[i.to] = dep[cur] + 1;
31                    q.push(i.to);
32                }
33            }
34        }
35    }
36    int dfs(int x, int fl){
37        if(x == t)
38            return fl;
39        for(int _ = iter[x] ; _ < G[x].size() ; _++){
40            auto &i = G[x][_];
41            if(i.cap > 0 && dep[i.to] == dep[x] + 1){
42                int res = dfs(i.to, min(fl, i.cap));
43                if(res <= 0)
44                    continue;
45                i.cap -= res;
46                G[i.to][i.rev].cap += res;
47                return res;
48            }
49            iter[x]++;
50        }
51        return 0;
52    }
53    int Dinic(){
54        int res = 0;
55        while(true){

```

```

56     fill(all(dep), -1);
57     fill(all(iter), 0);
58     bfs();
59     if(dep[t] == -1)
60         break;
61     int cur;
62     while((cur = dfs(s, INF)) > 0)
63         res += cur;
64 }
65 return res;
66 }
67 void init(int _n, int _s, int _t){
68     s = _s, t = _t;
69     G.resize(_n + 5);
70     dep.resize(_n + 5);
71     iter.resize(_n + 5);
72 }
73 };

```

6 Math

6.1 FastPow

```

1 long long qpow(long long x, long long powcnt, long
  ↳ tommod){
2     long long res = 1;
3     for(; powcnt ; powcnt >>= 1 , x = (x * x) %
  ↳ tommod)
4         if(1 & powcnt)
5             res = (res * x) % tommod;
6     return (res % tommod);

```

6.2 EXGCD

```

1 // ax + by = c
2 // return (gcd(a, b), x, y)
3 tuple<long long, long long, long long>exgcd(long
  ↳ long a, long long b){
4     if(b == 0)
5         return make_tuple(a, 1, 0);
6     auto[g, x, y] = exgcd(b, a % b);
7     return make_tuple(g, y, x - (a / b) * y);

```

6.3 EXCRT

```

1 long long inv(long long x){ return qpow(x, mod - 2,
  ↳ mod); }
2 long long mul(long long x, long long y, long long
  ↳ m){
3     x = ((x % m) + m) % m, y = ((y % m) + m) % m;
4     long long ans = 0;
5     while(y){
6         if(y & 1)
7             ans = (ans + x) % m;
8         x = x * 2 % m;
9         y >>= 1;
10    }
11    return ans;

```

```

12 }
13 pii ExCRT(long long r1, long long m1, long long r2,
  ↳ long long m2){
14     long long g, x, y;
15     tie(g, x, y) = exgcd(m1, m2);
16     if((r1 - r2) % g)
17         return {-1, -1};
18     long long lcm = (m1 / g) * m2;
19     long long res = (mul(mul(m1, x, lcm), ((r2 - r1)
  ↳ / g), lcm) + r1) % lcm;
20     res = (res + lcm) % lcm;
21     return {res, lcm};
22 }
23 void solve(){
24     long long n, r, m;
25     cin >> n;
26     cin >> m >> r; // x == r (mod m)
27     for(long long i = 1 ; i < n ; i++){
28         long long r1, m1;
29         cin >> m1 >> r1;
30         if(r != -1 && m != -1)
31             tie(r, m) = ExCRT(r, m, r1, m1);
32     }
33     if(r == -1 && m == -1)
34         cout << "no solution\n";
35     else
36         cout << r << '\n';
37 }

```

6.4 FFT

```

1 struct Polynomial{
2     int deg;
3     vector<int>x;
4     void FFT(vector<complex<double>>&a, bool invert){
5         int a_sz = a.size();
6         for(int len = 1; len < a_sz; len <= 1){
7             for(int st = 0; st < a_sz; st += 2 * len){
8                 double angle = PI / len * (invert ? -1 :
  ↳ 1);
9                 complex<double>wnow(1), w(cos(angle),
  ↳ sin(angle));
10                for(int i = 0; i < len; i++){
11                    auto a0 = a[st + i], a1 = a[st + len +
  ↳ i];
12                    a[st + i] = a0 + wnow * a1;
13                    a[st + i + len] = a0 - wnow * a1;
14                    wnow *= w;
15                }
16            }
17        }
18        if(invert)
19            for(auto &i : a)
20                i /= a_sz;
21    }
22    void change(vector<complex<double>>&a){
23        int a_sz = a.size();
24        vector<int>rev(a_sz);
25        for(int i = 1; i < a_sz; i++){
26            rev[i] = rev[i / 2] / 2;
27            if(i & 1)
28                rev[i] += a_sz / 2;
29        }

```



```

30     for(int i = 0; i < a_sz; i++)
31         if(i < rev[i])
32             swap(a[i], a[rev[i]]);
33     }
34     Polynomial multiply(Polynomial const&b){
35         vector<complex<double>>A(x.begin(), x.end()),
36         ↪ B(b.x.begin(), b.x.end());
37         int mx_sz = 1;
38         while(mx_sz < A.size() + B.size())
39             mx_sz <= 1;
40         A.resize(mx_sz);
41         B.resize(mx_sz);
42         change(A);
43         change(B);
44         FFT(A, 0);
45         FFT(B, 0);
46         for(int i = 0; i < mx_sz; i++)
47             A[i] *= B[i];
48         change(A);
49         FFT(A, 1);
50         Polynomial res(mx_sz);
51         for(int i = 0; i < mx_sz; i++)
52             res.x[i] = round(A[i].real());
53         while(!res.x.empty() && res.x.back() == 0)
54             res.x.pop_back();
55         res.deg = res.x.size();
56         return res;
57     }
58     Polynomial(): Polynomial(0) {}
59     Polynomial(int Size): x(Size), deg(Size) {}
60 };
```

6.5 Generating Functions

- Ordinary Generating Function $A(x) = \sum_{i \geq 0} a_i x^i$

$$\begin{aligned}
 & - A(rx) \Rightarrow r^n a_n \\
 & - A(x) + B(x) \Rightarrow a_n + b_n \\
 & - A(x)B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i} \\
 & - A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k} \\
 & - xA(x)' \Rightarrow n a_n \\
 & - \frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^n a_i
 \end{aligned}$$

- Exponential Generating Function $A(x) = \sum_{i \geq 0} \frac{a_i}{i!} x^i$

$$\begin{aligned}
 & - A(x) + B(x) \Rightarrow a_n + b_n \\
 & - A^{(k)}(x) \Rightarrow a_{n+k} \\
 & - A(x)B(x) \Rightarrow \sum_{i=0}^n n i a_i b_{n-i} \\
 & - A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} n i_1 i_2 \dots i_k a_{i_1} a_{i_2} \dots a_{i_k} \\
 & - xA(x) \Rightarrow n a_n
 \end{aligned}$$

- Special Generating Function

$$\begin{aligned}
 & - (1+x)^n = \sum_{i \geq 0} n i x^i \\
 & - \frac{1}{(1-x)^n} = \sum_{i \geq 0} i n - 1 x^i
 \end{aligned}$$

6.6 Numbers

- Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups. $S(n, k) = S(n-1, k-1) + kS(n-1, k)$, $S(n, 1) = S(n, n) = 1$, $S(n, k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} \binom{k}{i} i^n$, $x^n = \sum_{i=0}^n S(n, i)(x)_i$
- Catalan numbers $C_n = \frac{1}{n+1} 2n n = 2n n - 2n n + 1$, $\forall n \geq 0$
 $C_{n+1} = \sum_{i=0}^n C_i C_{n-i} = \frac{2(2n+1)}{n+2} C_n$, $C_0 = 1$

6.7 Theorem

- Cayley's Formula

- Given a degree sequence d_1, d_2, \dots, d_n for each *labeled* vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\dots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of *labeled* forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

- Erdős–Gallai theorem A sequence of nonnegative integers $d_1 \geq \dots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + \dots + d_n$ is even and $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$ holds for every $1 \leq k \leq n$.

- Gale–Ryser theorem A pair of sequences of nonnegative integers $a_1 \geq \dots \geq a_n$ and b_1, \dots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i, k)$ holds for every $1 \leq k \leq n$.

- Flooring and Ceiling function identity

$$\begin{aligned}
 & - \lfloor \frac{\lfloor \frac{a}{b} \rfloor}{c} \rfloor = \lfloor \frac{a}{bc} \rfloor \\
 & - \lceil \frac{\lceil \frac{a}{b} \rceil}{c} \rceil = \lceil \frac{a}{bc} \rceil \\
 & - \lceil \frac{a}{b} \rceil \leq \frac{a+b-1}{b} \\
 & - \lfloor \frac{a}{b} \rfloor \leq \frac{a-b+1}{b}
 \end{aligned}$$

- Möbius inversion formula

$$\begin{aligned}
 & - f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f\left(\frac{n}{d}\right) \\
 & - f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu\left(\frac{d}{n}\right) f(d) \\
 & - \sum_{d|n}^{\neq 1} \mu(d) = 1 \\
 & - \sum_{d|n} \mu(d) = 0
 \end{aligned}$$

- Spherical cap

- A portion of a sphere cut off by a plane.
- r : sphere radius, a : radius of the base of the cap, h : height of the cap, θ : $\arcsin(a/r)$.
- Volume $= \pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\cos\theta)^2/3$
- Area $= 2\pi r h = \pi(a^2+h^2) = 2\pi r^2(1-\cos\theta)$.