

# Trusted Execution – and how far you can trust it

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CIF Seminar, CiTiP, KU Leuven, 2020-02-07

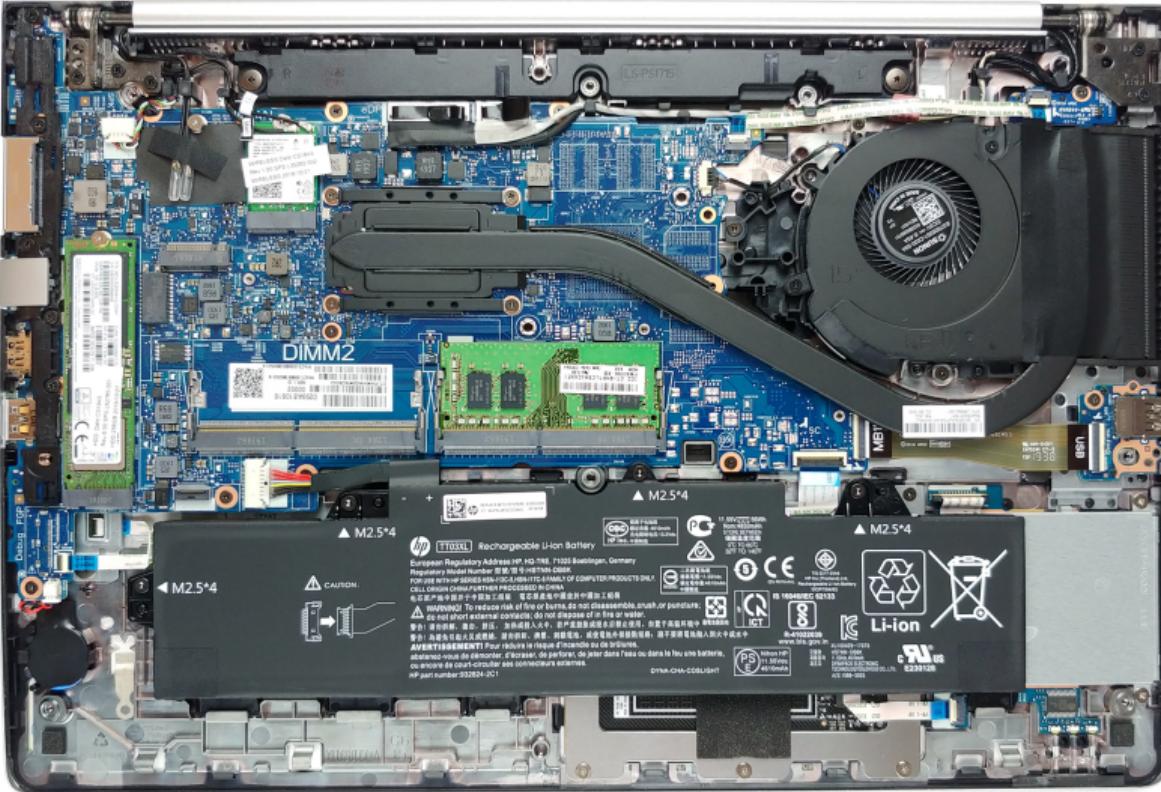
## **Trusted Computing...**

- Strong integrity protection and isolation for software components
- Software attestation: cryptographically bind a software to the executing hardware
- Sealed storage: bind data to attested software

## **...and how far you can trust it**

- Under which assumptions and attacker models?
- What about privacy?
- What are interesting use cases?

# Computers and how they get hacked



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**Memory** contains bits. Lots of them. Bits are grouped in bytes or words, which can be individually addressed.

# Computers and how they get hacked

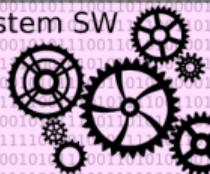
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|    |   |
|----|---|
| 00 | 1001000101011100011010100001100101000   |
| 90 | 110011110001010111000110101000011010101 |
| 70 | 0101000101011100011010100001101010101   |
| 50 | 110011110001010111000110101000011010100 |
| 30 | 0101000101011100011010100001101010101   |
| 10 | 110011110001010111000110101000011010100 |
| 11 | 0101000101011100011010100001101010101   |
| 12 | 110011110001010111000110101000011010101 |
| 13 | 0101000101011100011010100001101010101   |
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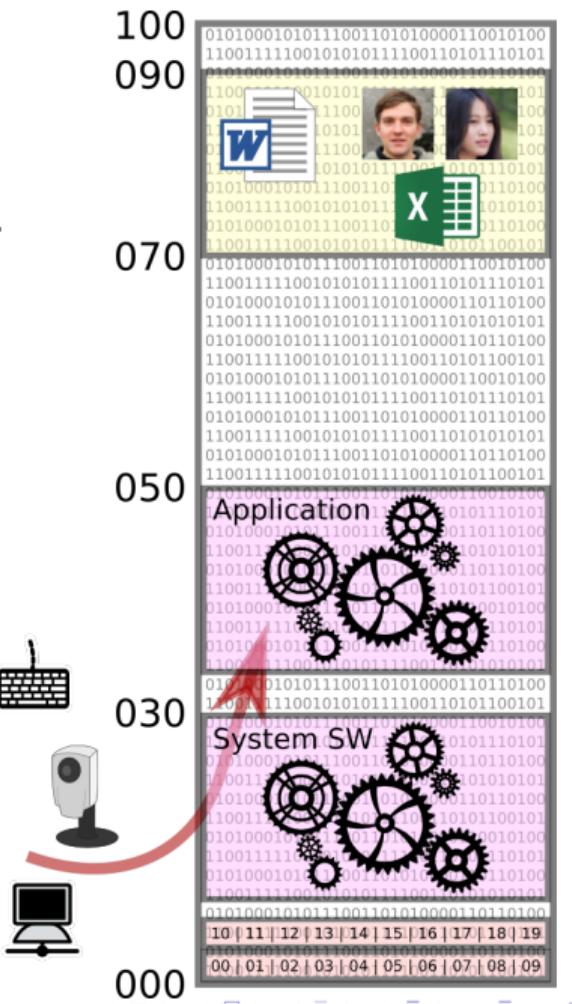
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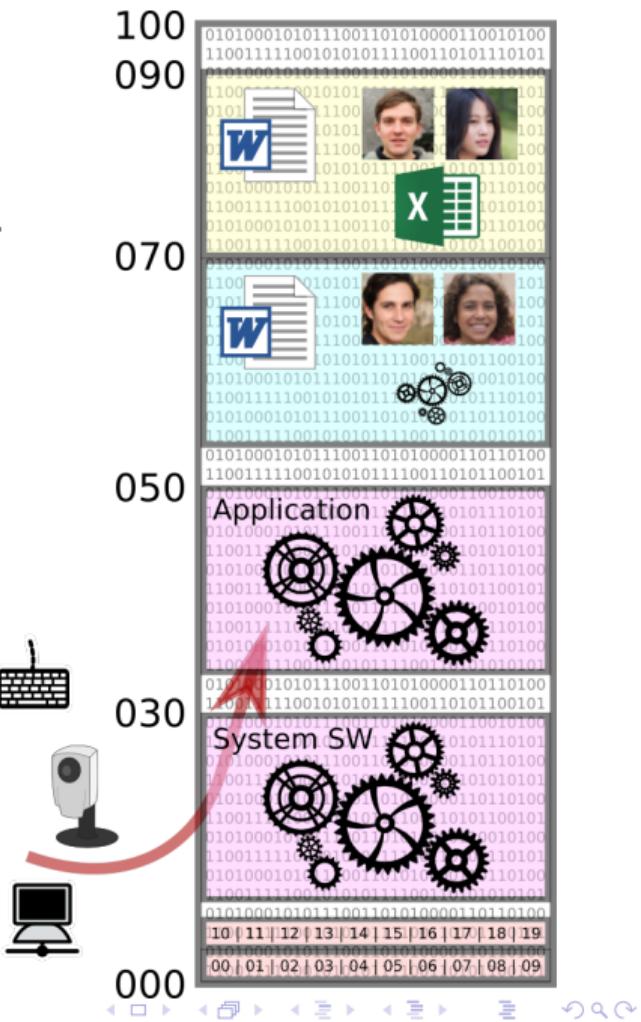
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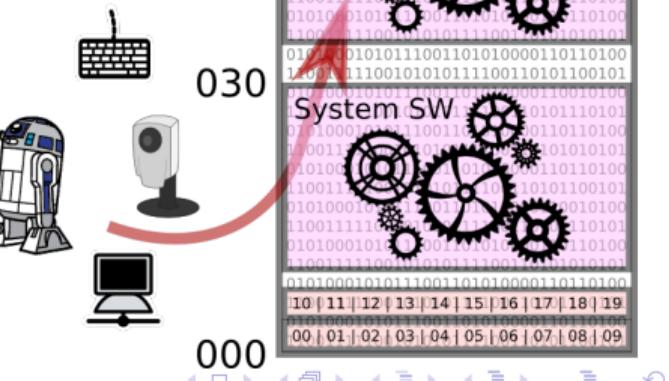
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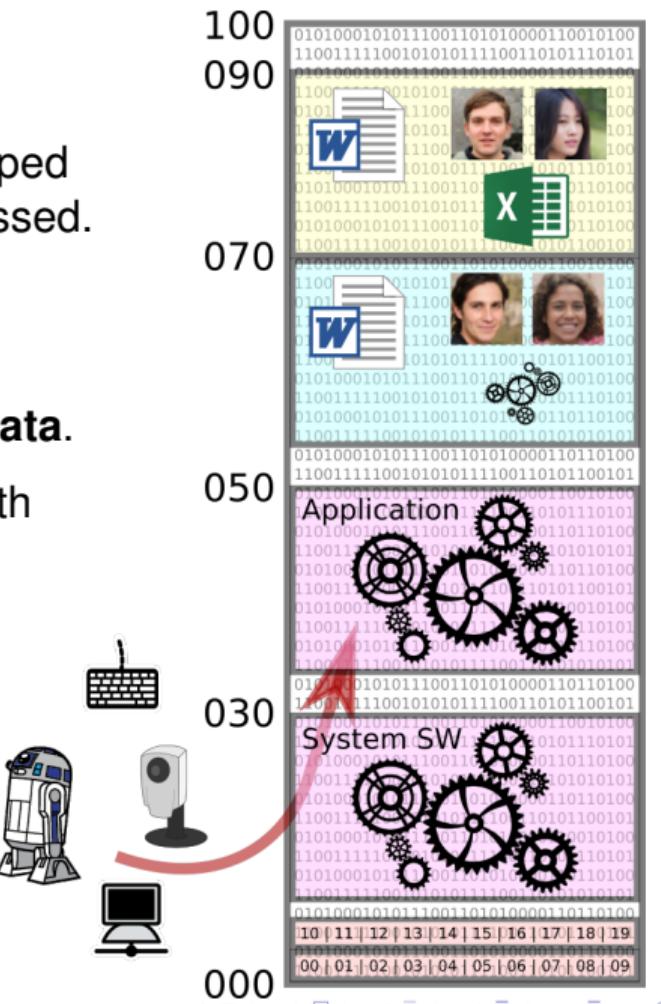
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**Software is modular!** Applications rely on a wealth of system software and shared code libraries to implement functionality.

**Security is based on assumptions.** What is trusted? What do we expect an attacker to do? What vulnerabilities are likely to exist in our code? Is there a moment when the system is supposedly secure?



# Computers and how they get hacked

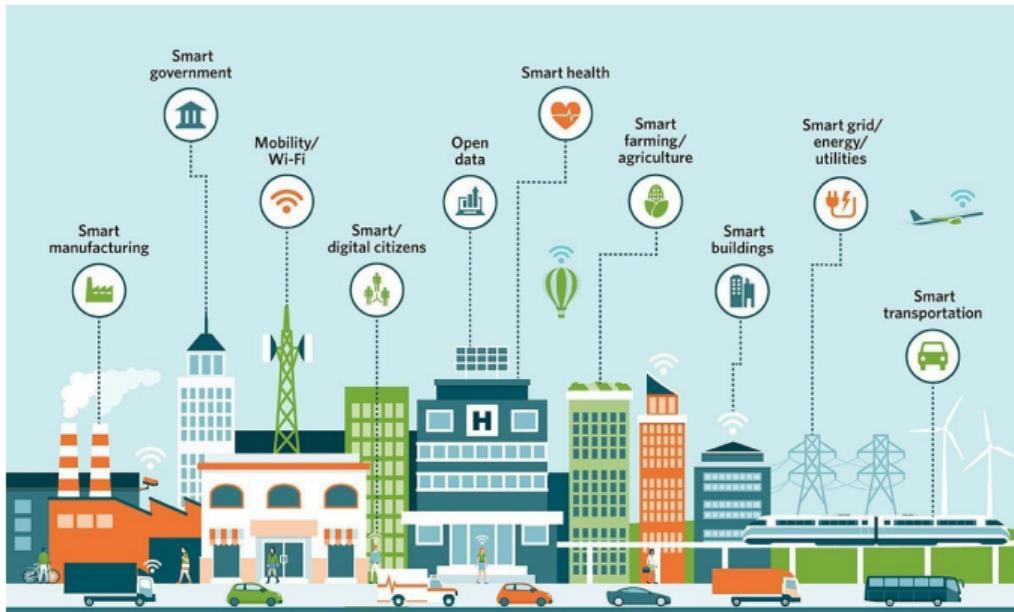
```
1  /* stack1.c; https://github.com/gerasdf/InsecureProgramming */
2
3 #include <stdio.h>
4
5 int main() {
6     int cookie;
7     char buf[80];
8
9     printf("buf: %08x cookie: %08x\n", &buf, &cookie);
10    gets(buf);
11
12    if (cookie == 0x41424344) {
13        printf("you win!\n");
14    }
15 }
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**Task: Compile and exploit to get “you win!”.**

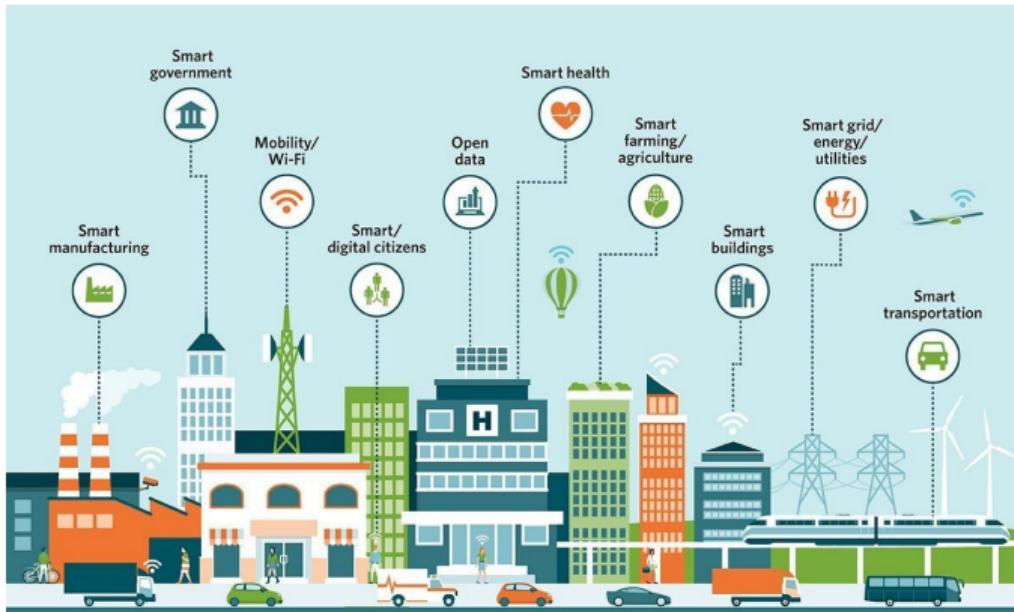
# Security in Smart Environments



**Vulnerabilities can hide anywhere:** There are 150M lines of code in a modern car. **Compartmentalisation** can help with managing complexity and bug containment.

**Image source:** <https://internetofthingsagenda.techtarget.com/definition/smart-city>

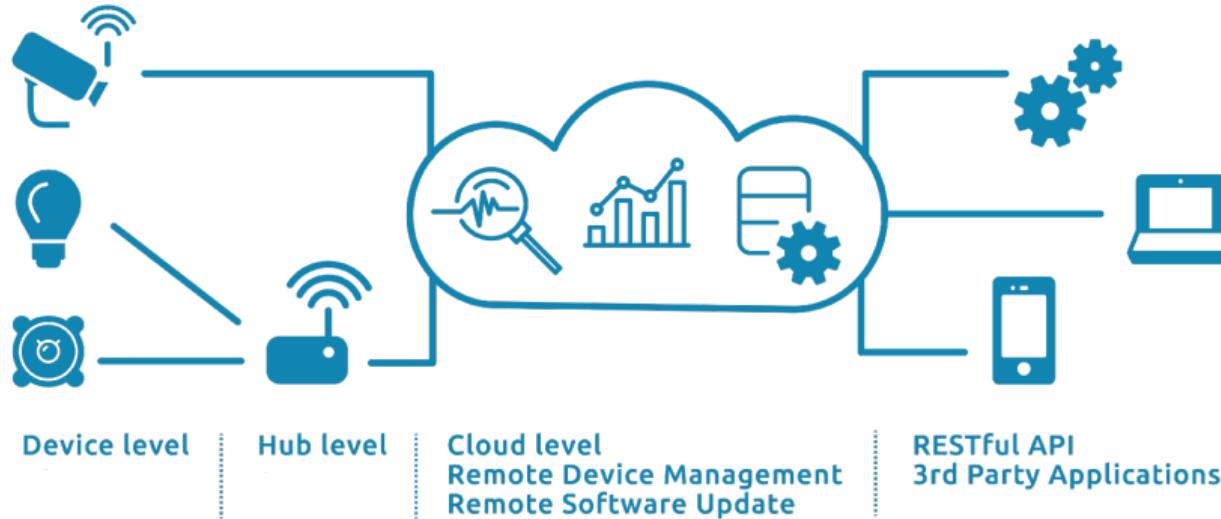
# Security in Smart Environments



**Infrastructure needs to be developed with safety, security and privacy in mind!** What is critical infrastructure? What is critical code? Where is personal data being processed? What's the impact of failure?

Image source: <https://internetofthingsagenda.techtarget.com/definition/smart-city>

# Security in Smart Environments



**Understanding can be really difficult:** What stake holders are involved? What are their objectives and abilities? What hardware and software is involved? Software quality? Data flows? Security requirements and guarantees?

**Image source:** <https://medium.com/connected-news/iot-foundation-what-is-an-iot-platform-c37c5e72d4a0>

# Security in Smart Environments

*Facebook Is Breached by Hackers,  
Putting 50 Million Users' Data at Risk*



One of the challenges for Facebook's chief executive Mark Zuckerberg is convincing users that the company handles their data responsibly.

**Source:** <https://www.nytimes.com/2018/09/28/technology/facebook-hack-data-breach.html>

**“The risks are about to get worse, because computers are being embedded into physical devices and will affect lives, not just our data.”**

— Bruce Schneier, [Sch18]

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Sex

# The looming deluge of connected dildos is a security nightmare

Just because the teledildonics patent has expired, sex tech companies shouldn't rush to bring connectivity to their products

**Source:** <https://www.wired.co.uk/article/teledildonics-hacking-sex-toys> (2017)

Meet us at WIRED Smarter this October

[BOOK TICKET](#)

# Smart dildos and vibrators keep getting hacked – but Tor could be the answer to safer connected sex

Connected sex toys are gathering huge amounts of data about our most intimate moments. Problem is, they're always getting hacked. Welcome to the emerging field of Onion Dildonics

**Source:** <https://www.wired.co.uk/article/sext-toy-bluetooth-hacks-security-fix> (2018)

# Security in Smart Environments

WIR ED

Inside the Cunning, Unprecedented Hack of Ukraine's Power Grid

KIM ZETTER SECURITY 03.03.16 07:00 AM

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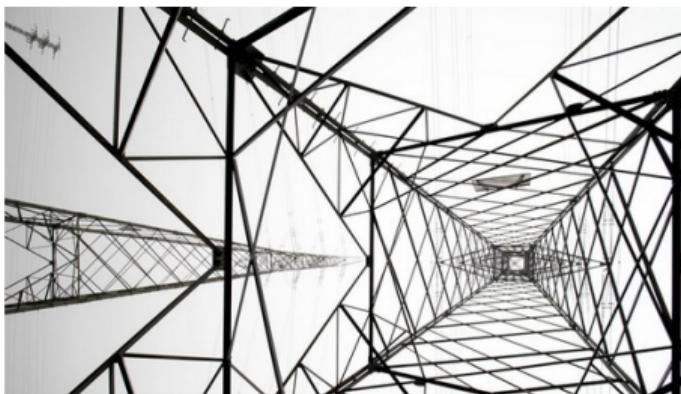
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t TWEET

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## INSIDE THE CUNNING, UNPRECEDENTED HACK OF UKRAINE'S POWER GRID



Source: <https://www.wired.com/2016/03/inside-cunning-unprecedented-hack-ukraines-power-grid/>

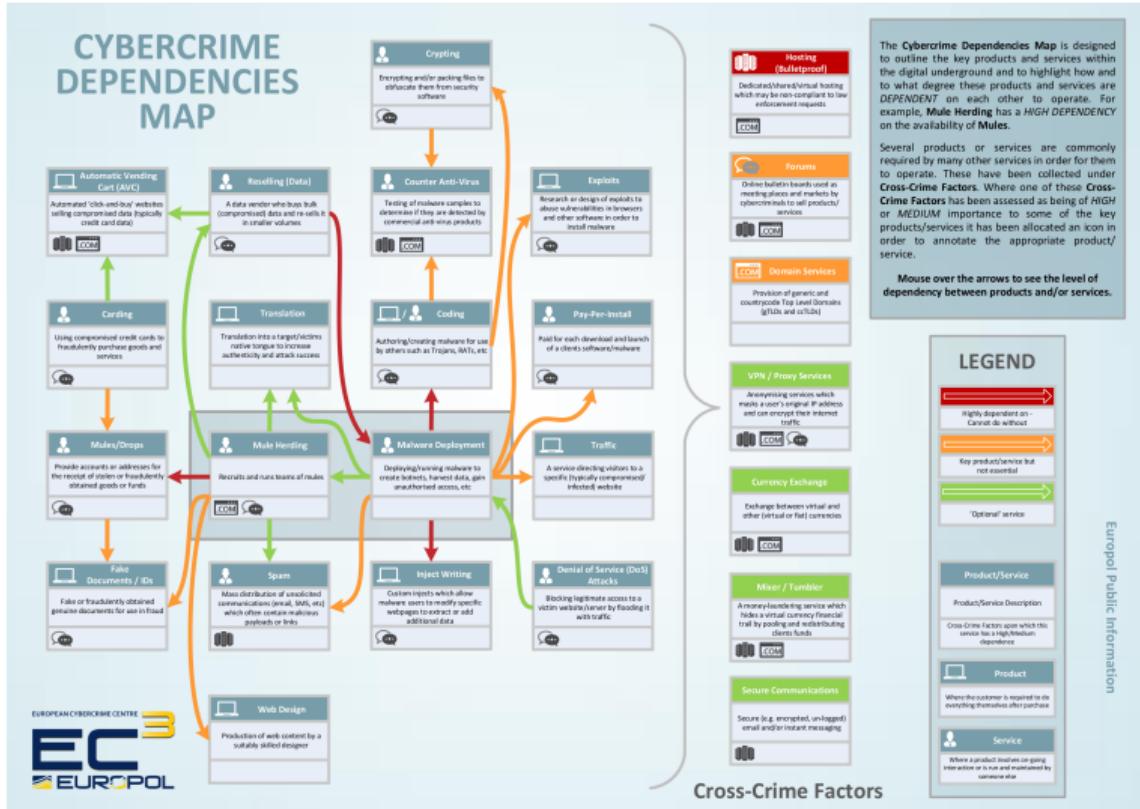
DO WE JUST SUCK  
AT... COMPUTERS?

YUP. ESPECIALLY SHARED ONES.



**Source:** <https://www.xkcd.com/1938/>

# Security in Smart Environments



Source: <https://www.europol.europa.eu/publications-documents/cybercrime-dependencies-map>

# What can we trust?

- **Reasoning about security is about setting boundaries**
  - Which parts are considered trusted, and which parts are not? **And why?**
  - How far do you want to go in defending your application?
  - What kind of security is economically viable?
- **Building secure systems requires rigorous security arguments**
  - Having a good idea about what you are building.
  - Determining **which attackers are considered** to be in scope.
  - Analysing potential vulnerabilities, and introducing appropriate countermeasures.
- **A security argument** is a rigorous argument that under a given system and adversary model, a countermeasure effectively counters a threat, or a security mechanism achieves a security goal.

# What can we trust?

Software?



Hardware?



Supply Chains?



People?



...

# What can we trust?

Menu Q, Search

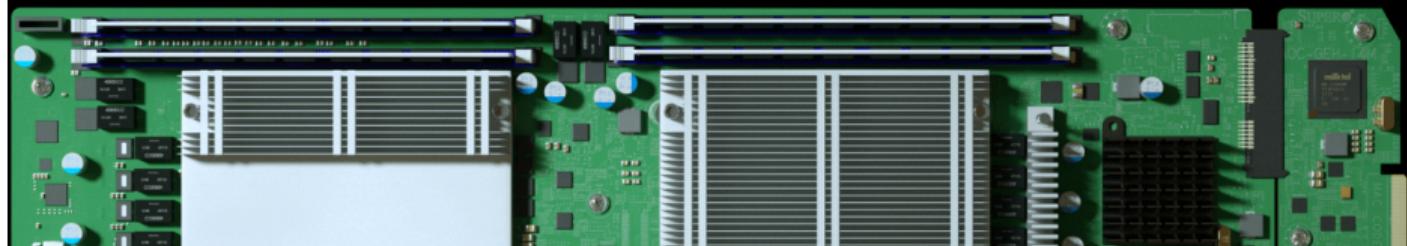
October 2018, 11:00 CEST

Bloomberg Businessweek

Sign In Subsc

## The Big Hack: How China Used a Tiny Chip to Infiltrate U.S. Companies

The attack by Chinese spies reached almost 30 U.S. companies, including Amazon and Apple, by compromising America's technology supply chain, according to extensive interviews with government and corporate sources.



Source: <https://www.bloomberg.com/news/features/2018-10-04/the-big-hack-how-china-used-a-tiny-chip-to-infiltrate-america...>

# Gathering Platform Requirements – A Thought Experiment



Sensors come from **different vendors**. Why would you trust them?

The cloud is “**other people’s computers**”. Why trust them?

Terminals may be used and managed by **health care professionals**...

There are **huge software and hardware stacks** with multiple vendors everywhere.

**Image source:** <https://medium.com/connected-news/iot-foundation-what-is-an-iot-platform-c37c5e72d4a0>

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- Thread modelling, risk assessment, etc.
- Anonymisation of data, if possible
- Zero Trust, micro-segmentation and granular perimeters

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- Encryption
- Isolation, Security Rings

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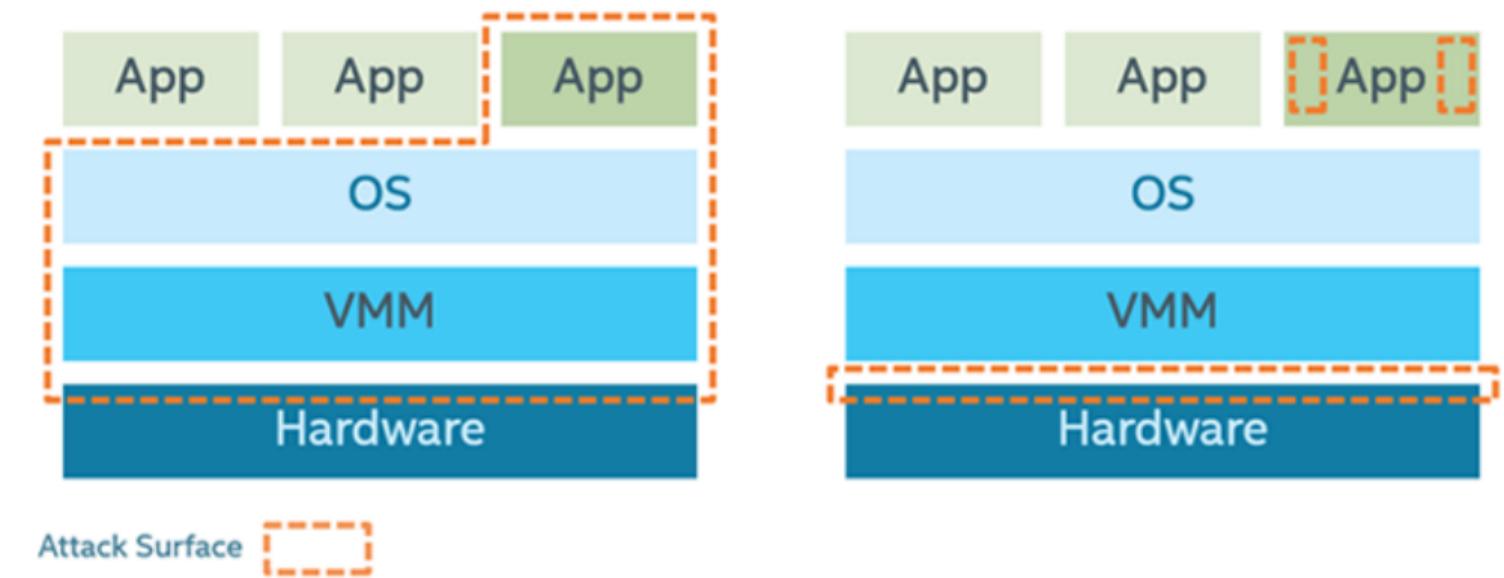
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## How can the execution environment (= hardware) help you?

- Encryption
- Isolation, Security Rings
- **Minimise Trusted Computing Base:**  
remove hypervisors, OSs, libraries from TCB  
**only trust hardware and your own code**

# Motivation: Application Attack Surface

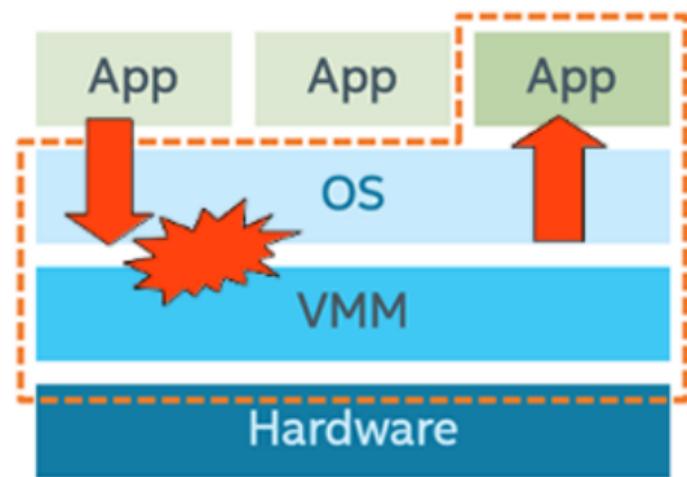
## Attack Surface Without Enclaves      Attack Surface With Enclaves



<https://software.intel.com/en-us/articles/intel-software-guard-extensions-tutorial-part-1-foundation>

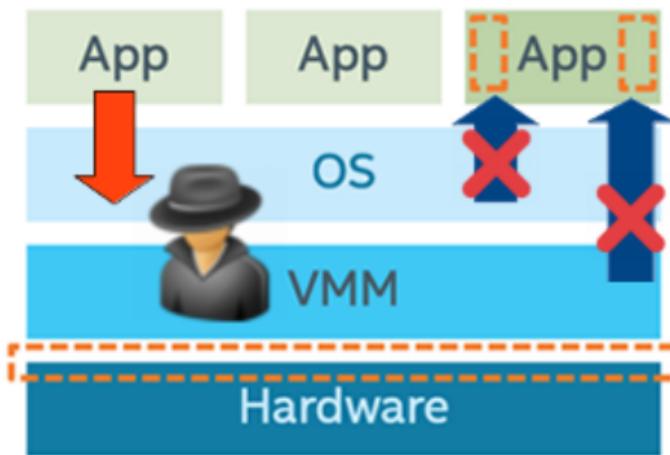
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Attack Surface

Attack Surface With Enclaves



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Layered architecture ↔ **hardware-only TCB**

## Gathering Platform Requirements – A Real System

**“We don’t want the Signal service to have visibility into the social graph of Signal users. Signal is always aspiring to be as ‘zero knowledge’ as possible, and having a durable record of every user’s friends and contacts on our servers would obviously not be privacy-preserving.”**



**Source:** <https://signal.org/blog/private-contact-discovery/>

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- ① Run a contact discovery service in a **secure SGX enclave**.
- ② Clients that wish to perform contact discovery negotiate a **secure connection** over the network all the way through the remote OS **to the enclave**.
- ③ Clients perform **remote attestation** to ensure that the code which is running in the *enclave is the same as the expected published open source code*.
- ④ **Clients transmit [...] their address book** to the enclave.
- ⑤ The **enclave looks up a client’s contacts** in the set of all registered users and **encrypts the results back** to the client.

**Source:** <https://signal.org/blog/private-contact-discovery/>

# Trusted Computing

**Source:** [https://en.wikipedia.org/wiki/Trusted\\_Computing](https://en.wikipedia.org/wiki/Trusted_Computing)

# Trusted Computing

**According to the *Trusted Computing Group***

Protect computing infrastructure at end points;

Hardware extensions to **enforce specific behaviour** and to provide cryptographic capabilities, protecting against unauthorised change and attacks

**Source:** [https://en.wikipedia.org/wiki/Trusted\\_Computing](https://en.wikipedia.org/wiki/Trusted_Computing)

# Trusted Computing

## According to the *Trusted Computing Group*

Protect computing infrastructure at end points;

Hardware extensions to enforce specific behaviour and to provide cryptographic capabilities, protecting against unauthorised change and attacks

- **Endorsement Key**, EK Certificate, Platform Certificate: Unique private key that never leaves the hardware, authenticate device identity
- **Memory curtaining**: provide isolation of sensitive areas of memory
- **Sealed storage**: Bind data to specific device or software
- **Remote attestation**: authenticate hardware and software configuration to a remote host
- **Trusted third party** as an intermediary to provide (ano|pseudo)nymity

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**In practice:** different architectures, subset of the above features, additions such as “enclaved” execution, memory encryption or secure I/O capabilities

**Source:** [https://en.wikipedia.org/wiki/Trusted\\_Computing](https://en.wikipedia.org/wiki/Trusted_Computing)

# Trusted Computing

## According to the *Trusted Computing Group*

Protect computing infrastructure at the hardware level by adding hardware extensions to enforce specific system capabilities, protecting against unauthorized access.

- **Endorsement Key, EK Certificate**: A key that never leaves the hardware
- **Memory curtaining**: provide isolation between memory regions
- **Sealed storage**: Bind data to specific hardware
- **Remote attestation**: authenticate the host to a remote host
- **Trusted third party** as an intermediary

**In practice:** different architectures, such as “enclaved” execution, memory encryption

## Possible Applications

### Digital rights management [edit]

Trusted Computing would allow companies to create a digital rights management (DRM) system that is difficult to bypass, though not impossible. An example is downloading a music file. Sealed storage could bind the file to a specific computer or player. If the file was played on an unauthorized player or computer, remote attestation could be used to authorize the device to play the music. The record company's rules. The music would be played from curtained memory, which would prevent the user from copying the file while it is playing, and secure I/O would prevent capturing what is being played. Any attempt to copy the file would be detected by the system, which would require either manipulation of the computer's hardware, capturing the audio signal before it reaches the recording device or a microphone, or breaking the security of the system.

New business models for use of software (services) over Internet may be boosted by the technology. One company could base a business model on renting programs for a specific time periods or "per play". For example, one could download a music file which could only be played a certain number of times before it becomes unusable. This could be used to prevent piracy only within a certain time period.

### Preventing cheating in online games [edit]

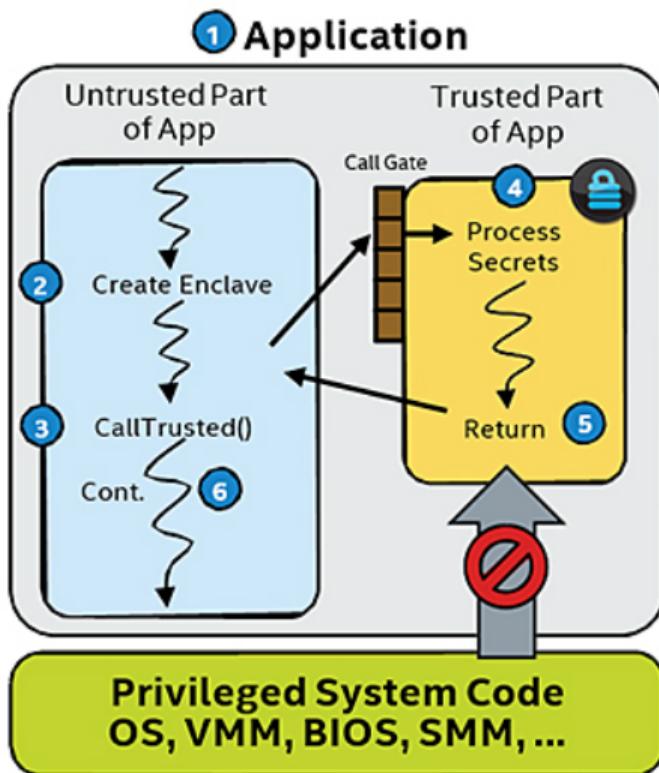
Trusted Computing could be used to combat [cheating in online games](#). Some players might be able to gain an advantage in the game; remote attestation, secure I/O and memory curtaining could be used to ensure that the player is running a server were running an unmodified copy of the software.<sup>[18]</sup>

### Verification of remote computation for grid computing [edit]

Trusted Computing could be used to guarantee participants in a [grid computing](#) system that they claim to be instead of forging them. This would allow large scale simulations to be run on multiple hosts using redundant computations to guarantee malicious hosts are not undermining the results.

Source: [https://en.wikipedia.org/wiki/Trusted\\_Computing](https://en.wikipedia.org/wiki/Trusted_Computing)

# Intel SGX Helicopter View



- Protected enclave in application's **virtual address space**
- Enclave can be entered through restrictive **call gate** only
- Provides **attestation** interface
- **Memory encryption** defends against untrusted system software and cold boot attacks

# Comparing Hardware-Based Trusted Computing Architectures

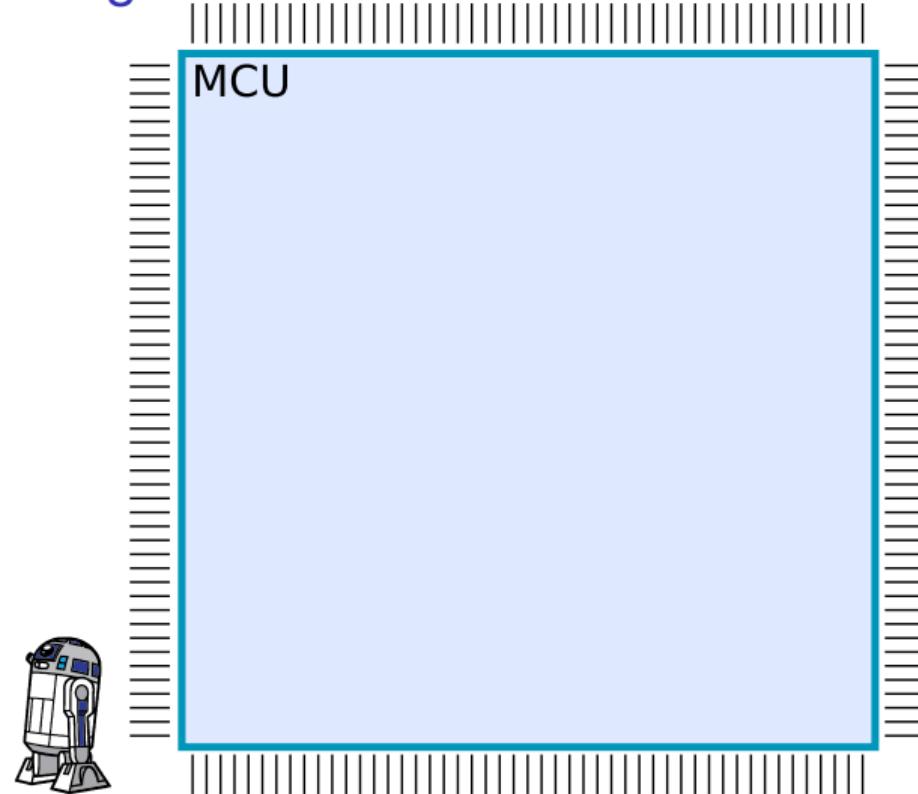
|              | Isolation | Attestation | Sealing | Dynamic RoT | Code Confidentiality | Side-Channel Resistance | Memory Protection | Lightweight | Coprocessor | HW-Only TCB | Preemption | Dynamic Layout | Upgradeable TCB | Backwards Compatibility | Open-Source | Academic | Target ISA    |
|--------------|-----------|-------------|---------|-------------|----------------------|-------------------------|-------------------|-------------|-------------|-------------|------------|----------------|-----------------|-------------------------|-------------|----------|---------------|
| AEGIS        | ●         | ●           | ●       | ●           | ●                    | ○                       | ●                 | ○           | ○           | ●           | ●          | ●              | ○               | ●                       | ○           | ●        | -             |
| TPM          | ○         | ●           | ●       | ○           | ●                    | -                       | ●                 | ○           | ●           | ●           | -          | -              | ○               | ●                       | ○           | ○        | -             |
| TXT          | ●         | ●           | ●       | ●           | ●                    | ●                       | ●                 | ○           | ●           | ●           | ○          | ●              | ●               | ●                       | ○           | ○        | x86_64        |
| TrustZone    | ●         | ○           | ○       | ●           | ●                    | ○                       | ○                 | ○           | ○           | ●           | ●          | ●              | ●               | ●                       | ○           | ○        | ARM           |
| Bastion      | ●         | ○           | ●       | ●           | ●                    | ●                       | ●                 | ○           | ○           | ●           | ●          | ●              | ●               | ●                       | ○           | ●        | UltraSPARC    |
| SMART        | ○         | ●           | ○       | ●           | ●                    | ○                       | -                 | ●           | ○           | ○           | -          | -              | ○               | ●                       | ○           | ●        | AVR/MSP430    |
| Sancus 1.0   | ●         | ●           | ○       | ●           | ●                    | ●                       | ●                 | ●           | ●           | ●           | ○          | ●              | ●               | ●                       | ●           | ●        | MSP430        |
| Soteria      | ●         | ●           | ●       | ●           | ●                    | ●                       | ●                 | ●           | ●           | ●           | ●          | ●              | ●               | ●                       | ●           | ●        | MSP430        |
| Sancus 2.0   | ●         | ●           | ●       | ●           | ●                    | ●                       | ●                 | ●           | ●           | ●           | ●          | ●              | ●               | ●                       | ●           | ●        | MSP430        |
| SecureBlue++ | ●         | ○           | ●       | ●           | ●                    | ●                       | ●                 | ○           | ○           | ●           | ●          | ●              | ●               | ●                       | ○           | ○        | POWER         |
| SGX          | ●         | ●           | ●       | ●           | ●                    | ●                       | ●                 | ○           | ○           | ○           | ●          | ●              | ●               | ●                       | ○           | ○        | x86_64        |
| Iso-X        | ●         | ●           | ○       | ●           | ●                    | ○                       | ●                 | ●           | ○           | ●           | ●          | ●              | ●               | ●                       | ○           | ●        | OpenRISC      |
| TrustLite    | ●         | ●           | ○       | ●           | ●                    | ●                       | ●                 | ●           | ●           | ●           | ●          | ●              | ●               | ●                       | ○           | ●        | Siskiyou Peak |
| TyTAN        | ●         | ●           | ●       | ●           | ●                    | ●                       | ●                 | ●           | ●           | ●           | ●          | ●              | ●               | ●                       | ○           | ●        | Siskiyou Peak |
| Sanctum      | ●         | ●           | ●       | ●           | ●                    | ●                       | ●                 | ●           | ○           | ○           | ○          | ●              | ●               | ●                       | ●           | ●        | RISC-V        |

● = Yes; ○ = Partial; ○ = No; - = Not Applicable

Adapted from  
“Hardware-Based  
Trusted Computing  
Architectures for  
Isolation and  
Attestation”, Maene et  
al., IEEE Transactions  
on Computers, 2017.  
[MGdC<sup>+</sup>17]

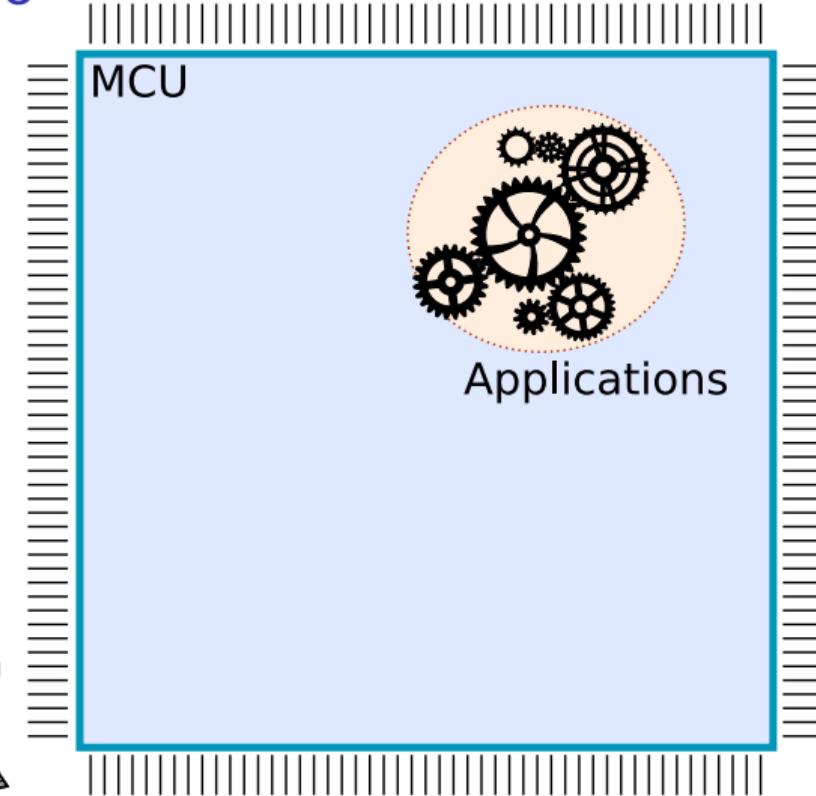
# Isolation and Attestation on Light-Weight MCUs

Many microcontrollers feature little security functionality



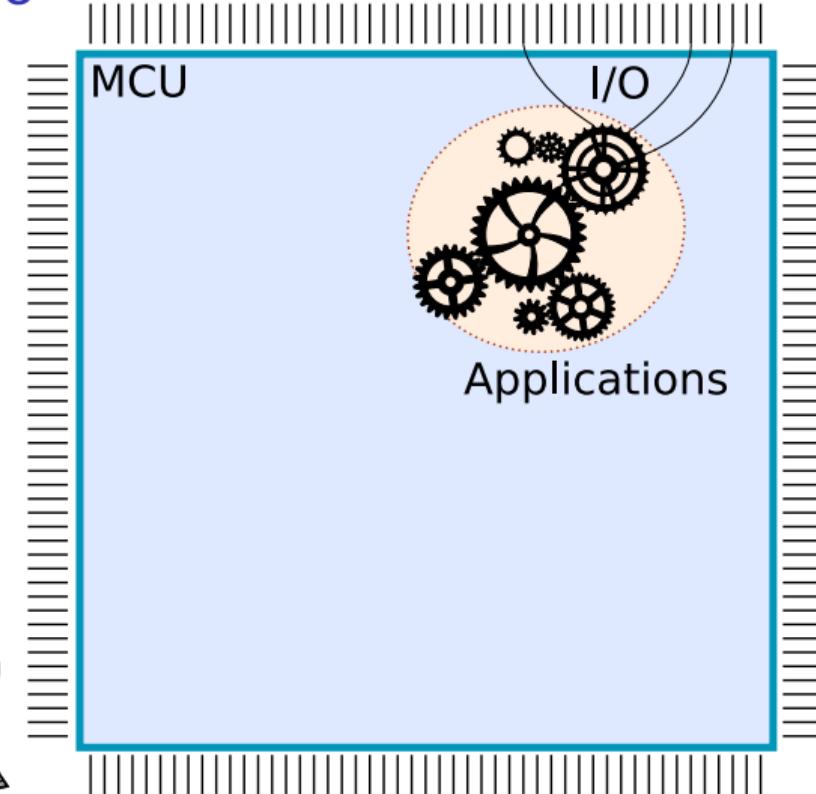
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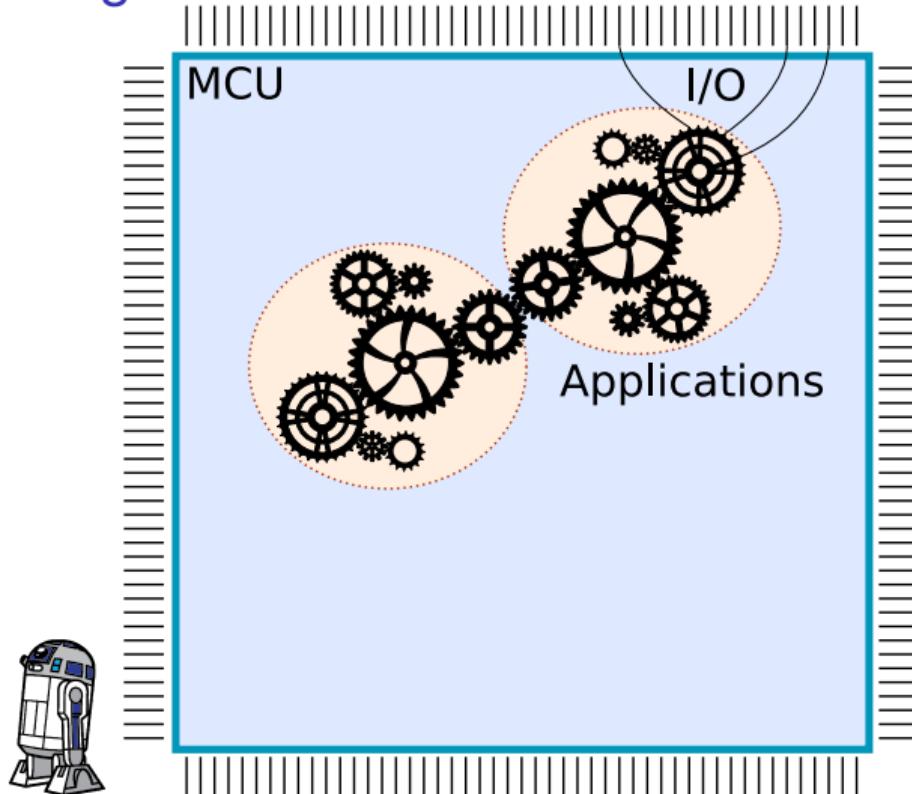
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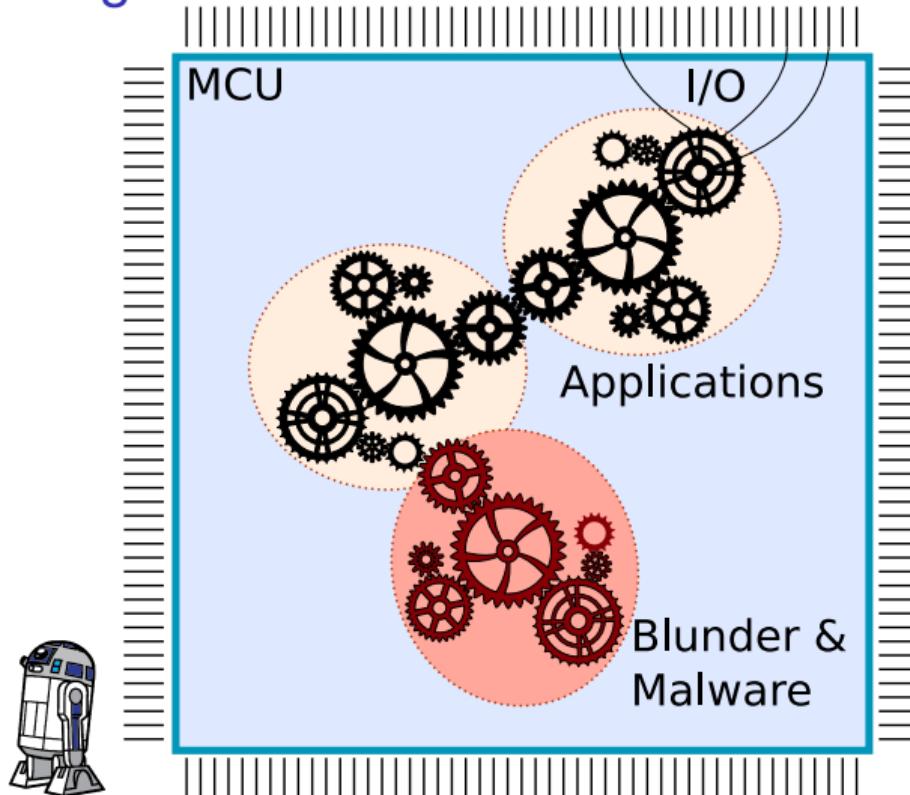
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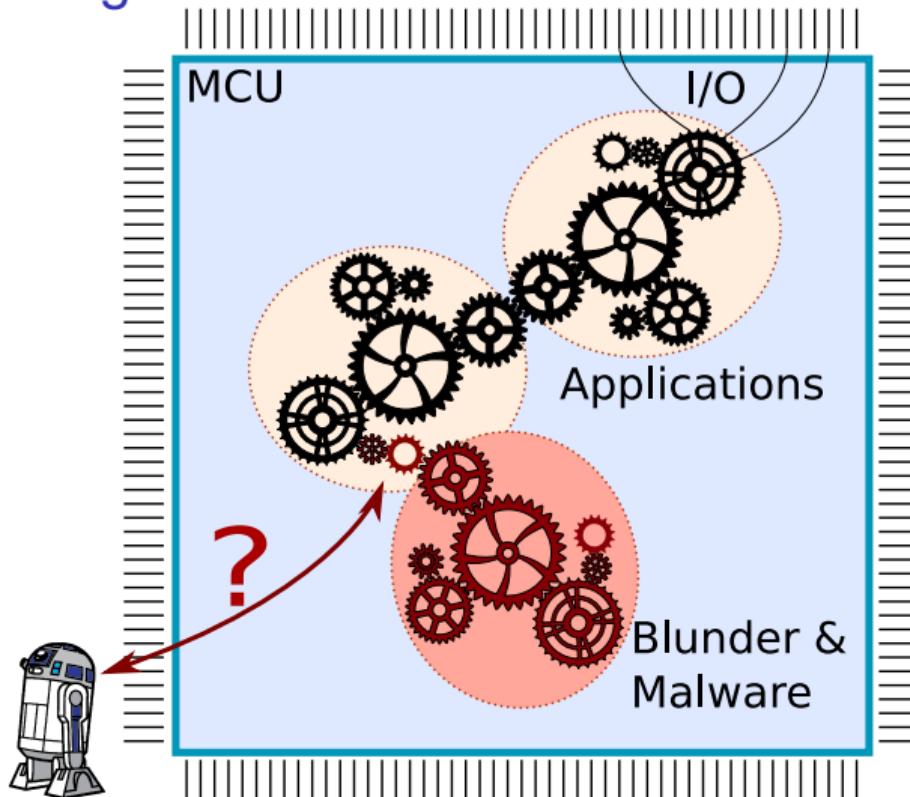
- Applications share address space
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Authenticity?



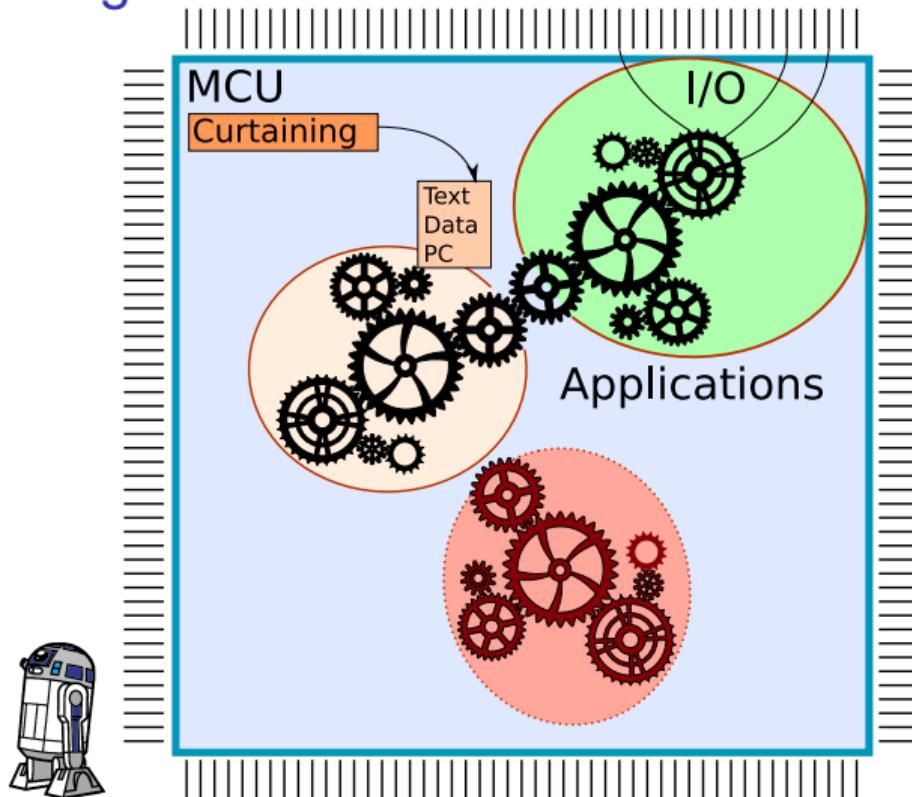
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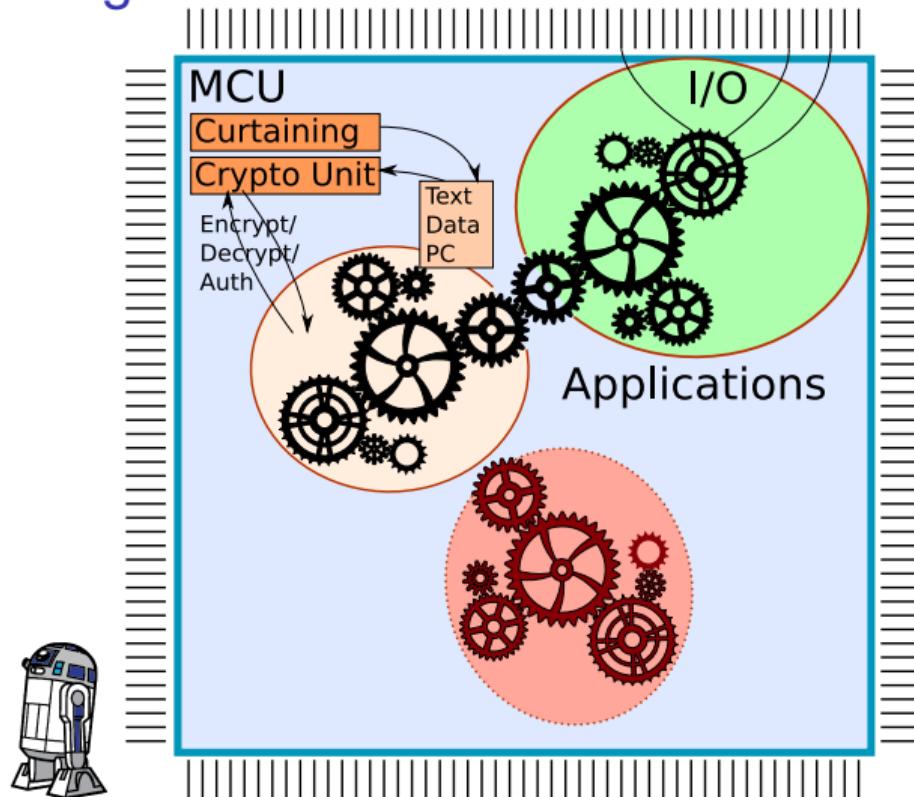
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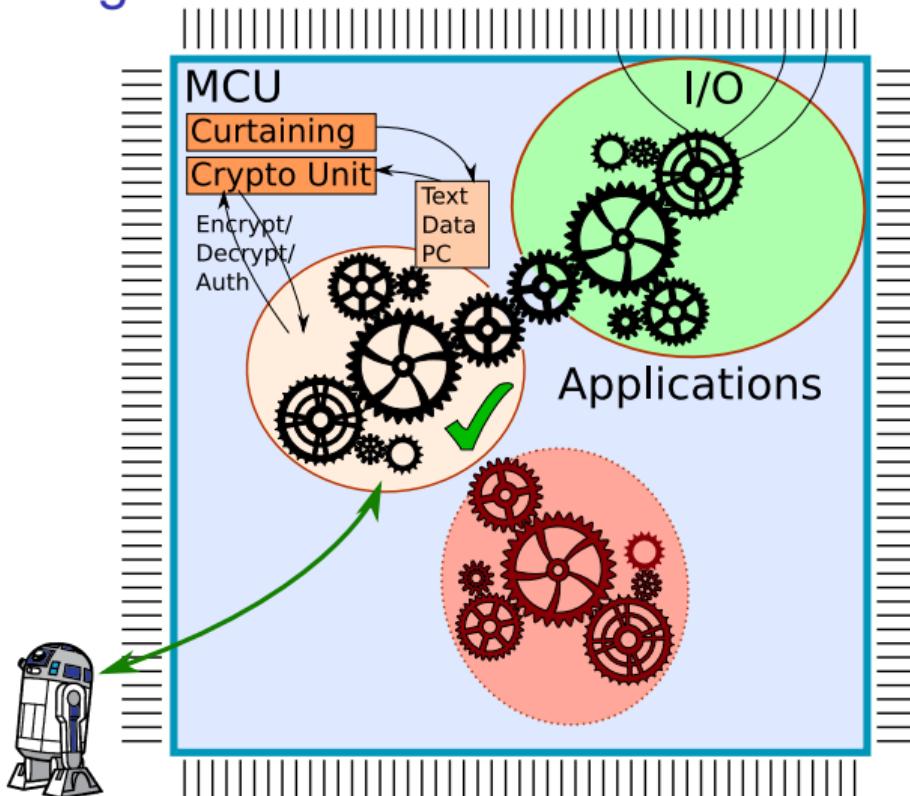
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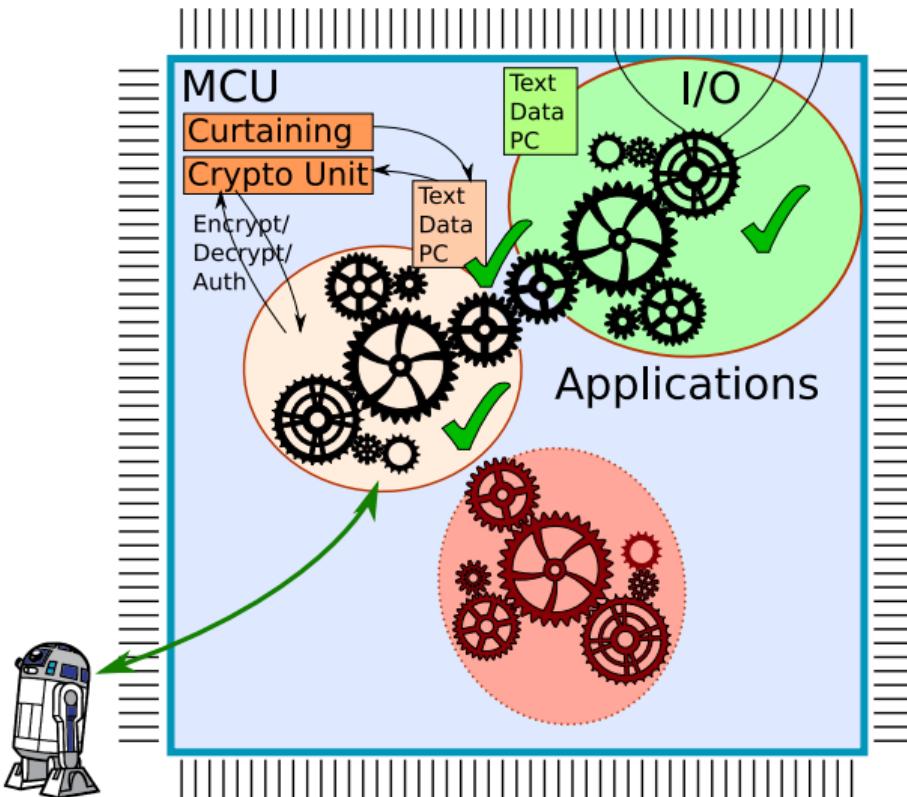
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# Sancus: Strong and Light-Weight Embedded Security [NVBM<sup>+</sup>17]

## Extends openMSP430 with strong security primitives

- Software Component Isolation
- Cryptography & Attestation
- Secure I/O through isolation of MMIO ranges

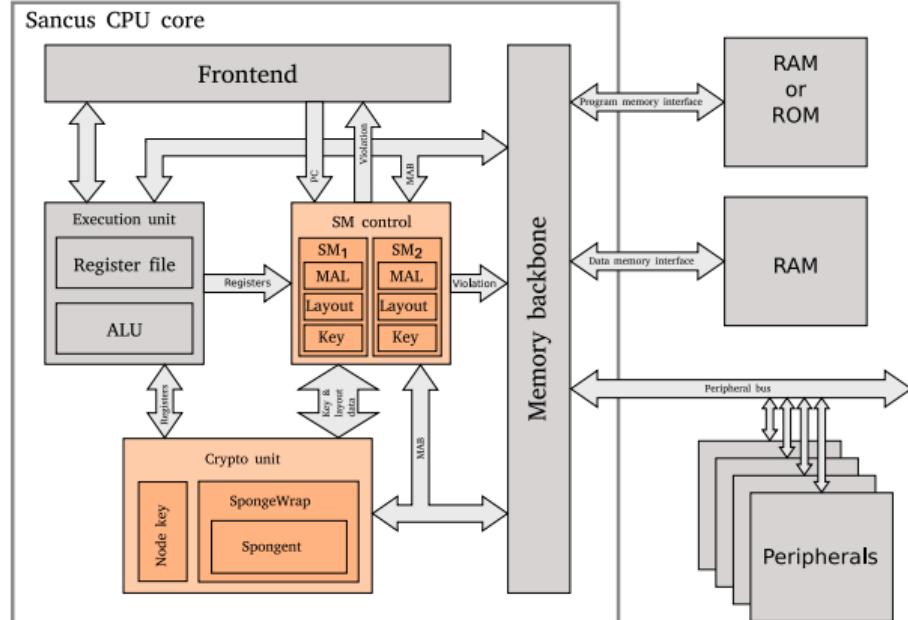
## Efficient

- Modular,  $\leq 2$  kLUTs
- Authentication in  $\mu$ s
- + 6% power consumption

## Cryptographic key hierarchy for software attestation

Isolated components are typically very small ( $< 1\text{kLOC}$ )

Sancus is Open Source: <https://distrinet.cs.kuleuven.be/software/sancus/>



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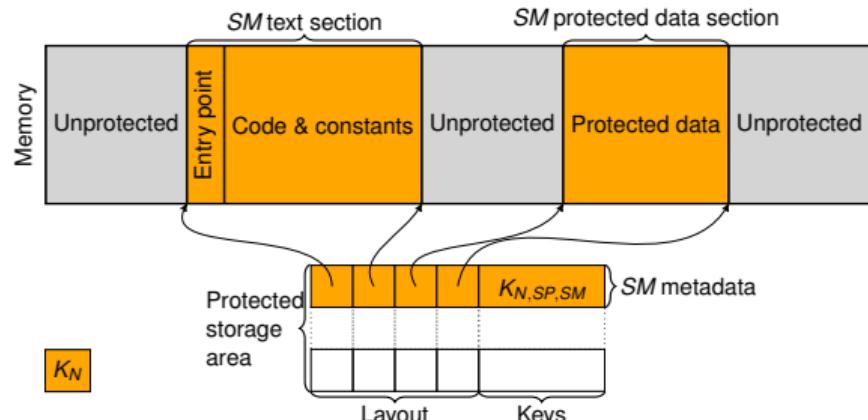
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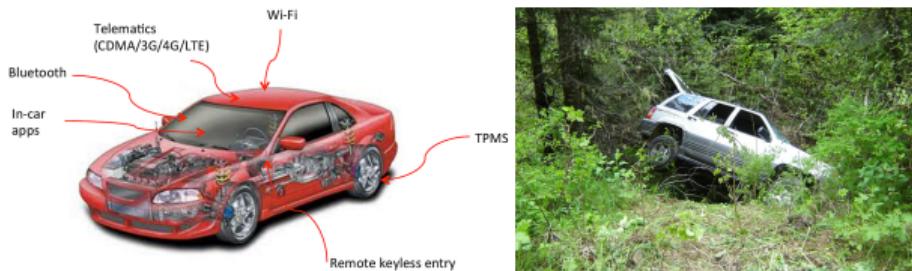
$N$  = Node;  $SP$  = Software Provider / Deployer  
 $SM$  = protected Software Module



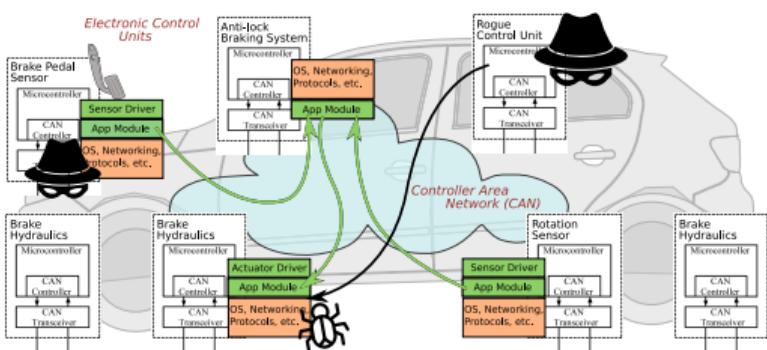
# Secure Automotive Computing with Sancus [VBMP17]

## Modern cars can be hacked!

- Network of more than 50 ECUs
- Multiple communication networks
- Remote entry points
- Limited built-in security mechanisms



Miller & Valasek, "Remote exploitation of an unaltered passenger vehicle", 2015



## Sancus brings strong security for embedded control systems:

- Message authentication
- Trusted Computing: software component isolation and cryptography
- Strong software security
- Applicable in automotive, ICS, IoT, ...

# Secure Automotive Computing with Sancus [VBMP17]



# Authentic Execution of Distributed Event-Driven Applications



“Authentic Execution of Distributed Event-Driven Applications with a Small TCB”,  
Noorman et al., STM 2017. [NMP17]

## When **not** to trust your TEE...

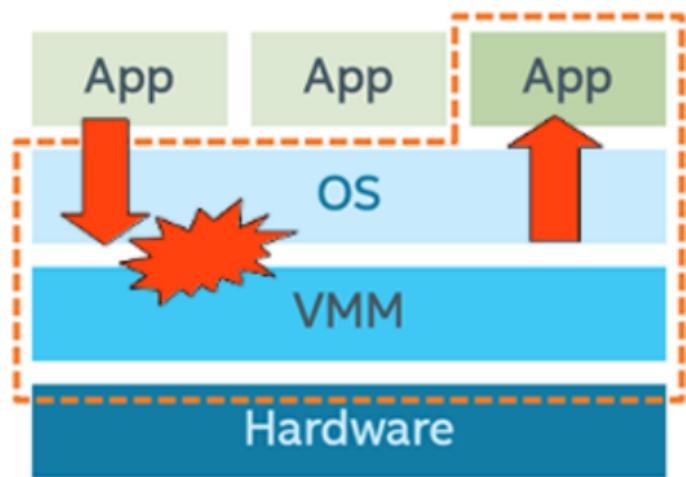
**Trusted Execution does not help you against bugs in your own (trusted) code.**

**Trusted Execution does not help you if you don't know what to protect.**

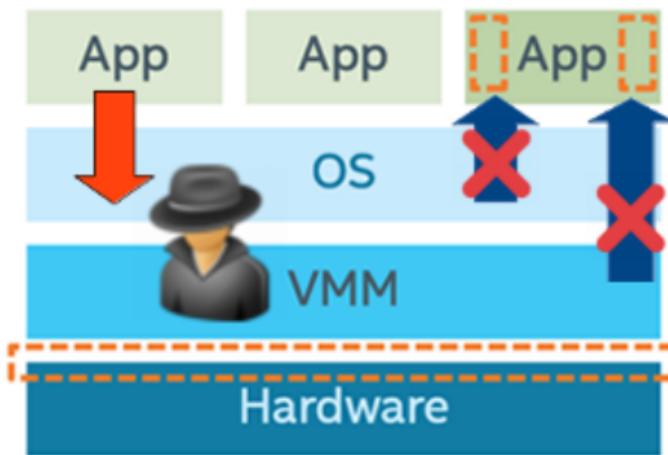
**(Trusted) Execution can be observed through indirect channels and may leak secrets through these channels.**

# Motivation: Application Attack Surface

Attack Surface Without Enclaves



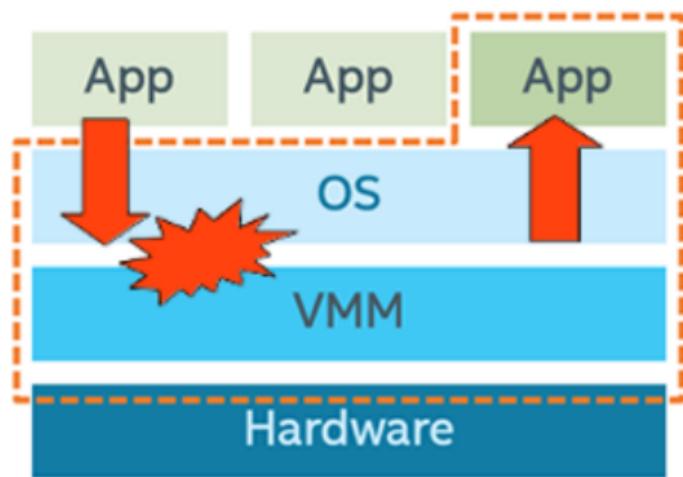
Attack Surface With Enclaves



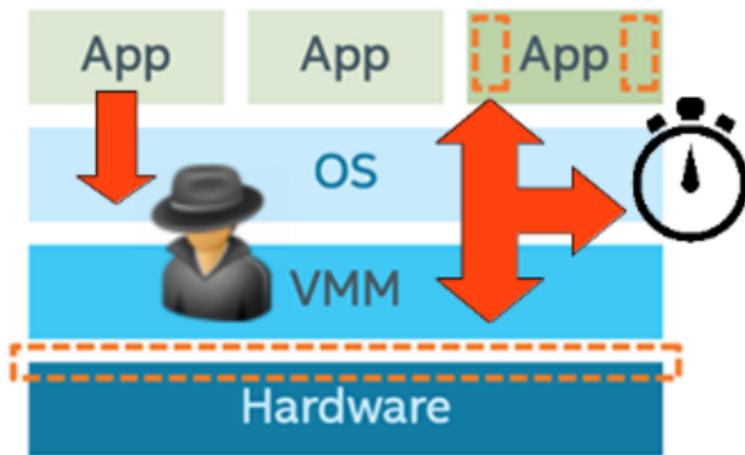
Layered architecture  $\leftrightarrow$  **hardware-only TCB**

# Motivation: Application Attack Surface

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Attack Surface With Enclaves



Attack Surface

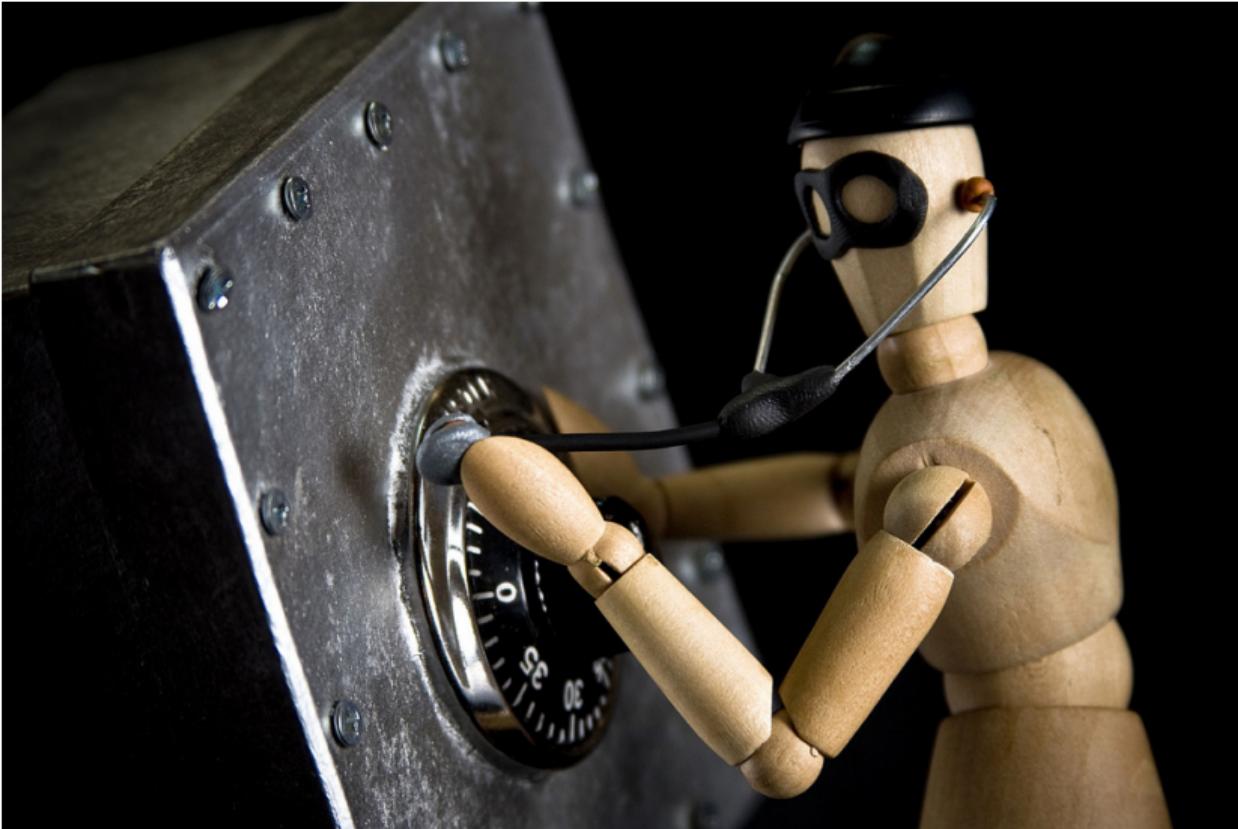
Untrusted OS → new class of powerful **side-channels**

# Side-Channel Attack Principle



Source: <https://commons.wikimedia.org/wiki/File:WinonaSavingsBankVault.JPG>

# Side-Channel Attack Principle



# Summary

## [Background]

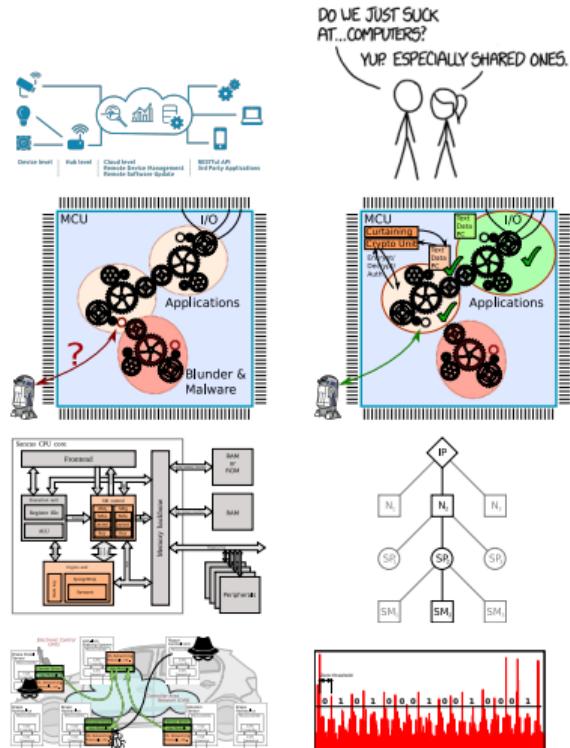
- ① Software vulnerabilities are hard to eliminate and can be exploited by attackers
- ② Even correct code needs protection against layer-below attacks!

## Trusted Execution Technology

- ① Strong application isolation and attestation: hardware-level security and taming complexity
- ② No protection against buggy software!
- ③ Potential for invasive use

## Sancus

- ① The Open-Source Trusted Computing Architecture
- ② Built upon openMSP430 16-bit MCU, applications in IoT and embedded control systems
- ③ Research prototype under active development!



# Trusted Execution for Everyone

**Fortanix** solves cloud security and privacy using runtime encryption technology build upon Intel SGX. <https://fortanix.com/>

**SCONE** enables secure execution of containers and programs using Intel SGX. <https://sconecontainers.github.io/>

**Graphene-SGX**: A practical library OS for unmodified applications on SGX. <https://github.com/oscarlab/graphene>

**Open Enclave** is an SDK for building enclave applications in C and C++. <https://github.com/Microsoft/openenclave>

**Our Tutorial:** Building distributed enclave applications with Sancus and SGX <https://github.com/sancus-pma/tutorial-dsn18>

# The Impact of ICT...

... and why the right choice of sustainable solutions really matters.



**Image sources:** Electronic waste recycling in Ghana, <https://en.wikipedia.org/>; Martin Falbisoner, Garzweiler surface mine, <https://en.wikipedia.org/>; Sebastian Meyer, "Blood, Sweat and Batteries", <https://www.sebmeyer.com/>

# Food For Thoughts

Trusted Computing in **public communications infrastructure**

Trusted Computing to **protect critical industrial infrastructure**

**Secret computations on secret data**, executing on a public clouds

**Hiding malware, computing cryptographic signatures on other people's computers**

Thank you!

**“The risks are about to get worse, because computers are being embedded into physical devices and will affect lives, not just our data.”**

— Bruce Schneier, [Sch18]

Thank you! Questions?

<https://distrinet.cs.kuleuven.be/>

<https://github.com/sancus-pma/tutorial-dsn18>

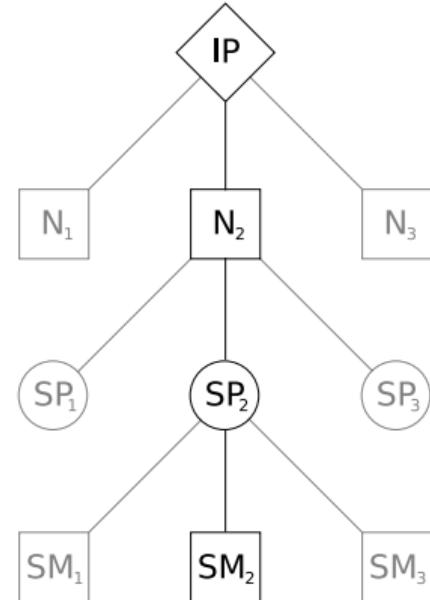
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# Attestation and Communication with Sancus

**Ability to use  $K_{N,SP,SM}$  proves the integrity and isolation  
of SM deployed by SP on N**

- Only  $N$  and  $SP$  can compute  $K_{N,SP,SM}$   
 $N$  knows  $K_N$  and  $SP$  knows  $K_{SP}$
- $K_{N,SP,SM}$  on  $N$  is computed after enabling isolation  
No isolation, no key; no integrity, wrong key
- Only  $SM$  on  $N$  is allowed to use  $K_{N,SP,SM}$   
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Through special instructions

**Remote attestation and secure communication by  
Authenticated Encryption with Associated Data**

- Confidentiality, integrity and authenticity
- Encrypt and decrypt instructions use  $K_{N,SP,SM}$  of the calling SM
- Associated Data can be used for nonces to get freshness

