

# MATH-458 Programming Concepts in Scientific Computing: Nonlinear Solvers

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## 1 Introduction

## 2 Algorithm

### 2.1 Bisection Method

The bisection method is based on Theorem 1.

**Theorem 1.** (*Zeros of a continuous functions*)

*Given a continuous function  $f : [a,b] \rightarrow \mathbb{R}$ , such that  $f(a)f(b) < 0$ , then  $\exists$  at least an  $\alpha \in (a,b)$  such that  $f(\alpha) = 0$ .*

Starting from  $I_0 = [a,b]$ , the bisection method generates a sequence of subintervals  $I_k = [a(k), b(k)]$ ,  $k \geq 0$ , with  $I_k \subset I_{k-1}$ ,  $k \geq 1$ , and enjoys the property that  $f(a(k))f(b(k)) < 0$ .

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**Algorithm 1** Bisection Method

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```
1:  $k \leftarrow 0$ ,  $a^{(0)} \leftarrow a$ ,  $b^{(0)} \leftarrow b$ ,  $x^{(0)} \leftarrow (a^{(0)} + b^{(0)})/2$ 
2: while  $x^{(k)} > \text{tol}$  and  $k < k_{\max}$  do
3:   if  $f(x^{(k-1)}) == 0$  then return  $\alpha = x^{(k-1)}$ 
4:   else
5:     if  $f(x^{(k-1)})f(a^{(k-1)}) < 0$  then
6:       set  $a^{(k)} = a^{(k-1)}$  and  $b^{(k)} = x^{(k-1)}$ 
7:     if  $f(x^{(k-1)})f(b^{(k-1)}) < 0$  then
8:       set  $a^{(k)} = x^{(k-1)}$  and  $b^{(k)} = b^{(k-1)}$ 
9:     set  $x^{(k)} = (a^{(k)} + b^{(k)})/2$ 
return  $\alpha = x^{(k)}$ 
```

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### 2.2 Newton/Quasi-Newton Methods

Assuming that  $f \in C^0(I)$  and is differentiable in the interval  $I = (a,b) \subseteq \mathbb{R}$ , the equation of a tangent line to the curve  $(x, f(x))$  at coordinate  $x^{(k)}$ , where  $x^{(k)} \in I$ , is  $y(x) = f(x^{(k)}) + df(x^{(k)})(x - x^{(k)})$ . If we assume  $y(x^{(k)}) = 0$ ,  $x^{(k+1)}$  can be computed from Eq. 1.

$$x^{(k+1)} = x^{(k)} - \frac{f(x^{(k)})}{q^{(k)}} \quad (1)$$

with  $q^{(k)}$  determining the method as Newton (Eq. 2) or Chord (Eq. 3):

$$q^{(k)} = \frac{df(x^{(k)})}{dx} \quad (2)$$

$$q^{(k)} = \frac{f(x^{(k)}) - f(x^{(k-1)})}{x^{(k)} - x^{(k-1)}} \quad (3)$$

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**Algorithm 2** General quasi-Newton Method

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```
1:  $k \leftarrow 0$ , set initial guess  $x^{(0)}$ 
2: while  $x^{(k)} > \text{tol}$  and  $k < k_{max}$  do
3:    $x^{(k+1)} = x^{(k)} - \frac{f(x^{(k)})}{(q^{(k)})}$ 
4:    $k = k + 1$ ;
   return  $\alpha = x^{(k)}$ 
```

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### 2.3 Fixed Point Iteration

For a given  $f : [a,b] \rightarrow \mathbb{R}$ , the problem can be transformed from  $f(x) = 0$  into an equivalent problem  $x - \phi(x) = 0$ . The auxiliary function  $\phi : [a,b] \rightarrow \mathbb{R}$  has to be chosen in such a way that  $\phi(\alpha) = \alpha$  whenever  $f(\alpha) = 0$ . Finding the fixed points of the mapping  $\phi$  thus results in finding the roots of the original equation  $f(x)$ , hence the name fixed point iteration.

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**Algorithm 3** Fixed point iterations

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```
1:  $k \leftarrow 0$ , set initial guess  $x^{(0)}$ 
2: while  $x^{(k)} - x^{(k-1)} > \text{tol}$  and  $k < k_{max}$  do
3:    $x^{(k+1)} = \phi(x^{(k)})$ 
4:    $k = k + 1$ ;
   return  $\alpha = x^{(k)}$ 
```

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### 2.4 Aitken Method

Aitken's method provides a way to accelerate the convergence of iterative methods for finding the roots of a function. The general algorithm is shown below.

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**Algorithm 4** Aitken Method

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```
1:  $k \leftarrow 0$ , set initial guess  $x^{(0)}$ 
2: while  $x^{(k)} - x^{(k-1)} > \text{tol}$  and  $k < k_{max}$  do
3:   Calculate  $x_{3i+1}$  and  $x_{3i+2}$  using any linear iterative method.
4:   Modify  $x_{3i+2}$  using  $x_{3i+2} = x_{3i} - \frac{(x_{3i+1} - x_{3i})^2}{(x_{3i} - 2x_{3i+1} + x_{3i+2})}$ 
5:    $k = k + 1$ ;
   return  $\alpha = x^{(k)}$ 
```

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## 3 Program Structure