Computer Architectures Exam of February 22, 2021 Part 1 – possible answers

Exe #1, Point #1

 Explain what Loop Unrolling is, stating who is in charge of applying it

 Loop unrolling is a static technique based on reducing the number of iterations a given loop is executed, modifying its body. It is normally implemented by the compiler.

Exe #1, Point #2

Describe the advantages and disadvantages it introduces

- Loop unrolling improves the performance by
 - Reducing the number of branches
 - Increasing the size of the loop body, often increasing the chances of identifying ILP in it
- Its main disadvantage lies in the increased code size.

Exe #1, Point #3

```
for (i=0;i<MAX;i++ )
         y[i] = x[i] + 5;
for (i=0;i<MAX;i+=4)
         y[i] = x[i] + 5;
         y[i+1] = x[i+1] + 5;
         y[i+2] = x[i+2] + 5;
         y[i+3] = x[i+3] + 5;
```

Exe #2, Point #1

Let consider a processor including a Branch Target Buffer (BTB).

Assuming that the processor uses 32 bit addresses, each instruction is 4 byte wide, and the BTB is composed of 4 entries, you are requested to

1. Describe the architecture of the BTB

The BTB is a table composed of 4 entries, each composed of 2 32-bit fields: Address and Target.

An additional 1 bit field may be present.

Exe #2, Point #2

2. Describe in details the behavior of a BTB, explaining when it is accessed, and which input and output information are involved with each access

The BTB is accessed each time an instruction is fetched.

Using the least significant log n bits of the instruction address (excluding the word alignment bits) an entry is selected. In this case n=4. The Address field of the entry is compared with the address of the fetched instruction: if they match, the Target field is uploaded in the PC.

When the instruction is completed, the BTB is possibly updated.

Exe #2, Point #3

- 3. Identify the final content of the BTB if
 - The BTB is initially empty (i.e., full of 0s)
 - The following instructions are executed in sequence

add.d f1,f2,f3	located at the address 0x00A50050
bnez r4,l1	located at the address 0x00A50054; the branch is taken, and the branch target address is 0x00A60050
mul.d f5,f6,f7	located at the address 0x00A60050
daddi r2,r2,-1	located at the address 0x00A60054
bez r2,12	located at the address 0x00A60058; the branch is not taken
j 13	located at the address 0x00A6005C;
	the branch target address is 0x00A60AA0

Exe ##2, Point #3: BTB initial content

address	target
0x00000000	0x0000000

Exe #2, Point #3: BTB content after add.d execution

add.d address	
0x00A50050	0000 0000 1010 0101 0000 0000 0101 00 00

address	target
0x00000000	0x0000000

Entry #0 is accessed

Exe #2, Point #3: BTB content after bnez execution

bnez address		
0x00A50054	0000 0000 1010 0101 0000 0000 0101 01 00	

address	target
0x00000000	0x0000000
0x00A50054	0x00A60050
0x00000000	0x0000000
0x00000000	0x0000000

Entry #1 is accessed

Exe #2, Point #3: BTB content after mul.d execution

mul.d address		
0x00A60050	0000 0000 1010 0110 0000 0000 0101 00 00	

address	target
0x00000000	0x0000000
0x00A50054	0x00A60050
0x00000000	0x0000000
0x00000000	0x0000000

Entry #0 is accessed

Exe #2, Point #3: BTB content after daddi execution

daddi address	
0x00A60054	0000 0000 1010 0110 0000 0000 0101 01 00

address	target
0x00000000	0x0000000
0x00A50054	0x00A60050
0x00000000	0x0000000
0x00000000	0x0000000

Entry #1 is accessed

Exe #2, Point #3: BTB content after bez execution

bez address	
0x00A60058	0000 0000 1010 0110 0000 0000 0101 10 00

address	target
0x00000000	0x000000
0x00A50054	0x00A60050
0x00000000	0x0000000
0x00000000	0x000000

Entry #2 is accessed

Exe #2, Point #3: BTB content after j execution

j address	
0x00A6005C	0000 0000 1010 0110 0000 0000 0101 11 00

address	target
0x00000000	0x0000000
0x00A50054	0x00A60050
0x00000000	0x0000000
0x00A6005C	0x00A60AA0

Entry #3 is accessed

Exe #2, Point #3: BTB final content

address	target
0x0000000	0x0000000
0x00A50054	0x00A60050
0x00000000	0x0000000
0x00A6005C	0x00A60AA0