

Applications of Mathematics in Computer Science (MACS)

LECTURE #2: CHECKSUMS AND HASH CODES

HEY, I LOST THE
SERVER PASSWORD.
WHAT IS IT, AGAIN?



IT'S— ...WAIT.
HOW DO I KNOW
IT'S REALLY YOU?



OOH, GOOD QUESTION!
I BET WE CAN CONSTRUCT A COOL
PROOF-OF-IDENTITY PROTOCOL. I'LL
START BY PICKING TWO RANDOM—

OH GOOD; IT'S YOU.
HERE'S THE PASSWORD...



The CS Problems:

- Redundancy in numbers
- Data integrity checking
- Authorization
- Hash tables

Redundancy in numbers

How can we quickly check that a number is correct?

- ID number

318892652*



3188926525

Why 5?

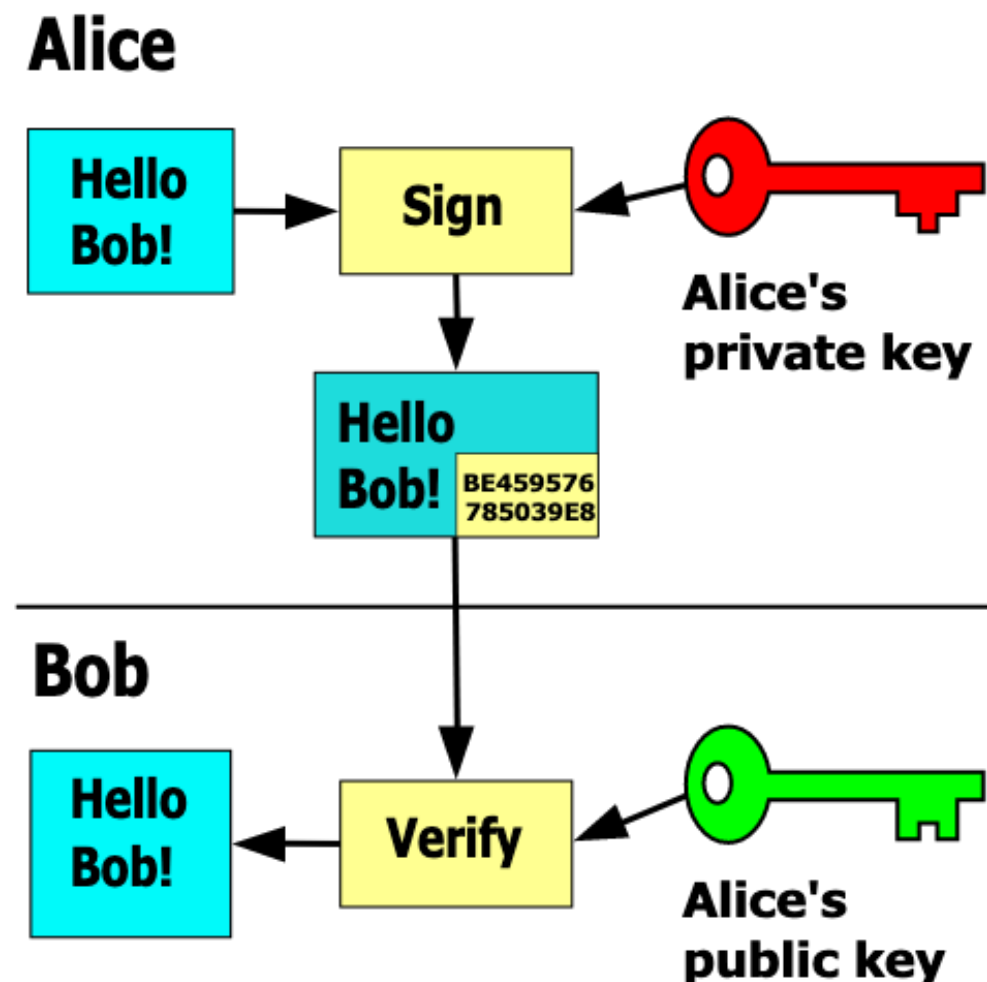
ספרת ביקורת

- “card number is invalid”

How do they know?

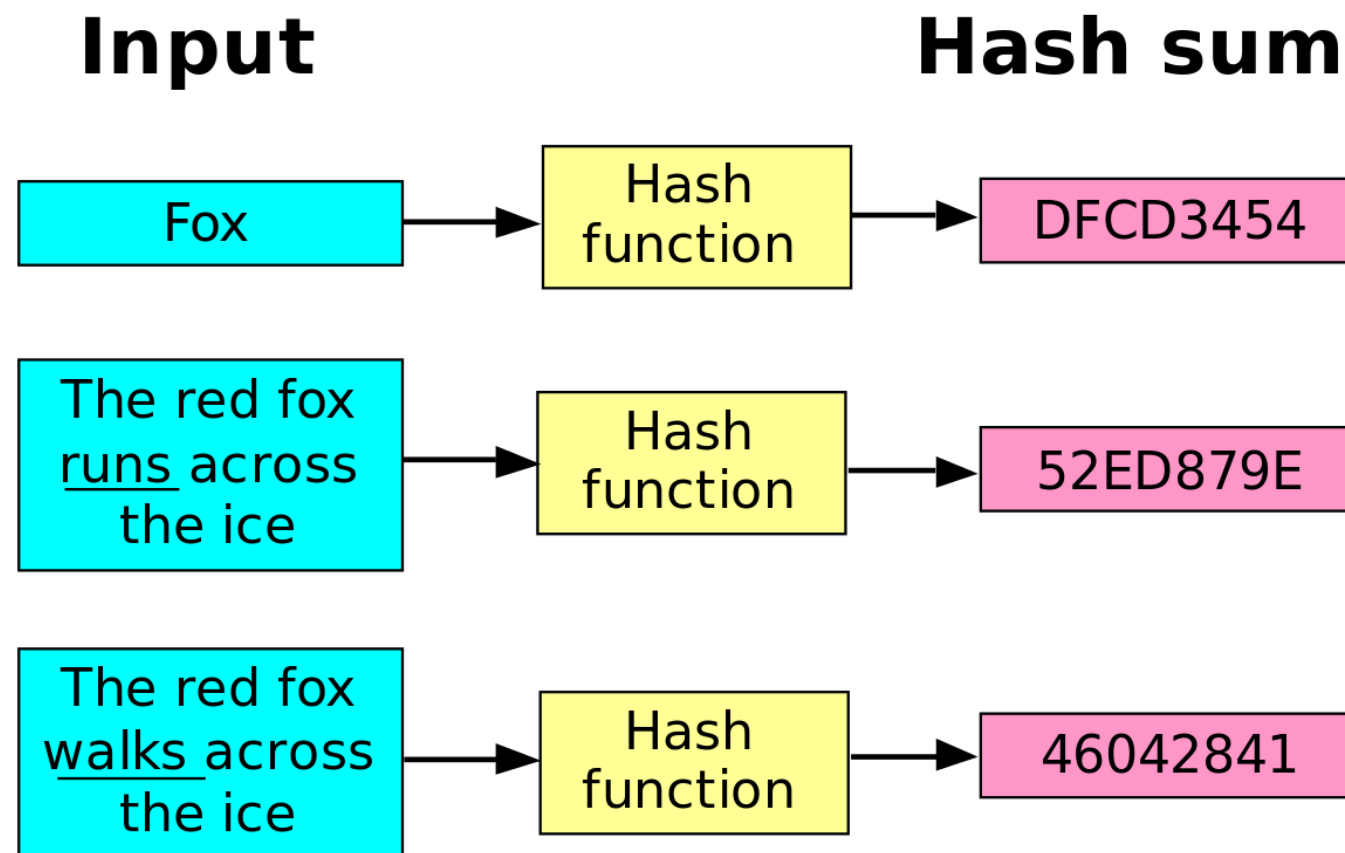
Data integrity checking

- Quickly check that a download is complete
- Verify a document:



Authorization

- Checking a person's identity
 - ID/password



Hash Tables (Dictionaries)

```
>>> tel = {'jack': 4098, 'sape': 4139}
>>> tel['guido'] = 4127
>>> tel
{'jack': 4098, 'sape': 4139, 'guido': 4127}
>>> tel['jack']
4098
>>> del tel['sape']
>>> tel['irv'] = 4127
>>> tel
{'jack': 4098, 'guido': 4127, 'irv': 4127}
>>> list(tel)
['jack', 'guido', 'irv']
>>> sorted(tel)
['guido', 'irv', 'jack']
>>> 'guido' in tel
True
>>> 'jack' not in tel
False
```

How can we construct
such a data type?

The math technique:
Modular Arithmetics

Modular Arithmetics

$$m = a \% b$$

$$0 \leq m < b$$



$$\exists n . a = n \cdot b + m$$

$$n \% 1 = 0$$

$$3 \% 2 = 1$$

$$10 \% 3 = 1$$

$$5 \% 3 = 2$$

Modular Equivalence

$$a \equiv b \pmod{c}$$



$$a \% c = b \% c$$

$$9 \equiv 5 \pmod{4}$$

$$4 \equiv 2 \pmod{2}$$

$$3 \equiv 111 \pmod{2}$$

- Reflexivity: $a \equiv a \pmod{n}$
- Symmetry: $a \equiv b \pmod{n} \Leftrightarrow b \equiv a \pmod{n}$
- Transitivity:

$$a \equiv b \pmod{n} \wedge b \equiv c \pmod{n} \Leftrightarrow a \equiv c \pmod{n}$$

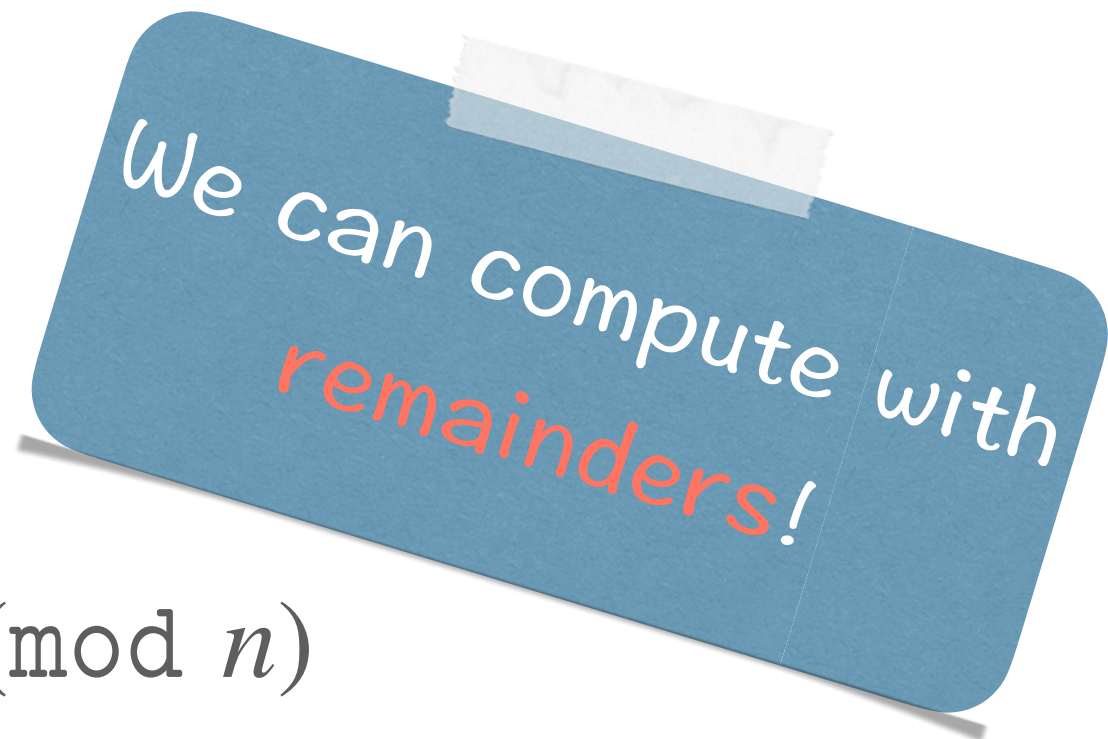
Properties

If $a \equiv b \pmod{n}$, then:

- $a + k \equiv b + k \pmod{n}$
- $ka \equiv kb \pmod{n}$ (for integer k)

If $a_1 \equiv b_1 \pmod{n}$ and $a_2 \equiv b_2 \pmod{n}$, then:

- $a_1 + a_2 \equiv b_1 + b_2 \pmod{n}$
- $a_1 - a_2 \equiv b_1 - b_2 \pmod{n}$
- $a_1 \cdot a_2 \equiv b_1 \cdot b_2 \pmod{n}$



Example

- S : 1000 integers sent
- R : 1000 integers received

$S=R?$

```
def checksum(A):  
    cksum = 0  
    for a in A:  
        cksum += a % 101  
    return cksum % 101
```

Will ~~overflow~~ ~~underflow~~ if integers are big

Let's try...