

Question 1

- a. $\{f : [T2 \rightarrow T3], g : [T1 \rightarrow T2], a : \text{Number}\} \vdash (f (g a)) : T3$
- b. $\{f : [T1 \rightarrow [T2 \rightarrow \text{Boolean}]], x : T1, y : T2\} \vdash (f x y) : \text{Boolean}$
- c. $\{f : [T1 \times T2 \rightarrow T3], y : T2\} \vdash (\text{lambda } (x) (f x y)) : [T1 \rightarrow T3]$
- d. $\{f : [T2 \rightarrow T1], x : T1, y : T3\} \vdash (f x) : T1$

- s' is applied to the type-expressions of s , i.e., for every variable T' for which $s'(T')$ is defined, occurrences of T' in type expressions in s are replaced by $s'(T')$.
- A variable T' in s' , for which $s(T)$ is defined, is removed from the domain of s' , i.e., $s'(T)$ is not defined on it anymore.
- The modified s' is added to s .
- Identity bindings, i.e., $s(T) = T$, are removed.
- If for some variable, $(s \circ s')(T)$ includes T , the combination fails.

https://bguppl.github.io/interpreters/class_material/3.3TypeInferenceSystem.html

a)
acc = {}
a = number
T1 = a
T2 = T2
f = [T2 -> T3]
g = [T1 -> T2]

—
acc = {a = number}
{T1 = a} o acc => {T1 = number }

acc o {T1 = number} = > acc = {a = number, T1 = number}
T2 = T2
f = [T2 -> T3]
g = [T1 -> T2]

—
{T2 = T2} o acc => {T2 = T2,...}
acc o {T2 = T2} => acc = {a = number, T1 = number, T2 = T2}

f = [T2 -> T3]
g = [T1 -> T2]

—

$\{f = [T2 \rightarrow T3]\} \circ acc \Rightarrow \{f = [T2 \rightarrow T3]\}$

$acc \circ \{f = [T2 \rightarrow T3]\} \Rightarrow acc = \{f = [T2 \rightarrow T3], a = \text{number}, T1 = \text{number}, T2 = T2\}$

$g = [T1 \rightarrow T2]$

—

$\{g = [T1 \rightarrow T2]\} \circ acc \Rightarrow \{g = [\text{Number} \Rightarrow T2]\}$

$acc \circ \{g = [T1 \rightarrow T2]\} \Rightarrow acc = \{f = [T2 \rightarrow T3], a = \text{number}, T1 = \text{number}, T2 = T2, g = [\text{number} \Rightarrow T2]\}$

—

Finally to calculate the inference:

MGU: $acc = \{f = [T2 \rightarrow T3], a = \text{number}, T1 = \text{number}, T2 = T2, g = [\text{number} \Rightarrow T2]\}$

$(f (g a))$

$(f ([\text{number} \rightarrow T2] \text{number}))$

$(f T2)$

$([T2 \rightarrow T3] T2)$

$T3$

b) This inference is wrong.

inferred equations:

f: $T3 \rightarrow \text{boolean}$

x: $T4$

y: $T5$

$T3: T4 * T5$

This clashes with environment substitutions:

f : $[T1 \Rightarrow [T2 \Rightarrow \text{boolean}]]$

x : $T1$

y: $T2$

solving the equations we reach

$T4 = T1$

$T5 = T2$

$T1 = T1 * T2$

which is impossible.

C) step 1 substitute:

$(\lambda x. (f x y)) \rightarrow_{\text{sub}} \lambda x. ([T1xT2 \rightarrow T3] x T2)$

stage 2: type variables for every equation:

$T0 = (\lambda x. (f x y))$

$T4 = (f x y)$

$T5 = x$

$T6 = y$

$T7 = f$

stage 3: generate type equations

$T0 = [T5 \rightarrow T4]$

$T7 = [T5xT6 \Rightarrow T4]$

$T6 = T2$ (from substitution)

$T7 = [T1 x T2 \rightarrow T3]$ (from substitution)

step 4: solve

$\text{acc} = \{\}$

$T0 = [[T5 \rightarrow T4]$

$T7 = [T5xT6 \Rightarrow T4]$

$T6 = T2$

$T7 = [T1 x T2 \rightarrow T3]$

—

$\text{acc} = \{ T0 = [[T5 \rightarrow T4]] \}$

$T7 = [T5xT6 \Rightarrow T4]$

$T6 = T2$

$T7 = [T1 x T2 \rightarrow T3]$

$\{T7 = [T5xT6 \Rightarrow T4]\} \circ \text{acc} \Rightarrow \{T7 = [T5xT6 \Rightarrow T4]\}$

$\text{acc} \circ \{T7 = [T5xT6 \Rightarrow T4]\} \Rightarrow \text{acc} = \{ T0 = [[T5 \rightarrow T4]] , T7 = [T5xT6 \Rightarrow T4] \}$

$T6 = T2$

$T7 = [T1 x T2 \rightarrow T3]$

—

$\{T6 = T2\} \circ \text{acc} \Rightarrow \{T6 = T2\}$

$\text{acc} \circ \{T6 = T2\} \Rightarrow \text{acc} = \{ T0 = [[T5 \rightarrow T4]] , T7 = [T5xT2 \Rightarrow T4], T6=T2 \}$

$T7 = [T1 x T2 \rightarrow T3]$

—

$\{T7 = [T1 x T2 \rightarrow T3]\} \circ \text{acc} \Rightarrow \{ [T1 x T2 \rightarrow T3] = [T5xT2 \rightarrow T4] \}$

$\text{acc} \circ \{ [T1 x T2 \rightarrow T3] = [T5xT6 \rightarrow T4] \} \Rightarrow$

$\text{acc} = \{ T0 = [[T5 \rightarrow T4]] , T6=T2, T7 = [T1 x T2 \rightarrow T3] \}$

$T1 = T5$

$T2 = T2$

$T3 = T4$

—

$\{T1 = T5\} \circ acc \Rightarrow \{T1 = T5\}$

$acc \circ \{T1 = T5\} \Rightarrow$

$acc = \{ T0 = [T5 \rightarrow T4], T6=T2, T7 = [T5 \times T2 \rightarrow T3], T1=T5 \}$

$T2 = T6$

$T3 = T4$

—

$\{T2 = T6\} \circ acc \Rightarrow \{T2 = T2\}$, not added.

$acc = \{ T0 = [T5 \rightarrow T4], T6=T2, T7 = [T5 \times T2 \rightarrow T3], T1=T5 \}$

$T3 = T4$

—

$\{T3 = T4\} \circ acc = \{ T3=T4\}$

$acc \circ \{T3 = T4\} \Rightarrow$

$acc = \{ T0 = [T5 \rightarrow T4], T6=T2, T7 = [T5 \times T2 \rightarrow T4], T1=T5, T3=T4 \}$

MGU: $acc = \{ T0 = [T5 \rightarrow T4], T6=T2, T7 = [T5 \times T2 \rightarrow T4], T1=T5, T3=T4 \}$

step 5: using the MGU as substitution

$T0 = (\lambda x) (f \times y)$

$T4 = (f \times y)$

$T5 = x$

$T6 = y$

$T7 = f$

$(\lambda (T5) T4)$

$(\lambda (T1) T3) =$

$[T1 \rightarrow T3]$

TRUE

d) running through inference, we will find $T2 = T1$ so f is $[T1 \Rightarrow T1]$ and $(f\ x)$ returns $T1$.

Question 2

2.1

- a) never
- b) string
- c) any
- d) number
- e) never
- f) boolean

2.2

```
(define (isBoolean : (any -> is? boolean))  
  (lambda ((x : any)) : is? boolean  
    (boolean? x)))
```

```
(define (good_in_L52 : ((union number boolean) -> boolean)  
  ( lambda ((z: ( union number boolean))) (if (isBoolean z)  
    z #f  
    ))
```

2.3)

```
(union string boolean)
```

The return type can be either a string if x is a number, or a boolean if x is a boolean. The function never reaches the “1” as x is either a number or a boolean, so a number is not returned.