

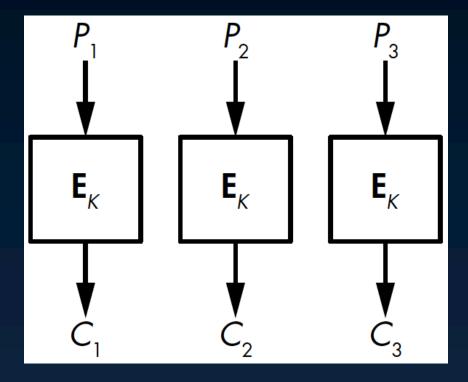


Applied Cryptography CPEG 472/672 Lecture 3B

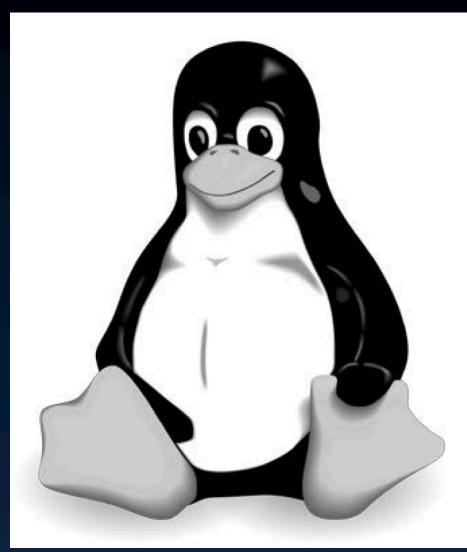
Instructor: Nektarios Tsoutsos

Modes of operation

- Electronic Codebook (ECB) Mode
 Never use not semantically secure
- Parallelizable



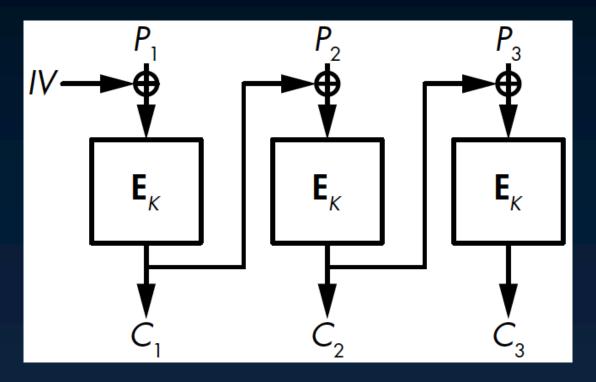
ECB insecurity example





Cipher Block Chaining (CBC) Mode

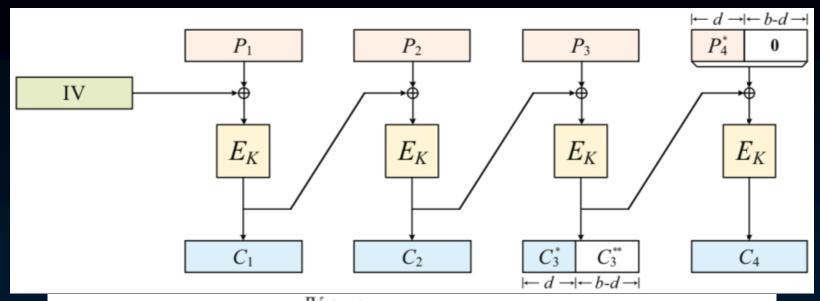
- Needs a random IV (not a secret)
 - Not semantic security otherwise
- Decryption can be parallelized



CBC for any message length

- What if the ptxt is not aligned with 16 byte length?
- Padding (append extra bytes)
 - ⊙ 01, 02 02, 03 03 03, ..., fifteen 0f
 - If ptxt aligns to blocks, pad with sixteen 10
 - Padding oracle attacks
 - ctxt size increases
- Ciphertext stealing
 - No increase in ctxt size

Ciphertext stealing



```
algorithm CBC-CS _K^{IV}(P)

n \leftarrow \lceil |P|/b \rceil

P_1 \cdots P_{n-1}P_n^* \leftarrow P where |P_1| = \cdots = |P_{n-1}| = b

P_n \leftarrow P_n^* \ 0^{b-d} where d \leftarrow |P_n^*|

C_0 \leftarrow IV;

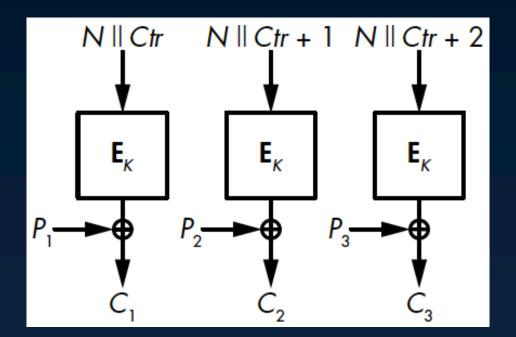
C_1 \cdots C_n \leftarrow \text{CBC }_K^{IV}(P_1 \cdots P_n) where |C_1| = \cdots = |C_n| = b

C_{n-1}^* \leftarrow \text{MSB}_d(C_{n-1})

return C_1 \cdots C_{n-2}C_{n-1}^*C_n
```

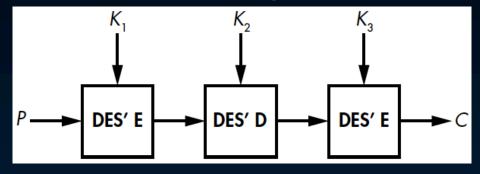
Counter (CTR) Mode

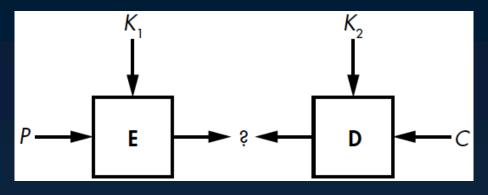
- Unique nonce per message
 - Nonce = number used once (non secret)
- Increase counter for each block
- Parallelizable



Attacks to Block Ciphers

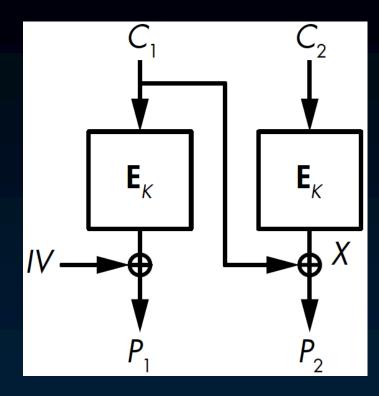
- Meet-in-the-Middle





Padding Oracle Attack

- Goal: Decrypt C2
- - Pick random C1
 - Brute force last byte until padding accepted
 - \odot C1[15] XOR X[15] = 01
 - Repeat for other bytes:
 - See how padding 02 02 can be accepted
 - ⊙ Find X[14] by setting C1[15] = X[15] XOR 02



Reading for next lecture

Aumasson: Chapter 3

Hands-on exercises

- Download tinyurl.com/my-blue-hen
- Encrypt it with AES in ECB mode
- Save the ciphertext as an image using PIL (Python) or Cimg (C++)
- Use AES-CBC to encrypt the same image
- Again, save the ctxt as an image
- Compare these images; what do you see?