

EL66 310

3/20/2018

Entropy

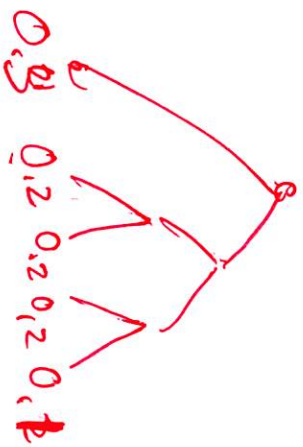
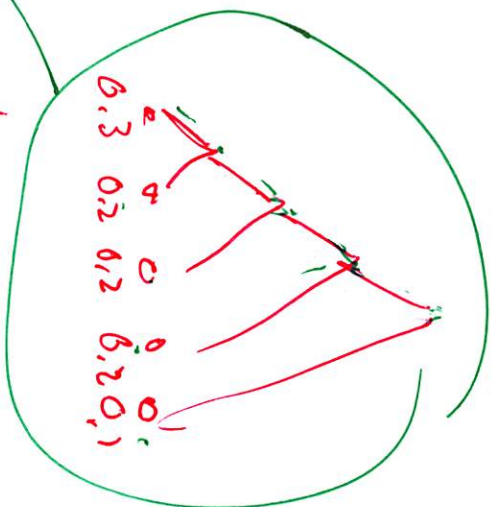
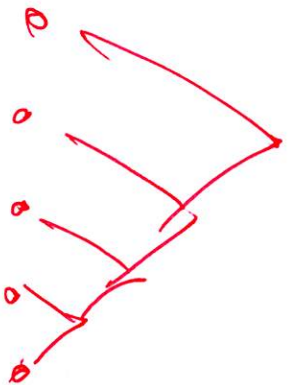
$$H(X) = - \sum_{k=1}^M P(k) \log_2 P(k) \geq 0$$

= measure of randomness (unpredictability)

Theorem:  $0 \leq H(X) \leq \log_2 M$

with  $H(X) = \log_2 M$  when  $P(k) = \frac{1}{M}$

↑  
Uniform distribution



Then  $H(X) + 1 > E[\text{Huff code}] \geq H(X)$

$$E(X) = 1 \times 0.3$$

$$+ 3 \times 0.7 = 2.1 \text{ bits}$$

$$EL = 4 \times 0.3 + 4 \times 0.2 + 3 \times 0.2 + 2 \times 0.2 + 1 \times 0.1$$

$$= 3.1 \text{ bits}$$

letter

Theorem: (Source Coding)

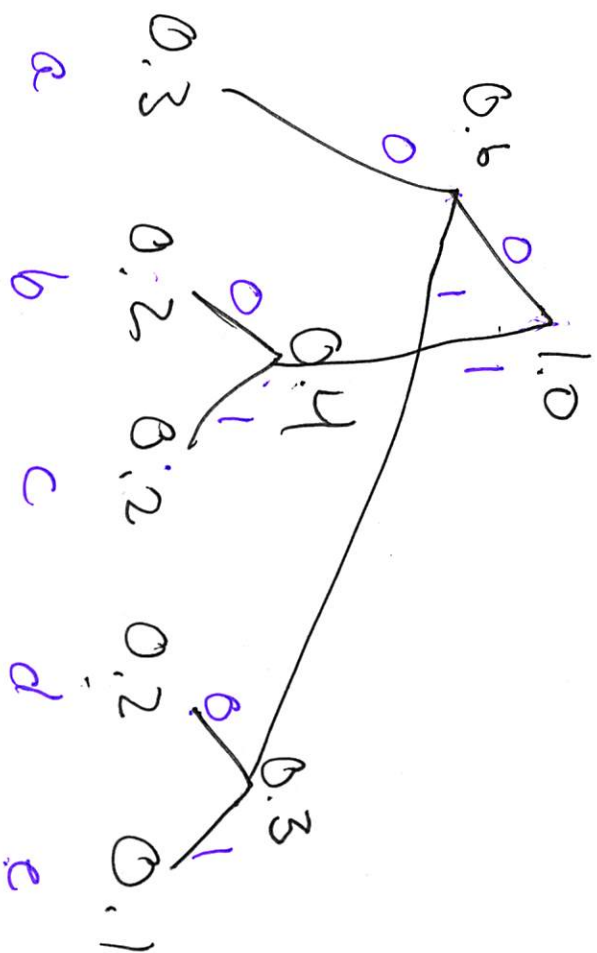
$$E L(X) \geq H(X)$$

$$2.3 \geq 2.24$$

where  $E L(X)$  = expected # of bits needed to encode  $X$

Huffman Code

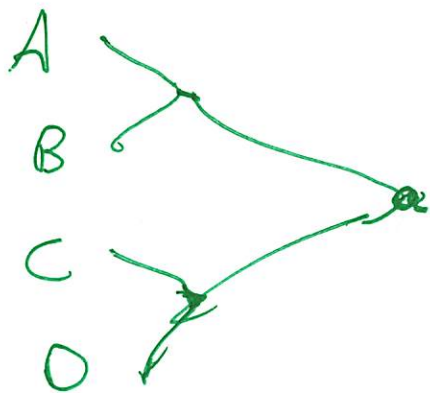
letter	p	Code	$L(X)$
a	0.3	00	2
b	0.2	10	2
c	0.2	11	2
d	0.2	010	3
e	0.1	011	3



$$E L(X) = 2 \times 0.3 + 2 \times 0.2 + 2 \times 0.2 + 3 \times 0.2 + 3 \times 0.1 = 2.3 \text{ bits}$$

$\approx 2.24$  bits/letter

$$H(X) = -0.3 \log 0.3 - 0.2 \log 0.2 - 0.2 \log 0.2 - 0.2 \log 0.2 - 0.1 \log 0.1$$



$P(A \text{ wins tournament})$

$= P(A \text{ wins 1st game})$

$\cap A \text{ wins 2nd game})$

$= P(A \text{ wins 1st game}) P(A \text{ wins 2nd game})$

$= 0.7 \times 0.7 = 0.49$

$P(A \text{ beats } X) = 0.7$

$P(B \text{ wins tournament}) = P(B \text{ wins 1st game})$

$\times P(B \text{ wins 2nd game})$

$= 0.3 \times 0.5 = 0.15$

$P(C \text{ wins tour}) = 0.5 \times \underline{\hspace{2cm}}$

$P(C \text{ wins Tour}) = P(C \text{ wins Tour} | A \text{ wins 1st game}) P(A \text{ wins 1st game})$   
 $+ P(C \text{ wins Tour} | B \text{ wins 1st game}) P(B \text{ wins 1st game})$

$= 0.5 \times 0.3 \times 0.7 + 0.5 \times 0.5 \times 0.3 = 0.18$

$X_1, X_2 \sim \text{Geometric}(p)$  ind

$S = X_1 + X_2$  pmf of  $S$ ?

$$P(X=k) = (1-p)^{k-1} p \quad k=1, 2, \dots, \infty$$

	$p$	$(1-p)p$	$(1-p)^2 p$	$(1-p)^3 p$	...
$p$	$p^2$	$(1-p)p^2$	$p^2 p^2$	$p^3 p^2$	...
$(1-p)p$	$q \cdot p^2$	$q^2 p^2$	$q^3 p^2$		
$q^2 p = (1-p)^2 p$		$q^2 p^2$	$q^3 p^2$		
$(1-p)^3 p$			$q^3 p^2$		
	$p^3$	$2qp^2$	$3q^2 p^2$	$4q^3 p^2$	