

SPRING 2023

INFORMATION TECHNOLOGY RESEARCH

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Lecture slides are based on the supplemental materials of the textbook: <https://algs4.cs.princeton.edu>



<http://algs4.cs.princeton.edu>

2.1 ELEMENTARY SORTS

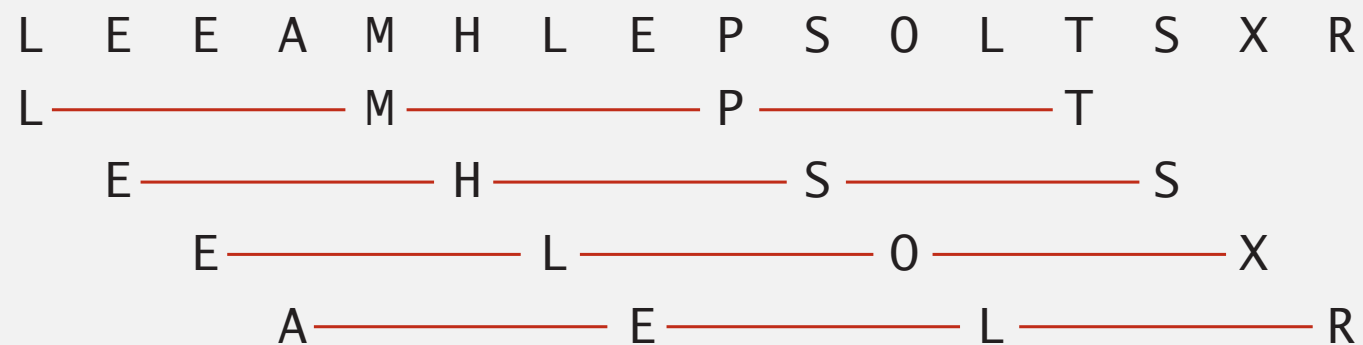
- *rules of the game*
- *selection sort*
- *insertion sort*
- *shellsort*
- *shuffling*

Shellsort overview

Idea. Move entries more than one position at a time by *h-sorting* the array.

an *h*-sorted array is *h* interleaved sorted subsequences

h = 4

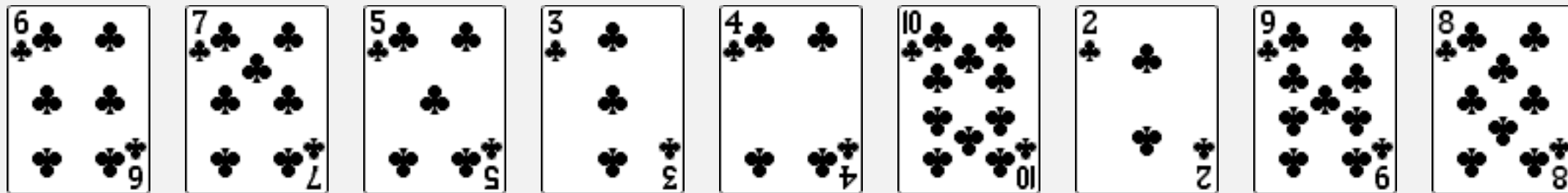


Shellsort. [Shell 1959] *h-sort* array for decreasing sequence of values of *h*.

input	S	H	E	L	L	S	O	R	T	E	X	A	M	P	L	E
13-sort	P	H	E	L	L	S	O	R	T	E	X	A	M	S	L	E
4-sort	L	E	E	A	M	H	L	E	P	S	O	L	T	S	X	R
1-sort	A	E	E	E	H	L	L	L	M	O	P	R	S	S	T	X

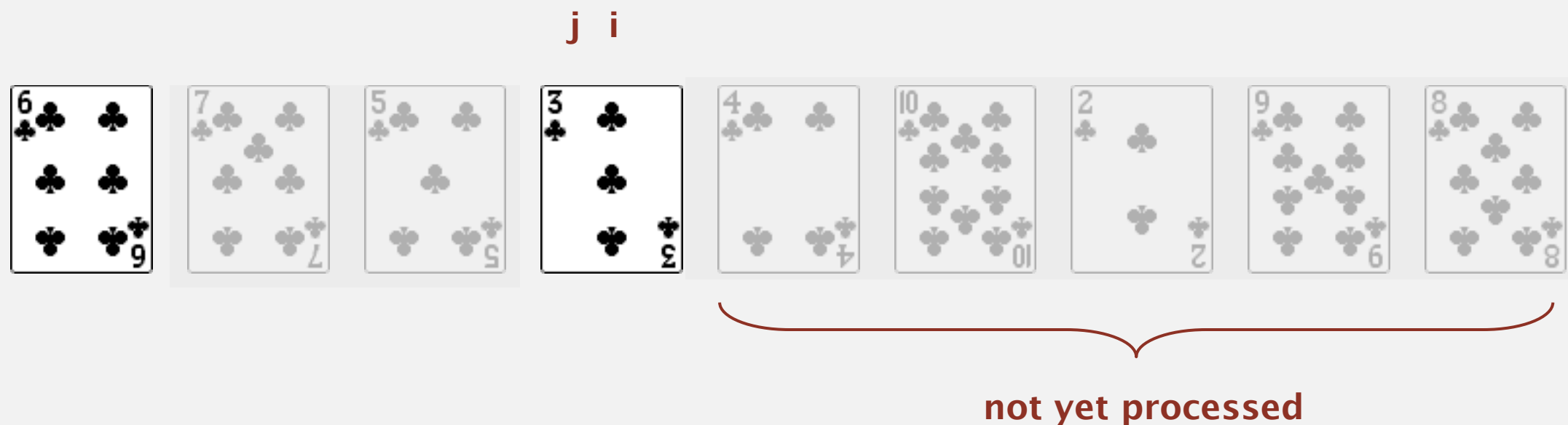
h-sorting demo

In iteration i , swap $a[i]$ with each larger entry h positions to its left.



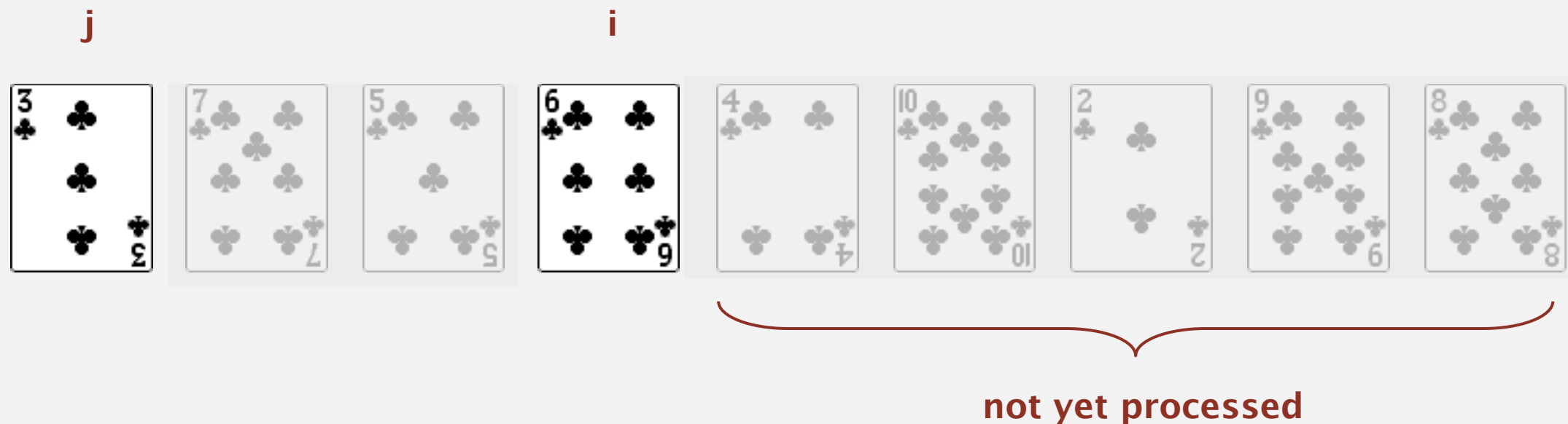
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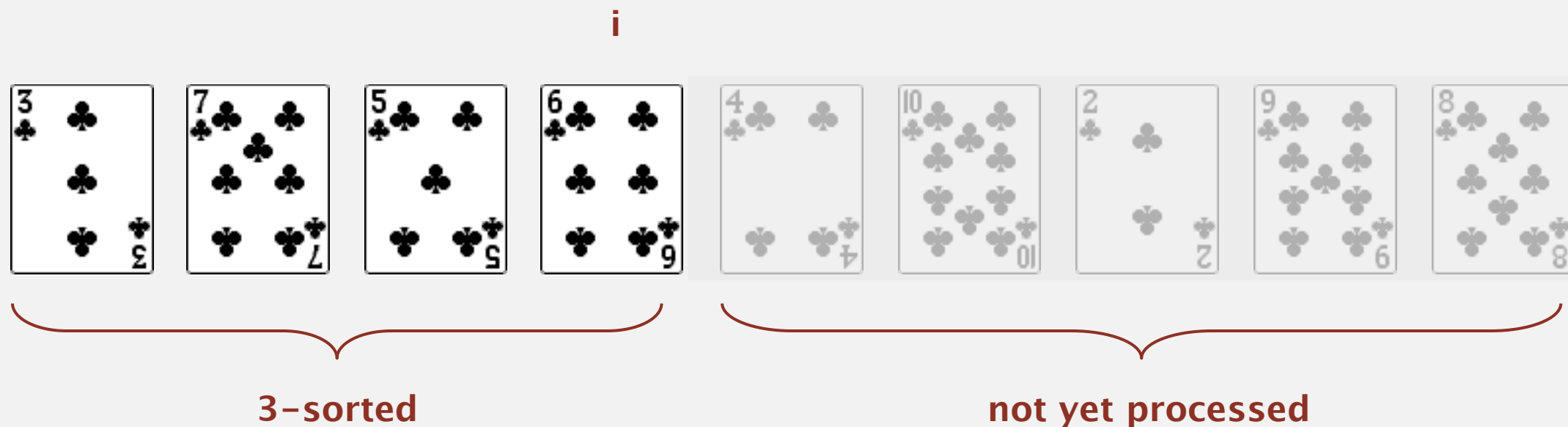
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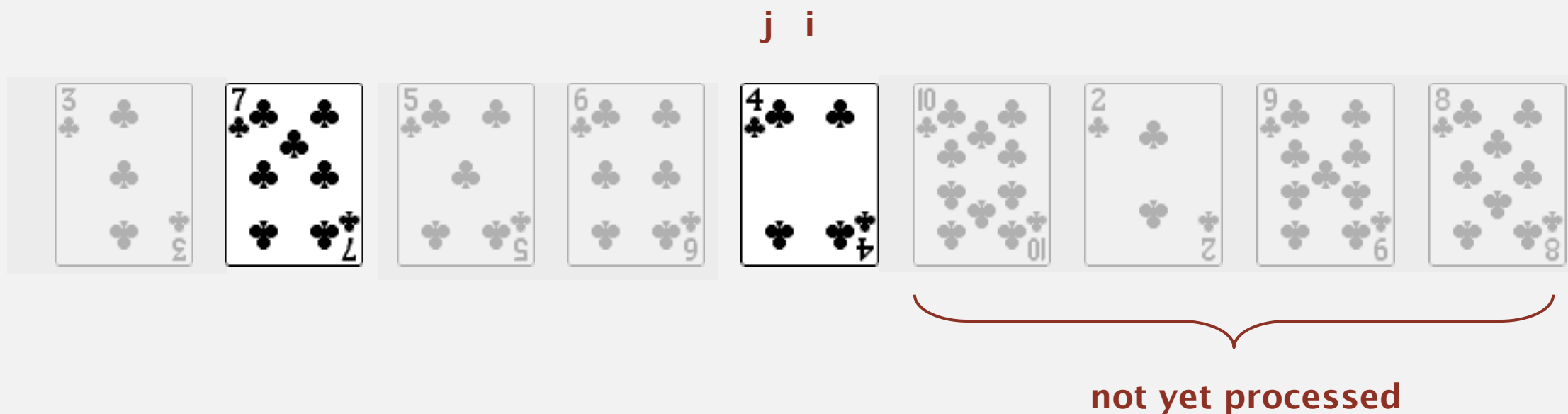
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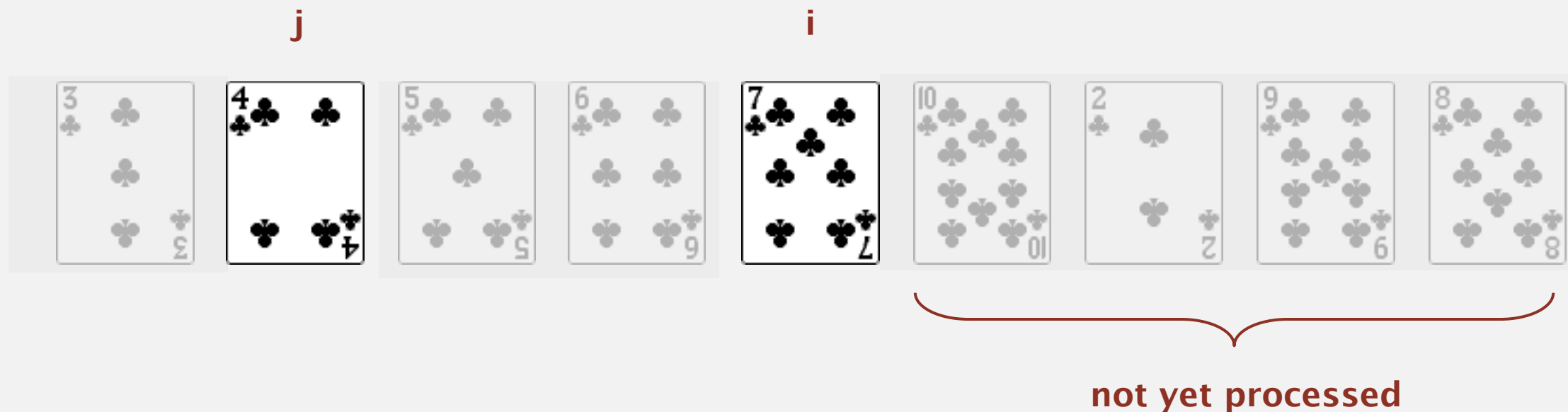
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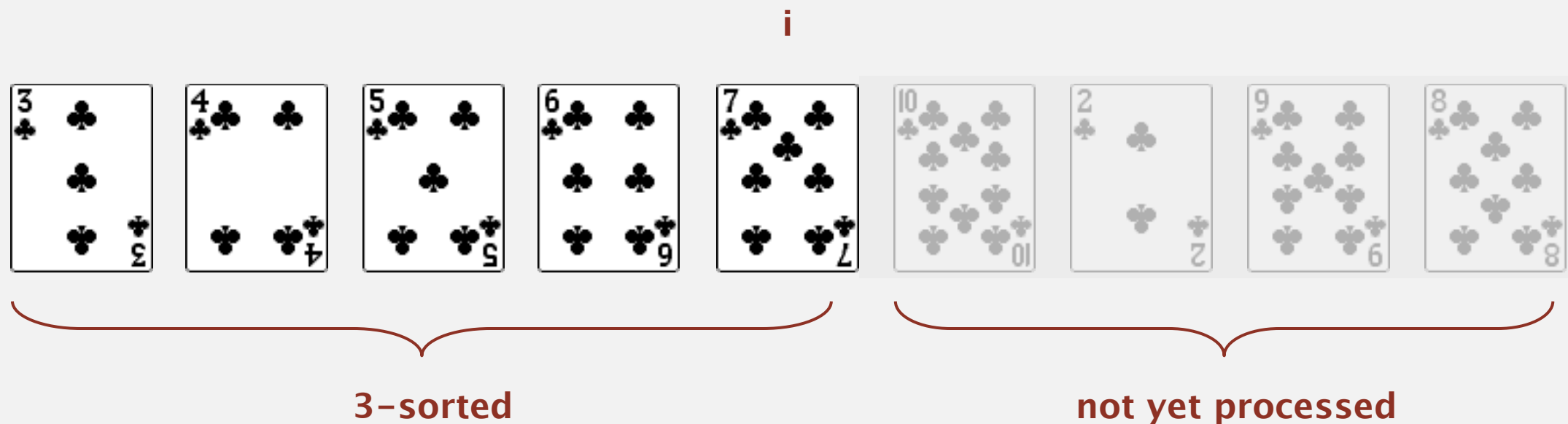
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h-sorting demo

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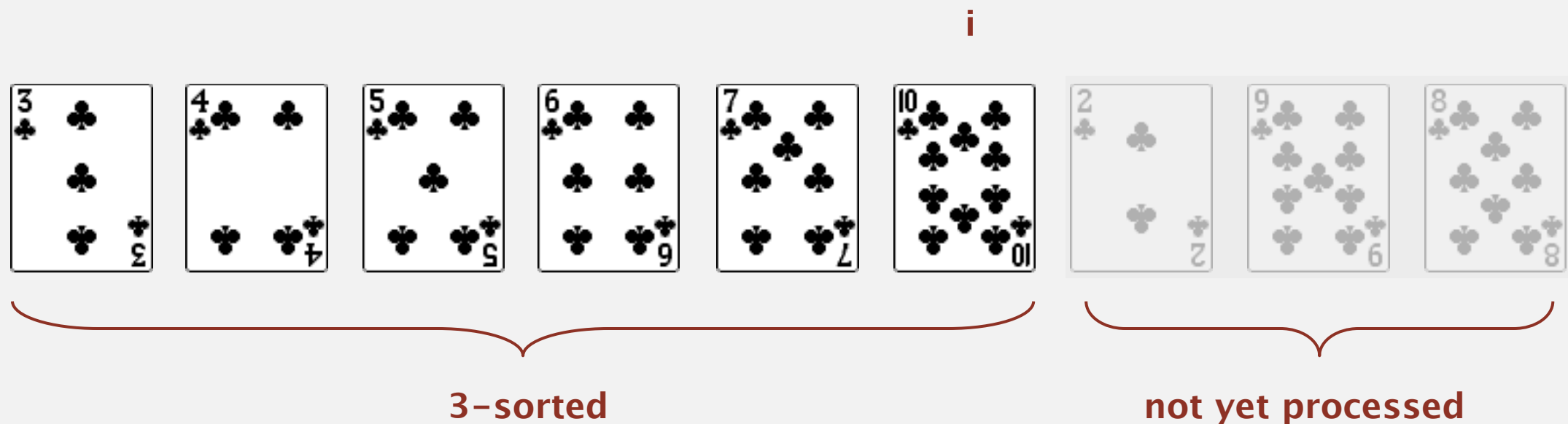
j i



not yet processed

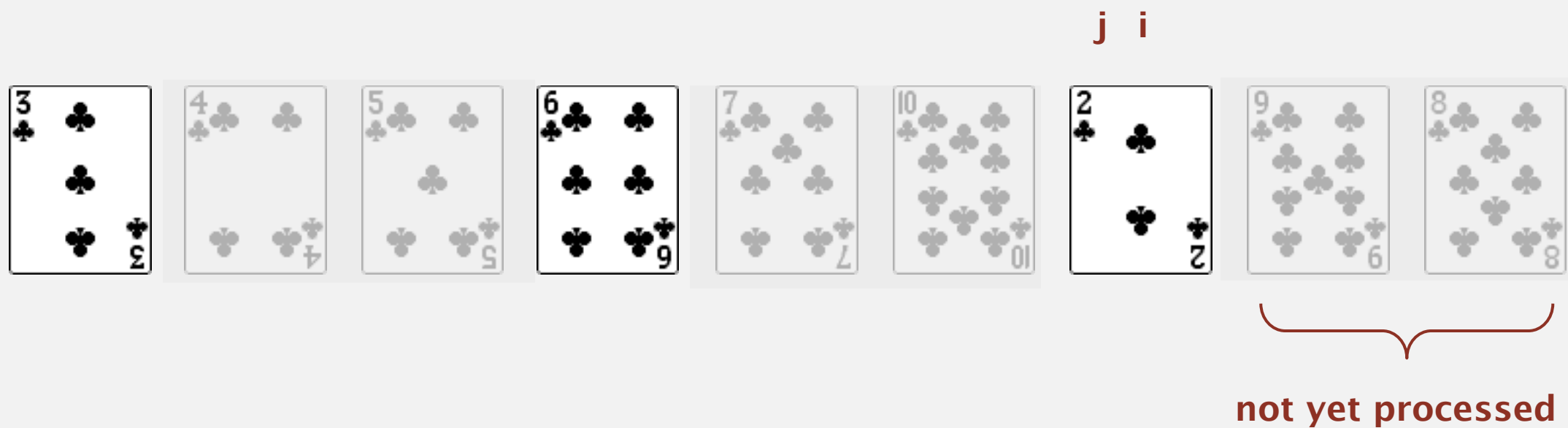
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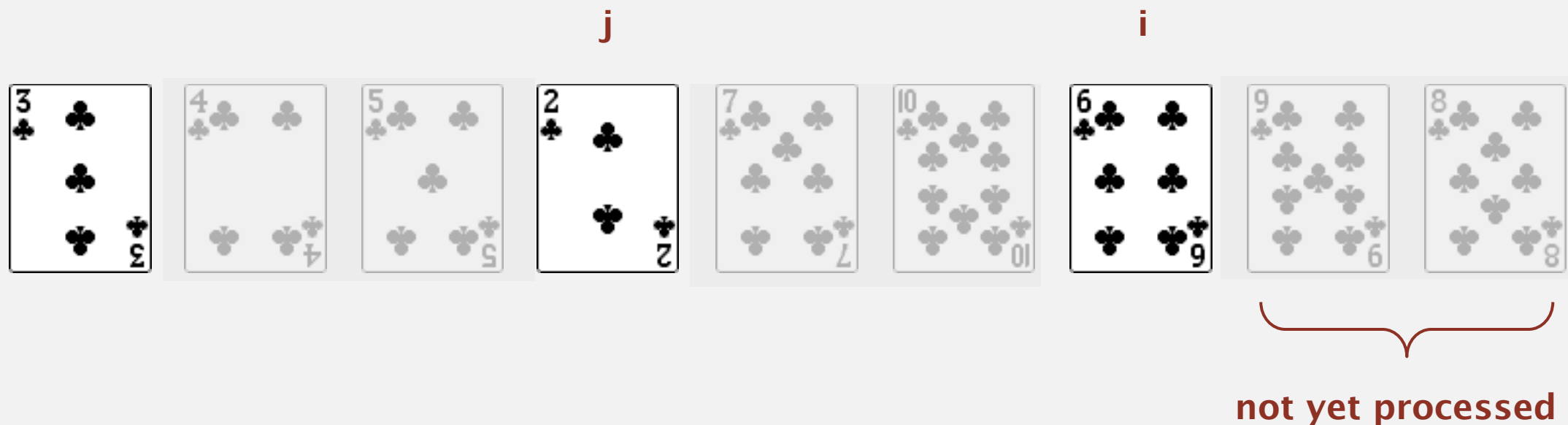
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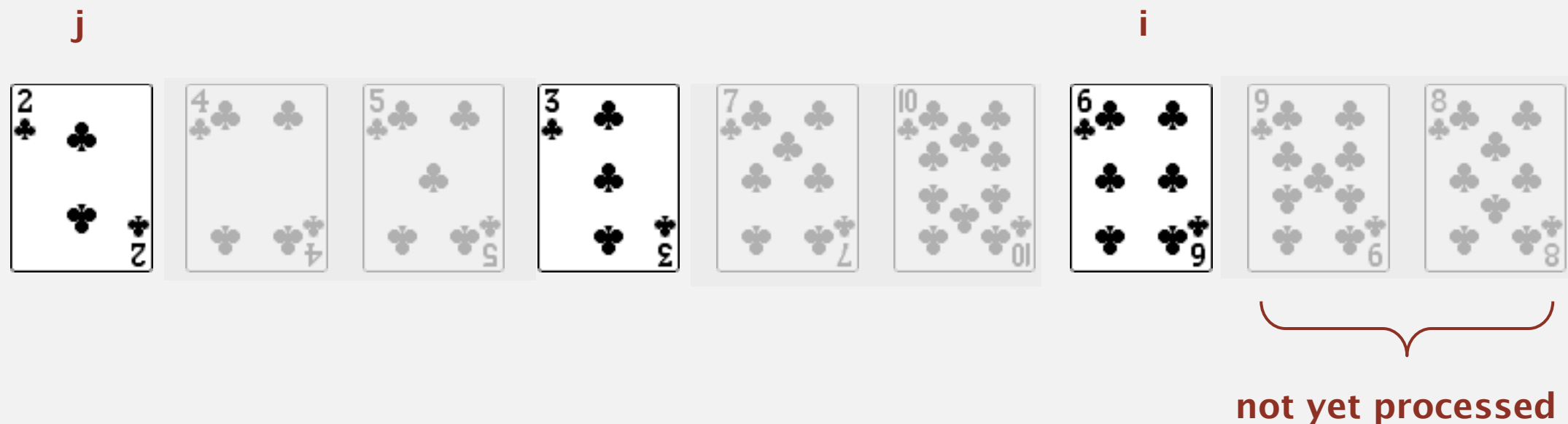
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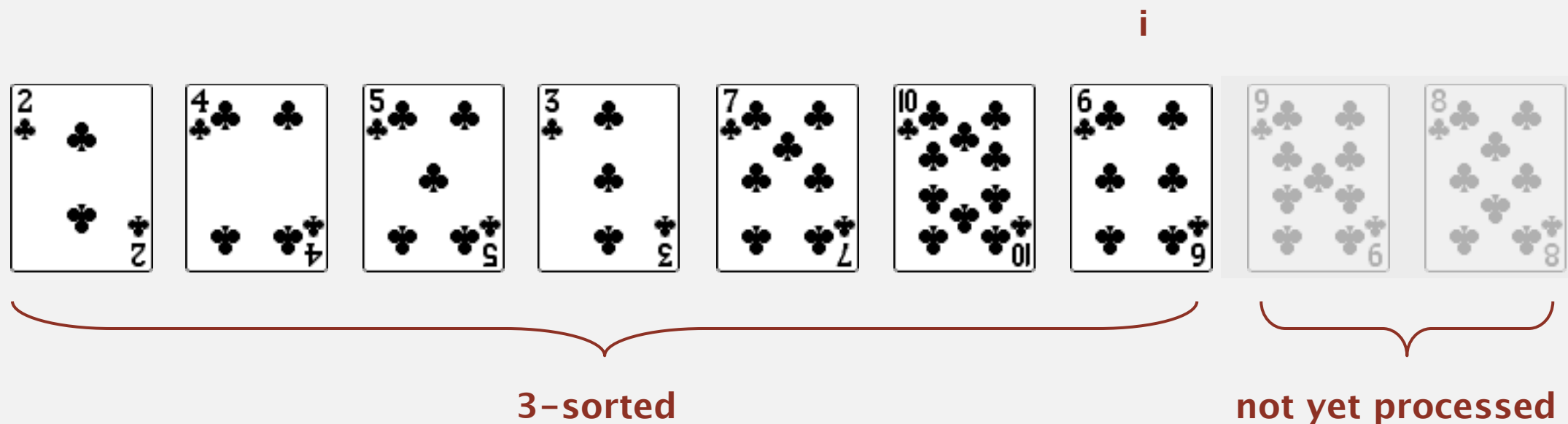
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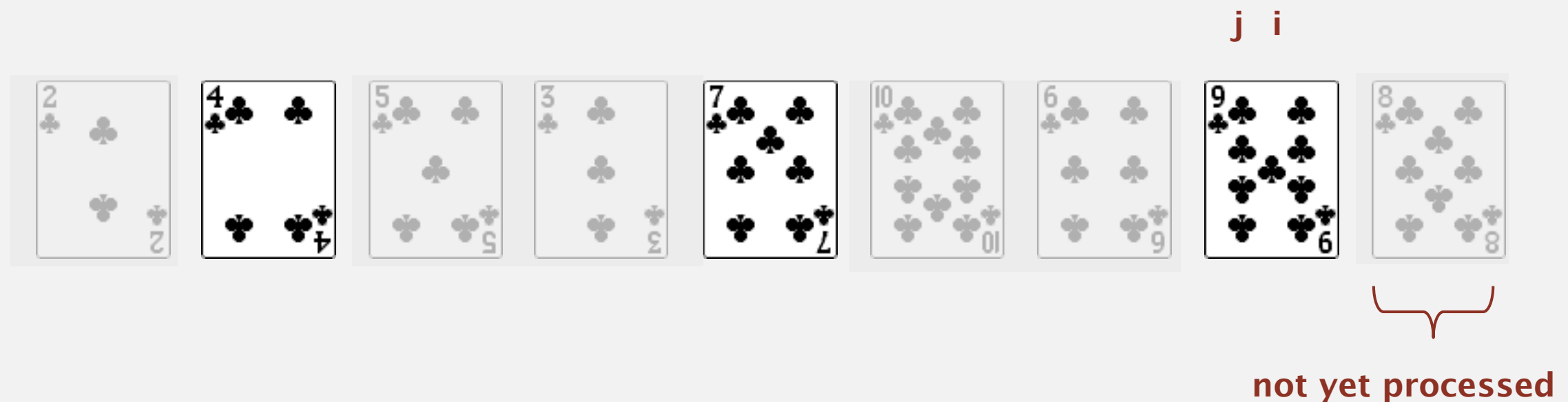
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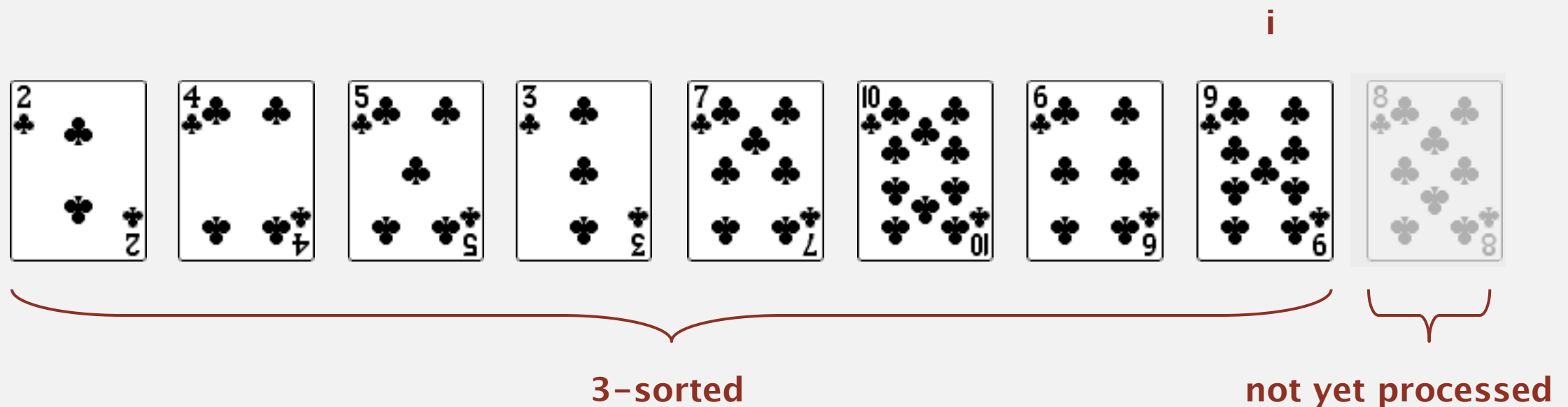
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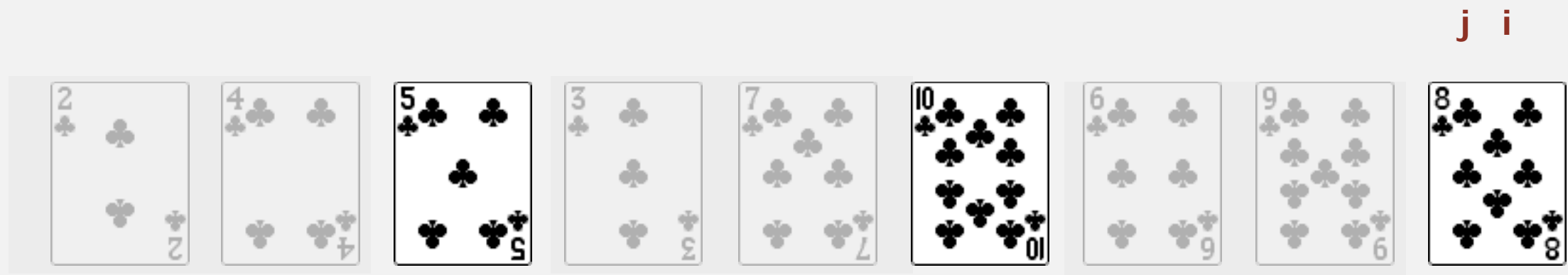
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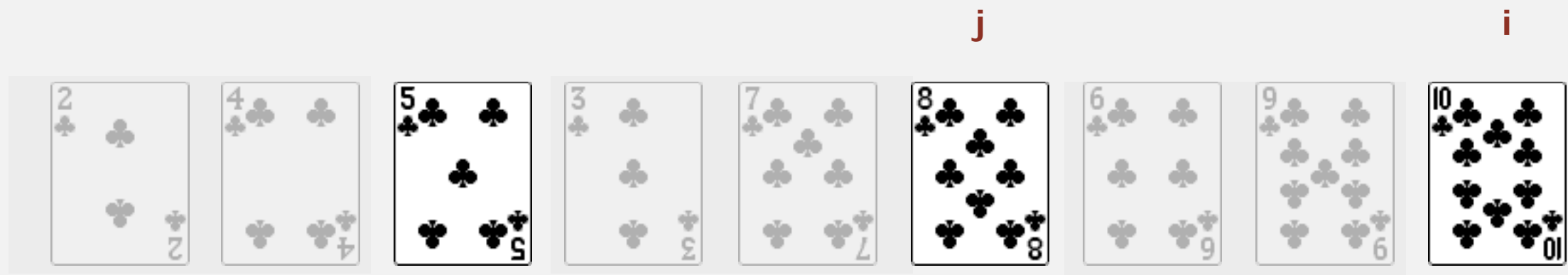
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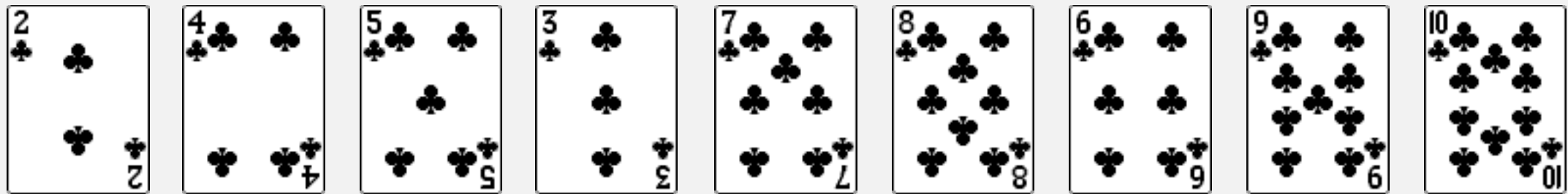
h-sorting demo

In iteration i , swap $a[i]$ with each larger entry h positions to its left.



h-sorting demo

An array is h-sorted if $a[i-h] \leq a[i]$ for each i .



3-sorted

h-sorting demo

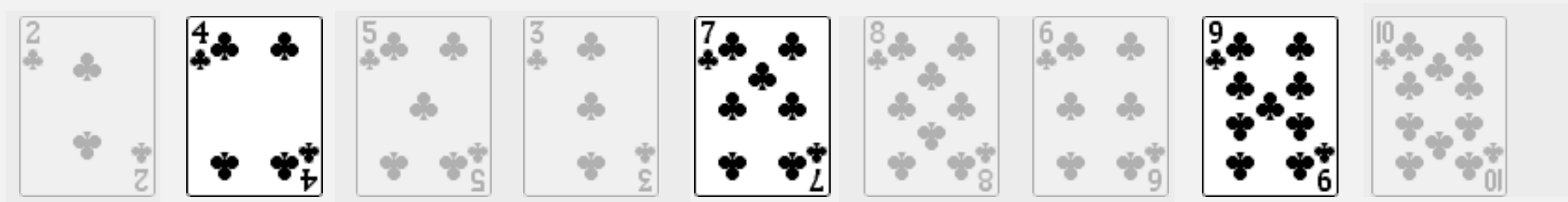
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3-sorted

h-sorting demo

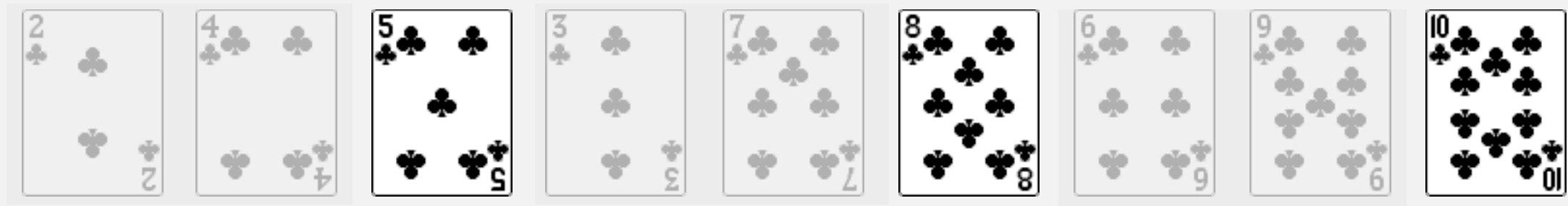
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3-sorted

h-sorting demo

An array is h-sorted if $a[i-h] \leq a[i]$ for each i .



3-sorted

h-sorting

How to h -sort an array? Insertion sort, with stride length h .

3-sorting an array

M	O	L	E	E	X	A	S	P	R	T
E	O	L	M	E	X	A	S	P	R	T
E	E	L	M	O	X	A	S	P	R	T
E	E	L	M	O	X	A	S	P	R	T
A	E	L	E	O	X	M	S	P	R	T
A	E	L	E	O	X	M	S	P	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T

Why insertion sort?

- Big increments \Rightarrow small subarray.
- Small increments \Rightarrow nearly in order. [stay tuned]

h-sort: how would you implement it from insertion sort?

```
public class Insertion
{
    public static void sort(Comparable[] a)
    {
        int N = a.length;
        for (int i = 0; i < N; i++)
            for (int j = i; j > 0; j--)
                if (less(a[j], a[j-1]))
                    exch(a, j, j-1);
                else break;
    }

    private static boolean less(Comparable v, Comparable w)
    { /* as before */ }

    private static void exch(Comparable[] a, int i, int j)
    { /* as before */ }
}
```

What to change here?
For h-sorted.
4 mins.

h-sort: how would you implement it from insertion sort?

```
public class Insertion
{
    public static void sort(Comparable[] a)
    {
        int N = a.length;
        for (int i = h; i < N; i++)
            for (int j = i; j > 0; j-=h)
                if (less(a[j], a[j-h]))
                    exch(a, j, j-h);
                else break;
    }

    private static boolean less(Comparable v, Comparable w)
    { /* as before */ }

    private static void exch(Comparable[] a, int i, int j)
    { /* as before */ }
}
```

Shellsort example: increments 7, 3, 1

input

S O R T E X A M P L E

7-sort

S	O	R	T	E	X	A	M	P	L	E
M	O	R	T	E	X	A	S	P	L	E
M	O	R	T	E	X	A	S	P	L	E
M	O	L	T	E	X	A	S	P	R	E
M	O	L	E	E	X	A	S	P	R	T

3-sort

M	O	L	E	E	X	A	S	P	R	T
E	O	L	M	E	X	A	S	P	R	T
E	E	L	M	O	X	A	S	P	R	T
E	E	L	M	O	X	A	S	P	R	T
A	E	L	E	O	X	M	S	P	R	T
A	E	L	E	O	X	M	S	P	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T

1-sort

A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	L	E	O	P	M	S	X	R	T
A	E	E	L	O	P	M	S	X	R	T
A	E	E	L	O	P	M	S	X	R	T
A	E	E	L	M	O	P	S	X	R	T
A	E	E	L	M	O	P	S	X	R	T
A	E	E	L	M	O	P	S	X	R	T
A	E	E	L	M	O	P	R	S	X	T
A	E	E	L	M	O	P	R	S	T	X

result

A E E L M O P R S T X

Shellsort: Java implementation

```
public class Shell
{
    public static void sort(Comparable[] a)
    {
        int N = a.length;
```

```
        int h = 1;
        while (h < N/3) h = 3*h + 1; // 1, 4, 13, 40, 121, 364, ...
```

← 3x+1 increment
sequence

What to write here? 5 mins.
Can use `exch` and `less` methods.

```
    }

    private static boolean less(Comparable v, Comparable w)
    { /* as before */ }
    private static void exch(Comparable[] a, int i, int j)
    { /* as before */ }
}
```

Shellsort: Java implementation

```
public class Shell
{
    public static void sort(Comparable[] a)
    {
        int N = a.length;

        int h = 1;
        while (h < N/3) h = 3*h + 1; // 1, 4, 13, 40, 121, 364, ...

        while (h >= 1)
        { // h-sort the array.
            for (int i = h; i < N; i++)
            {
                for (int j = i; j >= h && less(a[j], a[j-h]); j -= h)
                    exch(a, j, j-h);
            }

            h = h/3;
        }

        private static boolean less(Comparable v, Comparable w)
        { /* as before */ }
        private static void exch(Comparable[] a, int i, int j)
        { /* as before */ }
    }
}
```

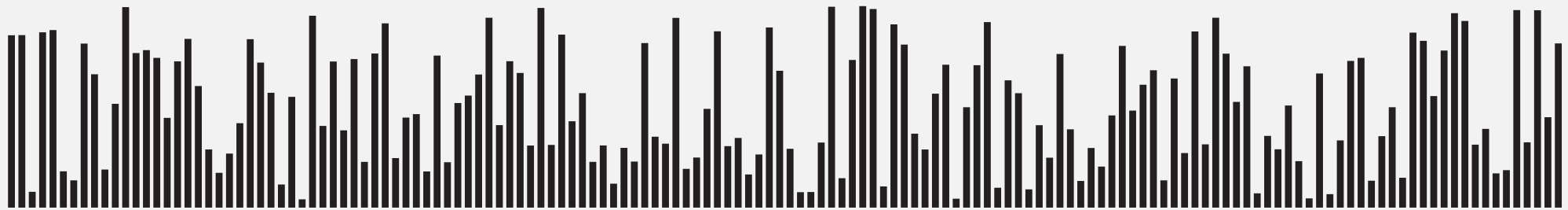
3x+1 increment sequence

insertion sort

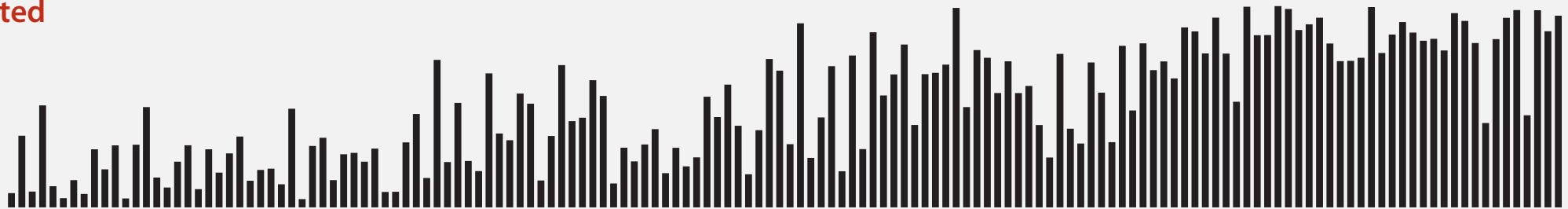
move to next increment

Shellsort: visual trace

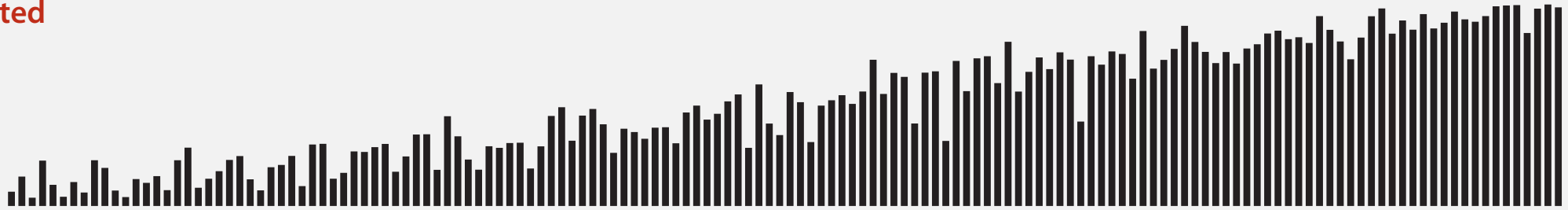
input



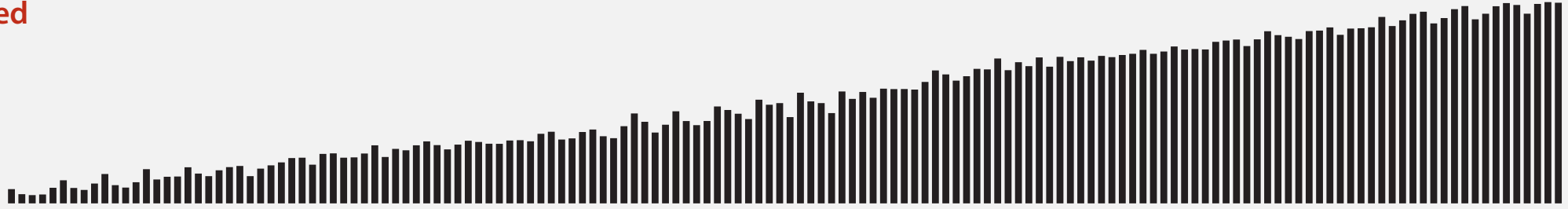
40-sorted



13-sorted



4-sorted

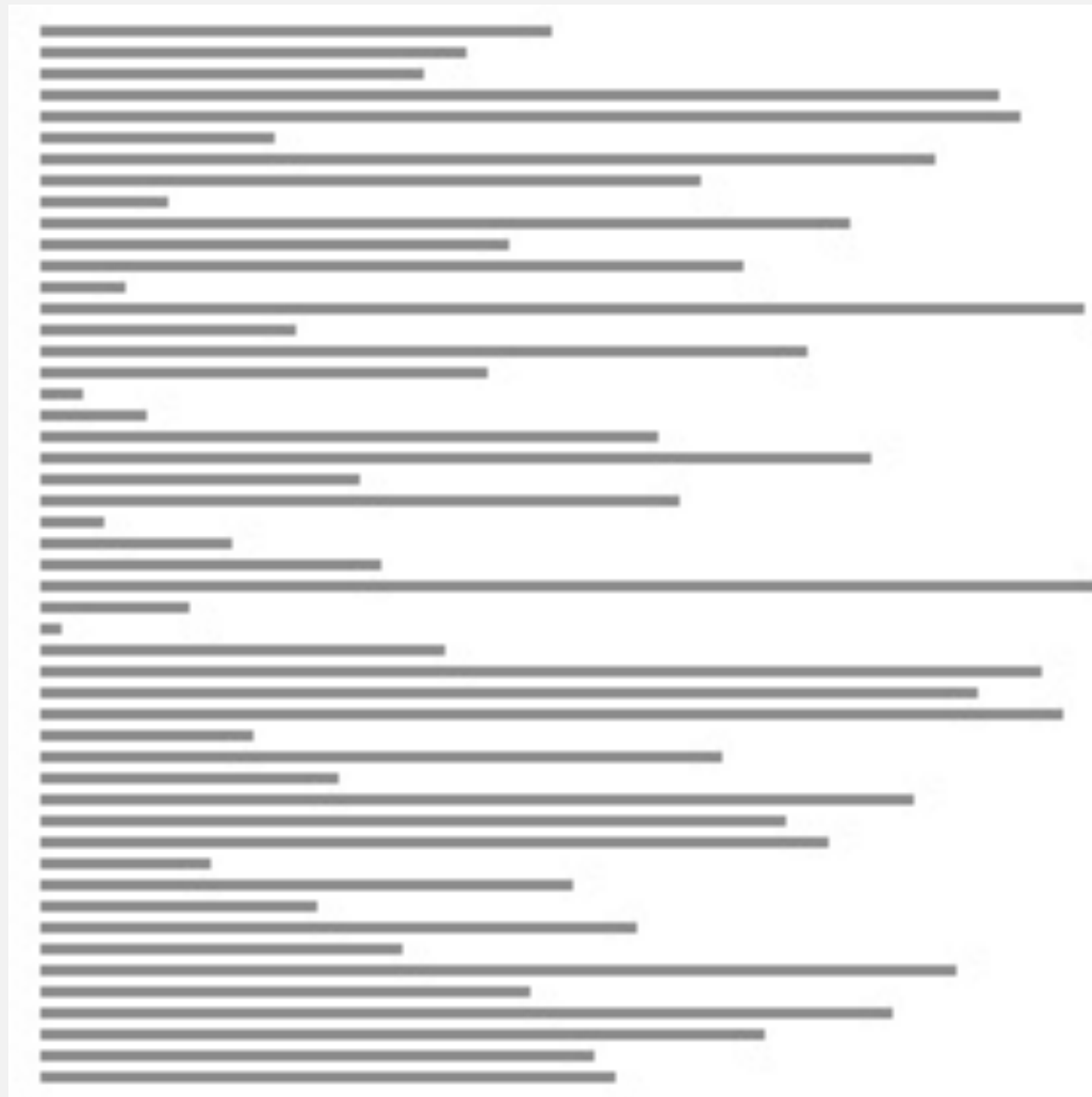


result



Shellsort: animation

50 random items

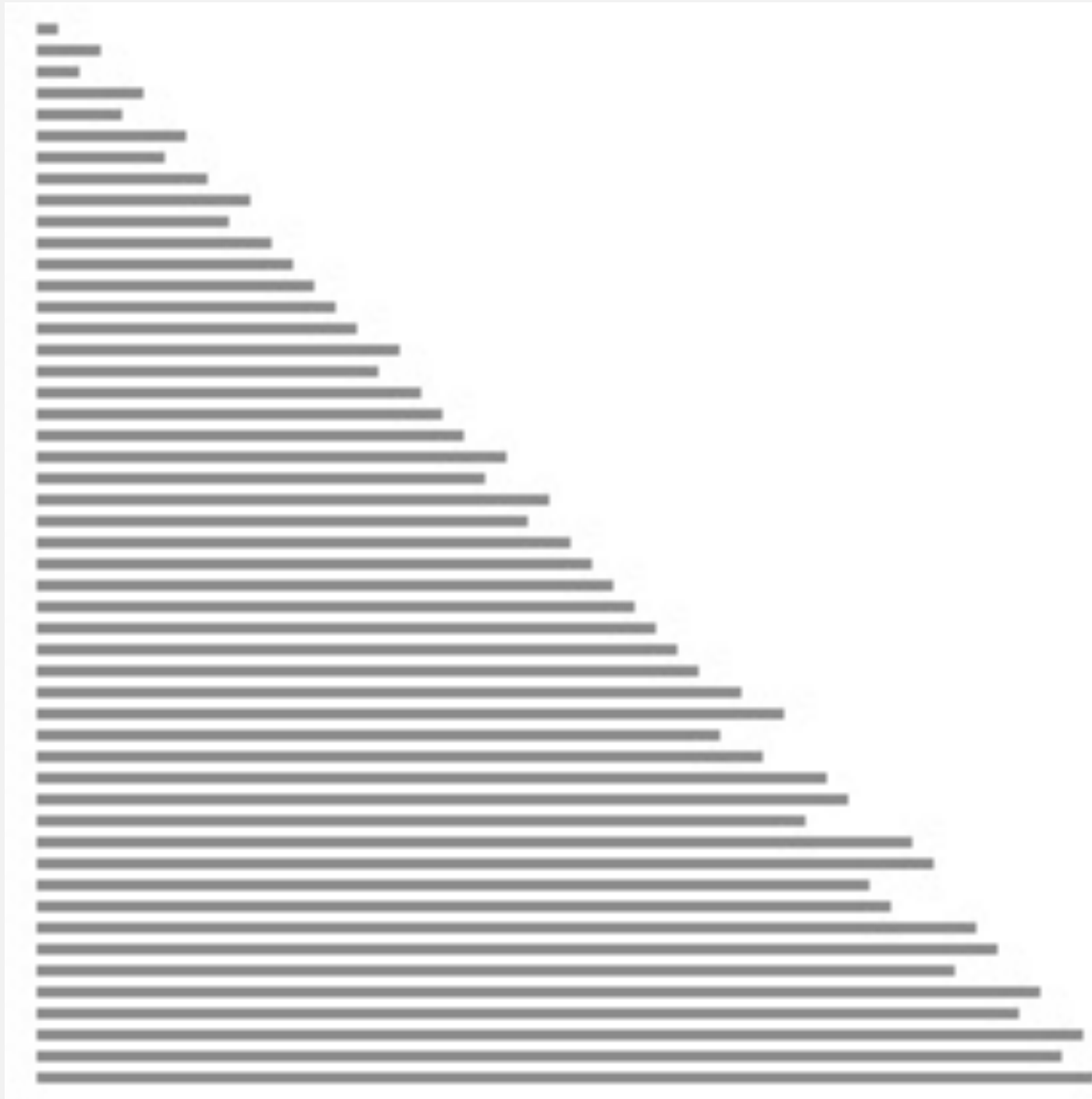


<http://www.sorting-algorithms.com/shell-sort>

- ▲ algorithm position
- h-sorted
- current subsequence
- other elements

Shellsort: animation

50 partially-sorted items



<http://www.sorting-algorithms.com/shell-sort>

- ▲ algorithm position
- h-sorted
- current subsequence
- other elements

Shellsort: analysis

Proposition. The order of growth of the worst-case number of compares used by shellsort with the $3x+1$ increments is $N^{3/2}$.

Property. The expected number of compares to shellsort a randomly-ordered array using $3x+1$ increments is....

N	compares	$2.5 N \ln N$	$0.25 N \ln^2 N$	$N^{1.3}$
5,000	93K	106K	91K	64K
10,000	209K	230K	213K	158K
20,000	467K	495K	490K	390K
40,000	1022K	1059K	1122K	960K
80,000	2266K	2258K	2549K	2366K

Remark. Accurate model has not yet been discovered (!)

Why are we interested in shellsort?

Example of simple idea leading to substantial performance gains.

Useful in practice.

- Fast unless array size is huge (used for small subarrays).
- Tiny, fixed footprint for code (used in some embedded systems).
- Hardware sort prototype.

Simple algorithm, nontrivial performance, interesting questions.

- Asymptotic growth rate?
- Best sequence of increments?  open problem: find a better increment sequence
- Average-case performance?

Lesson. Some good algorithms are still waiting discovery.



<http://algs4.cs.princeton.edu>

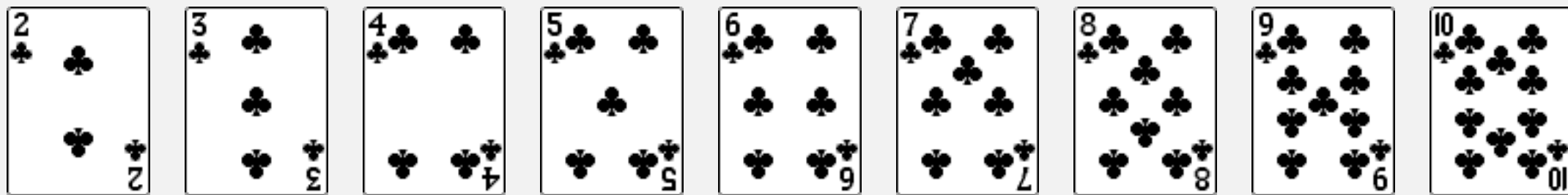
2.1 ELEMENTARY SORTS

- *rules of the game*
- *selection sort*
- *insertion sort*
- *shellsort*
- *shuffling*

How to shuffle an array

Goal. Rearrange array so that result is a uniformly random permutation.

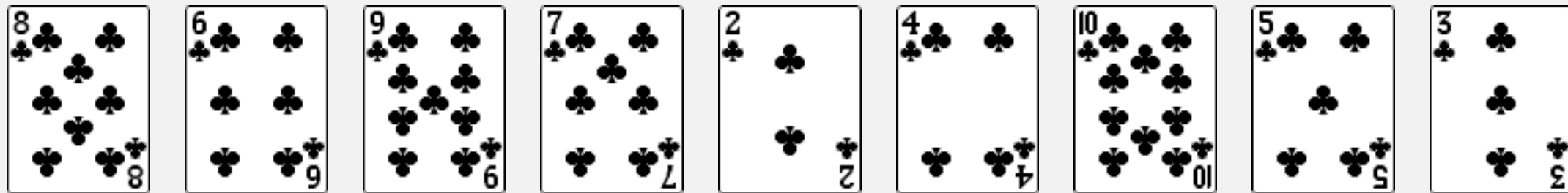
all permutations
equally likely



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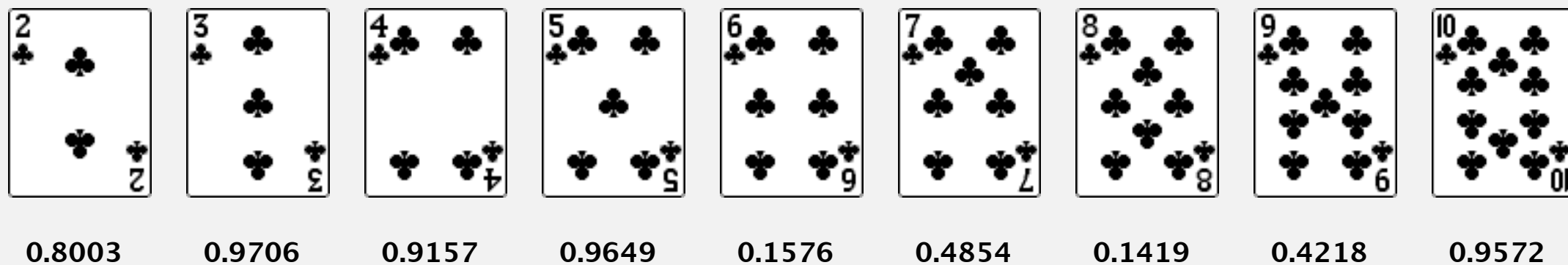
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Shuffle sort

- Generate a random real number for each array entry.
- Sort the array.

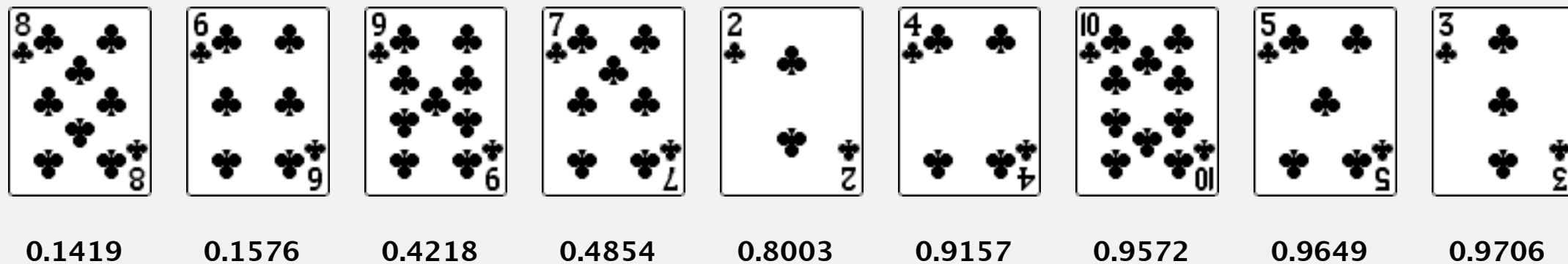
↑ useful for shuffling
columns in a spreadsheet



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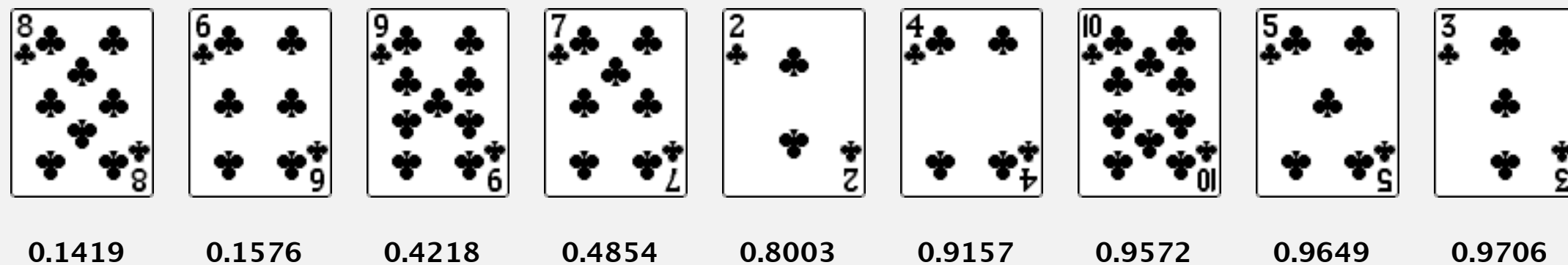
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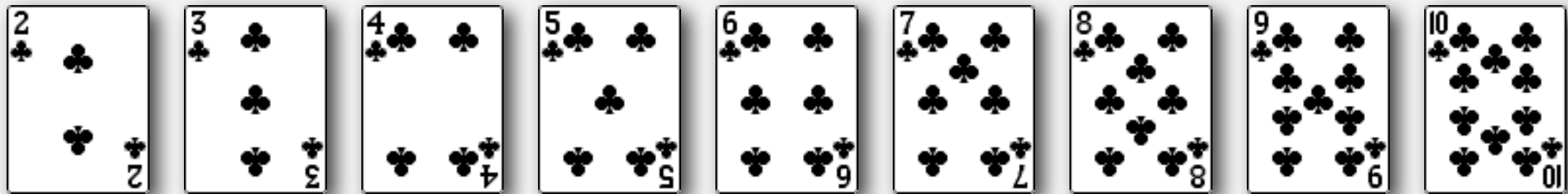


Proposition. Shuffle sort produces a uniformly random permutation.

↑ assuming real numbers
uniformly at random (and no ties)

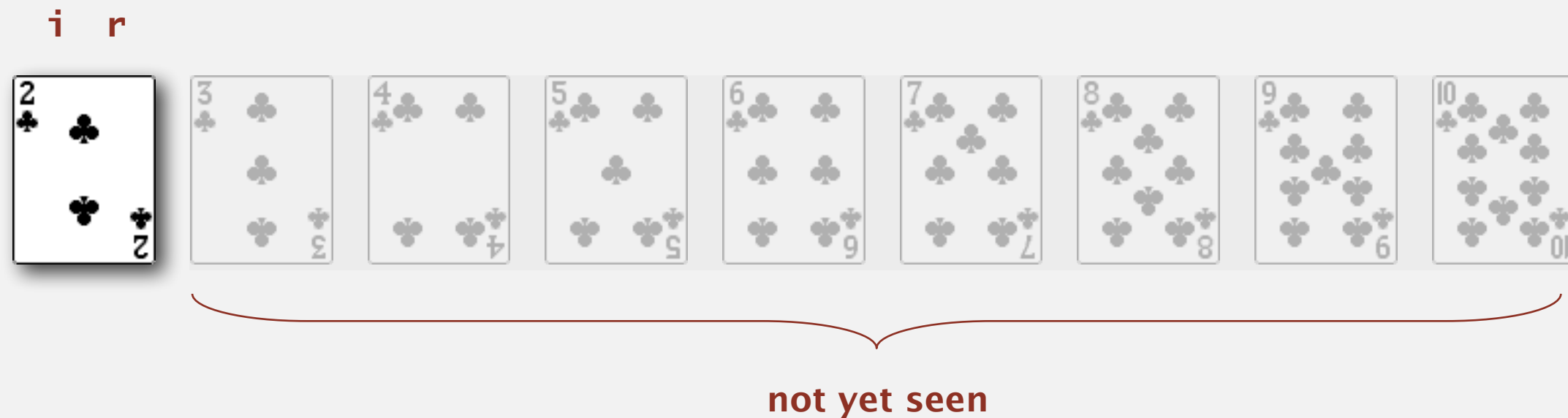
Knuth shuffle

- In iteration i , pick integer r between 0 and i uniformly at random.
- Swap $a[i]$ and $a[r]$.



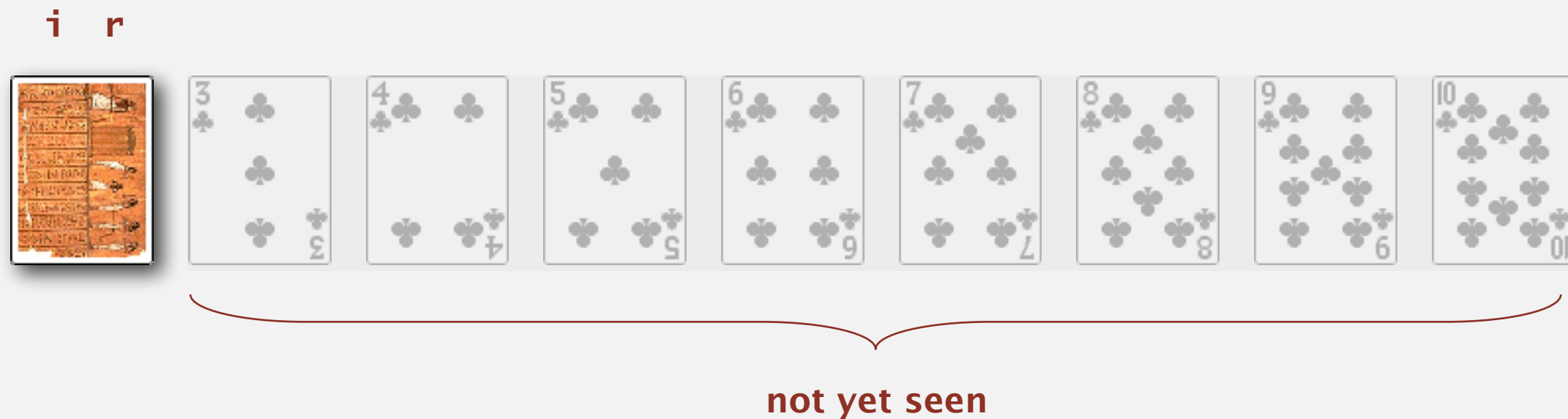
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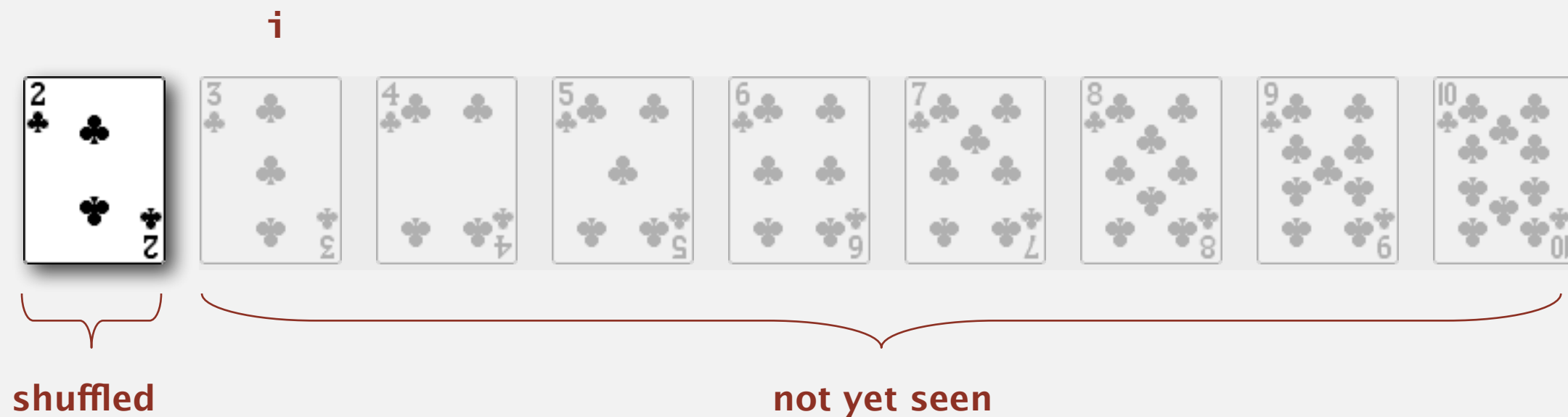
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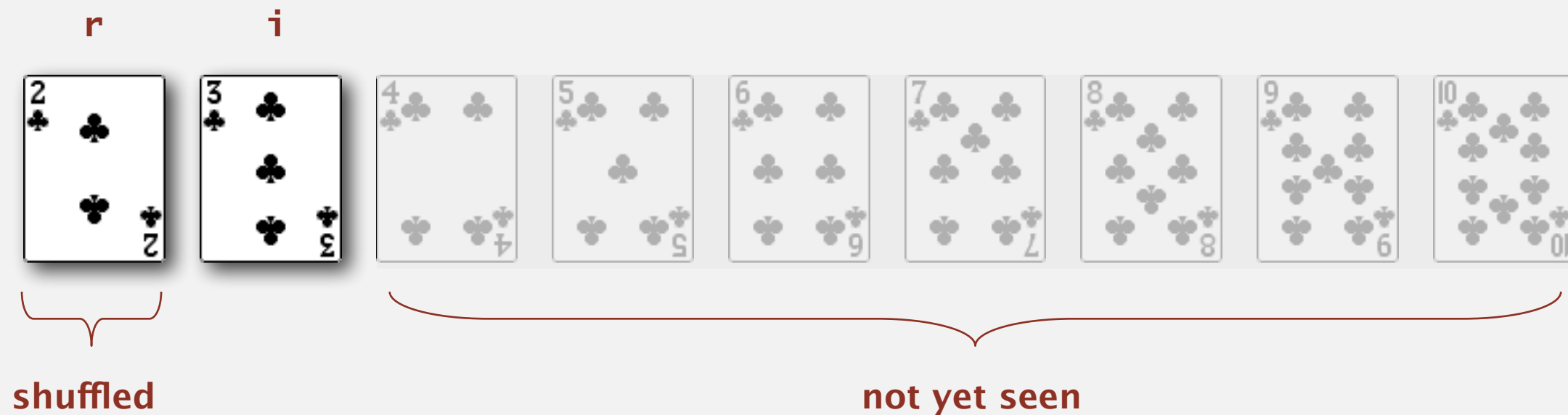
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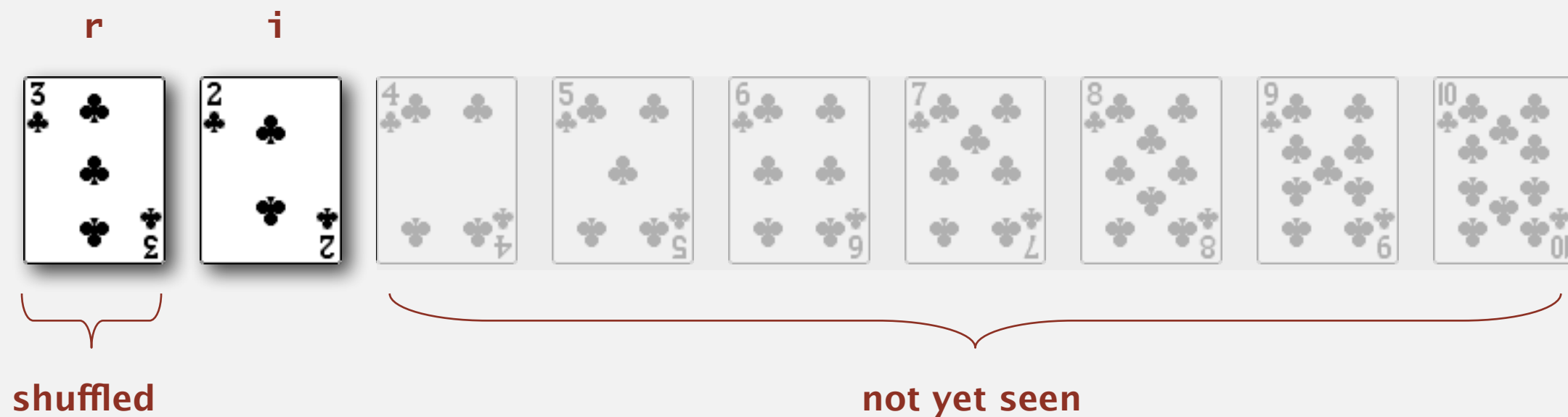
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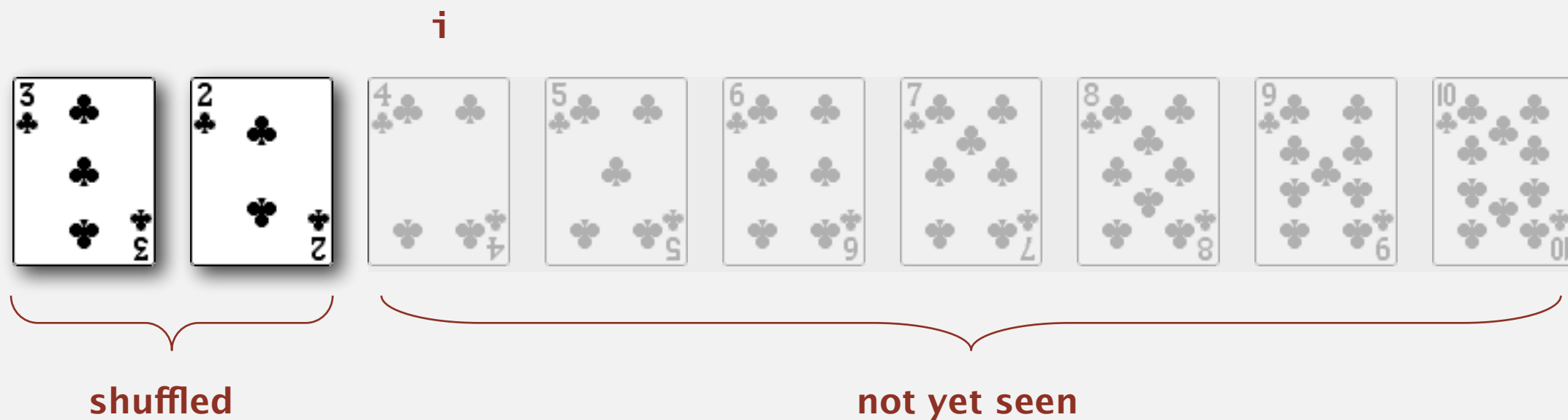
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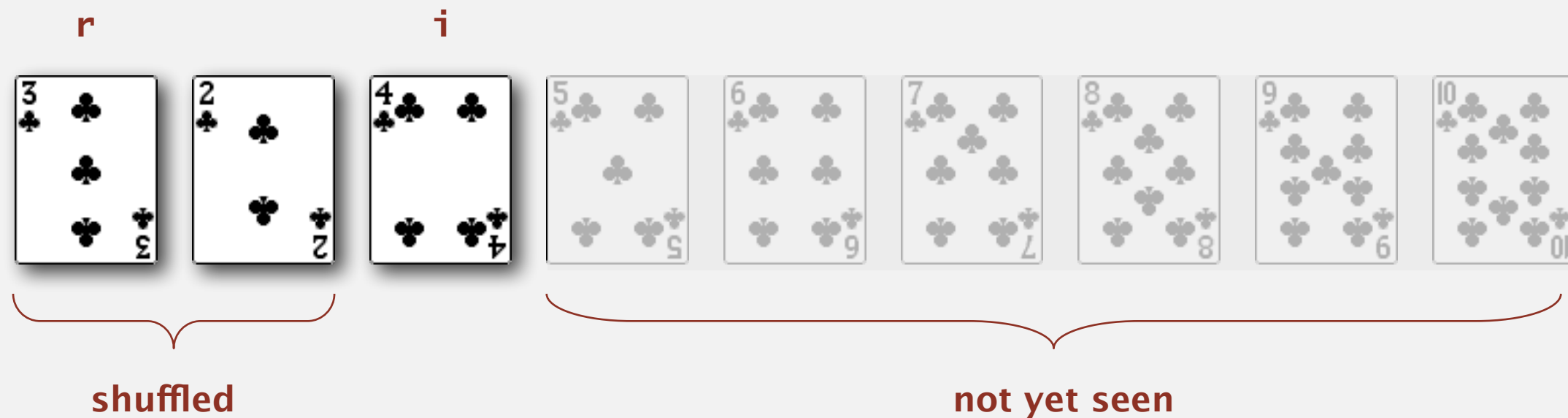
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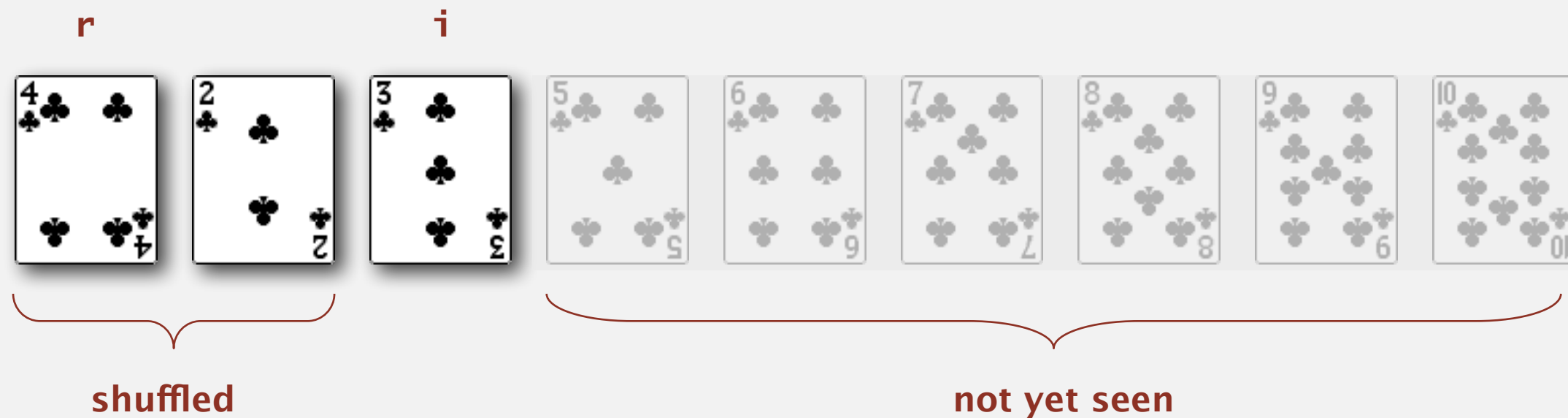
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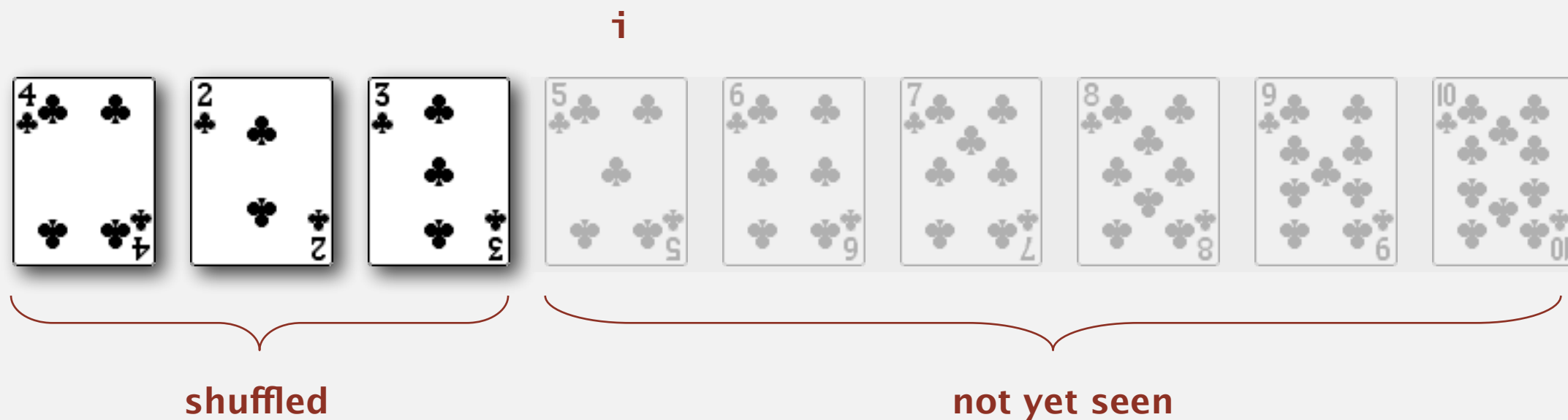
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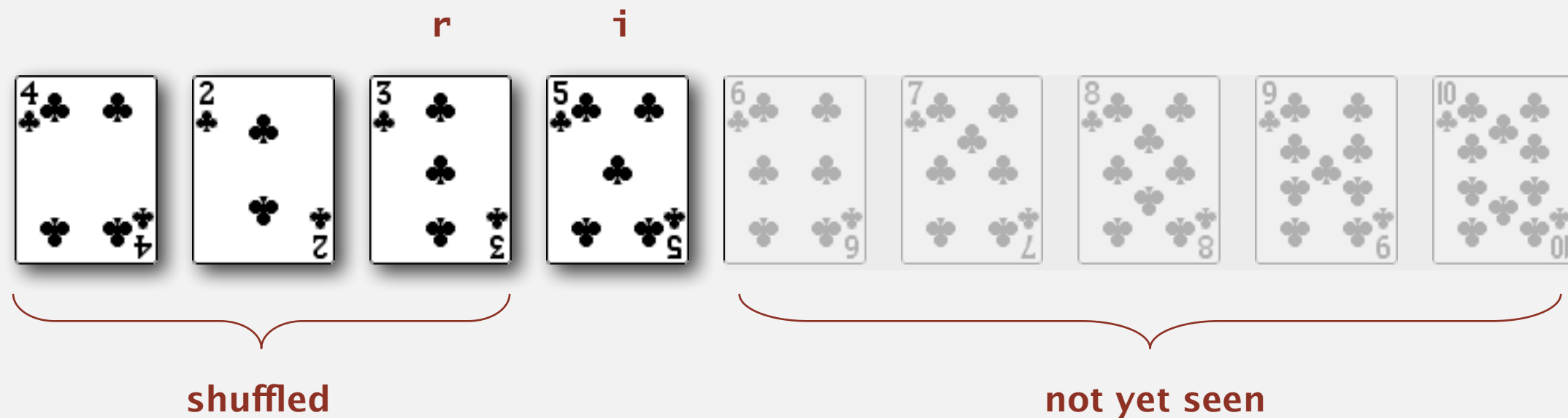
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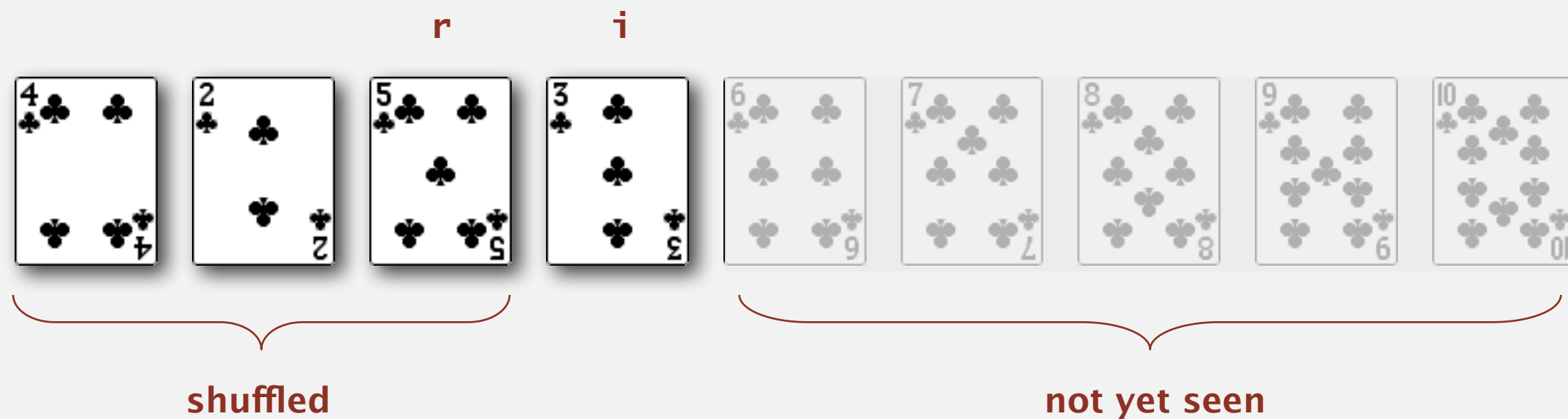
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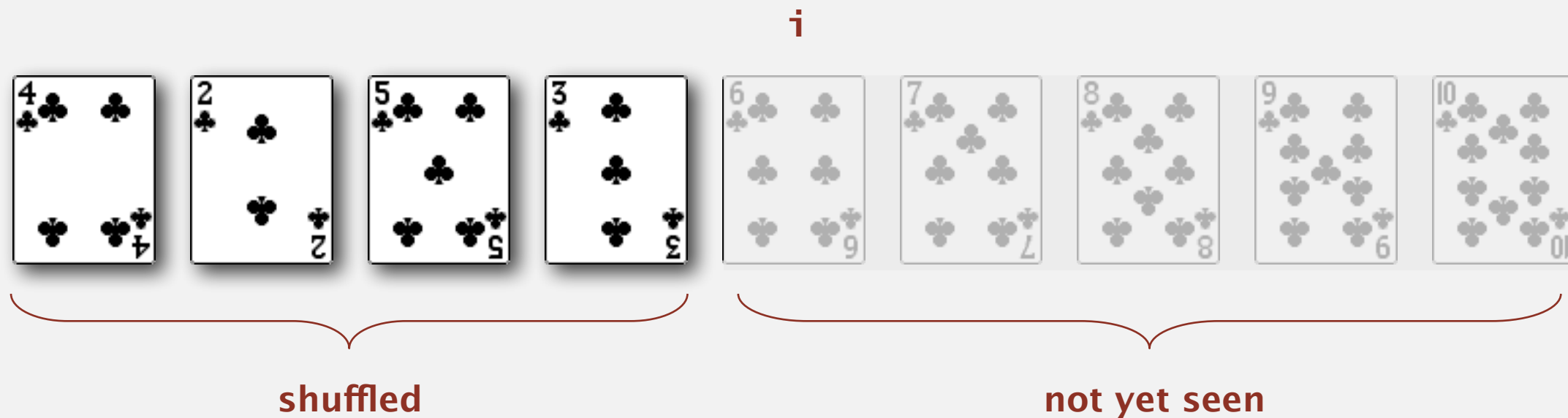
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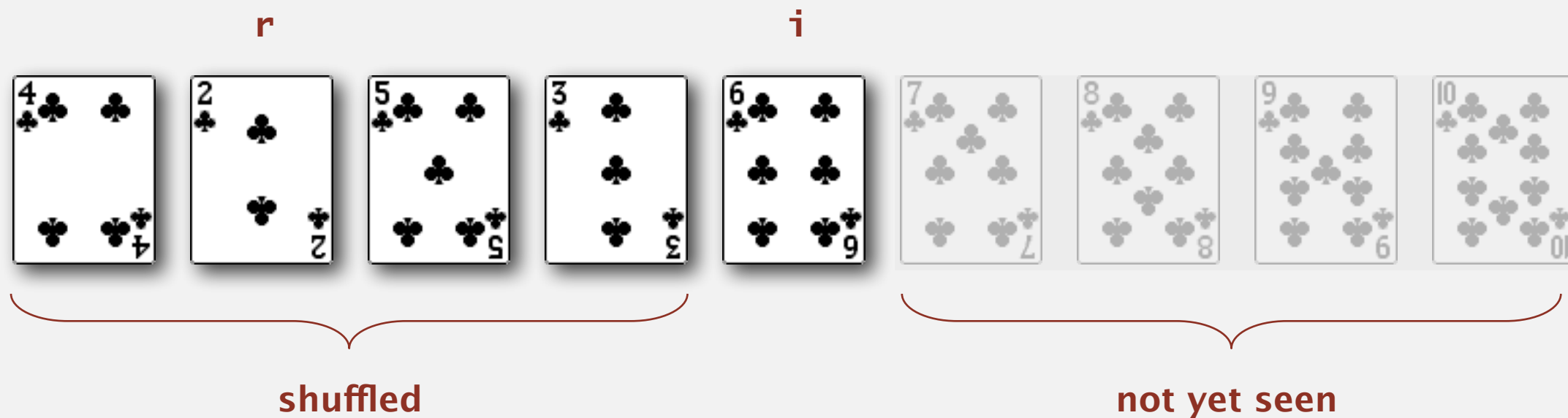
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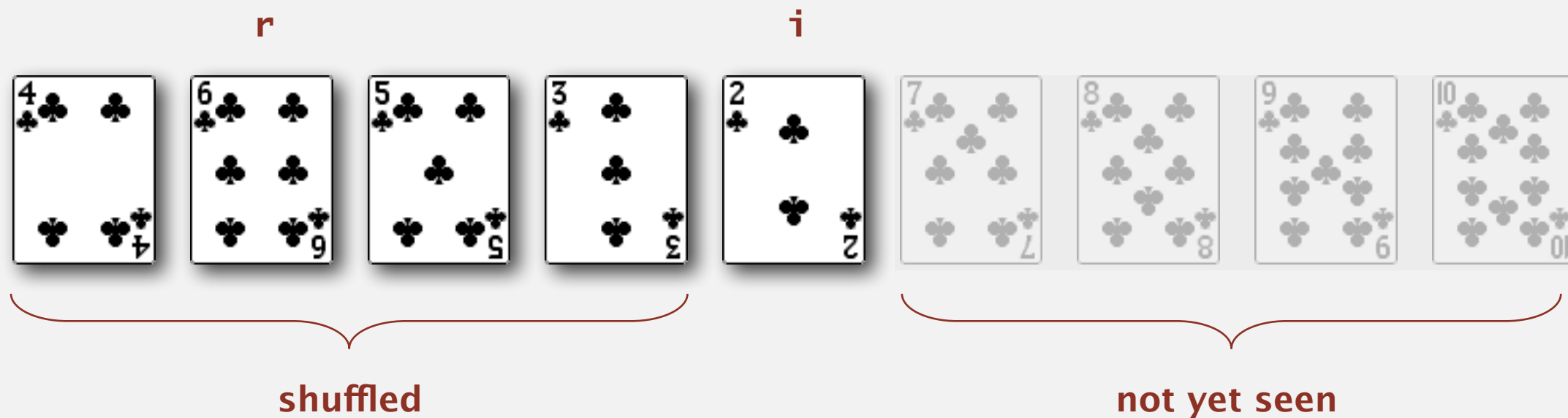
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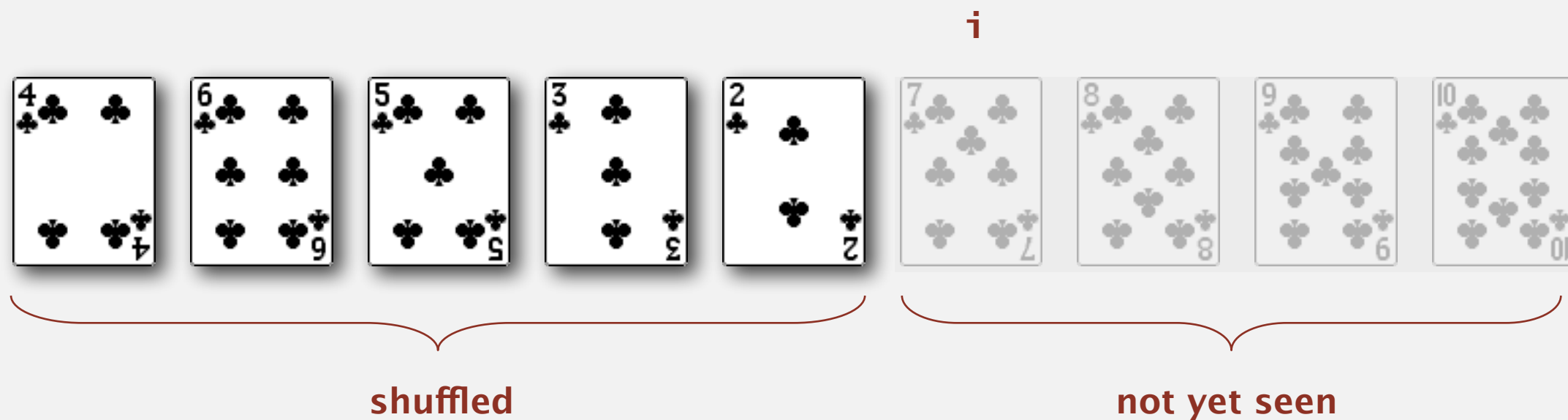
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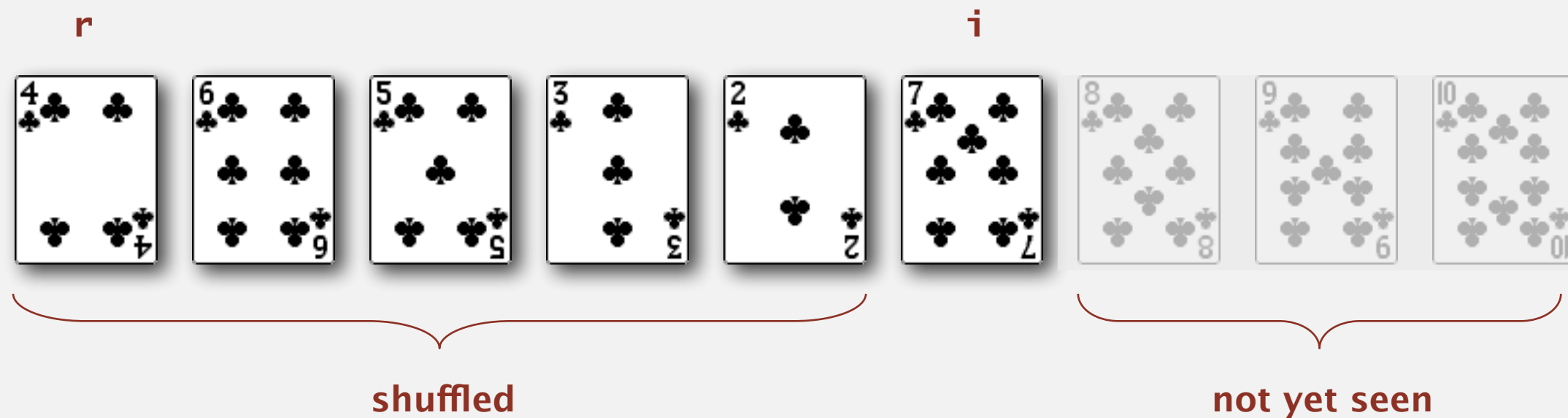
Knuth shuffle

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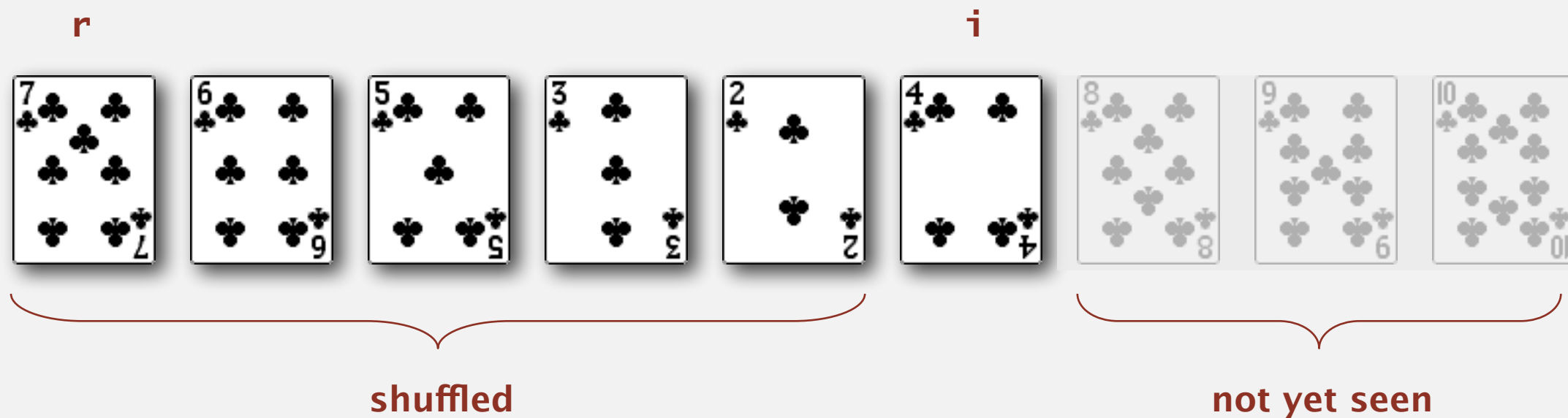
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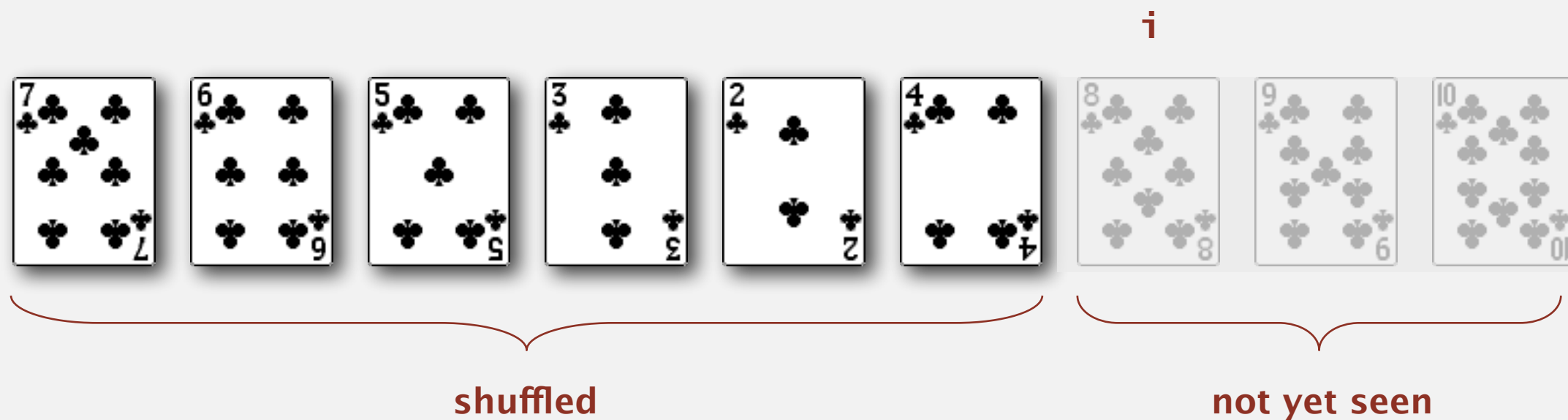
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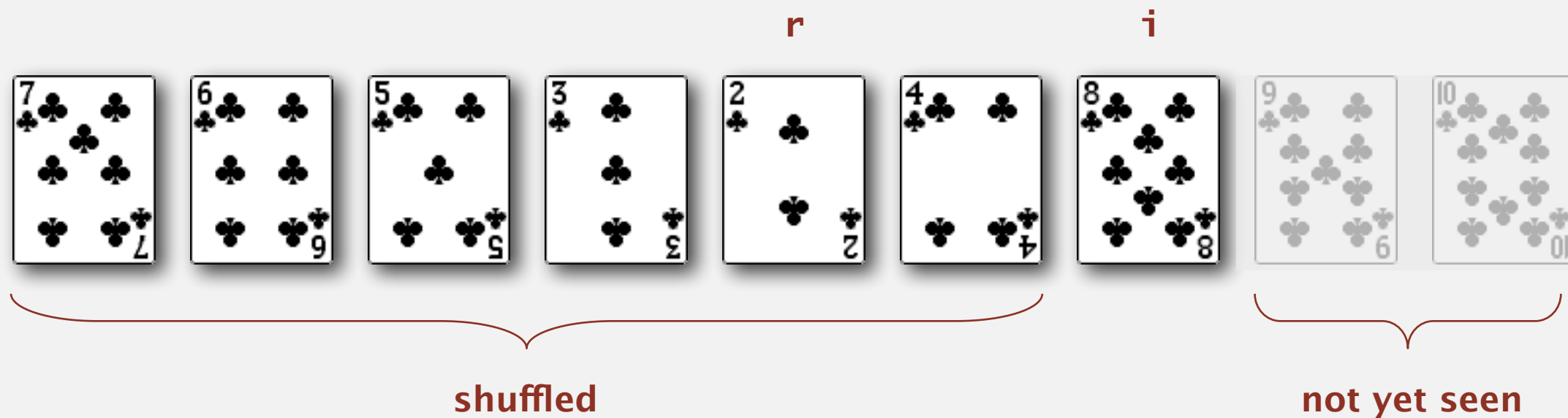
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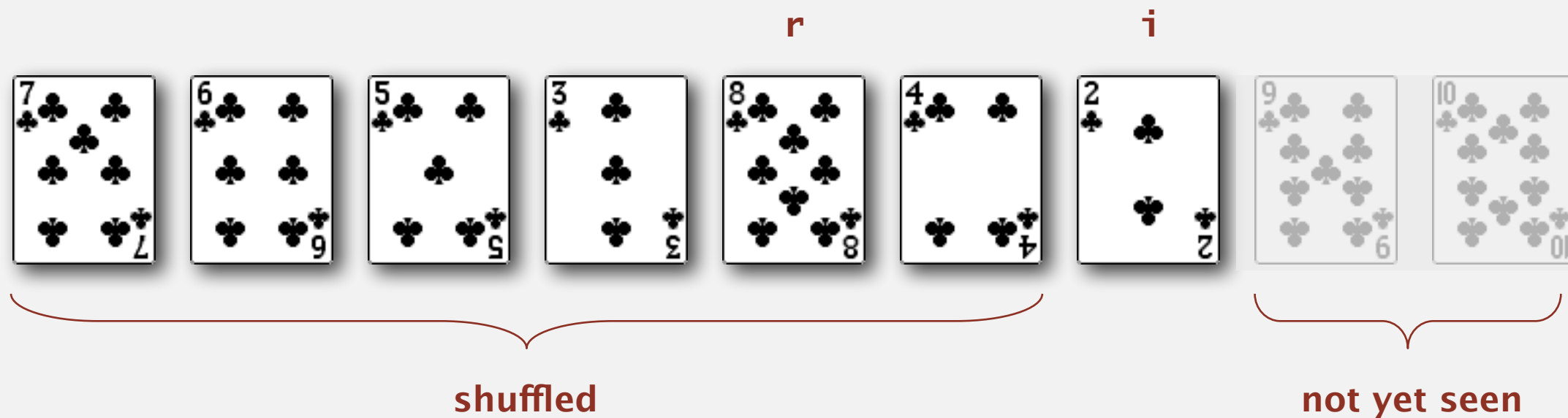
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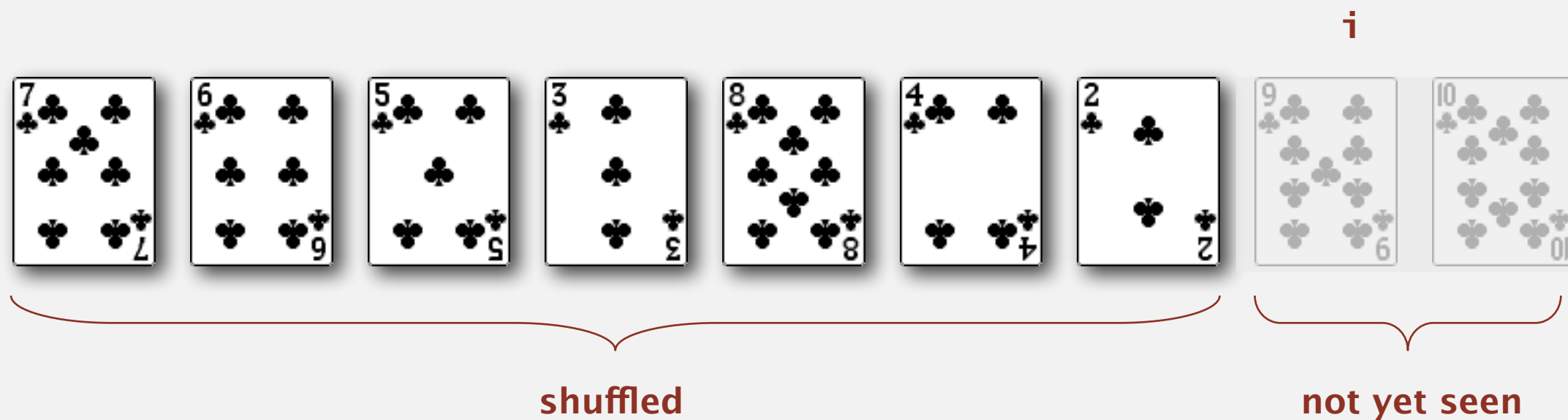
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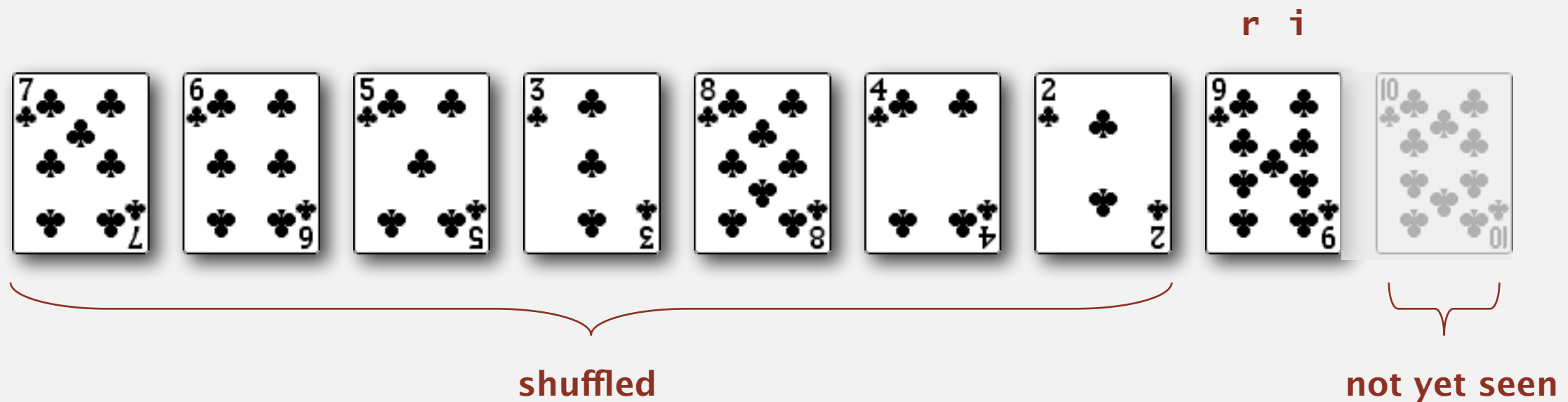
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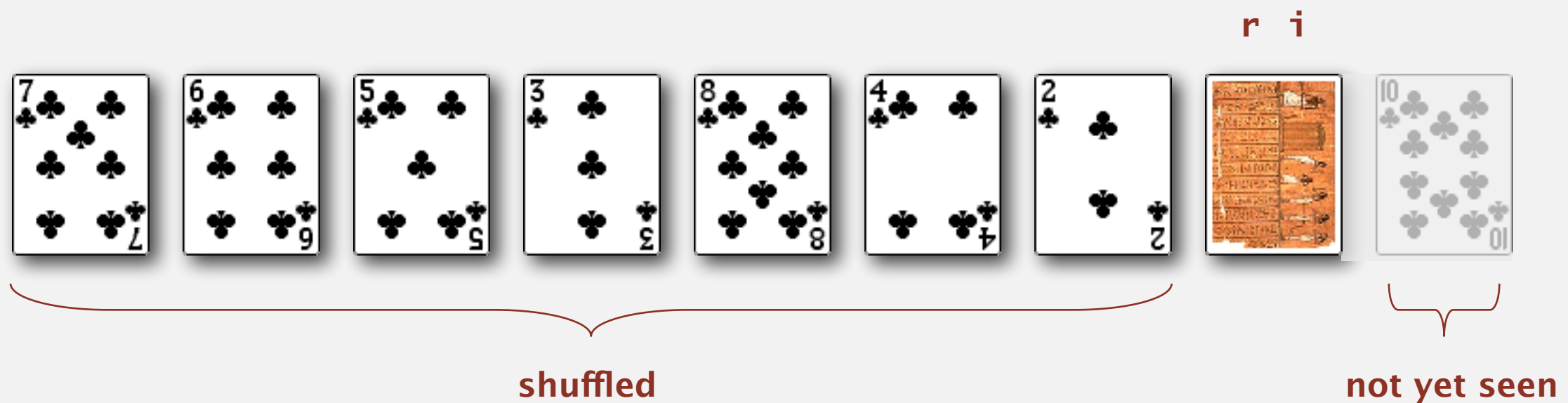
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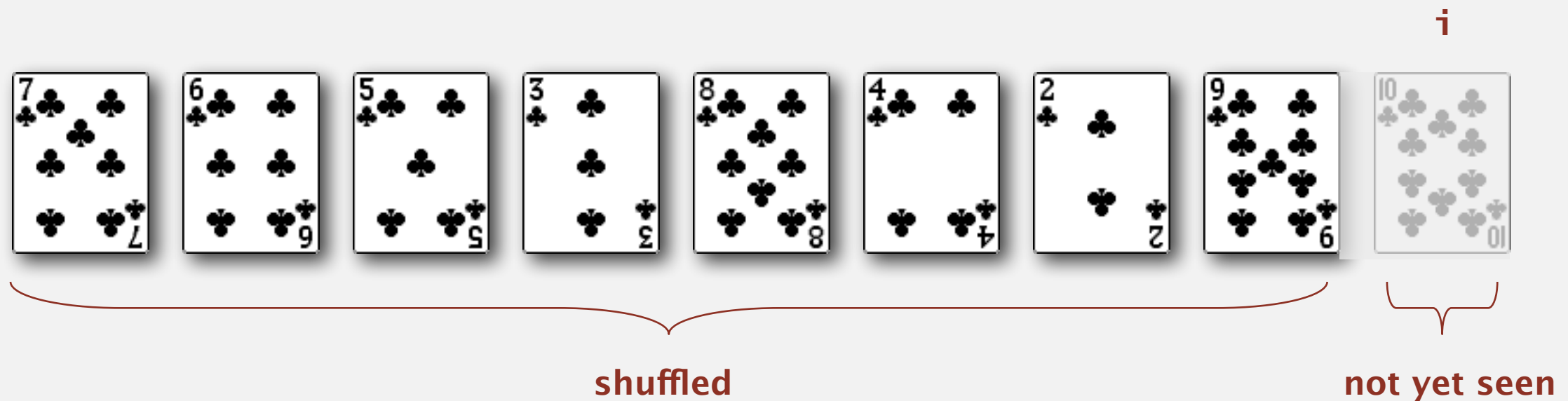
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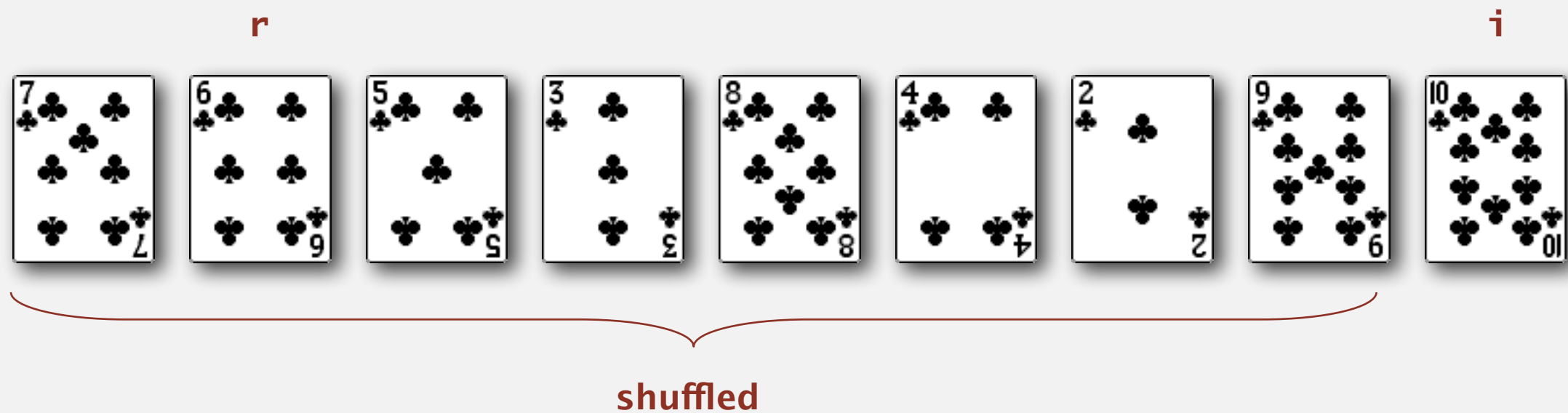
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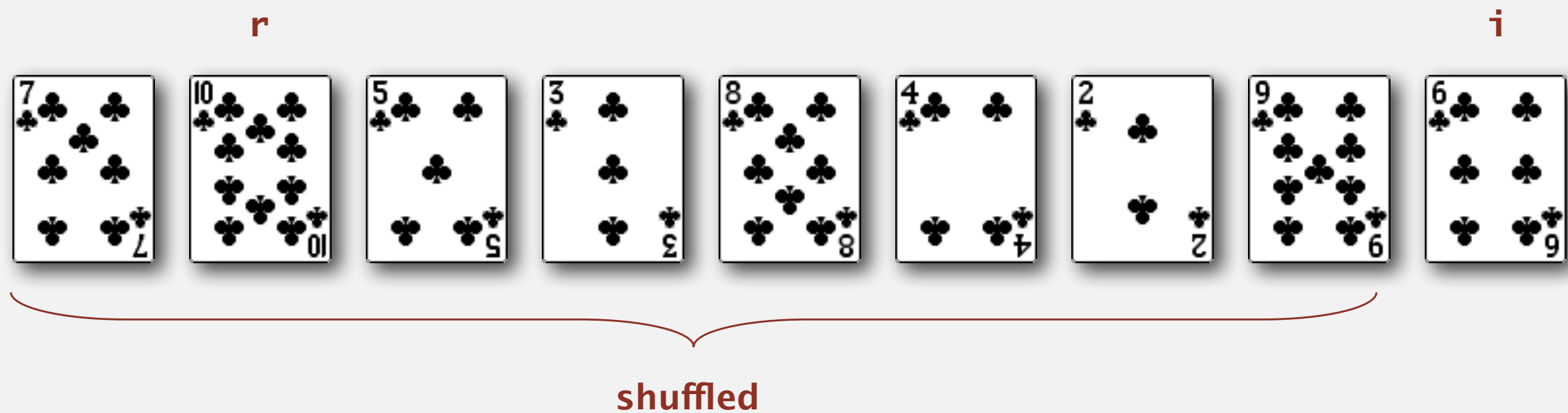
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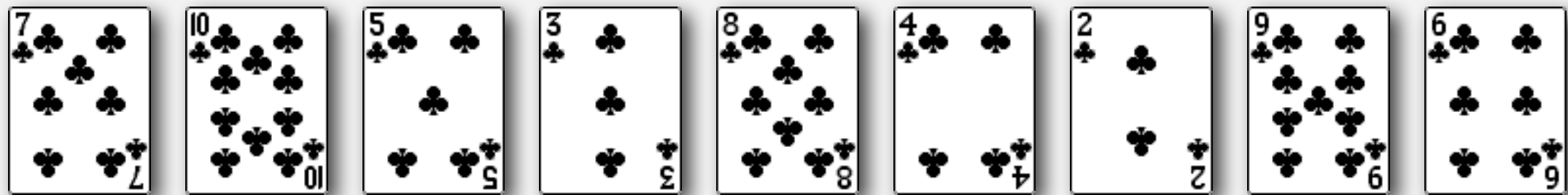
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Knuth shuffle

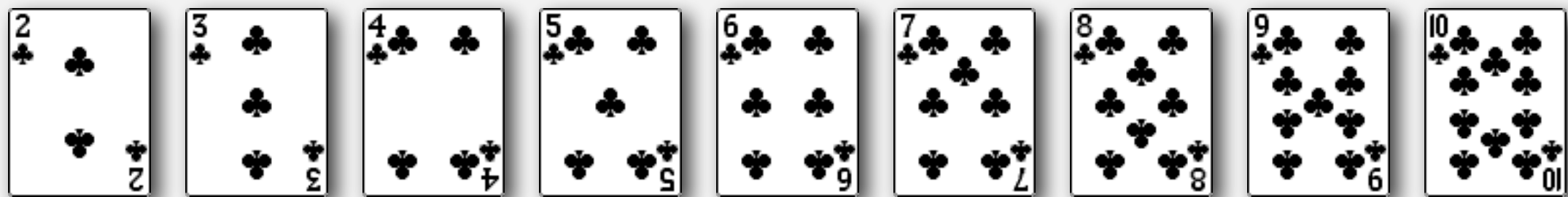
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shuffled

Knuth shuffle

- In iteration i , pick integer r between 0 and i uniformly at random.
- Swap $a[i]$ and $a[r]$.



Proposition. [Fisher-Yates 1938] Knuth shuffling algorithm produces a uniformly random permutation of the input array in linear time.

↖ assuming integers
uniformly at random

Knuth shuffle

- In iteration i , pick integer r between 0 and i uniformly at random.
- Swap $a[i]$ and $a[r]$.

common bug: between 0 and $N - 1$
correct variant: between i and $N - 1$

```
public class StdRandom
{
    ...
    public static void shuffle(Object[] a)
    {
        int N = a.length;
        for (int i = 0; i < N; i++)
        {
            int r = StdRandom.uniform(i + 1);
            exch(a, i, r);
        }
    }
}
```

between 0 and i , not including i

Broken Knuth shuffle

Q. What happens if integer is chosen between 0 and $N-1$?




instead of 0 and i

Broken Knuth shuffle

Q. What happens if integer is chosen between 0 and N-1 ?

A. Not uniformly random!

 instead of 0 and i

permutation	Knuth shuffle	broken shuffle
A B C	1/6	4/27
A C B	1/6	5/27
B A C	1/6	5/27
B C A	1/6	5/27
C A B	1/6	4/27
C B A	1/6	4/27

probability of each result when shuffling { A, B, C }

War story (online poker)

Texas hold'em poker. Software must shuffle electronic cards.



How We Learned to Cheat at Online Poker: A Study in Software Security

<https://www.developer.com/guides/how-we-learned-to-cheat-at-online-poker-a-study-in-software-security/>

War story (online poker)

Shuffling algorithm in FAQ at www.planetpoker.com

```
for i := 1 to 52 do begin
  r := random(51) + 1;
  swap := card[r];
  card[r] := card[i];
  card[i] := swap;
end;
```

Any problems with this shuffling algorithm?

War story (online poker)

Shuffling algorithm in FAQ at www.planetpoker.com

```
for i := 1 to 52 do begin
  r := random(51) + 1; ← between 1 and 51
  swap := card[r];
  card[r] := card[i];
  card[i] := swap;
end;
```

Bug 1. Random number r never 52 \Rightarrow 52nd card can't end up in 52nd place.

Bug 2. Shuffle not uniform (should be between 1 and i).

Seed = milliseconds since midnight \Rightarrow 86.4 million shuffles.

Exploit. After seeing 5 cards and synchronizing with server clock, can determine **all** future cards in real time.

“ The generation of random numbers is too important to be left to chance. ”

— Robert R. Coveyou

War story (online poker)

Best practices for shuffling (if your business depends on it).

- Use a hardware random-number generator that has passed both the FIPS 140-2 and the NIST statistical test suites.
- Continuously monitor statistic properties:
hardware random-number generators are fragile and fail silently.
- Use an unbiased shuffling algorithm.



RANDOM.ORG

Bottom line. Shuffling a deck of cards is hard!