

Comp3121 assignment 2 Question 3
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Question 3

[20 marks] Ryno needs your help! He has two directed acyclic graphs X and Y . Each of these graphs have n vertices, labelled $1, \dots, n$. He wants to know whether there is an ordering of $[1, \dots, n]$ which is a topological ordering of **both** graphs, and if so, what that ordering is.

A trivial example where this ordering does *not* exist is when $n = 2$, and X has an edge between $(1, 2)$ and Y has an edge between $(2, 1)$.

3.1 [5 marks] Provide another *non-trivial* example where this ordering does *not* exist. Non-trivial in this case means that an edge cannot appear in one direction in X , and in the opposite direction in Y , as per the example earlier.

3.2 [15 marks] Design an $O(n \log n)$ algorithm which solves Ryno's problem.

3.1

Assume graph X and Y contain three vertices, we can take the graph X as $3 \rightarrow 1 \rightarrow 2$ and graph Y as $1 \rightarrow 2 \rightarrow 3$.

3.2

We can adapt the topological sorting algorithm mentioned in the lecture slide (greedy p24) on both graphs X and Y to solve this problem because X and Y are both directed acyclic graphs.

We use an empty list L to record the ordered elements and put all the vertices of X with no incoming edge into a set S . Then we check if there is an element in S also a vertex of Y with no incoming edge. If there is then we add this vertex u (i.e. a number between 1 and n) into L and we remove all the edges satisfying the condition: $\{e = (u, v) \mid v \text{ has no incoming edge on both } X \text{ and } Y\}$. Nevertheless, if we cannot find such a vertex u , that means we cannot find a topological sorting on both graphs. We continue the step above until S is empty. If either graph has edges left, that also means we cannot find a topological sorting on both graphs. Otherwise, the list L (a topological sorted order) is the answer.

Time complexity

The worst case happens when we have to check every element in S and the maximum size of S is n . Also, for each iteration the size of L increases by 1. That is, the worst case time complexity is $O(n * n)$ which is $O(n^2)$.